

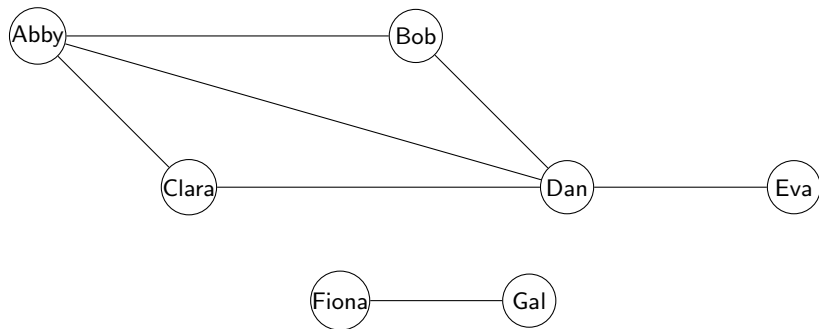
# Graph Theory and Map Coloring

Tal Berdichevsky and Corinne Mulvey

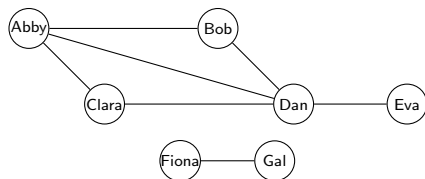
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May 18, 2019

# Modeling Facebook



# Modeling Facebook



## Definition

A *graph*  $G$  is a set of vertices  $V(G)$  and edges  $E(G)$ .

## Definition

The *order* and *size* of a graph  $G$  are the cardinality of its vertex set  $|V(G)|$  and edge set  $|E(G)|$ , respectively.

## Definition

For a graph  $G$  and vertex  $v \in V(G)$ , the *degree*  $\deg_G(v)$  of  $v$  is the number of edges incident with  $v$ .

# Fundamental Theorem of Graph Theory

## Theorem

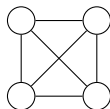
For any graph  $G$  of size  $m$ ,

$$\sum_{v \in V(G)} \deg_G(v) = 2m.$$

# Types of Graphs

## Definition

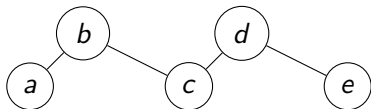
A *complete graph* is a graph  $G$  where every vertex  $v \in V(G)$  is adjacent to every other vertex  $u \in V(G)$ . A complete graph is denoted by  $K_n$  where  $n$  is the order of  $G$ .



# Types of Graphs

## Definition

A *path*  $P_n$  is a sequence of  $n$  vertices where every two consecutive vertices have an edge between them.

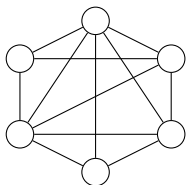


## Definition

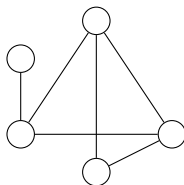
A graph  $H = (V', E')$  is a *subgraph* of the graph  $G = (V, E)$  if  $V' \subseteq V$ ,  $E' \subseteq E$ , and every edge  $e \in E'$  has its endpoints in  $V'$ .

# Types of Graphs

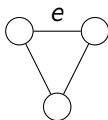
$H$ :



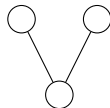
$H'$ :



$G$ :



$G - e$ :

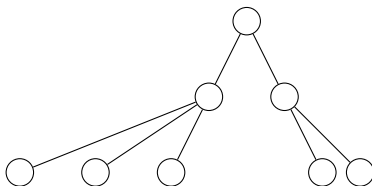


## Definition

A graph  $G$  is *connected* if for any pair of vertices in  $V(G)$  there is a path between them. A graph  $G$  is *disconnected* if there exists a pair of vertices in  $G$  which is not connected by a path.

## Definition

A *bridge* is an edge  $e$  of a graph  $G$  such the graph  $G - e$  is disconnected.



## Definition

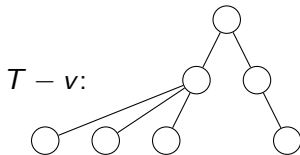
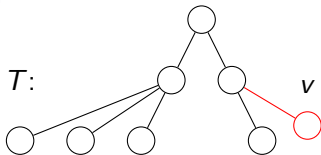
A *tree* is a connected graph where every edge is a bridge.



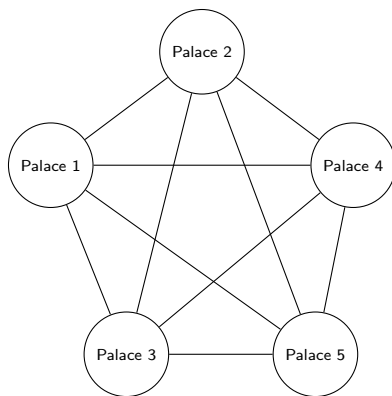
## Theorem

For every tree with order  $n$  and size  $m$ , we have  $m = n - 1$ .

## Proof



# Problem of the Five Palaces



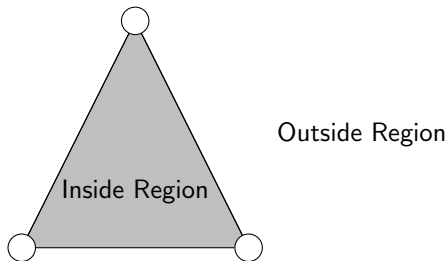
# Planarity

## Definition

A *planar graph* is a graph which can be drawn in a plane so that no edges intersect.

## Definition

A *region* in a plane graph  $G$  is a section of the plane enclosed by edges of  $G$ .



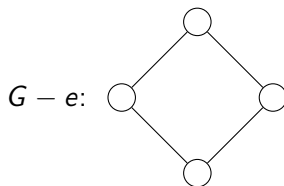
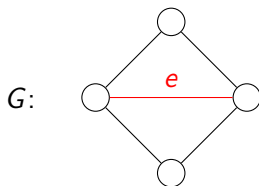
# Planarity

## The Euler Identity

If  $G$  is a connected plane graph of order  $n$ , size  $m$ , and  $r$  regions, then  $n - m + r = 2$ .

## Proof

Trees:  $k - (k - 1) + 1 = 2$



$$n - (m - 1) + (r - 1) = 2$$

$$n - m + r = 2$$

# Planarity

## Theorem

For any planar graph  $G$  of size  $m$  and order  $n \geq 3$ , we have

$$m \leq 3n - 6.$$

## Corollary

If  $G$  is a graph of order  $n \geq 3$  and size  $m$  where  $m > 3n - 6$ , then  $G$  is nonplanar.

# Planarity

## Corollary

Every planar graph contains a vertex of degree 5 or less.

## Theorem

The complete graph  $K_5$  is nonplanar.

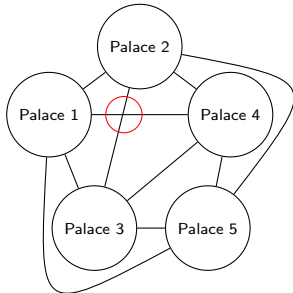
## Proof

The size of  $K_5$  is 10

The order of  $K_5$  is 5

$$m > 3n - 6$$

$$10 > 3(5) - 6 = 9$$



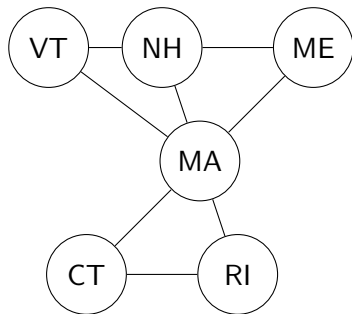
# Graph Colorings



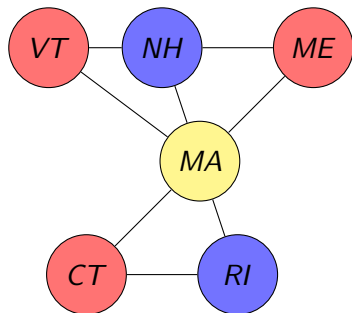
Figure: Blank Map of the New England States



# Graph Colorings



# Graph Colorings

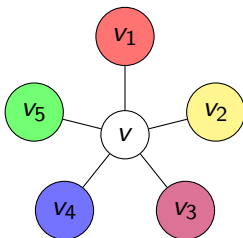


A 3-coloring of is the 3-chromatic graph  $G$  of the states  
( $\chi(G) = 3$ ).

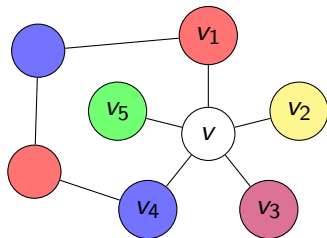
# Four and Five Color Theorems

## Theorem

Every planar graph is 5 - colorable



# Four and Five Color Theorems



# Four and Five Color Theorems

## Theorem

Every planar graph is 4 - colorable

- ▶ Proven in 1976 by Kenneth Appel and Wolfgang Haken
- ▶ 1,936 cases
- ▶ 1,200 hours of computer time