

RSA Public-Key Cryptography

The RSA Algorithm: It is a public key cryptography algorithm, which was proposed by Diffie and Hellman. RSA can be used for key exchange, digital signatures and the encryption of small blocks of data.

- RSA is primarily used to encrypt the session key used for secret key encryption or the message's hash value (digital signature).
- RSA's mathematical hardness comes from the ease in calculating large numbers and the difficulty in finding the prime factors of those large numbers.
- To create an RSA public/private key pair, here are the basic steps:
 - 1- Choose two prime numbers, p and q such that $p \neq q$.
 - 2- Calculate the modulus, $n = p \times q$.
 - 3- Calculate $\phi(n) = (p - 1) \times (q - 1)$.
 - 4- Select integer e such that $\gcd(\phi(n), e) = 1$ and $1 < e < \phi(n)$. (* gcd is greater common divisor)
 - 5- Calculate an integer d from the quotient $de \equiv 1 \pmod{\phi(n)} \Rightarrow de = 1 + k\phi(n) \Rightarrow d = (1 + k\phi(n)) / e$
- To encrypt a message, M , with the public key (e, n) , create the ciphertext, C , using the equation:
$$C = M^e \bmod n$$
- The receiver then decrypts the ciphertext with the private key (d, n) using the equation:
$$M = C^d \bmod n$$

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The RSA Example

1. Select two prime numbers, $p = 17$ and $q = 11$.
2. Calculate $n = p \times q = 17 \times 11 = 187$.
3. Calculate $\phi(n) = (p - 1)(q - 1) = 16 \times 10 = 160$.
4. Select e such that e is relatively prime to $\phi(n) = 160$ and less than $\phi(n)$; we choose $e = 7$.
5. Determine d such that $de \equiv 1 \pmod{160}$ and $d < 160$. $de = 1 + k \phi(n)$

The correct value is $d = 23$, because $23 \times 7 = 161 = 1 + (1 \times 160)$.

The resulting keys are public key $PU = \{7, 187\}$ and private key $PR = \{23, 187\}$.

Given a plaintext input of $M = 88$. For encryption, we need to calculate $C = 88^7 \pmod{187}$.
we can do this as follows.

$$88^7 \pmod{187} = [(88^4 \pmod{187}) * (88^2 \pmod{187}) * (88^1 \pmod{187})] \pmod{187}$$

$$88^1 \pmod{187} = 88$$

$$88^2 \pmod{187} = 7744 \pmod{187} = 77$$

$$88^4 \pmod{187} = 59,969,536 \pmod{187} = 132$$

$$88^7 \pmod{187} = (88 * 77 * 132) \pmod{187} = 894,432 \pmod{187} = 11$$

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For decryption, we calculate $M = 11^{23} \bmod 187$:

$$11^{23} \bmod 187 = [(11^1 \bmod 187) * (11^2 \bmod 187) * (11^4 \bmod 187) * (11^8 \bmod 187) * (11^8 \bmod 187)] \bmod 187$$

$$11^1 \bmod 187 = 11$$

$$11^2 \bmod 187 = 121$$

$$11^4 \bmod 187 = 14,641 \bmod 187 = 55$$

$$11^8 \bmod 187 = 214,358,881 \bmod 187 = 33$$

$$11^{23} \bmod 187 = (11 * 121 * 55 * 33 * 33) \bmod 187 = 79,720,245 \bmod 187 = 88$$

In the preceding example shows, we can make use of a property of modular arithmetic:

$$[(a \bmod n) * (b \bmod n)] \bmod n = (a * b) \bmod n$$

As another example, suppose we wish to calculate $x^{11} \bmod n$ for some integers x and n . Observe that $x^{11} = x^{1+2+8} = (x)(x^2)(x^8)$.

Public-Key Cryptography

Applications for Public-Key Cryptosystems:

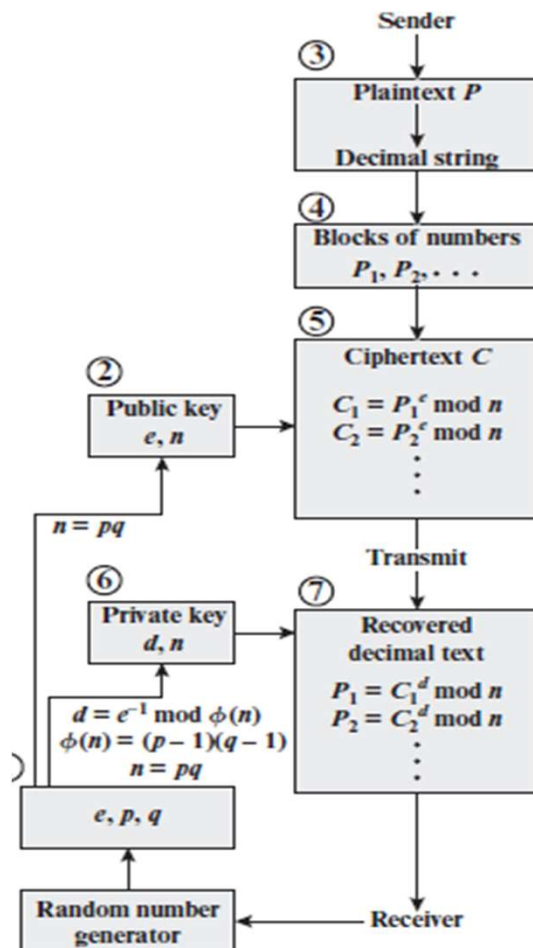
- **Encryption/decryption:** The sender encrypts a message with the recipient's public key, and the recipient decrypts the message with the recipient's private key.
- **Digital signature:** The sender "signs" a message with its private key.
- **Key exchange:** Two sides cooperate to exchange a session key.

The security of RSA:

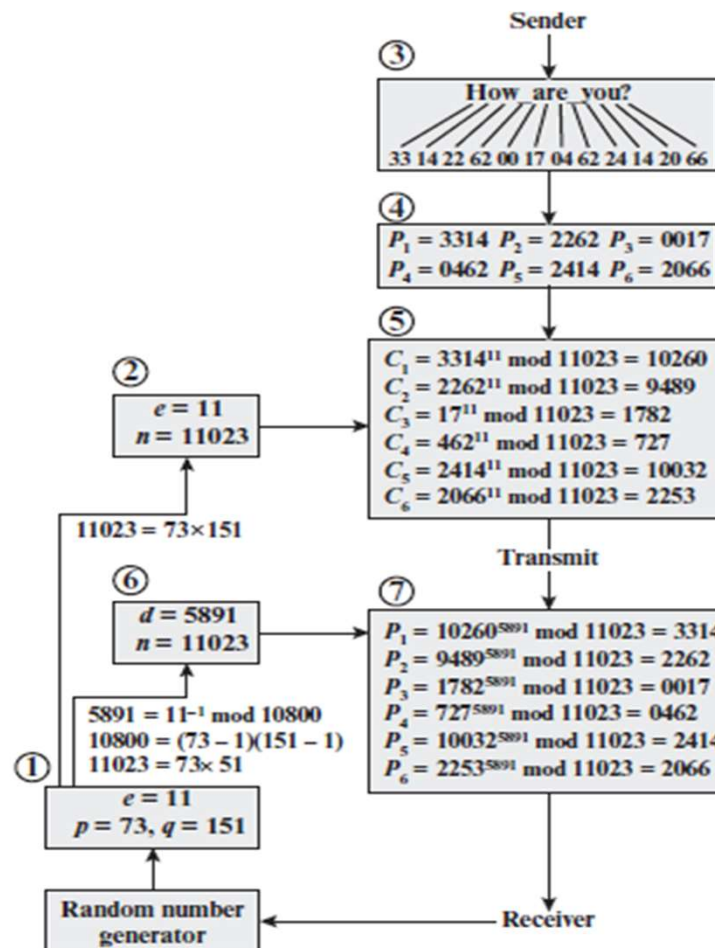
Five possible approaches to attacking the RSA algorithm are

- **Brute force:** This involves trying all possible private keys.
- **Mathematical attacks:** There are several approaches, all equivalent in effort to factoring the product of two primes.
- **Timing attacks:** These depend on the running time of the decryption algorithm.
- **Hardware fault-based attack:** This involves inducing hardware faults in the processor that is generating digital signatures.
- **Chosen ciphertext attacks:** This type of attack exploits properties of the RSA algorithm.

RSA processing of multiple blocks



(a) General approach



(b) Example