

Using OpenMP Tasking

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OpenMP Tasking



Introduced with OpenMP 3.0 in 2008

```
C/C++
#pragma omp task [clause]
... structured block ...
```

- Each encountering thread/task creates a new task
 - Code and data is being packaged up
 - Tasks can be nested
- Task barrier: taskwait
 - Encountering task is suspended until child tasks complete

```
C/C++
#pragma omp taskwait
```



Recursive computation of Fibonacci



On the following slides we will show three approaches to parallelize this recursive code with Tasking.



First version with Tasks (omp-v1)



```
int fib(int n)
int main (int argc,
         char* argv[])
                                          if (n < 2) return n;
                                       int x, y;
                                      #pragma omp task shared(x)
   [\ldots]
#pragma omp parallel
                                          x = fib(n - 1);
#pragma omp single
                                       #pragma omp task shared(y)
   fib(input);
                                          y = fib(n - 2);
   [...]
                                       #pragma omp taskwait
                                          return x+y;
```

- Only one Task / Thread enters fib() from main(),
 it is responsable for creating the two initial worker tasks
- Taskwait is required, as otherwise x and y inputs would be lost





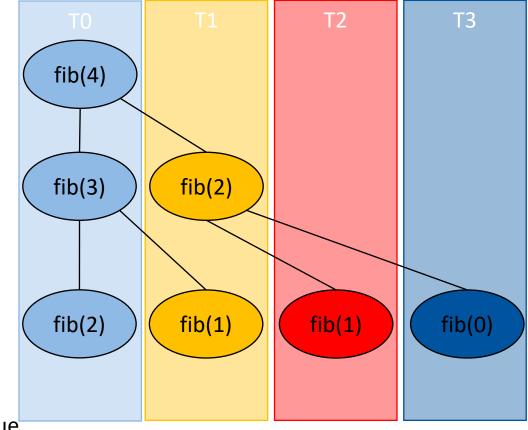
Dirk, how does that work in practice?



Illustration of Tasking



- T0 enters fib(4)
- T0 creates tasks for fib(3) and fib(2)
- T0 and T1 execute tasks from the queue
- T0 and T1 create 4 new tasks
- T0 T3 execute tasks









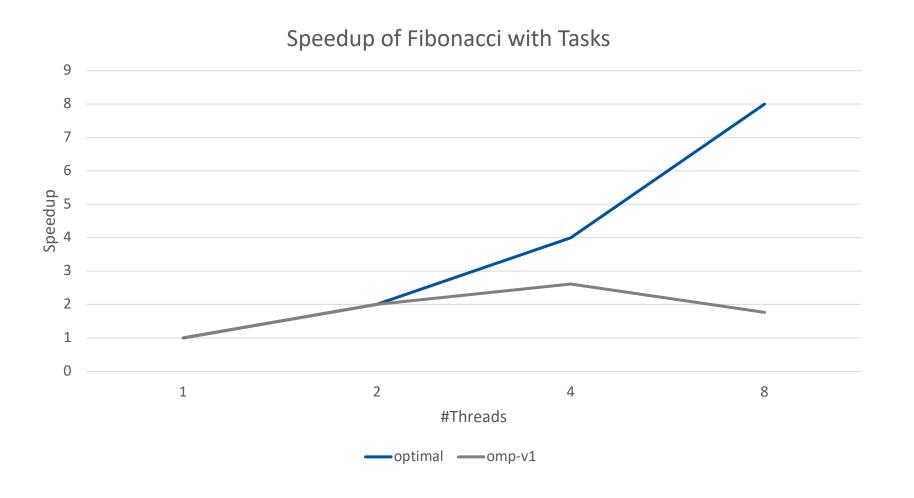
Got that.



Scalability Measurements (1/3)



Overhead of task creation prevents better scalability!







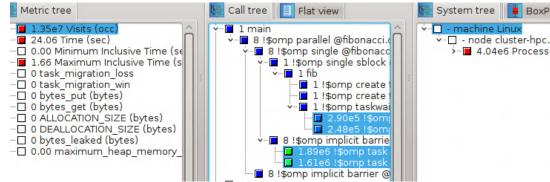
Christian, wait, let me explain!



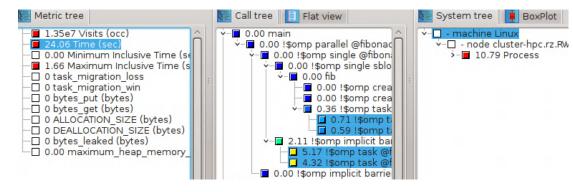
Performance Analysis



Event-based profiling gives a good overview :



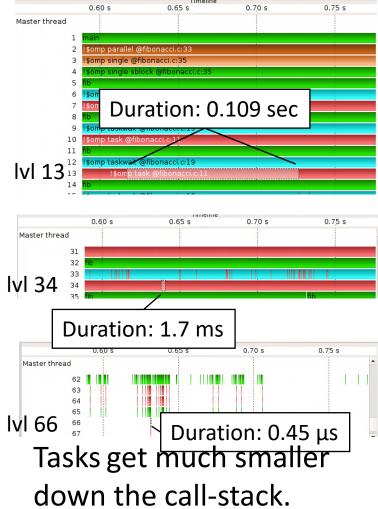
4.04 million threads are executed ...



... in ~10.79 seconds of CPU time.

=> average duration of a task is 2 2.6 µs

Tracing gives more details:





if Clause



- If the expression of an if clause on a task evaluates to false
 - The encountering task is suspended
 - The new task is executed immediately
 - The parent task resumes when the new task finishes
 - → Used for optimization, e.g., avoid creation of small tasks



Second version with Tasks (omp-v2)



Improvement: Don't create yet another task once a certain (small enough) n is reached

```
int main (int argc,
         char* argv[])
   [...]
#pragma omp parallel
#pragma omp single
   fib(input);
   [...]
```

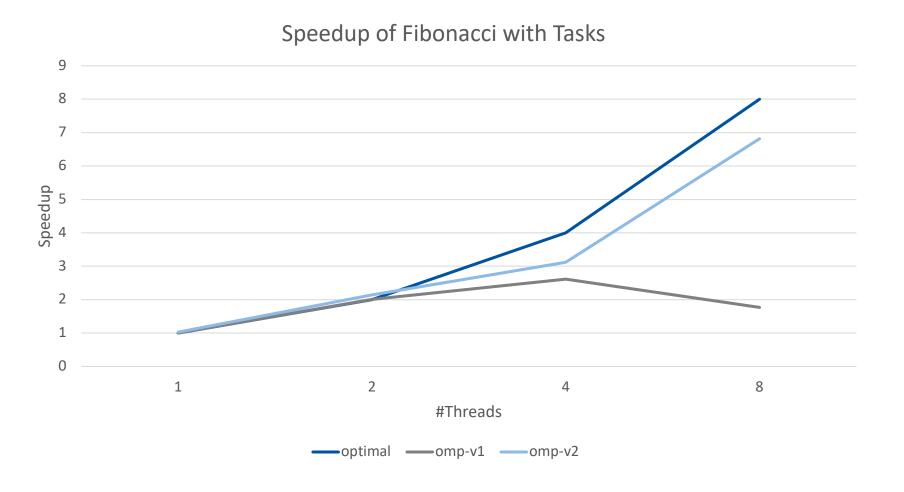
```
int fib(int n) {
   if (n < 2) return n;
int x, y;
#pragma omp task shared(x)
  \setminus if (n > 30)
  x = fib(n - 1);
#pragma omp task shared(y)
  y = fib(n - 2);
#pragma omp taskwait
   return x+y;
```



Scalability Measurements (2/3)



Speedup is ok, but we still have some overhead when running with 4 or 8 threads





Third version with Tasks (omp-v3)



Improvement: Skip the OpenMP overhead once a certain n is reached (no issue w/ production compilers)

```
int main (int argc,
         char* arqv[])
   [...]
#pragma omp parallel
#pragma omp single
   fib(input);
   [...]
```

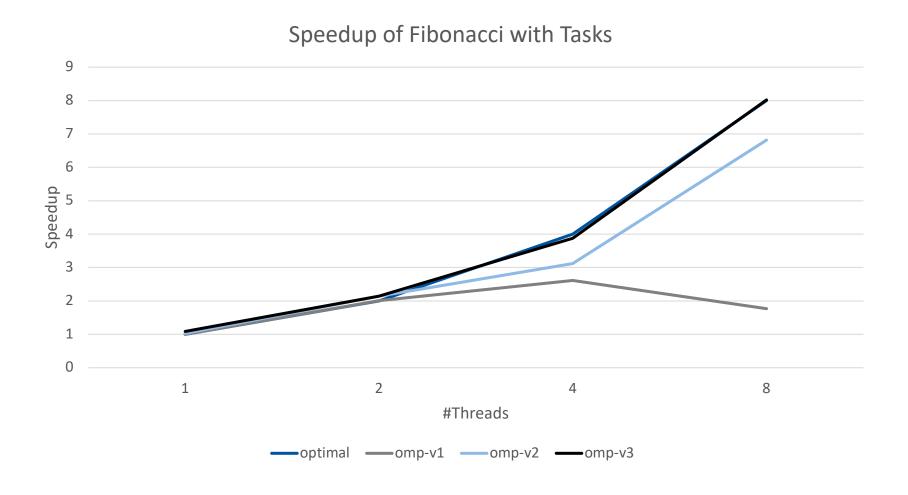
```
int fib(int n) {
   if (n < 2) return n;
   if (n \le 30)
      return serfib(n);
int x, y;
#pragma omp task shared(x)
   x = fib(n - 1);
#pragma omp task shared(y)
   y = fib(n - 2);
#pragma omp taskwait
   return x+y;
```



Scalability Measurements (3/3)



■ Everything ok now ^③







This also sounds interesting...



Example: taskloop



blocking

```
for ( i = 0; i<SIZE; i+=1) {
    A[i]=A[i]*B[i]*S;
}</pre>
```

```
taskloop
```

```
for ( i = 0; i < SIZE; i += TS) {
    UB = SIZE < (i + TS) ? SIZE: i + TS;
    for ( ii = i; ii < UB; ii + +) {
        A[ii] = A[ii] * B[ii] * S;
    }
}</pre>
```

```
for ( i = 0; i < SIZE; i += TS) {
    UB = SIZE < (i + TS) ? SIZE: i + TS;
    #pragma omp task private(ii) \
     firstprivate(i, UB) shared(S, A, B)
    for ( ii = i; ii < UB; ii + +) {
        A[ii] = A[ii] * B[ii] * S;
    }
}</pre>
```

```
#pragma omp taskloop grainsize(TS)
for ( i = 0; i<SIZE; i+=1) {
    A[i]=A[i]*B[i]*S;
}</pre>
```

- In manual transformation is difficult to determine grain
 - → 1 single iteration → to fine
 - → whole loop → no parallelism
- Apply blocking techniques
- taskloop: increase programmability



Example: task reductions



- Reduction operation
 - → perform some forms of recurrence calculations
 - → associative and commutative operators
- The (taskgroup) reduction clause

```
#pragma omp taskgroup task_reduction(op: list)
{structured-block}
```

- → Register a new reduction at [1]
- → Computes the final result after [3]
- The (task) in_reduction clause [participating]

```
#pragma omp task in_reduction(op: list)
{structured-block}
```

→ Task participates in a reduction operation

```
int res = 0:
node t* node = NULL;
#pragma omp parallel
 #pragma omp single
  #pragma omp taskgroup task_reduction(+: res)
  { // [1]
    while (node) {
     #pragma omp task in_reduction(+: res) \
              firstprivate(node)
     { // [2]
       res += node->value;
     node = node->next;
  } // [3]
                       OpenMP 5.0
```

Example: task dependencies



Task dependences as a way to define task-execution constraints

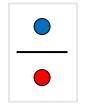
```
int x = 0;
#pragma omp parallel
#pragma omp single
{
    #pragma omp task
    std::cout << x << std::endl;

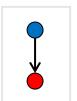
    #pragma omp taskwait

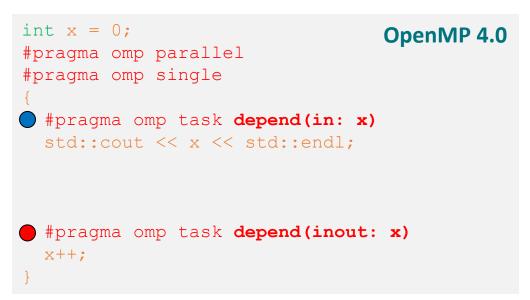
    #pragma omp task
    x++;
}</pre>
```

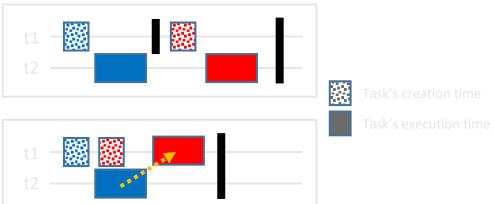
















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