

Lazy evaluation in Haskell

exploring some mental models and implementations

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Lazy,... zzz

..., It's fun.

NOTE

- Meaning of terms are different by communities.
- There are a lot of good documents. Please see also references.
- This is written for GHC's Haskell.

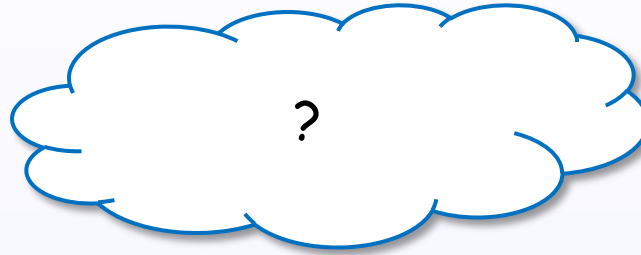
Contents

- Introduction
- Evaluations
- Expressions in Haskell
- Constructor
- Thunk
- let, case expression
- WHNF
- Evaluation in Haskell (GHC)
- Control the evaluation in Haskell
- Implementation in GHC
- Semantics
- References

Introduction

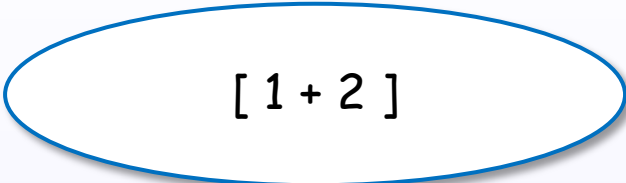
What is an expression?

An expression



An expression denotes a value

An expression



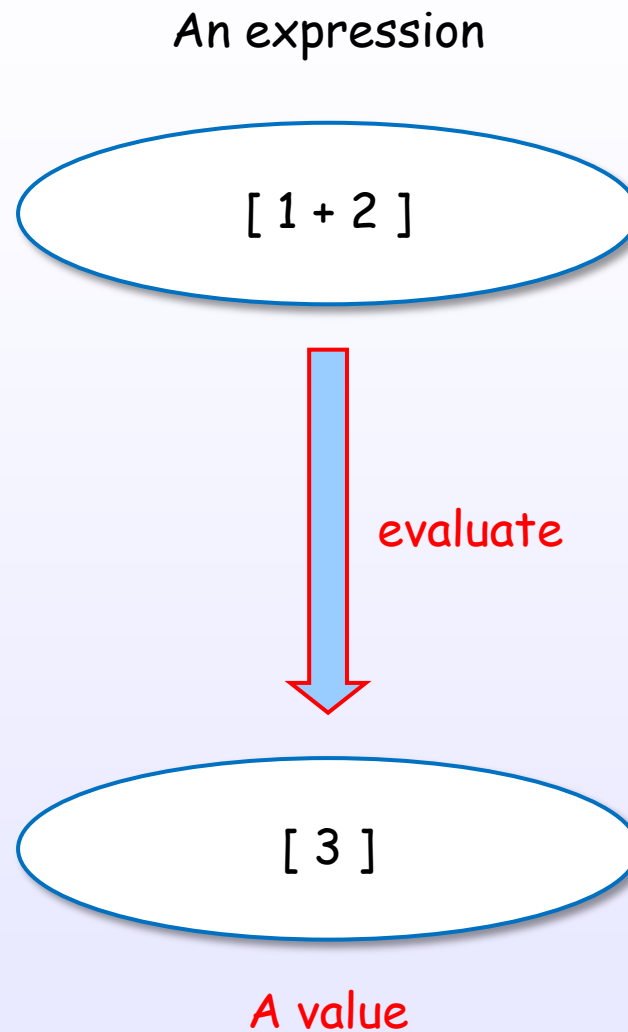
[1 + 2]

[HR2010]

[Bird, Chapter 2]

References : [1]

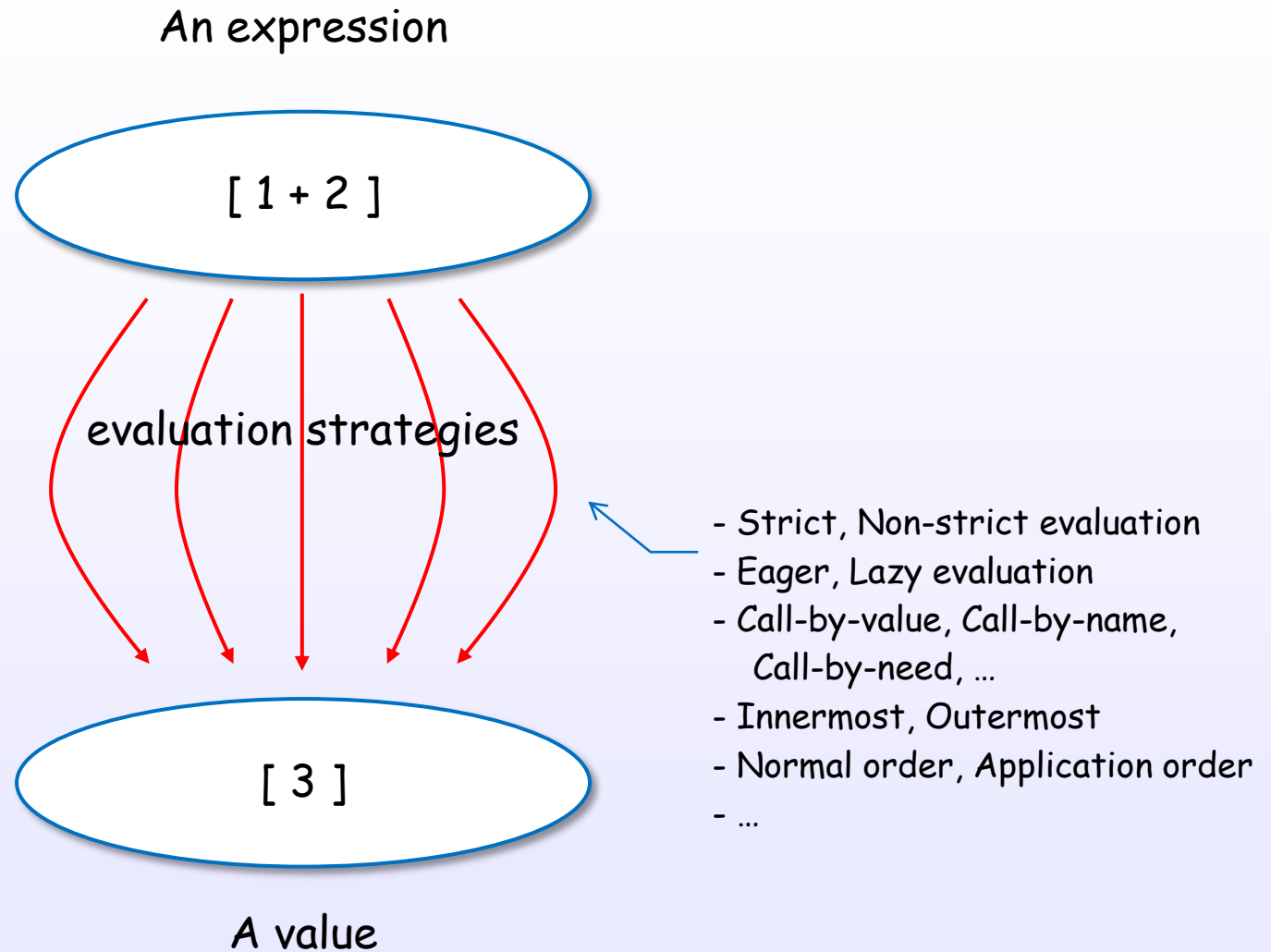
An expression evaluates to a value



[HR2010]

[Bird, Chapter 2]

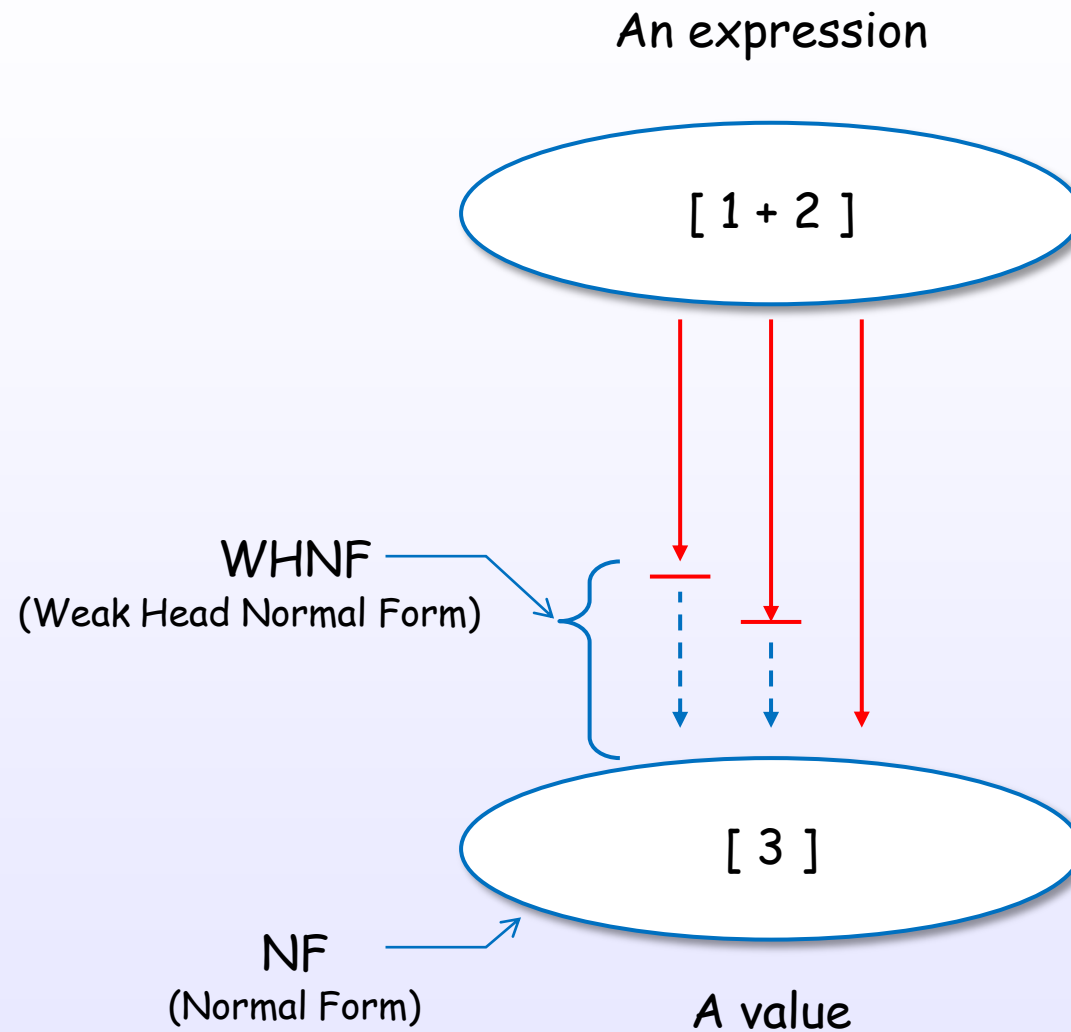
There are many evaluation approaches



[Bird, Chapter 2, 7]

[TAPL, Chapter 3]

There are some evaluation levels



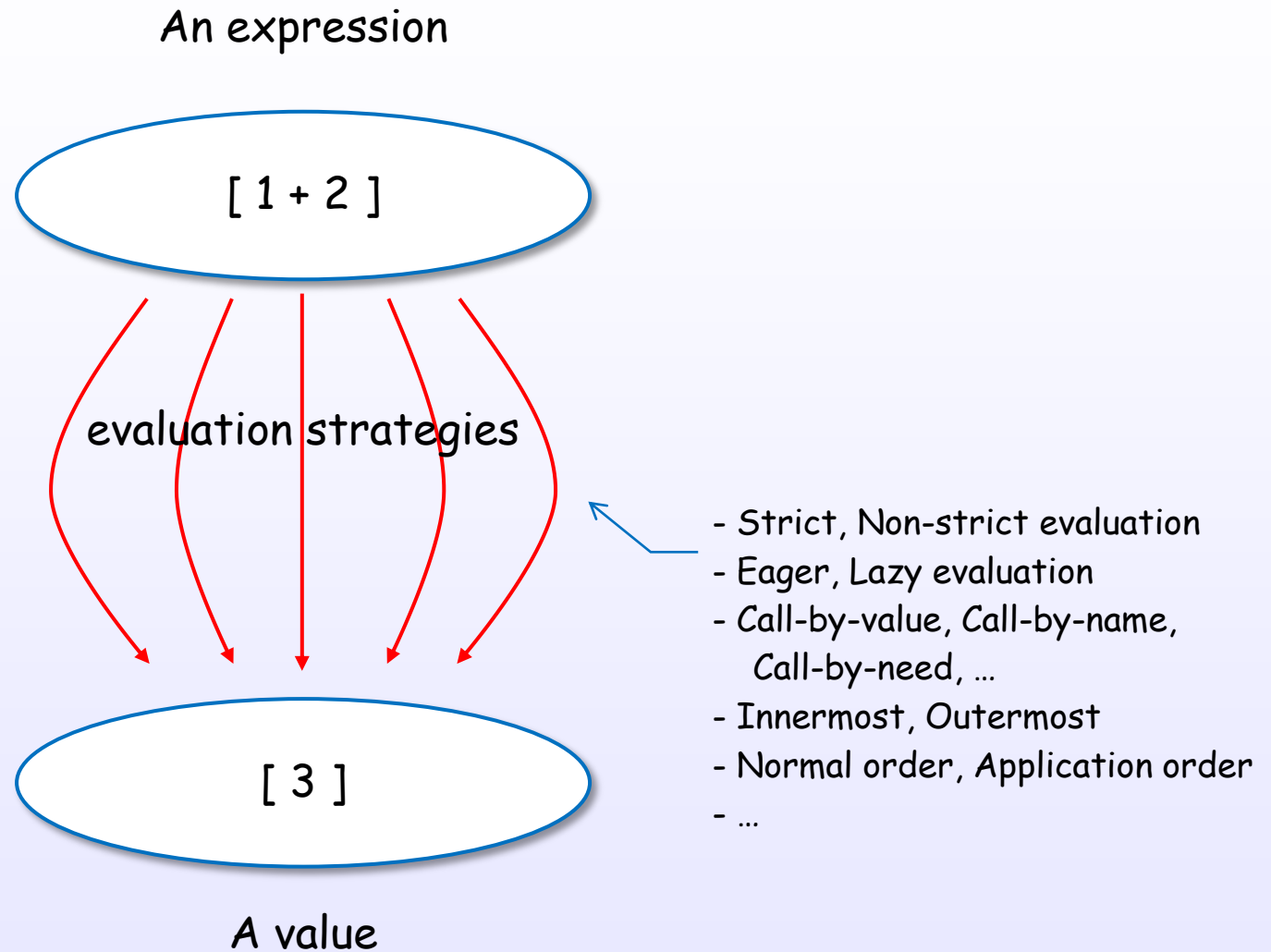
[Terei]

[Bird, Chapter 2, 7]

[TAPL, Chapter 3]

Evaluations

There are many evaluation approaches



[Bird, Chapter 2, 7]

[TAPL, Chapter 3]

Evaluation layers

denotational semantics

evaluation strategy

evaluation implementation

[Bird, Chapter 7]

[Hutton, Chapter 8]

[TAPL, Chapter 3]

References : [1]

Evaluation layers

denotational
semantics

Strict semantics

Non-strict semantics

evaluation
strategy

Eager evaluation
(Strict evaluation)

Nondeterministic
evaluation

Lazy evaluation
(Non-strict evaluation)

...

Call-by-Value

Call-by-Name

Call-by-Need

...

evaluation
implementation

Graph reduction

...

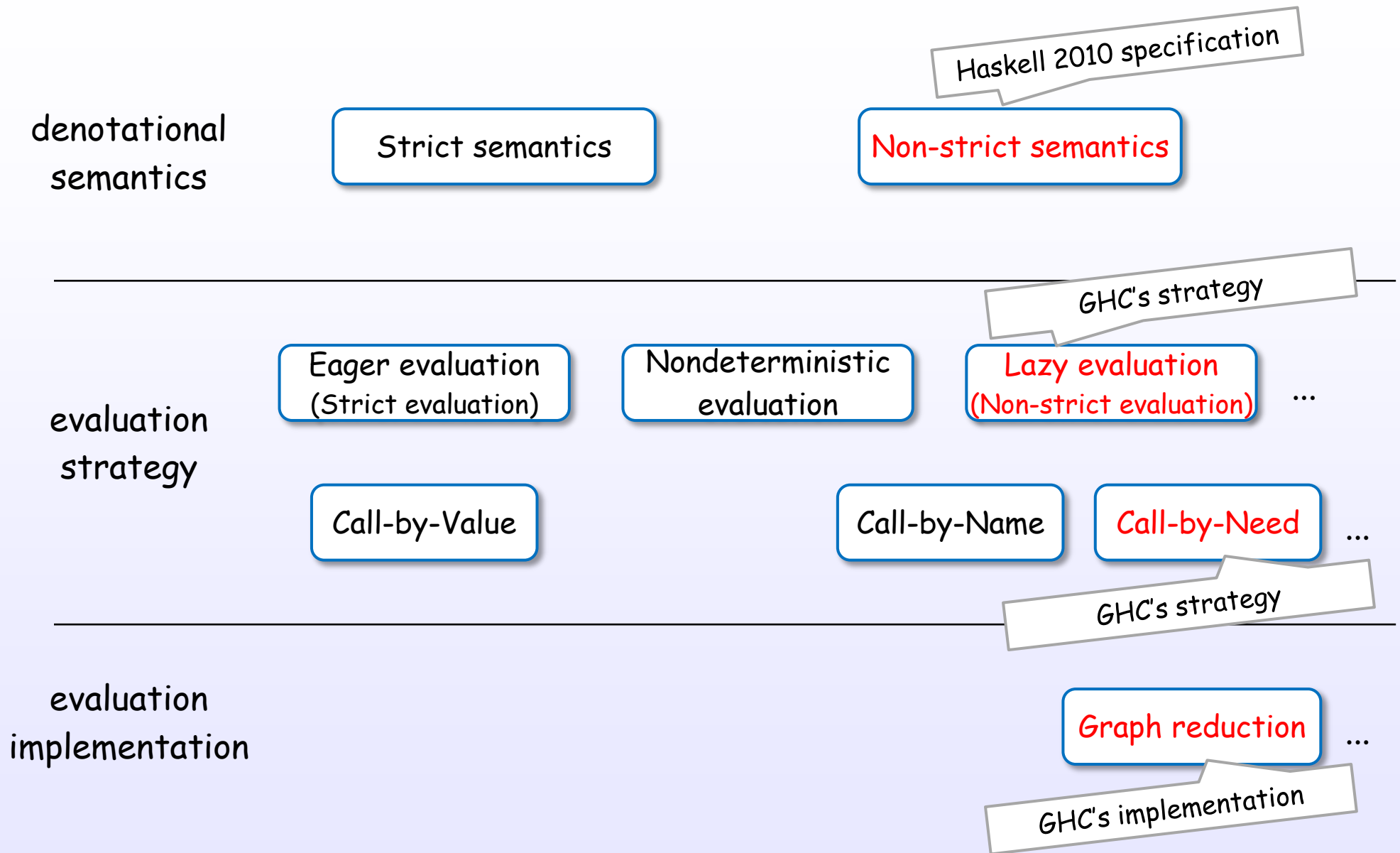
[Bird, Chapter 7]

[Hutton, Chapter 8]

[TAPL, Chapter 3]

References : [1]

Evaluation layers for GHC's Haskell



Simple example of both evaluations

Eager evaluation (Strict evaluation)

default
C, Java, JavaScript,
Python, OCaml, Scheme, ...

square (1 + 2)



argument
evaluation
first

square (3)



3 * 3



9

Lazy evaluation (Non-strict evaluation)

default
Haskell (GHC), ...

square (1 + 2)



apply
first

(1 + 2) * (1 + 2)



(3) * (3)



9

[Bird]
[Hutton]

Simple example of both evaluations

Eager evaluation
(Strict evaluation)

square (1 + 2)



square (3)



3 * 3



9

argument
evaluated

Lazy evaluation
(Non-strict evaluation)

square (1 + 2)



(1 + 2) * (1 + 2)



(3) * (3)



9

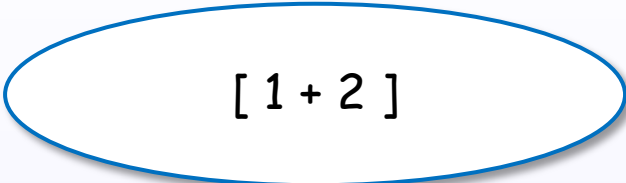
argument
evaluation
delayed !

[Bird]
[Hutton]

Expressions in Haskell

An expression denotes a value

An expression



[1 + 2]

[HR2010]

[Bird, Chapter 2]

References : [1]

There are many expressions in Haskell

Expressions

Just 5

$1 + 2$

$(1, 2)$

take 5 xs

$[1, 2, 3]$

let $x = 1$ in $x + y$

'a'

map f xs

if b then 1 else 0

7

$\forall x \rightarrow x + 1$

$x : xs$

fun arg

case x of $_ \rightarrow 0$

$(\forall x \rightarrow x + 1) 3$

do { $x \leftarrow$ get; put x}



categorizing

[HR2010]

[Bird, Chapter 2]

References : [1]

Expression categories in Haskell

lambda abstraction

$\forall x \rightarrow x + 1$

let expression

let $x = 1$ in $x + y$

conditional

if b then 1 else 0

case expression

case x of $_ \rightarrow 0$

do expression

do { $x \leftarrow \text{get}$; put x }

general constructor, literal and some forms

7

[1, 2, 3]

(1, 2)

'a'

$x : xs$

Just 5

function application

take 5 xs

$(\forall x \rightarrow x + 1)$ 3

1 + 2

map f xs

fun arg

[HR2010]
[Bird, Chapter 2]

Specification is defined in Haskell 2010 Language Report

Haskell 2010 Language Report, Chapter 3 Expressions [1]

<i>exp</i>	→	<i>infixexp</i> :: [context =>] type <i>infixexp</i>	(expression type signature)
<i>infixexp</i>	→	<i>lexp</i> <i>qop</i> <i>infixexp</i> - <i>infixexp</i> <i>lexp</i>	(infix operator application) (prefix negation)
<i>lexp</i>	→	\ <i>apat</i> ₁ ... <i>apat</i> _{<i>n</i>} -> <i>exp</i> let <i>decls</i> in <i>exp</i> if <i>exp</i> [<i>i</i>] then <i>exp</i> [<i>i</i>] else <i>exp</i> case <i>exp</i> of { <i>alts</i> } do { <i>stmts</i> } <i>fexp</i>	(lambda abstraction, $n \geq 1$) (let expression) (conditional) (case expression) (do expression)
<i>fexp</i>	→	[<i>fexp</i>] <i>aexp</i>	(function application)
<i>aexp</i>	→	<i>qvar</i> <i>gcon</i> <i>literal</i> (<i>exp</i>) (<i>exp</i> ₁ , ... , <i>exp</i> _{<i>k</i>}) [<i>exp</i> ₁ , ... , <i>exp</i> _{<i>k</i>}] [<i>exp</i> ₁ [, <i>exp</i> ₂] .. [<i>exp</i> ₃]] [<i>exp</i> <i>qual</i> ₁ , ... , <i>qual</i> _{<i>n</i>}] (<i>infixexp</i> <i>qop</i>) (<i>qop</i> { - } <i>infixexp</i>) <i>qcon</i> { <i>fbind</i> ₁ , ... , <i>fbind</i> _{<i>n</i>} } <i>aexp</i> _{<i>qcon</i>} { <i>fbind</i> ₁ , ... , <i>fbind</i> _{<i>n</i>} }	(variable) (general constructor) (parenthesized expression) (tuple, $k \geq 2$) (list, $k \geq 1$) (arithmetic sequence) (list comprehension, $n \geq 1$) (left section) (right section) (labeled construction, $n \geq 0$) (labeled update, $n \geq 1$)

Constructor

Constructor

Constructor is one of the key elements to understand WHNF and lazy evaluation.

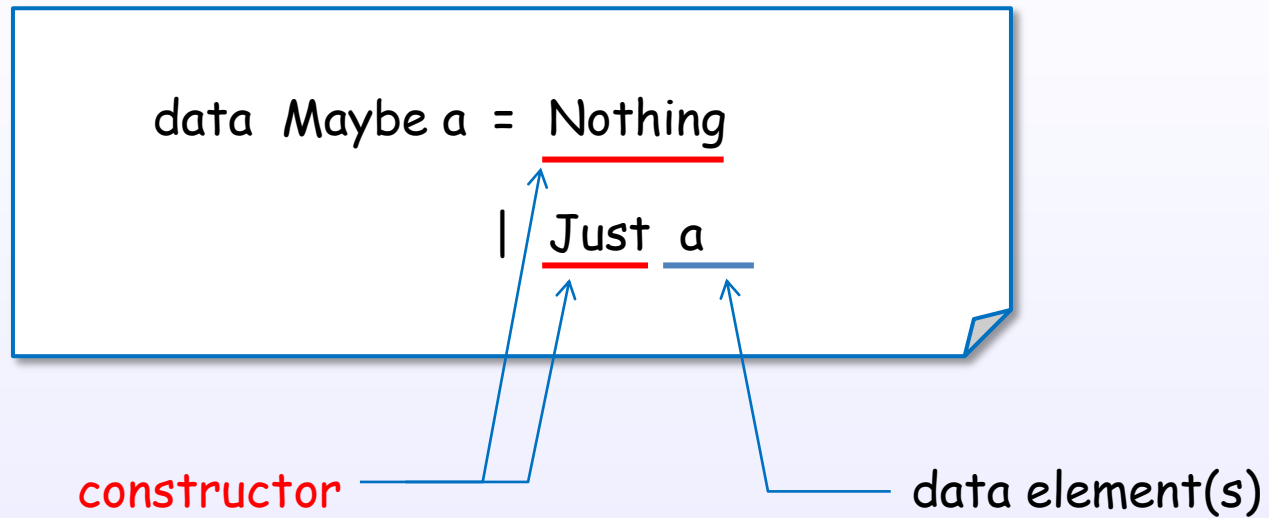
data文で宣言する代数的データ型とその値

```
data Maybe a = Nothing  
              | Just a
```

Algebraic Data Type

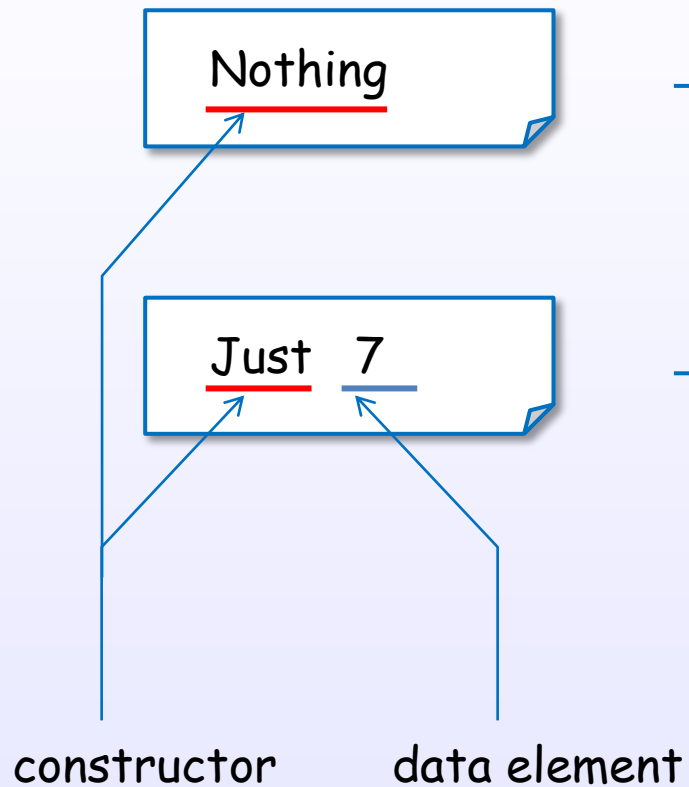
Algebraic Data **Value**

Constructorはdata文で宣言する代数的データ値

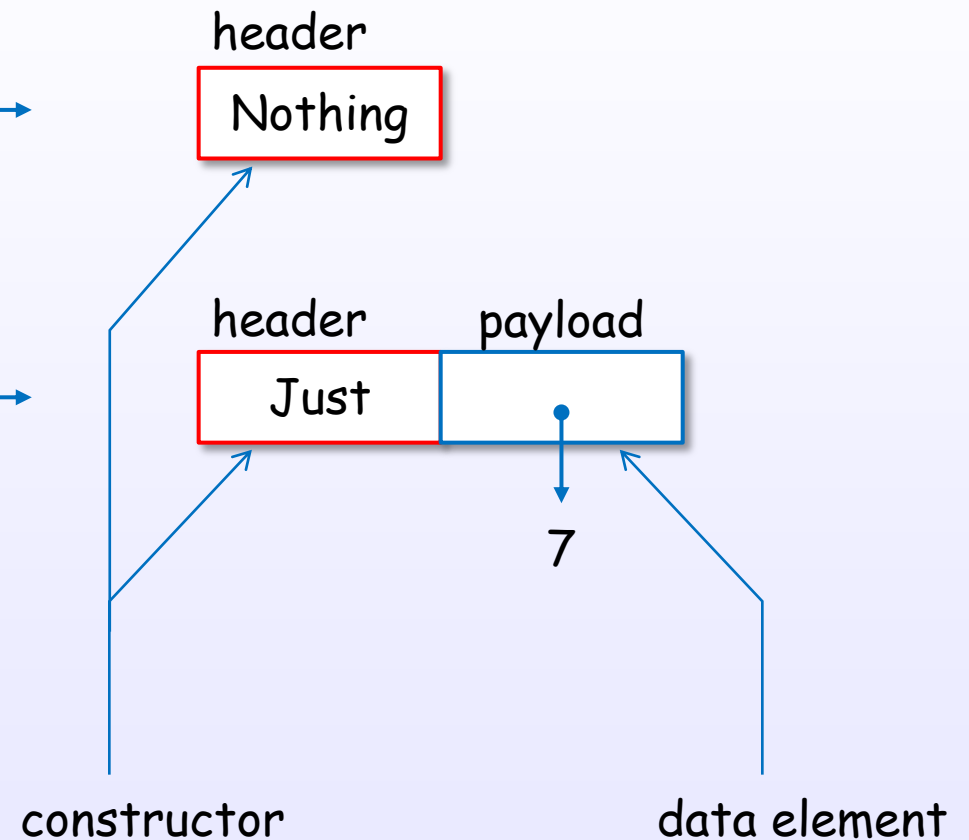


Constructorの内部表現

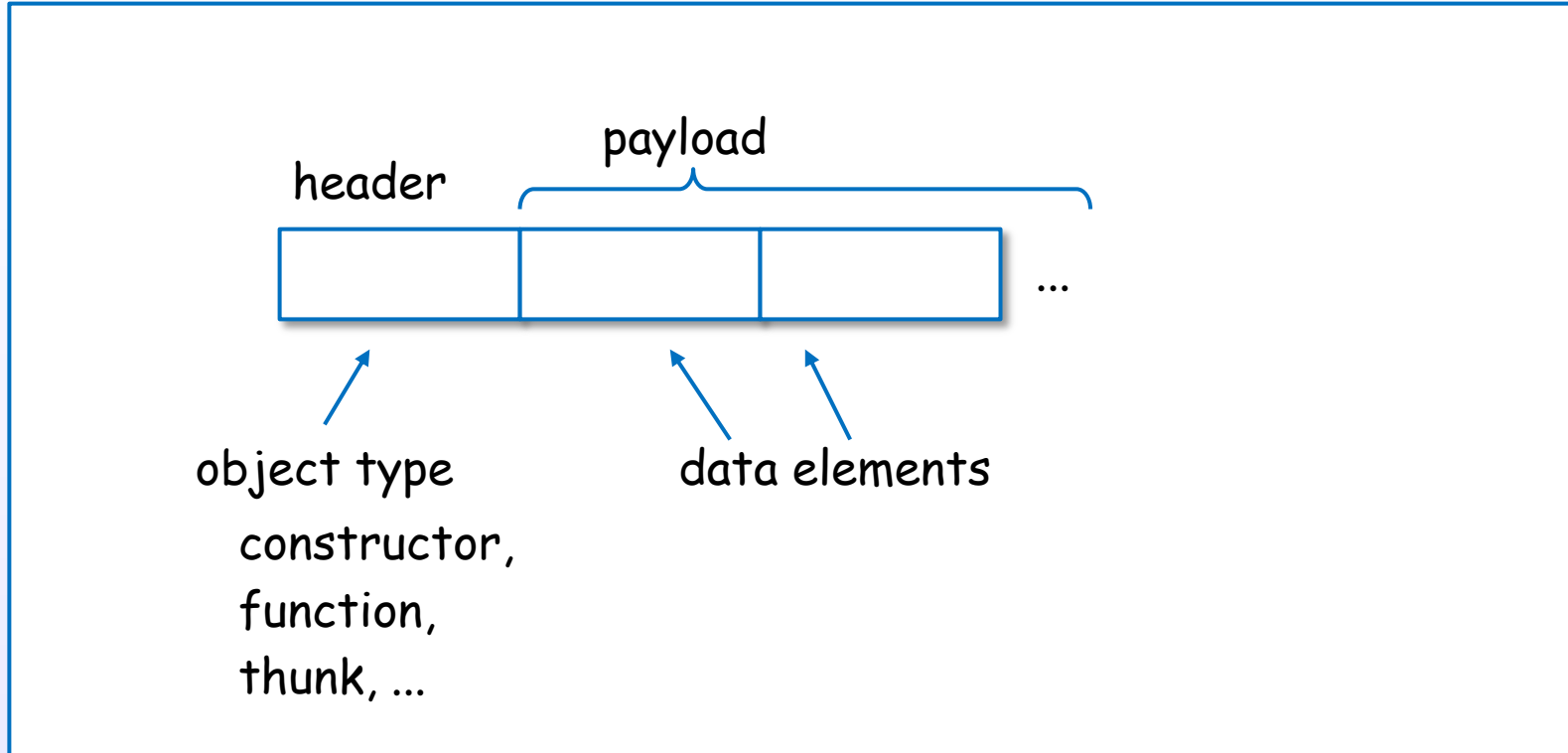
Haskell code



GHC's internal representation
in heap memory



Constructorは統一内部表現で表現される



in heap memory, stack, registers or static memory

いろいろなコンストラクタと内部表現

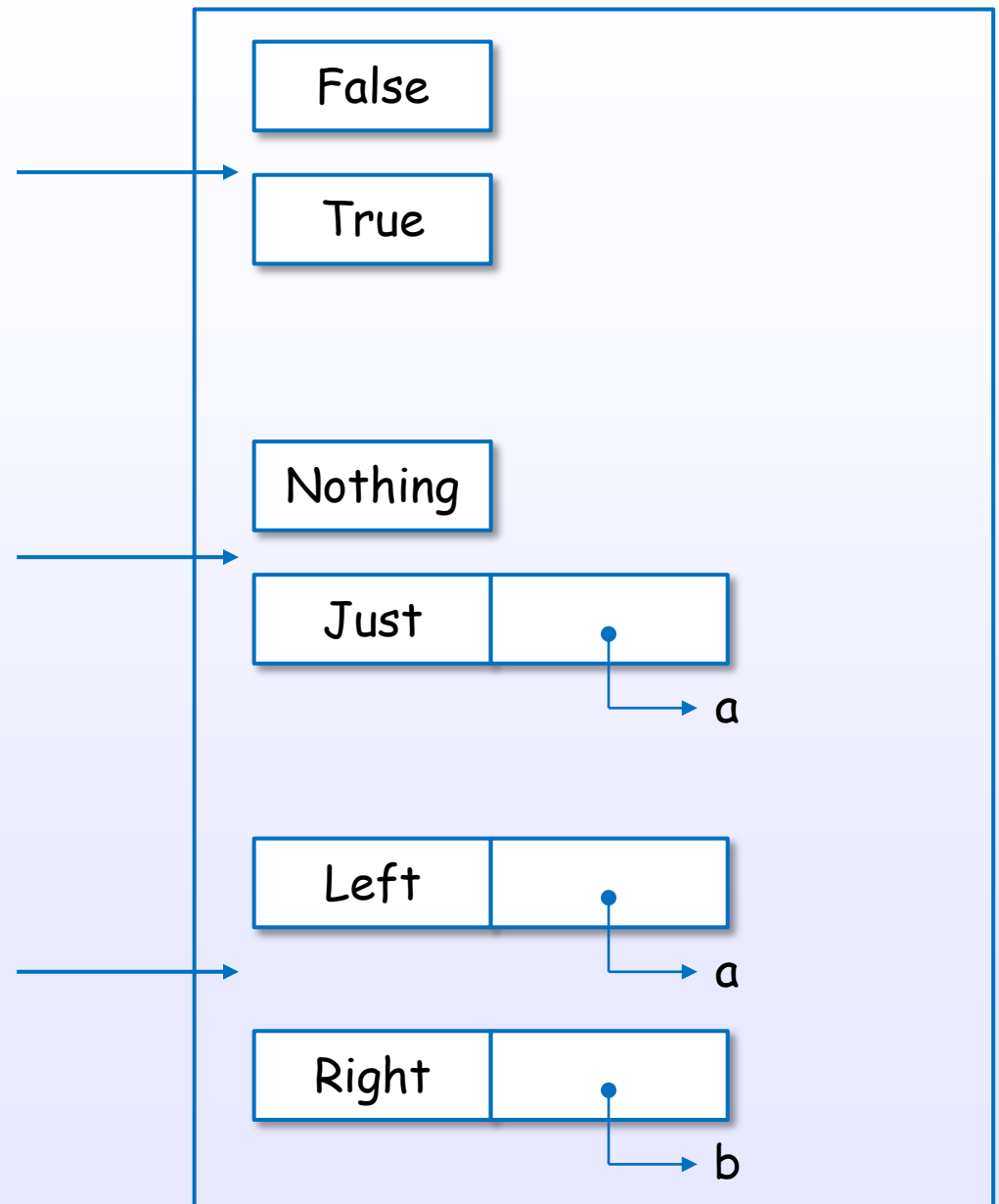
Haskell code

```
data Bool = False  
          | True
```

```
data Maybe a = Nothing  
              | Just a
```

```
data Either a b = Left a  
                 | Right b
```

GHC's internal representation



基本データ型も実はコンストラクタで構成されている

Haskell code

```
data Int = Int# 0#  
        | Int# 1#  
        | ...
```

```
data Char = Char# 'a'#  
          | Char# 'b'#  
          | ...
```

GHC's internal representation

→

Int#	0#
Int#	1#
⋮	

→

Char#	'a'#
Char#	'b'#
⋮	

リストも実はコンストラクタで構成されている

List

[1, 2, 3]

syntactics desugar

1 : (2 : (3 : []))

prefix notation by section

(:) 1 ((:) 2 ((:) 3 []))

equivalent data constructor

Cons 1 (Cons 2 (Cons 3 Nil))

constructor

リストも実はコンストラクタで構成されている

List

```
[ 1, 2, 3 ]
```

syntactics desugar

```
1 : ( 2 : ( 3 : [] ) )
```

prefix notation by section

```
(:) 1 ( (:) 2 ( (:) 3 [] ) )
```

equivalent data constructor

```
Cons 1 ( Cons 2 ( Cons 3 Nil ) )
```

** pseudo code*

```
data List a = []  
             | : a (List a)
```

```
data List a = Nil  
             | Cons a (List a)
```


リストも実はコンストラクタで構成されている

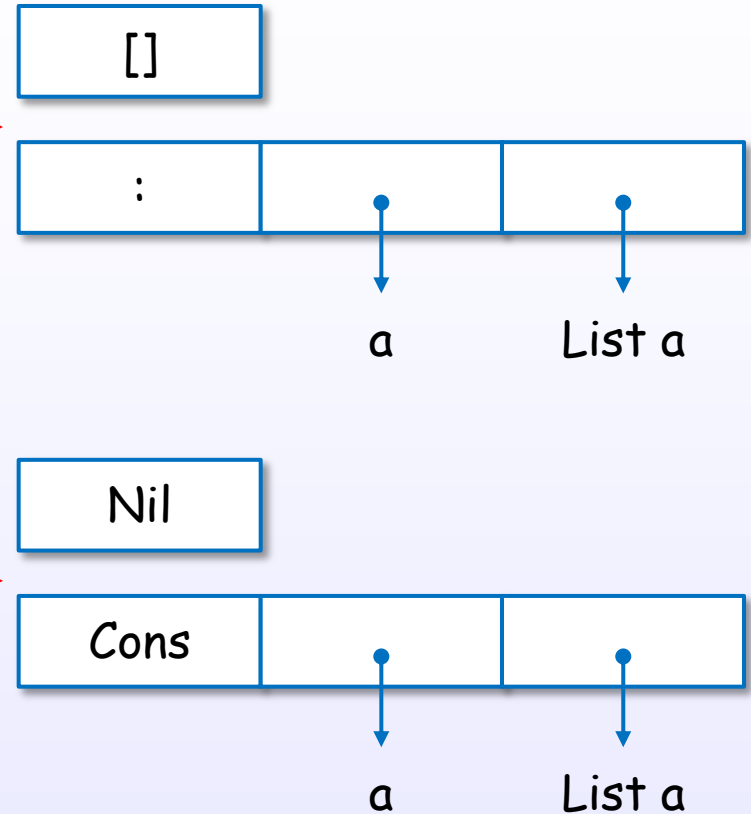
Haskell code

```
data List a = []  
           | : a (List a)
```

↕ equivalent data constructor

```
data List a = Nil  
           | Cons a (List a)
```

GHC's internal representation



リストも実はコンストラクタで構成されている

Haskell code

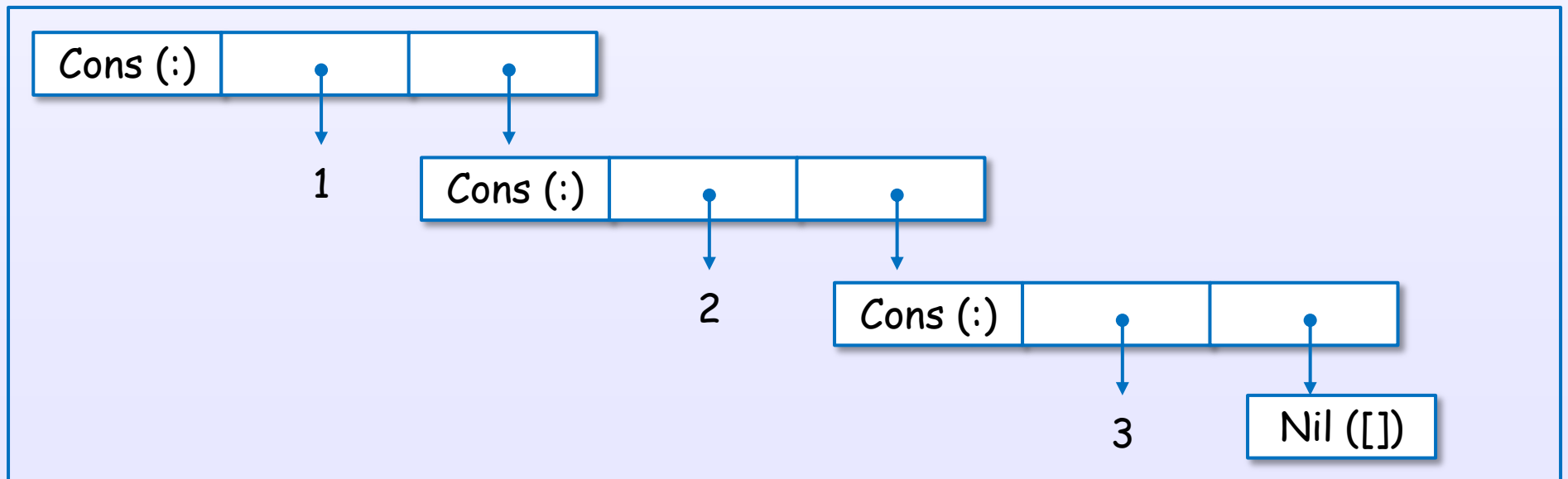
```
[ 1, 2, 3 ]
```

```
1 : ( 2 : ( 3 : [] ) )
```

```
(:) 1 ( (:) 2 ( (:) 3 [] ) )
```

```
Cons 1 ( Cons 2 ( Cons 3 Nil ) )
```

GHC's internal representation



タプルも実はコンストラクタで構成されている

Tuple (Pair)

(7 , 8)

prefix notation by section

(,) 7 8

equivalent data constructor

Pair2 7 8

constructor

** pseudo code*

data Pair2 a = (,) a a

data Pair2 a = Pair2 a a

タプルも実はコンストラクタで構成されている

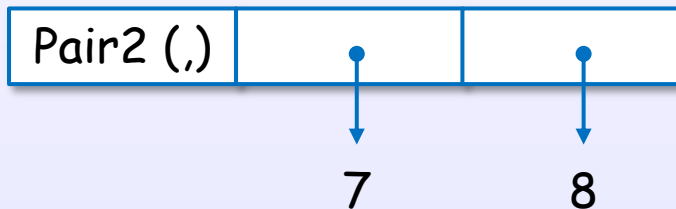
Haskell code

```
( 7 , 8 )
```

```
(,) 7 8
```

```
Pair2 7 8
```

GHC's internal representation



Thunk

Thunk

think
(unevaluated expression/
suspended computation)

A thunk is an **unevaluated** expression in heap memory.

Thunk

An unevaluated expression

take x ys



create/allocate

thunk
(unevaluated expression/
suspended computation)

in heap memory

A thunk is created for an unevaluated expression.

Thunkの内部表現

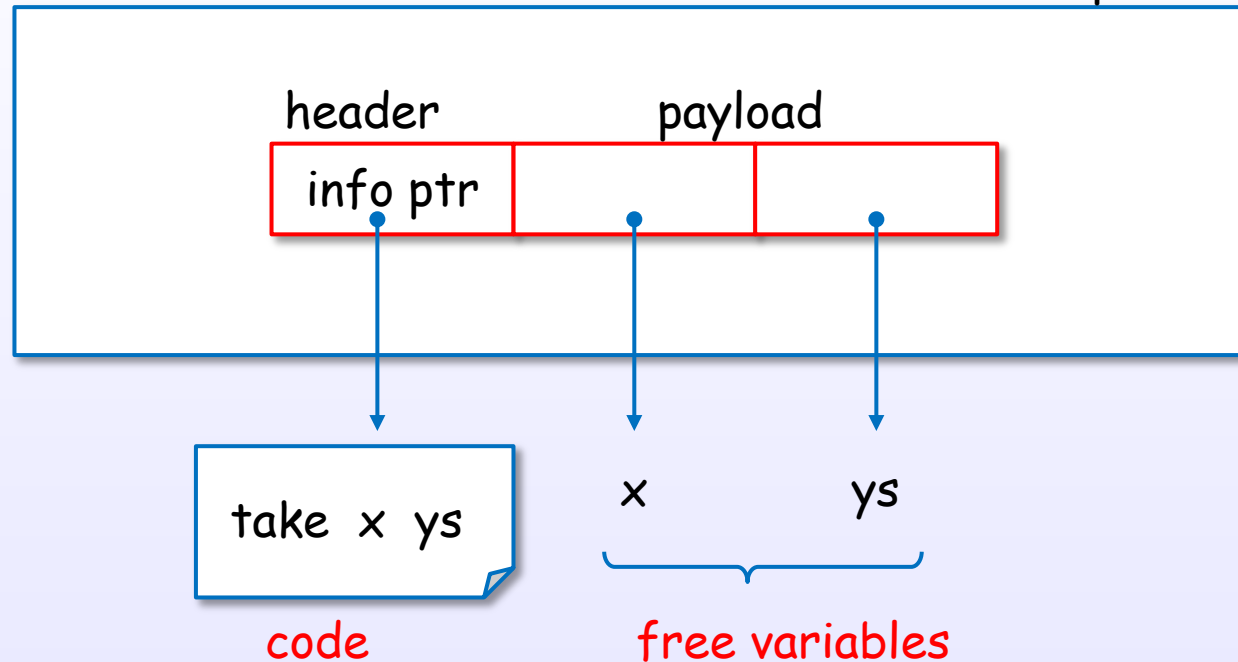
An unevaluated expression

take x ys



thunk

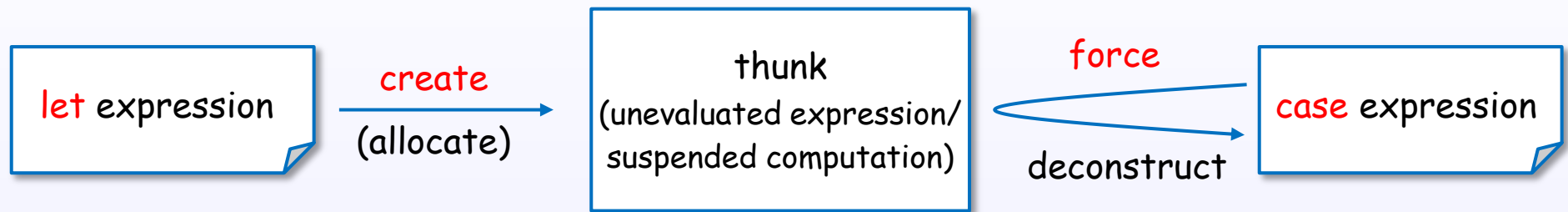
GHC's internal representation



A thunk is represented with header(code) + payload(free variables).

let, case expression

let/case expressions and thunk



A let expression may create a thunk.

A case expression forces and deconstructs the thunk.

A let expression creates a thunk

let expression

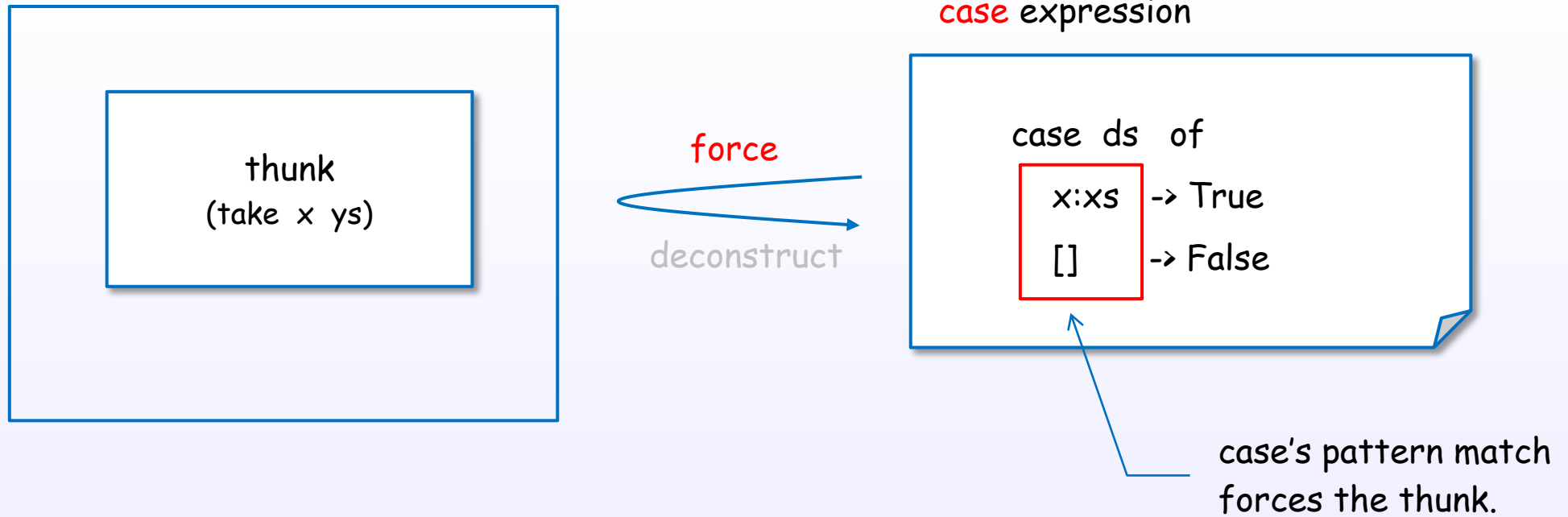
```
let ds = take x ys
```

create
(allocate)

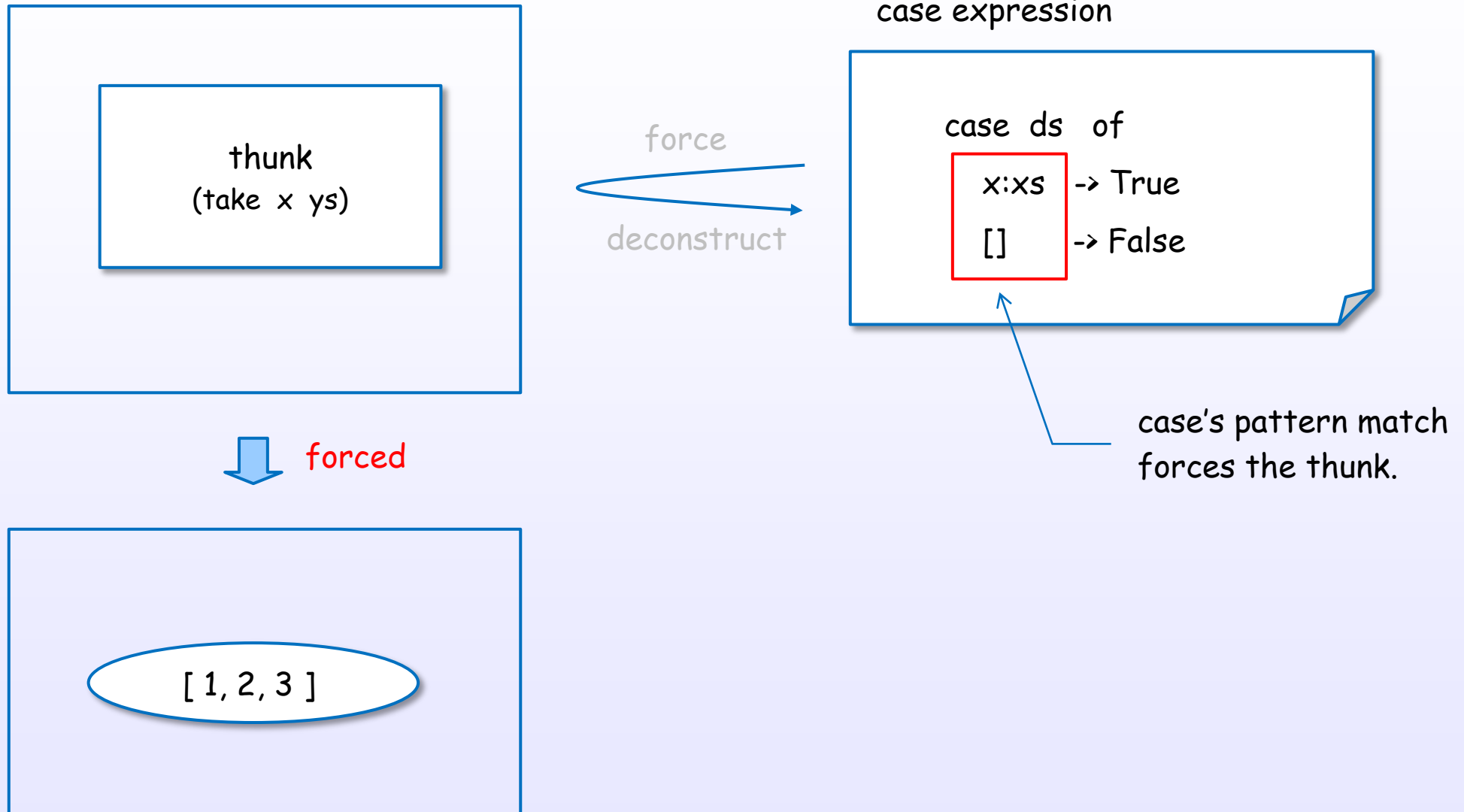
```
thunk  
(take x ys)
```

heap memory

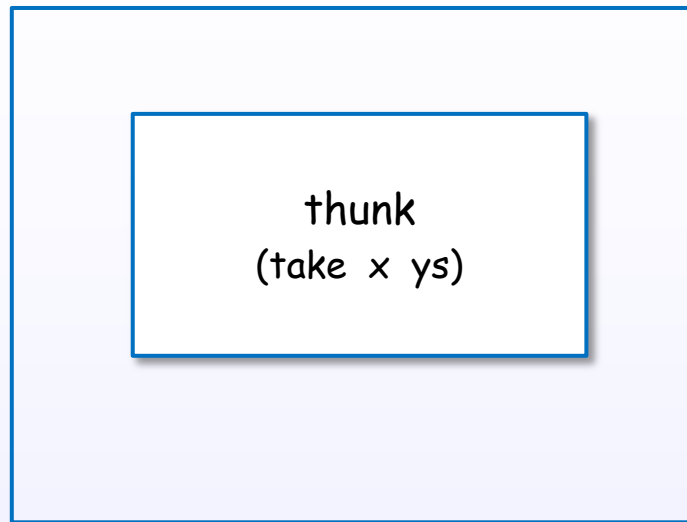
A case expression forces a thunk



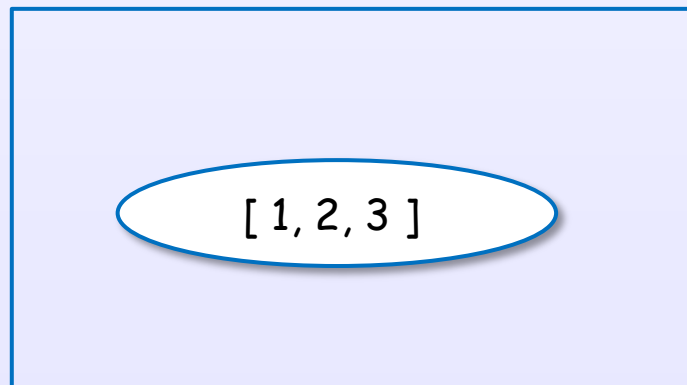
A case expression forces a thunk



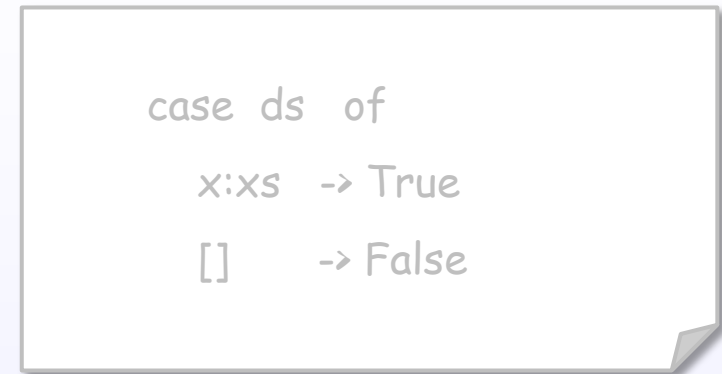
A case expression forces a thunk



↓ forced

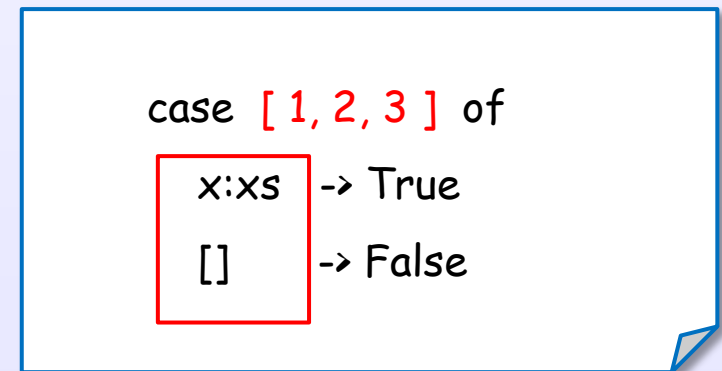


case expression



force
deconstruct

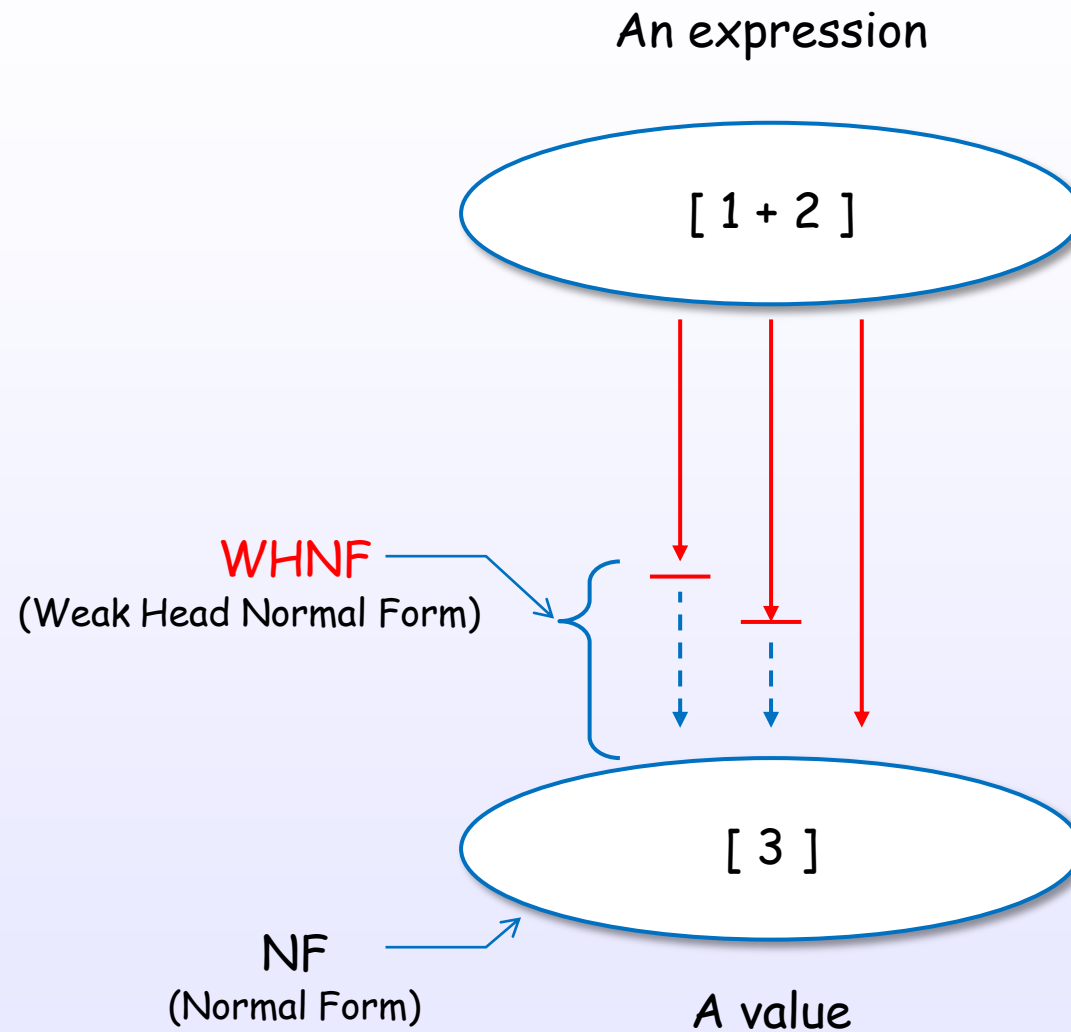
case expression



force
deconstruct

WHNF

There are some evaluation levels



[Terei]

[Bird, Chapter 2, 7]

[TAPL, Chapter 3]

[4]

normal form:

an expression without an redexes

head normal form:

an expression where the top level (head) is neither a redex NOR
a lambda abstraction with a reducible body

weak head normal form:

an expression where the top level (head) isn't a redex

[Terei]

[4]

evaluation strategies:

call-by-value: arguments evaluated before function entered (copied)

call-by-name: arguments passed unevaluated

call-by-need: arguments passed unevaluated but an expression is only evaluated once (sharing)

no-strict evaluation Vs. lazy evaluation:

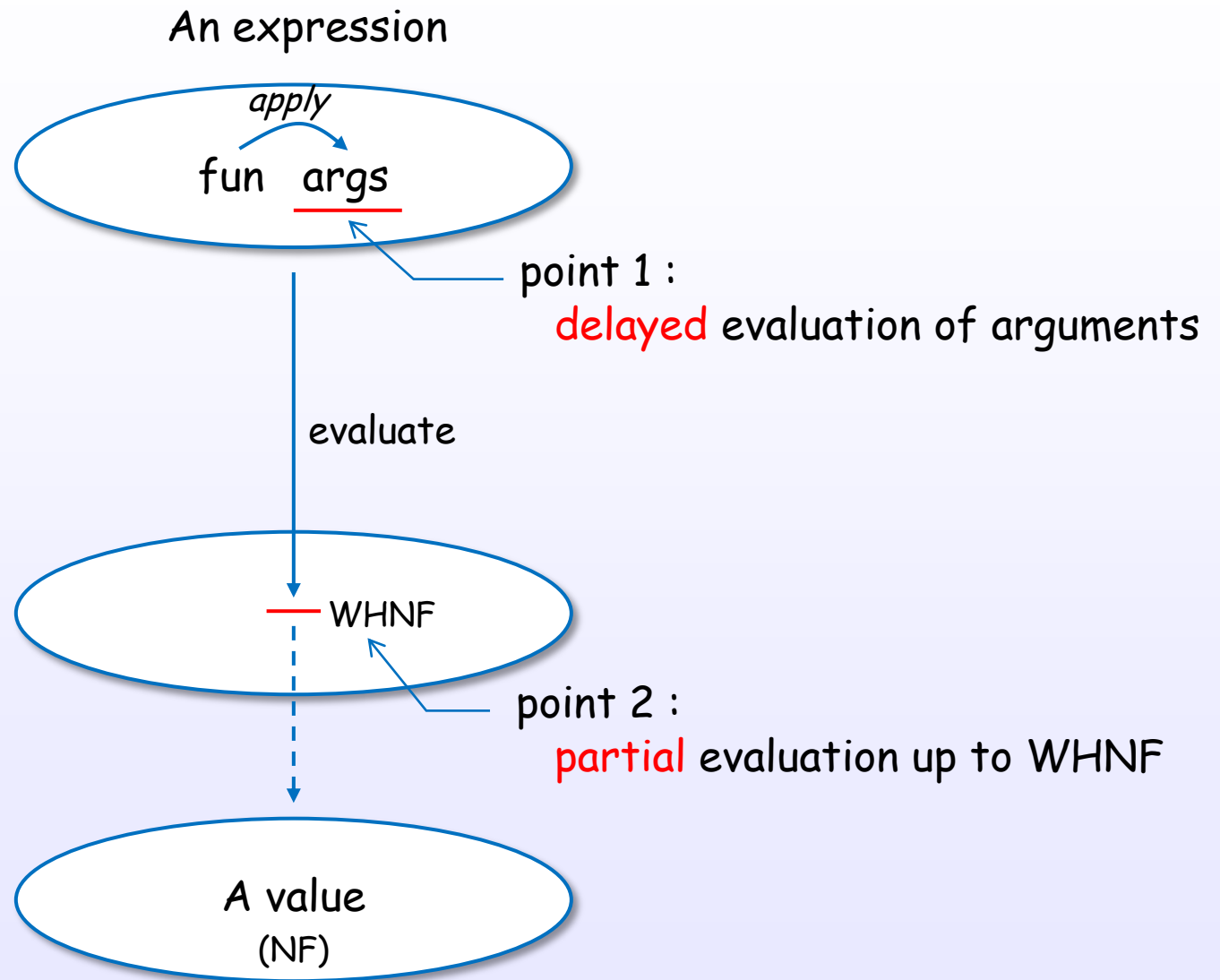
non-strict: Includes both call-by-name and call-by-need, general term for evaluation strategies that don't evaluate arguments before entering a function

lazy evaluation: Specific type of non-strict evaluation. Uses call-by-need (for sharing).

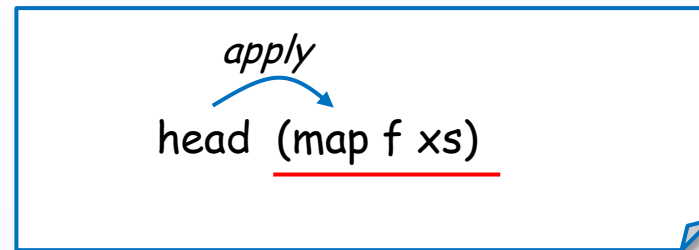
[Terei]

Evaluation in Haskell (GHC)

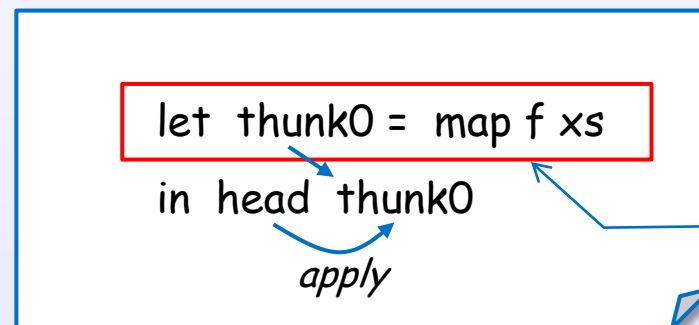
Key concept of Haskell's lazy evaluation



point 1 : delayed evaluation of arguments



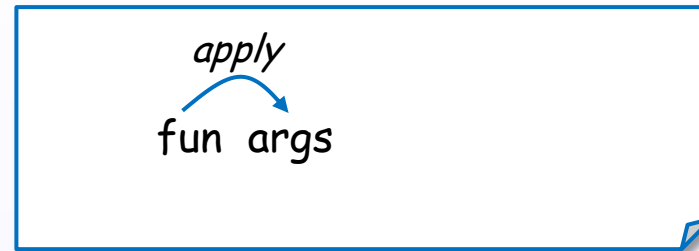
↓ internal transformation by *GHC*



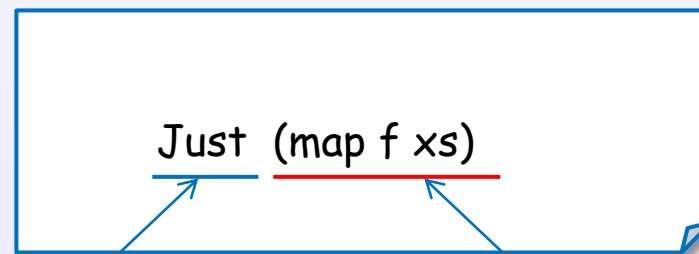
create a **thunk**
in heap memory

GHC implements lazy evaluation using the thunk.
Evaluation of arguments is delayed with the thunk.

point 2 : partial evaluation up to WHNF



↓ evaluation up to WHNF



evaluated part
(head constructor)

unevaluated part
(thunk)

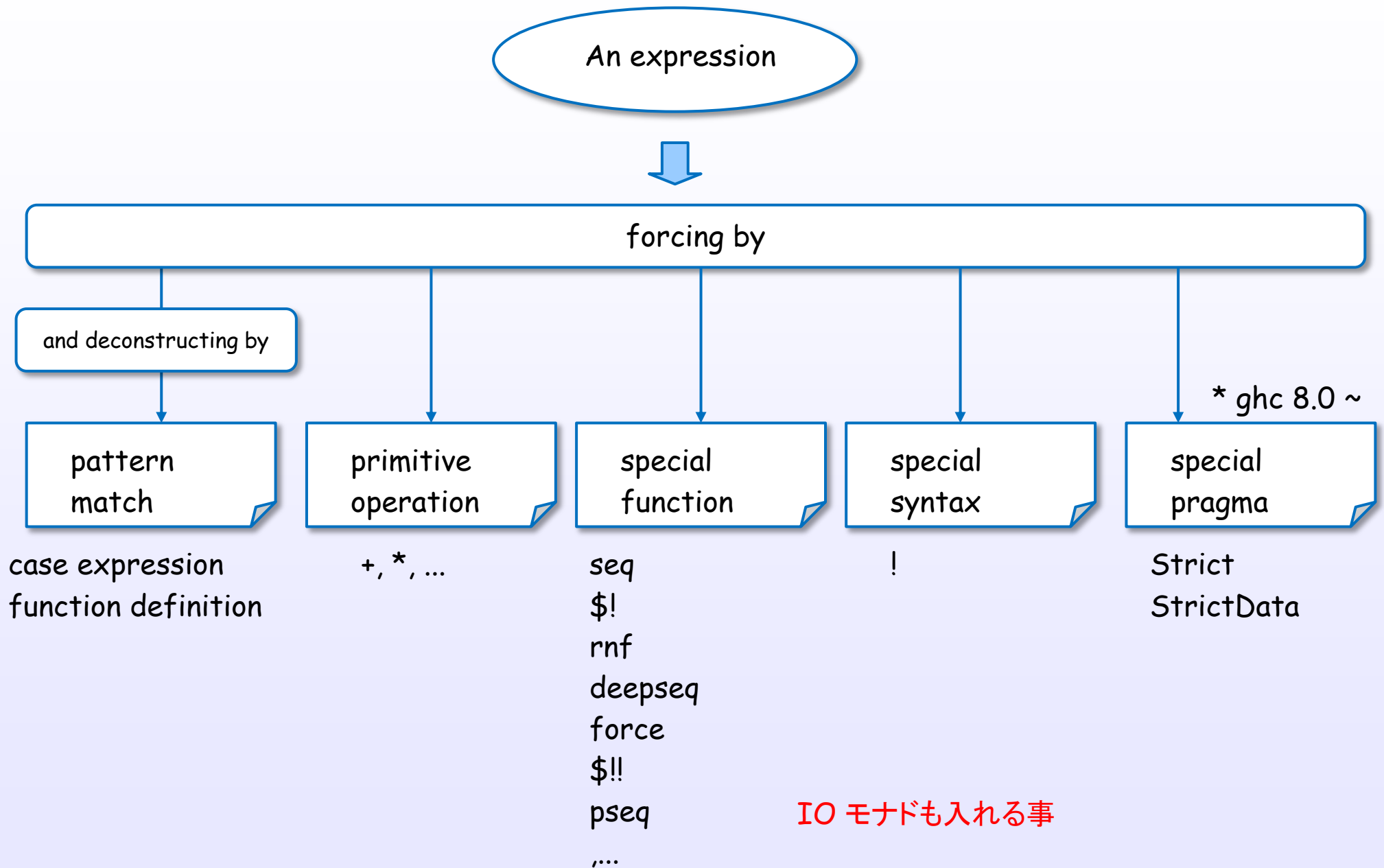
GHC can partially evaluate an expression.
Constructor can hold an unevaluated expression (a thunk).

Pattern match

[CIS194]

Control the evaluation in Haskell

How to drive evaluation



Implementation in GHC

Tree, Graph

a expression

AST

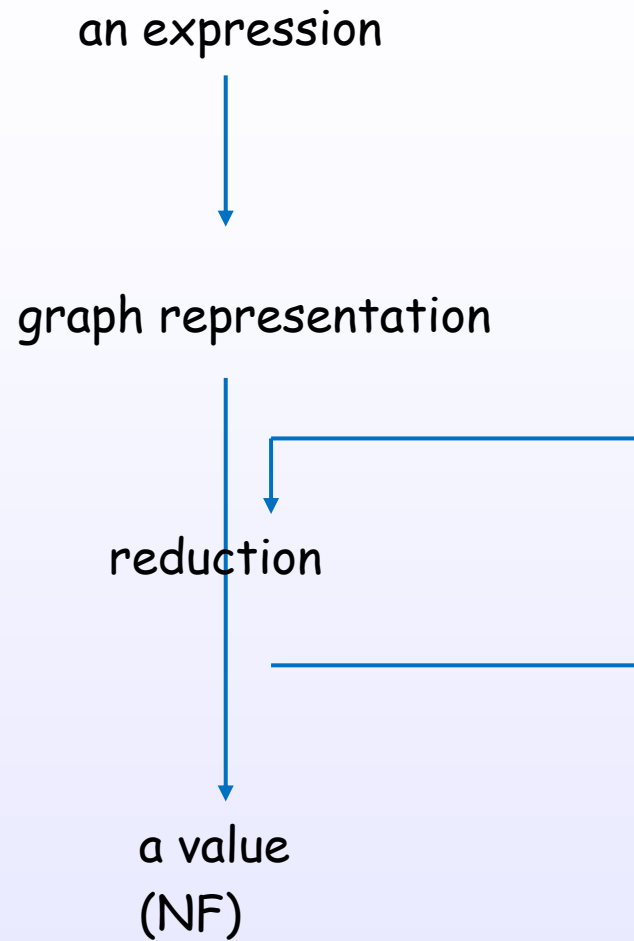
Tree

Graph

Shared Term

Lazy

evaluation, reduction



Expressions examples

STG heap objects

language

Just 5

implementation

heap object

Layer

Non-strictness

$$f \perp = \perp$$

Lazy evaluation

Graph reduction

STG machine

Layer

Haskell semantics

take 5 [1..10]

internal representation

graph

STG semantics

heap object

STG machine

Semantics

Bottom

domain

co-domain

defined

undefined

$$f \perp = \perp$$

[Bird, Chapter 2]

Strictness, Bottom

[Bird, Chapter 2]

References : [1]

References

References

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[http://www.scs.stanford.edu/11au-cs240h/notes/ghc-slides.html#\(11\)](http://www.scs.stanford.edu/11au-cs240h/notes/ghc-slides.html#(11))
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<https://ghc.haskell.org/trac/ghc/wiki/Commentary/Compiler/GeneratedCode>
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http://takenobu-hs.github.io/downloads/haskell_ghc_illustrated.pdf

Lazy,... ^{!!!}