CC Lab Final Report

Lexor code:

```
import java.io.*;
       keywords.add("int");
       keywords.add("float");
       keywords.add("while");
       keywords.add("main");
       keywords.add("if");
       keywords.add("else");
       keywords.add("new");
       Pattern operators = Pattern.compile("^{(-,*,+,/,>,<,\&,|,=]}\subseteq");
           Scanner read = new Scanner(file);
```

```
flag = false;
```

```
if (operators.matcher(temp).find()) { // sub
String data;
public Node(String data, String type, int line) {
```

```
}
```

Lexor Output:

Lexem	Туре	Line num	Repetition
int	Keyword	1	2
main	Keyword	1	1
	Brackets	1	3
	Brackets	1	3
	Brackets	1	1
float	Keyword	2	1
test	Variable	2	1
	Operator	2	1
	Constant	2	1
printf	Variable		1
"Sum of "	String		1
"and"	String		1
"is "	String		1
	0perator		1
	Brackets		1
Process fini	shed with exit	code 0	

Parser code:

```
for rule in rules:
closureSet = []
```

```
in rule not in tempClosureSet:
                        tempClosureSet.append(in rule)
              closureSet.append(rule)
return closureSet
              generateStatesFor.append(dotPointsHere)
              shiftedRule[1][indexOfDot] = \
              newState.append(shiftedRule)
    indexDot = rule[1].index('.')
# check that rule is not "Handle"
if rule[1][-1] != '.':
```

```
closureRes = \
                    addClosureRules.append(rule)
def generateStates(statesDict):
               compute GOTO(key)
```

```
return fres
def follow(nt):
```

```
res = follow(curNT)
```

```
tempRow.append('')
                  + f"S{stateMap[entry]} "
tempRule = copy.deepcopy(rule)
```

```
# *** MAIN *** - Driver Code
```

```
nonterm_userdef = ['E', 'T', 'F']
term_userdef = ['id', '+', '*', '(', ')']
print("\nGrammar after Augmentation: \n")
separatedRulesList = \
printResult(separatedRulesList)
start_symbol = separatedRulesList[0][0]
print("\nCalculated closure: IO\n")
I0 = findClosure(0, start symbol)
printResult(I0)
statesDict = {}
stateMap = {}
statesDict[0] = I0
stateCount = 0
generateStates(statesDict)
print("Result of GOTO computation:\n")
printAllGOTO(stateMap)
diction = {}
```

Parser Output:

```
Original grammar input:

E -> E + T | T
T -> T * F | F
F -> (E) | id

Grammar after Augmentation:

E' -> . E
E -> . E + T
E -> . T
T -> . T * F
T -> . (E)
F -> . id

Calculated closure: I0

E' -> . E
E -> . E + T
E -> . T
T -> . T * F
F -> . (E)
F -> . id
```

States Generated: State = I0 E' -> . E E -> . E + T E -> . T T -> . T * F T -> . F F -> . (E) F -> . id State = I1 E' -> E . E -> E . + T State = I2 E -> T . T -> T . * F State = I3

```
State = I4
F -> ( . E )
E -> . E + T
E -> . T
T -> . T * F
T -> . F
F -> . ( E )
F -> . id

State = I5
F -> id .

State = I6
E -> E + . T
T -> . F
F -> . ( E )
F -> . id

State = I7
T -> . T * F
F -> . ( E )
F -> . id
```

```
State = I9
State = I10
State = I11
Result of GOTO computation:
GOTO (IO, E) = I1
GOTO (IO, T) = I2
GOTO (IO, F) = I3
GOTO (I1 , + ) = I6
GOTO (I2, *) = I7
GOTO (I4, E) = I8
GOTO ( I4 , T ) = I2
GOTO (I4, F) = I3
GOTO (I4, () = I4
```

```
Result of GOTO computation:
GOTO ( IO , E ) = I1
GOTO (IO, T) = I2
GOTO (IO, F) = I3
GOTO (IO, id) = I5
GOTO (I1 , +) = I6
GOTO (I2, *) = I7
GOTO (I4, E) = I8
GOTO (I4, T) = I2
GOTO (I4, F) = I3
GOTO (I4, () = I4
GOTO (I4, id) = I5
GOTO (I6, T) = I9
GOTO (I6, F) = I3
GOTO (I6, id) = I5
GOTO (I7, F) = I10
GOTO (I7, () = I4
GOTO (I7, id) = I5
GOTO ( I8 , ) ) = I11
GOTO ( I9 , * ) = I7
```

SLR(1)	SLR(1) parsing table:										
	id			()	\$	Е	Т	F		
10	S 5			S 4			1	2	3		
I1		S6				Accept					
12		R2	S7		R2	R2					
13		R4	R4		R4	R4					
I4	S5			S4			8	2	3		
15		R6	R6		R6	R6					
I6	S 5			S4				9	3		
17	S5			\$4					10		
18		S6			S11						
I9		R1	S7		R1	R1					
I10		R3	R3		R3	R3					
I11		R5	R5		R5	R5					
Proces	ss finish	ed with	exit cod	e 0							

LL1 Parser code:

```
def removeLeftRecursion(rulesDiction):
    # for rule: A->Aa|b
    # result: A->bA',A'->aA'|#

# 'store' has new rules to be added
store = {}
    # traverse over rules
for lhs in rulesDiction:
    # alphaRules stores subrules with left-recursion
    # betaRules stores subrules without left-recursion
    alphaRules = []
    betaRules = []
    # get rhs for current lhs
    allrhs = rulesDiction[lhs]
    for subrhs in allrhs:
        if subrhs[0] == lhs:
            alphaRules.append(subrhs[1:])
        else:
            betaRules.append(subrhs)
# alpha and beta containing subrules are separated
# now form two new rules
if len(alphaRules) != 0:
    # to generate new unique symbol
    # add ' till unique not generated
    lhs_ = lhs + "'"
    while (lhs_ in rulesDiction.keys()) \
            or (lhs_ in store.keys()):
            lhs_ += "'"
# make beta rule
for b in range(0, len(betaRules)):
            betaRules[b].append(lhs_)
```

```
alphaRules[a].append(lhs )
        alphaRules.append(['#'])
def LeftFactoring(rulesDiction):
           temp[subrhs[0]].append(subrhs)
        if len(allStartingWithTermKey) > 1:
              ex rules.append(g[1:])
```

```
# if result is not already returned
# - control reaches here
# lastly if eplison still persists
# - keep it in result of first
```

```
ansNew = follow(curNT)
```

```
if res is not None:
def computeAllFirsts():
  diction = LeftFactoring(diction)
```

```
def computeAllFollows():
def createParseTable():
      row = []
for y in terminals:
```

```
mat.append(row)
               firstFollow.append(u)
         res = firstFollow
      ttemp.append(res)
```

```
entry = parsing table[x][y]
buffer = buffer[:-1]
```

```
rules=["S -> A k O",

"A -> A d | a B | a C",

"C -> c",

"B -> b B C | r"]
nonterm_userdef=['A','B','C']
term_userdef=['k','O','d','a','c','b','r']
firsts = {}
follows = {}
computeAllFirsts()
start symbol = list(diction.keys())[0]
computeAllFollows()
```

LL1 Parser Output:

```
Rules:
S->[['A', 'k', '0']]
A->[['A', 'd'], ['a', 'B'], ['a', 'C']]
C->[['c']]
B->[['b', 'B', 'C'], ['r']]
After elimination of left recursion:
S->[['A', 'k', '0']]
A->[['a', 'B', "A'"], ['a', 'C', "A'"]]
C->[['c']]
B->[['b', 'B', 'C'], ['r']]
A'->[['d', "A'"], ['#']]
After left factoring:
S->[['A', 'k', '0']]
A->[['a', "A''"]]
A''->[['B', "A'"], ['C', "A'"]]
C->[['c']]
B->[['b', 'B', 'C'], ['r']]
A'->[['d', "A'"], ['#']]
```

```
Calculated firsts:
first(S) => {'a'}
first(A) => {'a'}
first(A'') => {'r', 'b', 'c'}
first(C) => {'c'}
first(B) => {'r', 'b'}
first(A') => {'#', 'd'}

Calculated follows:
follow(S) => {'$'}
follow(A) => {'k'}
follow(A'') => {'k'}
follow(C) => {'k', 'c', 'd'}
follow(B) => {'k', 'c', 'd'}
follow(A'') => {'k'}
```

```
Non-T FIRST FOLLOW
S {'a'} {'$'}
A {'a'} {'k'}
A'' {'r', 'b', 'c'} {'k', 'c', 'd'}
B {'r', 'b'} {'k', 'c', 'd'}
A' {'#', 'd'} {'k', 'c', 'd'}
Generated parsing table:

k 0 d a c b r $
S S->A k 0
A A A A->A A''

A'' A'-># A'-># A'->d A'
```

```
Validate String => a r k 0
             Buffer
                                   Stack
                                                       Action
                                     S $
                                                T[S][a] = S->A k 0
          $ 0 kra
                                                T[A][a] = A->a A''
          $ 0 k r a
                                 A k 0 $
          $ 0 k r a
                             a A'' k 0 $
                                                    Matched:a
                                             T[A''][r] = A''->B A'
            $ 0 k r
                               A'' k 0 $
            $ 0 k r
                                                    T[B][r] = B->r
            $ 0 k r
                              r A' k 0 $
                                                    Matched:r
              $ 0 k
                               A' k 0 $
                                                  T[A'][k] = A'->#
              $ 0 k
                                   k 0 $
                                                    Matched:k
                                     0 $
                $ 0
                                                    Matched:0
                  $
                                       $
                                                        Valid
Valid String!
Process finished with exit code 0
```

Semantic Analyser Code:

```
import pandas as pd
import copy
try:
    a=pd.read_csv("input.csv")
    print("\a One thing in this program is that it takes an input of csv
file.\n\aThe formate of the csv file is in following manner: ")
    print(a)
    print("\a The output of this program is based on above csv file.\n\aYou
can take the different input csv file for diffrent required output.")
    print("\a One thing is to be noticed that in the csv file, there are
two columns one is left and other is right for left and right values
respectively.")
    print("\a There is only one left side variable for each equation and it
may be possible that more than one varible in right side.")
    print("\a The operators and operands which are used in right side must
be space separated from each other.")
    print("\a This program is case sensitive, this means that 'd*10' and
'10*d' are treated as different equation.\nIf you want to solve this
problem then you can use the 'CODE OPTIMIZATION TECHNIQUE'.")
    print('\t')
    c=a.shape# It will gives an tuple of numbers of rows and columns
    #print(c)
    l=[]
    o=list("+-*/")#If you want to add more operator youn can use that as
well
    o1=[]
    r=[]
    for i in range(c[0]):# Here c[0] is Oth element of tuple c, which is
a.shape (c=a.shape)
```

```
l=l+[a['left'][i]]
d=a['right'][i]
     li.remove(x)
     akm=akm+[r[li.index(l[ak])]]
        akm.clear()
```

```
akm.clear()
ACounter+=2
akm.pop(len(akm)-2)
akm.clear()
```

Code Optimization and generation:

```
import pandas as pd
import copy
try:
    a=pd.read_csv("input.csv")
    print("\a One thing in this program is that it takes an input of csv
file.\n\aThe formate of the csv file is in following manner: ")
    print(a)
    print("\a The output of this program is based on above csv file.\n\aYou
can take the different input csv file for diffrent required output.")
    print("\a One thing is to be noticed that in the csv file, there are
two columns one is left and other is right for left and right values
respectively.")
    print("\a There is only one left side variable for each equation and it
may be possible that more than one varible in right side.")
    print("\a The operators and operands which are used in right side must
be space separated from each other.")
    print("\a This program is case sensitive, this means that 'd*10' and
'10*d' are treated as different equation.\nIf you want to solve this
problem then you can use the 'CODE OPTIMIZATION TECHNIQUE'.")
    print('\t')
    c=a.shape# It will gives an tuple of numbers of rows and columns
    #print(c)
    l=[]
    o=list("+-*/")#If you want to add more operator youn can use that as
well
    o1=[]
    r=[]
    for i in range(c[0]):# Here c[0] is 0th element of tuple c, which is
a.shape (c=a.shape)
    l=!+[a('left'][i]]
    d=a('right')[i]
    x=d.split()
```

```
li.remove(x)
   akm.clear()
akm.pop(len(akm)-2)
```

```
akm.pop(len(akm)-2)
       akm.clear()
akm.pop(len(akm)-2)
ACounter+=2
akm.pop(len(akm)-2)
ACounter+=1
```

```
print("\t")
    elif((l[ak] not in o1)or (l[ak] not in string.ascii_lowercase)):
        print("\f Error!\n\f Please enter valid syntax for three
address code.\n\f Check your csv file...")
        print(f"\f Error description...\nError in line number {z} and
place number {ak}.")
        print(f"\f Error element is {a['right'][z]}.")
        break
except (FileNotFoundError):
    print("Please check you input file. It may possible that file doesn't
exist.")
    print("Also check the file name that is given in input section at the
starting place.")
except(ArithmeticError):
    print("An arithmetic error is caused due to which program is not
proceed futher.Please check for the solution.")
except(IndexError):
    print("List index out of range.")
except:
    print("An exceptions occurred.")
```

Code Optimization and generation Output:

```
MOV b , R0

MOV c , R7

MOV d , R5

ADD R7 , R5

STOR R5 , R0

MOV f , R8

ADD R5 , R8

STOR R8 , R0

MOV b , R0

MOV b , R0

ADD R0 , R0

STOR R0 , R8

MOV r , R12

MOV f , R8

STOR R8 , R12

MOV a , R11

MOV 10 , R10

STOR R10 , R11
```

MOV s , R4

MOV a , R11

MOV 10 , R10

ADD R11 , R10

STOR R10 , R4

MOV d , R5

MOV s , R4

MOV 10 , R10

SUB R4 , R10

STOR R10 , R5

MOV g , R1

MOV 10 , R10 MOV d , R5 MUL R10 , R5 STOR R5 , R1

MOV j , R2 MOV d , R5 MOV 10 , R10 MUL R5 , R10 STOR R10 , R2