

ASSIGNMENT #4: SORTING AND COMPLEXITY

DUE DATE WITH BRIGHTSPACE: FEBRUARY 1 AT 11:50 PM

THE LEARNING OUTCOMES

- to follow an algorithm from a description and to implement it in C++ (a sorting algorithm)
- to do a theoretical analysis of some code and to express it in big O notation
- to do a theoretical analysis of code written by the student and to express it in big O notation
- to write a recurrence relation with initial value if the function being analyzed is recursive
- to analyze code for best, average, worst cases

READINGS

- class notes
- §3.9 and §4.4 on Generating Random Numbers
- §7.10 Sorting Arrays
- chapter 17 on Recursion
 - not covered §17.3, §17.6, §17.7
- chapter 18 on Developing Efficient Algorithms
 - §18.1 - §18.4
 - not covered: Towers of Hanoi analysis §18.4.3
 - not covered: §18.5, (§18.6, §18.7), §18.8, §18.9, §18.10
 - yes Chapter Summary
- chapter 19 on Sorting §19.1 - §19.5
 - not covered §19.6, §19.7, §19.8

THEORETICAL ANALYSIS ON PAPER

Do a theoretical analysis of the following C++ function `multiply`.

Count the (exact) number of **multiplications** done in `multiply` (and its helper functions) as a function of n .

In addition to giving your answer as a function of n , express the answer in big O notation.

Show your workings.

```
int I(int i, int j, int n)
{
    return n * i + j;
}

int dotProduct(const int A[], const int B[], int i, int j, int n)
{
    int t = 0;
    for (int k = 0; k < n; k++) {
        t += A[I(i,k,n)] * B[I(k,j,n)] * A[I(i,k,n)] * B[I(k,j,n)];
    }
    return t;
}

void multiply( const int A[], const int B[], int C[], int n )
{
    for (int i = 0; i < n; i++ ) {
        for (int j = 0; j < n; j++ ) {
            C[I(i,j,n)] = dotProduct(A, B, i, j, n );
        }
    }
}
```

DESCRIPTION OF A SORTING ALGORITHM

Suppose we want to sort n integers of an array A in **ascending** order using the algorithm described below.

For the **following explanation given below**, assume that the number of integers is $n > 5$ and that there are no duplicates in the array. A is an array with n integers.

- We place the largest of the array A of n integers into $A[n-1]$ by comparing pairs of numbers:
 - we compare $A[0]$ with $A[1]$ and if $A[0] > A[1]$ then $A[0]$ and $A[1]$ are swapped
 - we compare $A[1]$ with $A[2]$ and if $A[1] > A[2]$, then $A[1]$ and $A[2]$ are swapped
 - we continue comparing in pairs
 - ..
 - we compare $A[n-2]$ with $A[n-1]$ and if $A[n-2] > A[n-1]$, then $A[n-2]$ and $A[n-1]$ are swapped
- We place the smallest of the first $n-1$ integers into $A[0]$ by comparing pairs of numbers “in the opposite direction” from the previous pass:
 - we compare $A[n-2]$ with $A[n-3]$ and if $A[n-2] < A[n-3]$, then $A[n-2]$ and $A[n-3]$ are swapped
 - we compare $A[n-3]$ with $A[n-4]$ and if $A[n-3] < A[n-4]$, then $A[n-3]$ and $A[n-4]$ are swapped
 - we continue comparing in pairs
 - ..
 - we compare $A[1]$ with $A[0]$ and if $A[1] < A[0]$, then $A[0]$ and $A[1]$ are swapped
- We place the largest of the integers from $A[1]$ up to and including $A[n-2]$ into $A[n-2]$ by comparing pairs of numbers “in the opposite direction” from the previous pass
 - we compare $A[1]$ with $A[2]$ and if $A[1] > A[2]$ then $A[1]$ and $A[2]$ are swapped
 - we compare $A[2]$ with $A[3]$ and if $A[2] > A[3]$, then $A[2]$ and $A[3]$ are swapped
 - we continue comparing in pairs
 - ..
 - we compare $A[n-3]$ with $A[n-2]$ and if $A[n-3] > A[n-2]$, then $A[n-3]$ and $A[n-2]$ are swapped
- We place the smallest of the integers from $A[1]$ up to and including $A[n-3]$ into $A[1]$ by comparing pairs of numbers “in the opposite direction” from the previous pass:
 - we compare $A[n-3]$ with $A[n-4]$ and if $A[n-3] < A[n-4]$, then $A[n-2]$ and $A[n-3]$ are swapped
 - we compare $A[n-4]$ with $A[n-5]$ and if $A[n-4] < A[n-5]$, then $A[n-4]$ and $A[n-5]$ are swapped
 - we continue comparing in pairs
 - ..
 - we compare $A[2]$ with $A[1]$ and if $A[2] < A[1]$, then $A[2]$ and $A[1]$ are swapped
- We continue comparing in pairs in both directions until the whole array is sorted in ascending order.

For your assignment and in general, do **not** assume that $n > 5$.

For your assignment and in general, do **not** assume that there are no duplicates. Duplicates are allowed in the array A .

I just made those two simplifications to explain the algorithm above.

IMPLEMENT

1. Implement the algorithm described in the previous section and use the function prototype

```
void sort(int A[], int n);
```

You may overload the function `sort` but you must still use the function prototype above to call the overloaded function.

Do not keep track of the swaps. No need to keep a Boolean variable that would help you optimize whenever any integers were not swapped in a pass.

2. Now using the `rand()` library function, generate an array of 19 random integers in the range
 $0 \leq \text{random number} \leq 999$

Print onto standard output the array before you `sort` it and after you `sort` it. Print 5 integers per line (use a named constant for 5) with the last line having at most 5 numbers. Format your output so that you can tell quickly whether the sorting is working.

3. For `sort`
 - a. Give the time complexity in big O notation: no need to show your workings or how you came up with the answer.
 - b. What is the worst case and the best case and the average case?
 - c. Is the running time complexity different for the different cases? Why yes or why not?
4. Repeat 1. and 2. but this time write a **recursive function** that implements the sorting algorithm. Use the function prototype

```
void sortR(int A[], int n);
```

Yes, you may overload the function.

Like in the class notes for selection sort and insertion sort, recursive functions *can* have loops

5. Given that n is the size of the array (i.e. n is the number of integers in the array), do a theoretical analysis for the running time of the function `sortR`.
In your **theoretical** analysis, count the number of comparisons that `sortR` does to sort n values. Count comparisons between the array values the way we have done in class.
Show your workings.
In addition give the complexity in big O notation.
Do not modify your code for the theoretical analysis. The assignment is different from the lab.
6. For `sortR`
 - a. What is the worst case and the best case and the average case?
 - b. Is the running time complexity different for the different cases? Why yes or why not?

SUBMIT WITH BRIGHTSPACE AS A SINGLE ZIP FILE: PART A) TO PART F) AS LISTED BELOW

- A) The source code of the complete sorting programs: one program with an iterative version `sort` function and one program with a recursive version `sortR` function (you may put both versions in a single program if you want). Document clearly your functions.
- B) Some “screen captures” or “print screens” that show that your iterative version `sort` works:
1. Show the input array of 19 random integers (at most 5 integers per line)
 2. Show the sorted output array of 19 random integers (at most 5 integers per line)
- C) Some “screen captures” or “print screens” that show that your recursive version `sortR` works:
1. Show the input array of 19 random integers (at most 5 integers per line)
 2. Show the sorted output array of 19 integers (at most 5 integers per line)

Submit as a pdf file (or as raster files from handwritten notes)

- D) The theoretical analysis of the recursive version `sortR` as described in 5 and the answers to the questions 6a. 6b.
- Give a recurrence relation with initial value.
 - Solve the recurrence relation with initial value and give your answer in closed-form as a function of n .
 - Write the resulting expression in big O notation
 - Show your work.
- E) As per 3.a give the complexity of `sort` in big O notation (no need to show your work) and give the answer to the questions 3b. and 3c.
- F) The theoretical analysis of multiply

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