# Algorithms and Data Structures

Spring 2019

## Assignment 9

Date: April 23, 2019

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### Problem 9.1 Hash Tables

(a) Given the sequence 3, 10, 2, 4, apply the double-hashing strategy for open addressing to store the sequence in the given order in a hash table of size m=5 with hash functions  $h_1(k)=k \mod 5$  and  $h_2(k)=7k \mod 8$ . Document all collisions and how they are resolved. Write down your computations.

Double hashing uses a hash function of the form:

$$h(k,i) = (h_1(k) + ih_2(k)) \mod m$$

The initial probe goes to position  $T[h_1(k)]$ ; successive probe positions are offset from previous positions by the amount  $h_2(k)$ , modulo m.

#### Insertion by double-hashing:

#### 1. Insert 3

First, we calculate  $h_1(3) = 3 \mod 5 = 3$ 

Since the slot 3 in the hash table is empty, we insert the key 3 into it.

2. Insert 10

Again, we start with calculating  $h_1(k) = 10 \mod 5 = 0$ 

Since the slot 0 in the hast table is empty, we insert the key 10 into it.

3. Insert 2

 $h_1(k) = 2 \mod 5 = 2$ 

Since the slot 2 in the hash table is not occupied, we insert the key 2 into it.

4. Insert 4

 $h_1(k) = 4 \mod 5 = 4$ 

Since the slot 4 in the hash table is not occupied, we insert the key 4 into it.

There are 0 collisions. That's how the final hash table looks like:



(b) Implement a hash table that supports insertion and querying with open addressing using linear probing. Select an h' function and explain why your selected h' is well-suited for your test data. The implementation should be consistent with the following or equivalent class specifications:

Implementation of the HashTable class. hashTable.cpp

```
Node::Node(int key, int value) {
    this->key = key;
    this->value = value;
}
    Obrief constructor
    by default maxSize is 100
HashTable::HashTable(int maxSize) {
    this->maxSize = maxSize;
    currentSize = 0;
    // allocate memory
    arr = new Node*[maxSize];
    for (int i = 0; i < maxSize; i++) {</pre>
        arr[i] = nullptr;
}
    Obrief destructor
HashTable::~HashTable() {
    for (int i = 0; i < maxSize; i++) {</pre>
        delete arr[i];
   delete[] arr;
}
    Obrief Hash Function
int HashTable::hashCode(int key) {
    return key % maxSize;
}
    Obrief Insert element at a key
void HashTable::insertCode(int key, int value) {
    int hashValue = (hashCode(key)) % maxSize;
    int init = -1;
    int i = 1;
    while (hashValue != init && arr[hashValue]!= nullptr && arr[hashValue]->key != key)
        if (init == -1) {
            init = hashValue;
        hashValue = (hashCode(key) + i) % maxSize;
    }
    if (arr[hashValue] == nullptr)
        arr[hashValue] = new Node(key, value);
```

```
currentSize++;
    }
}
    Obrief Search for an element with the given key
    Oreturn the position of the element we are looking for
int HashTable::get(int key) {
    for (int i = 0; i < maxSize; i++) {</pre>
        int hashValue = (hashCode(key) + i) % maxSize;
        if (arr[hashValue] == nullptr) {
            return -1;
        else if (arr[hashValue]->key == key) {
            return hashValue;
    }
    return -1;
}
    Obrief check whether the HashTable is empty
bool HashTable::isEmpty() {
    return currentSize == 0;
```

To execute run make

### Problem 9.2 Greedy Algorithms

(a) Show that a greedy algorithm for the activity-selection problem that makes the greedy choice of selecting the activity with shortest duration may fail at producing a globally optimal solution.

Our lemma is that the greedy choice of picking activity with the shortest duration as a first choice is optimal choice. Let's prove that it is wrong. Let's consider the following example. We have the set S containing 3 activities. Each of these activities has its corresponding start and end times.  $S = \{[2,6], [5,7], [6,13]\}$ . The activity with the shortest duration is [5,7]. The following one is [2,6]. Since it overlaps with [5,7], we cannot pick it. The same applies for [6,13]. Consequently, our result is only one activity, which is definitely not the most optimal one. If we used another approach with the greedy choice of picking  $a_1$  as first choice, it would lead to the optimal choice, which is indeed  $\{[2,6], [6,13]\}$ .

(b) Assuming an unsorted sequence of activities, derive a greedy algorithm for the activity-selection problem that selects the activity with the latest starting time. Your solution should not simply sort the activities and then select the activity.

```
/*
Algorithms and Data Structures
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Assignment #8
Ofile ActivitySelectionProblem.cpp
Qauthor Taiyr Begeyev
Quersion 1.0 23/04/19
*/
```

```
/*
           \__ \ / / < / / \_ __ \
            (____ / | ___ | / ____ | | | | |
                         \/
*/
#include <iostream>
#include <algorithm>
#include <vector>
using namespace std;
struct Activity {
    int startTime;
    int finishTime;
};
/*
    Obrief print the elements
void print(vector<Activity> S) {
    int n = S.size();
    for (int i = 0; i < n; i++) {
        cout << "Activity " << i + 1<< ":" << endl;</pre>
        cout << S[i].startTime << " " << S[i].finishTime << endl << endl;</pre>
    }
}
    Obrief Greedy algorithm that solves the activity selection problem
    As our greedy choice we pick the activitu that has the latest start time
    Oparams the structure with our activities
    Oreturn the structure with the solution
vector <Activity> ActivitySelection(vector <Activity> S) {
    // this struct will contain the solution
    vector<Activity> ActivitySolution;
    int latestStartTime;
    int latestStartTimeIndex;
    while (!S.empty()) {
        latestStartTime = 0;
        // select activity with the latest start time
        for (int i = 0; i < S.size(); i++) {</pre>
            if (S[i].startTime > latestStartTime) {
                latestStartTime = S[i].startTime;
                latestStartTimeIndex = i;
            }
        }
        // check if it overlaps with activities in the ActivitySolution
        bool overLaps = false;
        for (int i = 0; i < ActivitySolution.size(); i++) {</pre>
            if (S[latestStartTimeIndex].finishTime > ActivitySolution[i].startTime) {
                overLaps = true;
            }
```

```
}
        // push it to the ActivitySolution
        if (!overLaps)
            ActivitySolution.push_back(S[latestStartTimeIndex]);
        // delete this activity from S
        S.erase(S.begin() + latestStartTimeIndex);
    return ActivitySolution;
}
int main() {
    vector <Activity> myActivities1, myActivities2, myActivities3;
    myActivities1 = \{\{1, 2\}, \{2, 3\}, \{4, 10\}\};
    myActivities2 = {
        {1, 4}, {3, 5},
        {0, 6}, {5, 7},
        {3, 8}, {5, 9},
        {6, 10}, {8, 11},
        {8, 12}, {2, 13}, {2, 14}
    };
    myActivities3 = {
        {2, 14}, {2, 13},
        {8, 12}, {8, 11},
        {5, 9}, {3, 8}
    };
    myActivities1 = ActivitySelection(myActivities1);
    myActivities2 = ActivitySelection(myActivities2);
    myActivities3 = ActivitySelection(myActivities3);
    // print the result
    print(myActivities1);
    print(myActivities2);
    print(myActivities3);
    return 0;
}
```