File Systems

2024 Semester 2 COMPSCI 340: Operating Systems Talia Xu

Lecture 1 1.0.0

To read a file

The high level idea

- Perform a path lookup
- Traverse the path until we get the file's inode
- The root inode is typically 2 in a Unix file system
- Once we have the inode of the file, we can start to read its data

Example: We want to read an 8 KB file /foo/bar

To read a file

Example: We want to read an 8 KB file /foo/bar

First we need to open the file.

Operation	data bitmap	inode Bitmap	foo's inode		foo data	bar data[0]	bar data[1]
open(bar)							

To read a file

Example: We want to read an 8 KB file /foo/bar

Now we have the file open.

Each block is 4 KB, we must perform read twice.

Operation	data bitmap	inode Bitmap	foo's inode		foo data	bar data[0]	bar data[1]
read()							
read()							

To write a file

The high level idea

- Similar to read need to traverse the path to the file
- Unlike read may need to allocate a file/block
- Block allocation updates multiple structures

Example: We want to create and write to an 8 KB file /foo/bar

To write a file

Example: We want to create and write to an 8 KB file /foo/bar

Operation	data bitmap	inode Bitmap	root inode	foo's inode	bar's inode	root data	foo data	bar data[0]	bar data[1]
create(bar)									
create(bal)									

To write a file

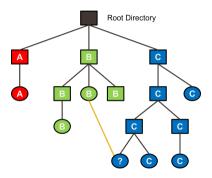
Example: We want to create and write to an 8 KB file /foo/bar

File now exists but it is empty.

We need to perform 2 writes of 4 KB each.

Operation	data bitmap	inode Bitmap	root inode	foo's inode	bar's inode	root data	foo data	bar data[0]	bar data[1]
write()									
write()									

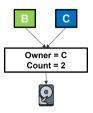
Sharing a file



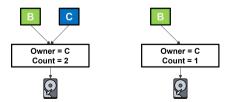
What are the different ways User B and User C can share a file?

Sharing a file

Hard Links are pointers to inodes.

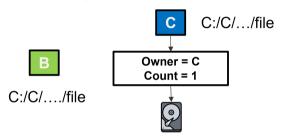


Deleting a file only removes one hard link. The file can still be accessed



Sharing a file

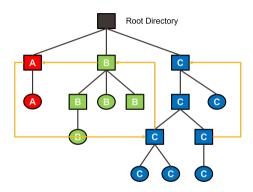
Soft links are paths to another file.



When resolving the file, the file system is redirected somewhere else, so:

- Soft link targets do not need to exist
- Soft link targets can be deleted without notice of the soft link
- Unresolvable soft links lead to an exception

Cycles in Directory Graph



Why are cyclic loops a problem? How do we solve this?

Filesystem Example Problem

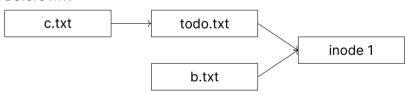
```
touch todo.txt
ln todo.txt b.txt
ln -s todo.txt c.txt
mv todo.txt d.txt
rm b.txt
```

How does the FS look like before and after the mv and rm commands?

Filesystem Example Solution (1)

```
touch todo.txt
In todo.txt b.txt
In -s todo.txt c.txt
mv todo.txt d.txt
rm b.txt
```

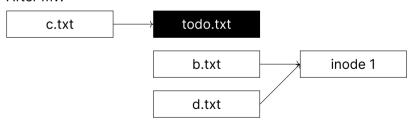
Before my:



Filesystem Example Solution (2)

touch todo.txt
In todo.txt b.txt
In -s todo.txt c.txt
mv todo.txt d.txt
rm b.txt

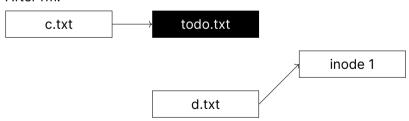
After mv:



Filesystem Example Solution (3)

touch todo.txt
In todo.txt b.txt
In -s todo.txt c.txt
mv todo.txt d.txt
rm b.txt

After rm:



What Data is Stored in an inode?

- a Filename
- b Containing Directory name
- c File Size
- d File type
- e # of soft links to file
- f location of soft links
- g # of hard links to file
- h location of hard links
- i access rights
- j timestamps
- k file contents
- I ordered list of data blocks

What Data is Stored in an inode? Solution

- a Filename No. Names are stored in directories
- b Containing Directory name No. File can be in multiple dirs
- c File Size Yes
- d File type Yes
- e # of soft links to file No (they are unknown)
- f location of soft links No (they are unknown)
- g # of hard links to file Yes (to know when to erase the file, check stat)
- h location of hard links No (they are unknown to the inode)
- i access rights Yes
- j timestamps Yes
- k file contents Sometimes
- I ordered list of data blocks Yes, by definition

Symbolic and Hard Links - More Practise

What is the output?