

Formula to Calculate Weight for Low-Weight Weight 3 Inputs and Proof

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0.1 Equation and Proof

Proof. The polynomial representation of a weight-3 RTZ input is given by

$$Q(x) = x^{h\tau+t}(1 + x^{\beta\tau+1} + x^{\gamma\tau+2})$$

The impulse response of the RSC encoder is

$$(1 \ 1 \ 1 \ 0 \ 1 \ 1 \ 0 \ 1 \ 1 \ 0 \ \dots)$$

and using the impulse response, we can calculate the parity weight as well as Hamming weight of the turbo codeword.

Let

$$\begin{aligned} \phi_1 &= (0 \ 0 \ 1), \ \phi'_1 = (0 \ 1 \ 0), \ \phi''_1 = (1 \ 0 \ 0), \\ \phi_2 &= (0 \ 1 \ 1), \ \phi'_2 = (1 \ 1 \ 0), \ \phi''_2 = (1 \ 0 \ 1), \\ \phi_3 &= (1 \ 1 \ 1). \end{aligned}$$

Then, the weight-3 RTZ occurs since $\phi_2 + \phi'_2 + \phi''_2 = \mathbf{0}_\tau = \mathbf{0}_3$. Now, we consider the weight of the vector derived by the summation of the followings vectors.

$$\begin{aligned} &(\mathbf{0}_{3(\gamma+h)} \ \phi_1 \ \phi'_2 \ \dots) \\ &(\mathbf{0}_{3(\beta+h)} \ \phi_2 \ \phi''_2 \ \dots) \\ &(\mathbf{0}_{3h} \ \phi_3 \ \phi_2 \ \dots) \end{aligned}$$

Without loss of generality, we can assume that all weight-3 RTZ inputs begin at the 0th position, ie $h = t = 0$. This is because the case where $h > 0$ or $t > 0$ is just a right-shifted version of the weight-3 RTZ. With this assumption, we only need to consider cases where $h = 0$, $\gamma \geq h$. To simplify calculation, we have included an addition table for all the vectors which is shown in Table 1

	ϕ_1	ϕ'_1	ϕ''_1	ϕ_2	ϕ'_2	ϕ''_2	ϕ_3
ϕ_1	$\mathbf{0}_3$	—	—	—	—	—	—
ϕ'_1	ϕ_2	$\mathbf{0}_3$	—	—	—	—	—
ϕ''_1	ϕ'_2	ϕ'_2	$\mathbf{0}_3$	—	—	—	—
ϕ_2	ϕ'_1	ϕ_1	ϕ_3	$\mathbf{0}_3$	—	—	—
ϕ'_2	ϕ_3	ϕ''_1	ϕ'_1	ϕ''_2	$\mathbf{0}_3$	—	—
ϕ''_2	ϕ''_1	ϕ_3	ϕ_1	ϕ'_2	ϕ_2	$\mathbf{0}_3$	—
ϕ_3	ϕ'_2	ϕ''_2	ϕ_2	ϕ''_1	ϕ_1	ϕ'_1	$\mathbf{0}_3$

Table 1: Truth Table

Furthermore, we consider 4 general cases for all possible values of i, j, k where $i \geq k$. These cases are $(= \ =)$, $(= \ <)$, $(< \ =)$ and $(< \ <)$

Case 0: $\gamma = \beta = h$

For this case, the vectors to sum will be

$$\begin{aligned} &(\phi_1 \ \phi'_2 \ \dots) \\ &(\phi_2 \ \phi''_2 \ \dots) \\ &(\phi_3 \ \phi_2 \ \dots) \\ &\hline &(\phi''_2 \ \mathbf{0}_3 \ \dots) \end{aligned}$$

and the derived vector will be $(\phi''_2 \ \mathbf{0}_3 \ \dots)$ with a weight of $w_p = 2$

Case 1a: $\gamma = h < \beta$
vectors to sum:

$$\begin{array}{c}
 (\phi_1 \ \phi'_2 \ \phi'_2 \ \phi'_2 \ \phi'_2 \ \cdots) \\
 (\mathbf{0}_3 \ \cdots \ \mathbf{0}_3 \ \phi_2 \ \phi''_2 \ \cdots) \\
 +(\phi_3 \ \phi_2 \ \phi_2 \ \phi_2 \ \phi_2 \ \cdots) \\
 \hline
 (\phi'_2 \ \phi''_2 \ \phi''_2 \ \phi'_2 \ \mathbf{0}_3 \ \cdots)
 \end{array}$$

derived vector : $(\phi'_2 \ (\phi''_2)_{\beta-h-1} \ \phi'_2 \ \mathbf{0}_3 \ \cdots)$

Parity weight:

$$w_p = 2(\beta - h) + 2 = 2\beta + 2 \quad (0-1)$$

Case 1b: $\beta = h < \gamma$
vectors to sum:

$$\begin{array}{c}
 (\mathbf{0}_3 \ \cdots \ \cdots \ \mathbf{0}_3 \ \phi_1 \ \phi'_2 \ \cdots) \\
 (\phi_2 \ \phi''_2 \ \cdots \ \phi''_2 \ \phi''_2 \ \phi''_2 \ \cdots) \\
 +(\phi_3 \ \phi_2 \ \cdots \ \phi_2 \ \phi_2 \ \phi_2 \ \cdots) \\
 \hline
 (\phi''_1 \ \phi'_2 \ \cdots \ \phi'_2 \ \phi_3 \ \mathbf{0}_3 \ \cdots)
 \end{array}$$

derived vector : $(\phi''_1 \ (\phi_2)_{\gamma-h-1} \ \phi_3 \ \mathbf{0}_3 \ \cdots)$

Parity weight:

$$w_p = 2(\gamma - h) + 2 = 2\gamma + 2 \quad (0-2)$$

Case 2a: $h < \gamma = \beta$
vectors to sum:

$$\begin{aligned} & (\mathbf{0}_3 \cdots \cdots \mathbf{0}_3 \phi_1 \phi'_2 \cdots) \\ & (\mathbf{0}_3 \cdots \cdots \mathbf{0}_3 \phi_2 \phi''_2 \cdots) \\ & + (\phi_3 \phi_2 \cdots \phi_2 \phi_2 \phi_2 \cdots) \\ & \hline & (\phi_3 \phi_2 \cdots \phi_2 \phi_1 \mathbf{0}_3 \cdots) \end{aligned}$$

derived vector : $(\phi_3 (\phi_2)_{\gamma-h-1} \phi_1 \mathbf{0}_3 \cdots)$

Parity weight:

$$w_p = 2(\gamma - h) + 2 = 2\gamma + 2 \quad (0-3)$$

Case 3a: $h < \gamma < \beta$
vectors to sum:

$$\begin{aligned} & (\mathbf{0}_3 \cdots \cdots \mathbf{0}_3 \phi_1 \phi'_2 \cdots \phi'_2 \phi'_2 \phi'_2 \cdots) \\ & (\mathbf{0}_3 \cdots \cdots \cdots \cdots \cdots \mathbf{0}_3 \phi_2 \phi''_2 \cdots) \\ & + (\phi_3 \phi_2 \cdots \phi_2 \phi_2 \phi_2 \cdots \phi_2 \phi_2 \phi_2 \cdots) \\ & \hline & (\phi_3 \phi_2 \cdots \phi_2 \phi'_1 \phi'_2 \cdots \phi'_2 \phi'_2 \mathbf{0}_3 \cdots) \end{aligned}$$

derived vector : $(\phi_3 (\phi_2)_{\gamma-h-1} \phi'_1 (\phi'_2)_{\beta-\gamma-1} \phi'_2 \mathbf{0}_3 \cdots)$

Parity weight:

$$\begin{aligned} w_p &= 2(\gamma - h) + 2 + 2(\beta - i) \\ &= 2(\beta - h) + 2 \\ &= 2\beta + 2 \end{aligned} \quad (0-4)$$

Case 3b: $h < \beta < \gamma$

$$\begin{aligned} & (\mathbf{0}_3 \cdots \cdots \cdots \cdots \cdots \mathbf{0}_3 \phi_1 \phi'_2 \cdots) \\ & (\mathbf{0}_3 \cdots \cdots \mathbf{0}_3 \phi_2 \phi''_2 \cdots \phi''_2 \phi''_2 \phi''_2 \cdots) \\ & + (\phi_3 \phi_2 \cdots \phi_2 \phi_2 \phi_2 \cdots \phi_2 \phi_2 \phi_2 \cdots) \\ & \hline & (\phi_3 \phi_2 \cdots \phi_2 \mathbf{0}_3 \phi'_2 \cdots \phi'_2 \phi_3 \mathbf{0}_3 \cdots) \end{aligned}$$

derived vector : $(\phi_3 (\phi_2)_{j-k-1} \mathbf{0}_3 (\phi'_2)_{i-j-1} \phi_3 \mathbf{0}_3 \cdots)$

Parity weight:

$$\begin{aligned} w_p &= 2(\beta - h) + 1 + 2(\gamma - \beta) + 1 \\ &= 2(\gamma - h) + 2 \\ &= 2\gamma + 2 \end{aligned} \quad (0-5)$$

From all the above cases we can conclude that the parity weight for a weight-3 RTZ sequence may be calculated as

$$w_p = 2l + 2 \quad (0-6)$$

where $l = \max\{\gamma, \beta\} - k = \max\{\gamma, \beta\}$ since $k = 0$

Assuming that after interleaving, another weight-3 RTZ input is produced. Let γ', β', h', l' and w'_p be similarly defined. Then the Hamming weight w_H of the turbo codeword produced can be calculated as

$$w_H = 7 + 2(l + l') \quad (0-7)$$

□