

Q1 (15 Marks)

Solve the following recurrence, with all the steps and details and find the complexity of the recurrence. Show all detail.

$$T(N) = 5 * T\left(\frac{N}{5}\right) + N^{1/2} \quad T(1) = 1$$

$$\begin{aligned}
 T(N) &= 5 * T\left(\frac{N}{5}\right) + N^{1/2} \\
 &= 5 \left[5 * T\left(\frac{N}{25}\right) + \left(\frac{N}{5}\right)^{1/2} \right] + N^{1/2} \\
 &= 25 T\left(\frac{N}{25}\right) + 5 \frac{N^{1/2}}{5^{1/2}} + N^{1/2} \\
 &= 25 \left[5 * T\left(\frac{N}{125}\right) + \left(\frac{N}{25}\right)^{1/2} \right] + 5^{1/2} N^{1/2} + N^{1/2} \\
 &= 125 T\left(\frac{N}{125}\right) + 25 \left(\frac{N^{1/2}}{25^{1/2}}\right) + 5^{1/2} N^{1/2} + N^{1/2} \\
 &= 5^3 T\left(\frac{N}{5^3}\right) + \frac{25}{5} N^{1/2} + 5^{1/2} N^{1/2} + N^{1/2} \\
 &= 5^3 T\left(\frac{N}{5^k}\right) + 5 N^{1/2} + 5^{1/2} N^{1/2} + N^{1/2} \\
 &= 5^k T\left(\frac{N}{5^k}\right) + N^{1/2} [5^1 + 5^{1/2} + 5^0] \\
 &= 5^k T\left(\frac{N}{5^k}\right) + N^{1/2} \sum_{i=0}^{k-1} (5^{1/2})^i \\
 &= \frac{\log_5 N}{5^k} T(1) + N^{1/2} \left[\sum_{i=0}^{k-1} (5^{1/2})^i \right] \\
 &= N^{\frac{\log_5 N}{5^k}} * 1 + N^{1/2} \left[\sum_{i=0}^{k-1} (5^{1/2})^i \right] \\
 &= N * 1 + N^{1/2} [N^{1/2} - 1] \\
 &= N * 1 + N^{1/2} * N^{1/2} \\
 &= O(N)
 \end{aligned}$$

$$\begin{aligned}
 \frac{N}{5^k} &= 1 \Rightarrow N = 5^k \\
 \log_5 N &= \log_5 5^k \\
 \log_5 N &= k \log_5 5 \\
 \frac{\log_5 N}{\log_5 5} &= k
 \end{aligned}$$

$$\begin{aligned}
 &\therefore \frac{(5^{1/2})^k - 1}{5^{1/2} - 1} \\
 &\therefore \frac{N^{\frac{\log_5 N}{5^k}} (5^{1/2})^k - 1}{5^{1/2} - 1} \quad \therefore 5^{1/2} - 1 = C_1 \\
 &\therefore \frac{N^{1/2} - 1}{C_1}
 \end{aligned}$$

Mid Solution

Q2 (25 Marks)

Given a sorted array, **Arr**, of size **N**, which contains only two distinct elements **x** & **y**, where **x < y**. It is required to find the count of **y** in the array **Arr**. **The Time Complexity of the algorithm shall be $O(\log_2 n)$. Keep in mind, function, naming convention, complexity, correct implementation and desired solution.**

- a) Briefly describe your algorithm in words. (5 Marks)
- b) Write a function that returns the count of **y** elements. (20 Marks)

Example:

Arr:	2 2 2 2 2 2 5 5 5	Returned Value:	4
Arr:	-12 -12 -12 7 7 7 7 7 7 7	Returned Value:	9

conceptual solution:

Function run with 2 arguments, Arr and N

We have two pointers one at start = 0 and one at the end = N-1

we will assigned x to the value

we will loop until start < end

we will calculate mid

if mid index value is equal to value move start to mid + 1

if mid index value is greater than value and mid - 1 index value is value **then return the count**

else end to mid -1

or return count at the end

```
int number_count_v1(int Arr[], int N){
    int count= 0, start = 0, end = N-1, mid = -1;
    int value = Arr[0];

    while(start < end){

        cout<<"At start "<<start<<" "<<end;
        mid = start + (end - start)/2;
        if(Arr[mid] == value){
            start = mid +1;
        }else if(Arr[mid] != value && Arr[mid-1] == value){
            return N - mid;
        }else{
            end = mid -1;
        }
        cout<<" at end "<<start<<" "<<end<<endl;
    }
    return N - end;
}
```

conceptual solution:

Function run with 2 arguments, Arr and N

We have two pointers one at start = 0 and one at the end = N-1

we will assigned x to the value

we will loop until start < end

- we will calculate mid

- if mid index value is equal to value move start to mid + 1

- else end to mid

or return count at the end

```
int number_count_v2(int Arr[], int N){
    int count= 0, start = 0, end = N-1, mid = -1;
    int value = Arr[0];

    while(start < end){

        cout<<"At start "<<start<<" "<<end;
        mid = start + (end - start)/2;
        if(Arr[mid] == value){
            start = mid + 1;
        }else{
            end = mid;
        }
        cout<<" at end "<<start<<" "<<end<<endl;
    }
    return N - end;
    return 0;
}
```