



Heap

DATASTRUCTURES & ALGORITHMS



Definition

The (**binary**) heap is:

- Complete binary tree
- Has heap property

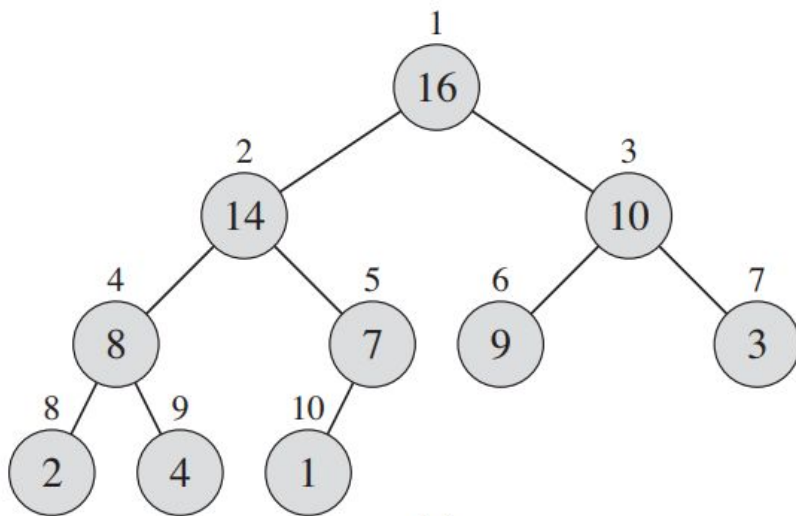
There are two kinds of binary heaps:

- Max-heaps
- Min-heaps

max-heap property is that for every node i other than the root:

$$n_i.parent.key \geq n_i.key$$

Max-heap example



(a)

PARENT(i)

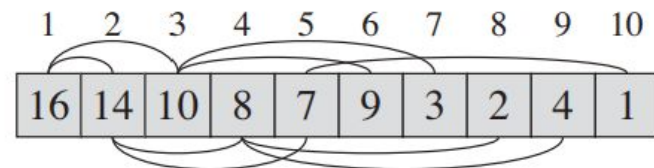
1 **return** $\lfloor i/2 \rfloor$

LEFT(i)

1 **return** $2i$

RIGHT(i)

1 **return** $2i + 1$



(b)



Exercise

6.1-1

What are the minimum and maximum numbers of elements in a heap of height h ?

6.1-2

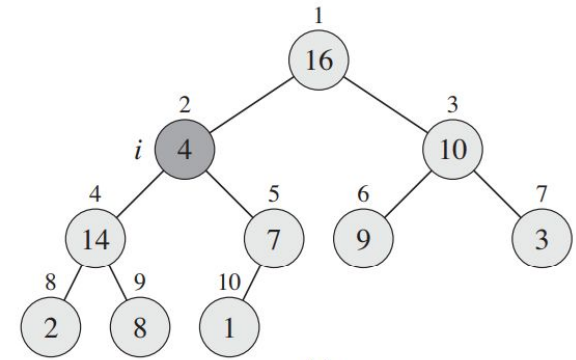
Show that an n -element heap has height $\lceil \lg n \rceil$.

Maintaining the heap property

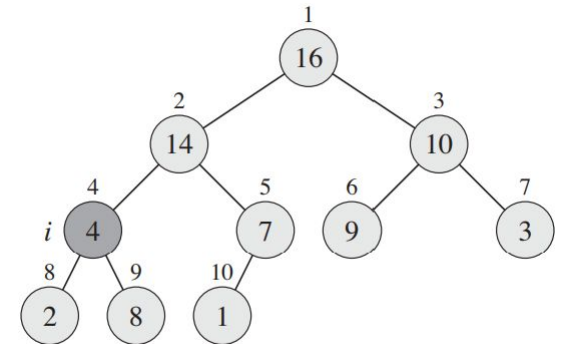
MAX-HEAPIFY(A, i)

```

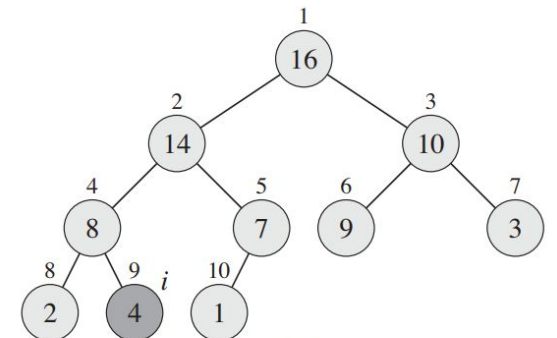
1   $l = \text{LEFT}(i)$ 
2   $r = \text{RIGHT}(i)$ 
3  if  $l \leq A.\text{heap-size}$  and  $A[l] > A[i]$ 
4       $\text{largest} = l$ 
5  else  $\text{largest} = i$ 
6  if  $r \leq A.\text{heap-size}$  and  $A[r] > A[\text{largest}]$ 
7       $\text{largest} = r$ 
8  if  $\text{largest} \neq i$ 
9      exchange  $A[i]$  with  $A[\text{largest}]$ 
10     MAX-HEAPIFY( $A, \text{largest}$ )
    
```



(a)



(b)



(c)



Running Time of Max_Heapify

What is the Worst case?

$$T(n) \leq T(2n/3) + \Theta(1)$$

$$T(n) = O(\lg n) = O(h)$$

Master Theorem:

$$T(n) = aT(n/b) + f(n)$$

$$\text{If } f(n) = \Theta(n^{\log_b a}), \text{ then } T(n) = \Theta(n^{\log_b a} \lg n)$$



Building a heap

BUILD-MAX-HEAP(A)

```
1   $A.heap-size = A.length$ 
2  for  $i = \lfloor A.length/2 \rfloor$  downto 1
3      MAX-HEAPIFY( $A, i$ )
```

Exercise:

6.3-3

Show that there are at most $\lceil n/2^{h+1} \rceil$ nodes of height h in any n -element heap.



Running time of building heap

$$\sum_{h=0}^{\lfloor \lg n \rfloor} \left\lceil \frac{n}{2^{h+1}} \right\rceil O(h) = O \left(n \sum_{h=0}^{\lfloor \lg n \rfloor} \frac{h}{2^h} \right)$$

$$\begin{aligned} \sum_{h=0}^{\infty} \frac{h}{2^h} &= \frac{1/2}{(1 - 1/2)^2} \\ &= 2. \end{aligned}$$

$$\begin{aligned} O \left(n \sum_{h=0}^{\lfloor \lg n \rfloor} \frac{h}{2^h} \right) &= O \left(n \sum_{h=0}^{\infty} \frac{h}{2^h} \right) \\ &= O(n). \end{aligned}$$



Heap sort

HEAPSORT(A)

```
1  BUILD-MAX-HEAP( $A$ )
2  for  $i = A.length$  downto 2
3      exchange  $A[1]$  with  $A[i]$ 
4       $A.heap-size = A.heap-size - 1$ 
5      MAX-HEAPIFY( $A, 1$ )
```

Exercise:

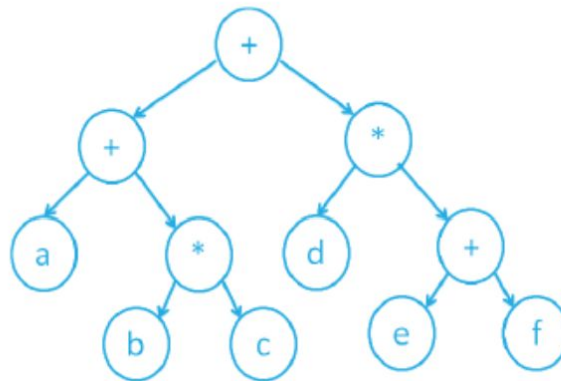
6.4-1

Using Figure 6.4 as a model, illustrate the operation of HEAPSORT on the array $A = \langle 5, 13, 2, 25, 7, 17, 20, 8, 4 \rangle$.

Expression Tree

Expression trees are binary trees with each leaf node serving as an operand and each internal (non-leaf) node serving as an operator.

$a + (b * c) + d * (e + f)$





Building expression tree

- Convert expression to postfix
- If a character is an operand, add it to the stack.
- If a character is an operator, pop both values from the stack and make both its children and push the current node again.