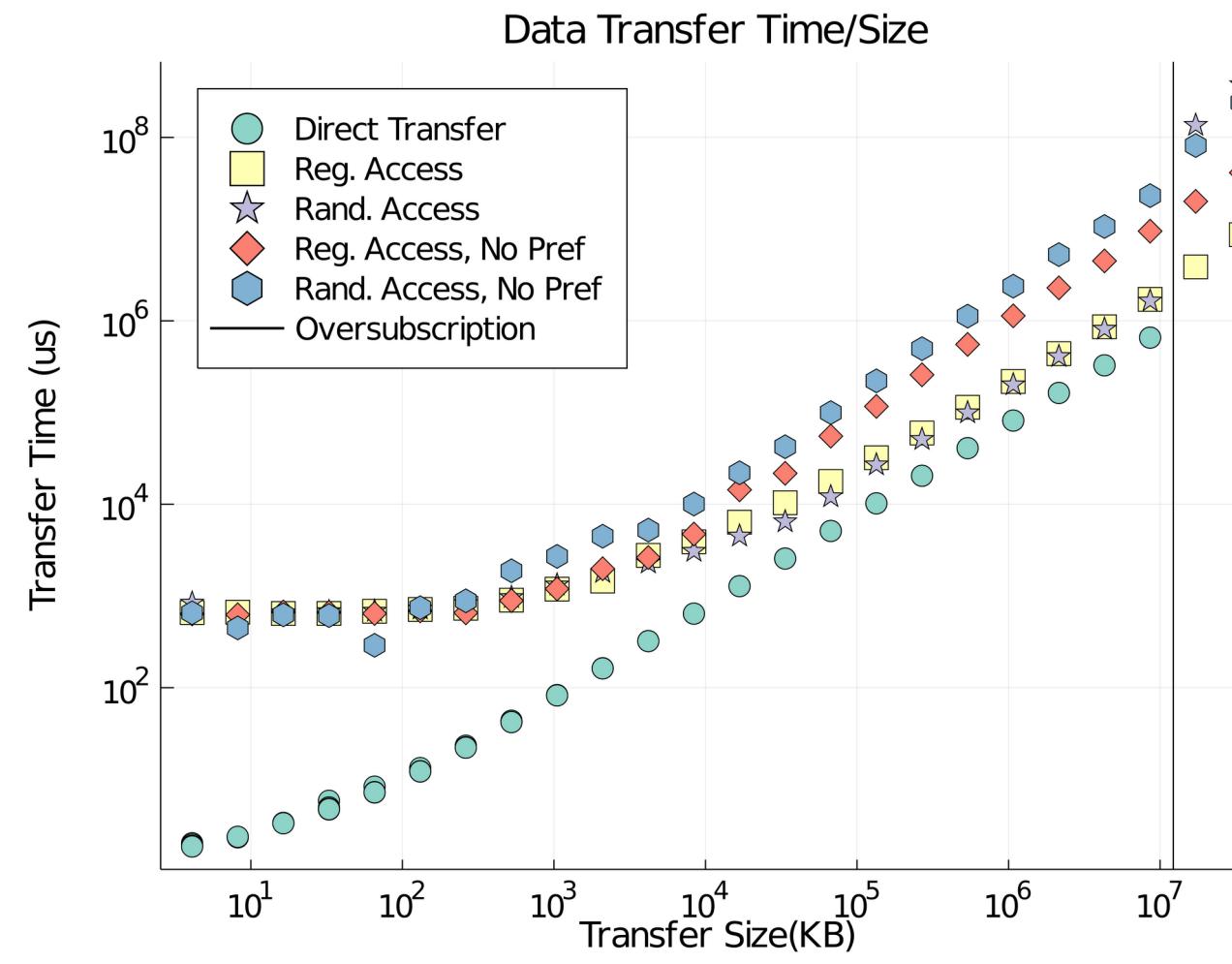


Holistic Performance Analysis and Optimization of Unified Virtual Memory

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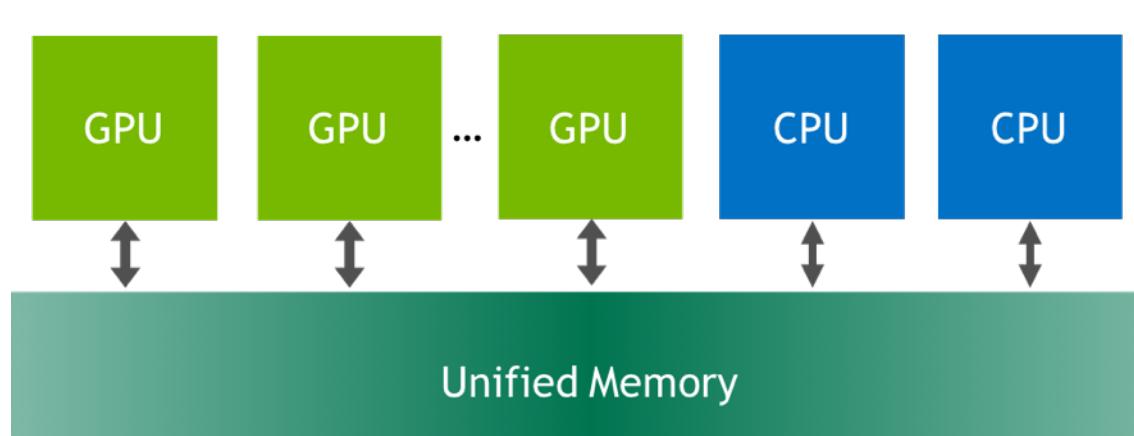
Problem and Motivation

- Problem:** UVM improves programmer productivity at the cost of performance. Can we have both?
- Big Picture:** UVM is increasingly popular - used in HPC applications and frameworks such as Kokkos, Raja, and Trilinos.
- SotA:** UVM performance loss not well studied at **systems software, system architecture** level.
- Our Approach:** Holistic Analysis, then Optimization
 - Study full system – systems software, workload generation, applications.
 - Include advanced features – prefetching and oversubscription - for real-world use-cases.
 - Apply optimizations at *method* level, applicable to other implementations besides NVIDIA UVM.



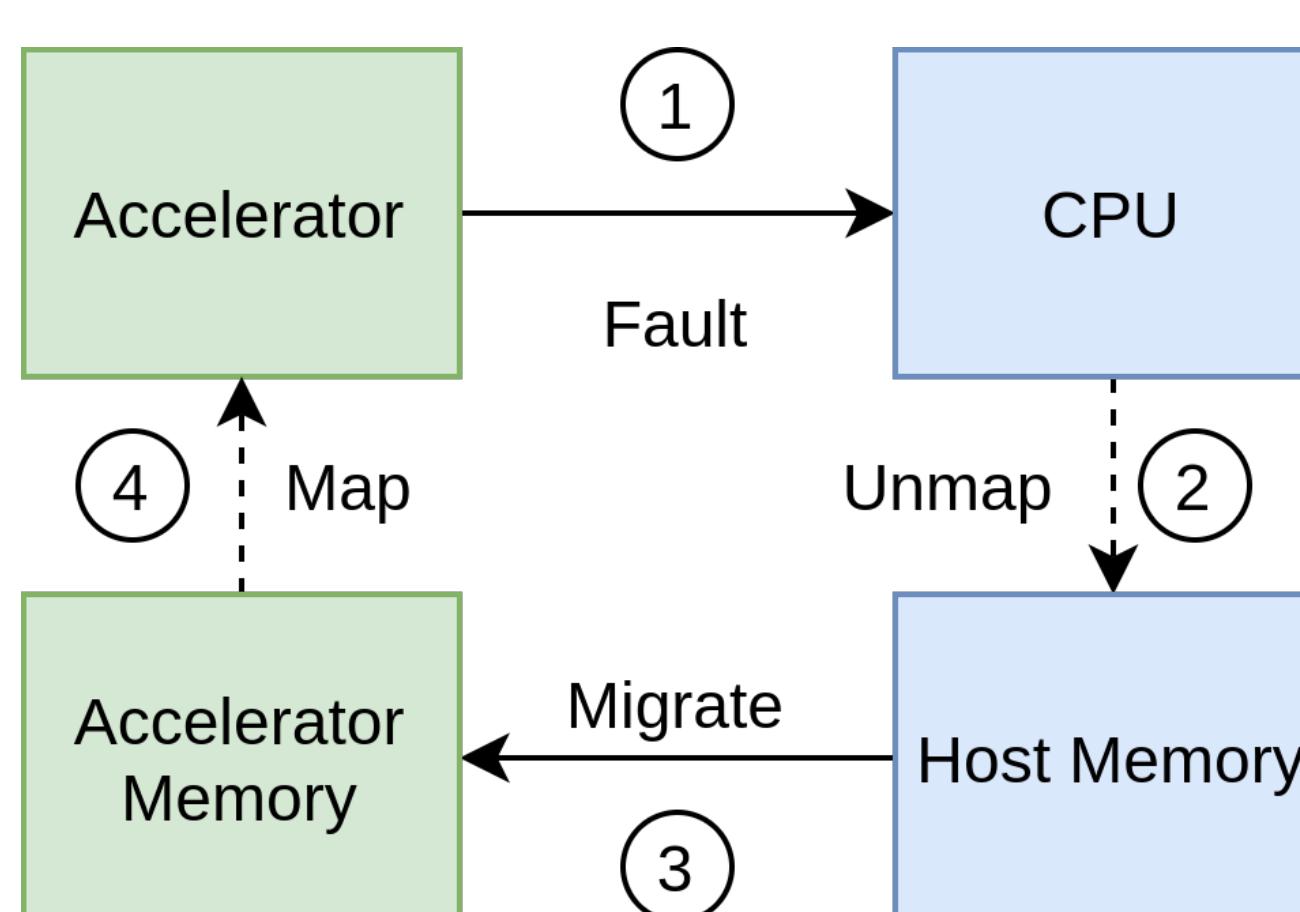
Background

- UVM is an example of *heterogeneous shared memory*.
- Heterogeneous shared memory systems fundamentally use *on-demand page migration* as the backing technology for physical memory placement.
- Same programming view of memory on all devices.
- UVM uses software *density prefetching* because on-demand migration is prohibitively slow.
- Enables device memory *oversubscription*, expanding capacity to host memory – similar to Swap storage.



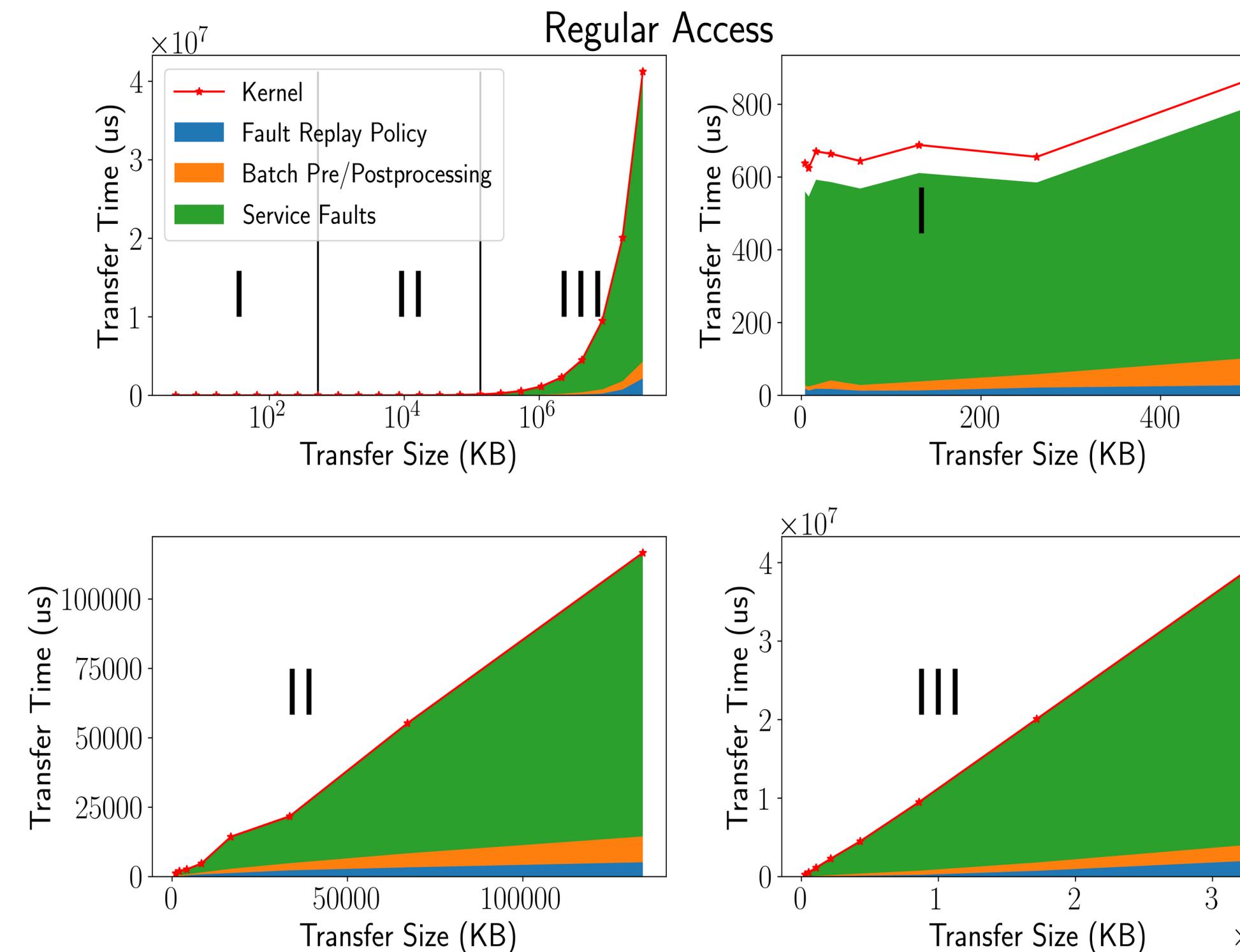
<https://developer-blogs.nvidia.com/wp-content/uploads/2017/11/Unified-Memory-MultiGPU-FI.png>

- All heterogeneous shared memory systems are fundamentally 4-steps: **fault** → **unmap** → **migrate** → **map**



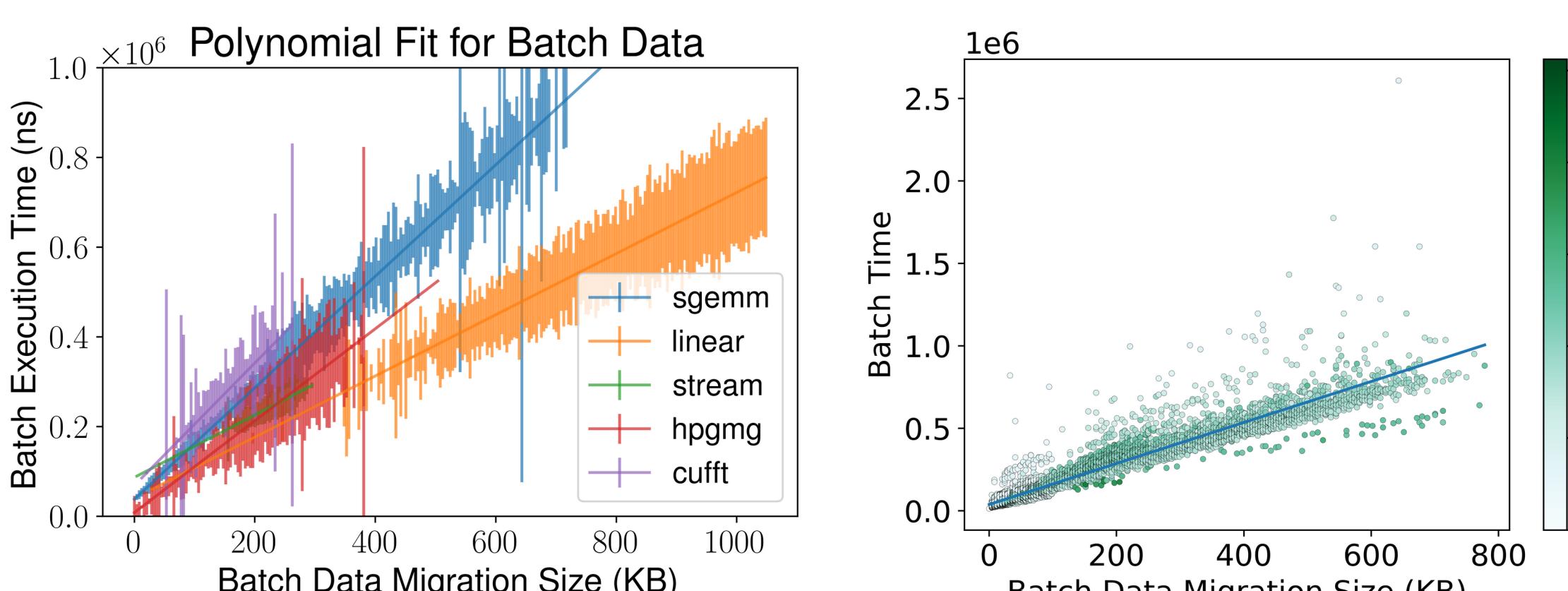
High-Level Performance

- The majority of cost is the “service faults” category.
- “Service faults” is primarily the **unmap** → **migrate** → **map** process – we focus further study here.



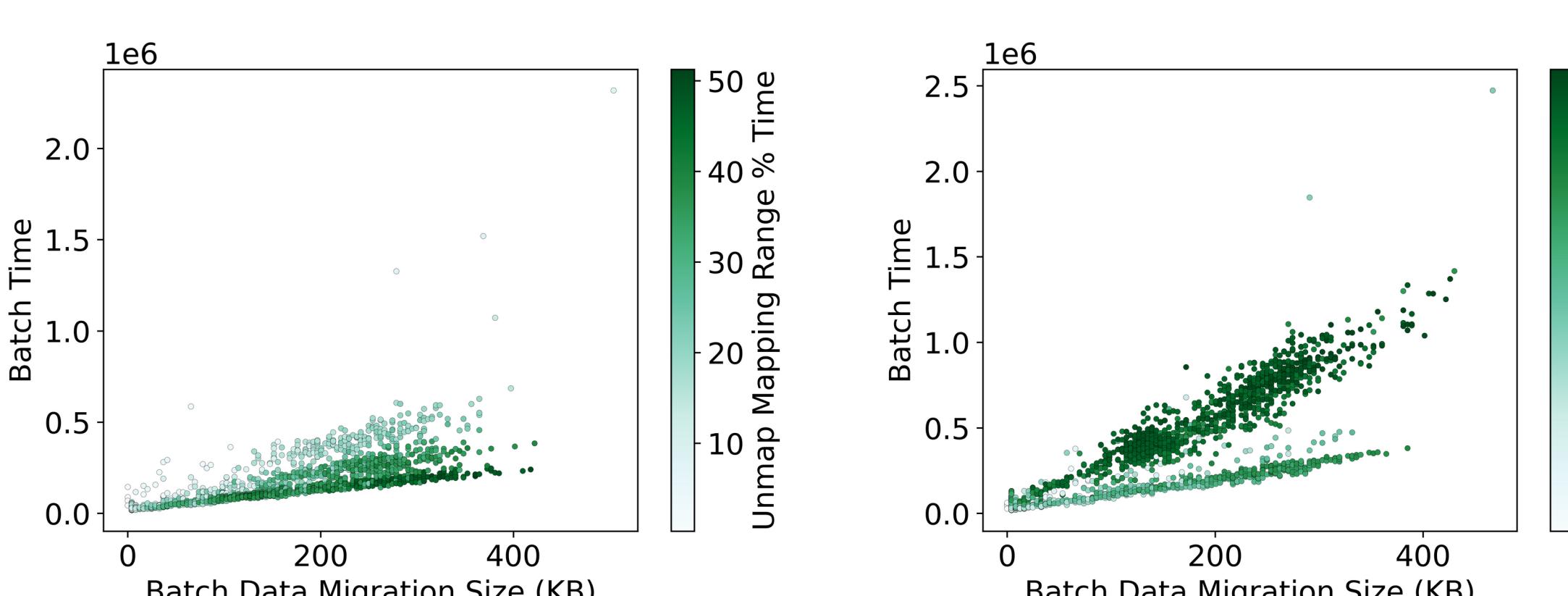
Impact of Data Migration

- Key Insight:** Most of UVM cost is software overhead.
- The amount of data migrated is highly correlated with the time required to process, but (left):
 - Applications unexpectedly show different trends for migration costs.
 - Fault batches within the same application have varied costs.
 - Data transfer is not the majority of cost (right).
 - Interconnect hardware bandwidth is not the main problem.**



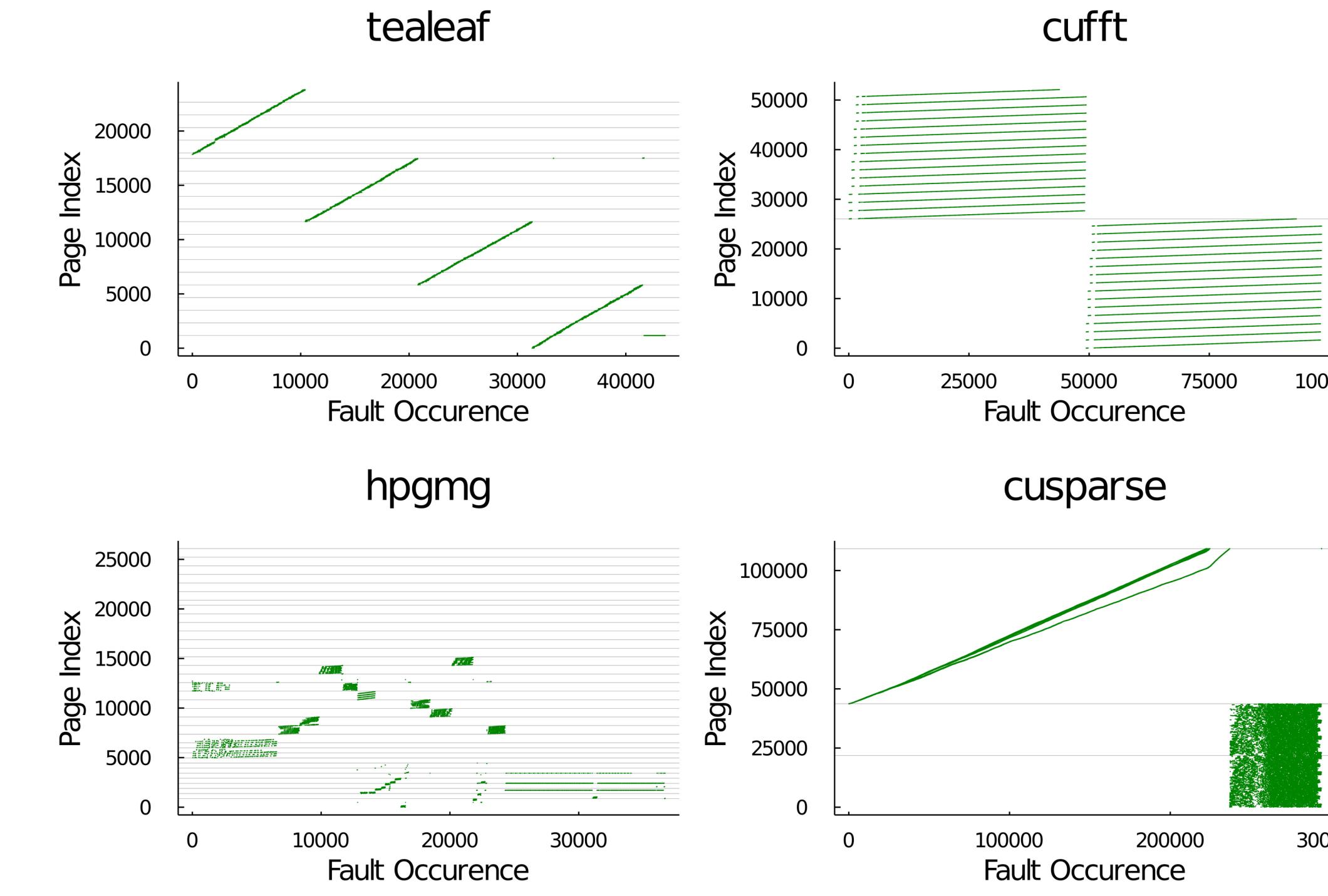
Page Unmapping

- Key Insight:** A primary source of overhead is CPU page unmapping.
- This overhead can account for up to 50% of batch execution time (left), and over 80% in certain scenarios with CPU parallelization (right).
- Tracing indicates that this overhead is likely attributed to TLB shootdowns.
 - Optimizing TLB shootdowns for migration has broad impacts.

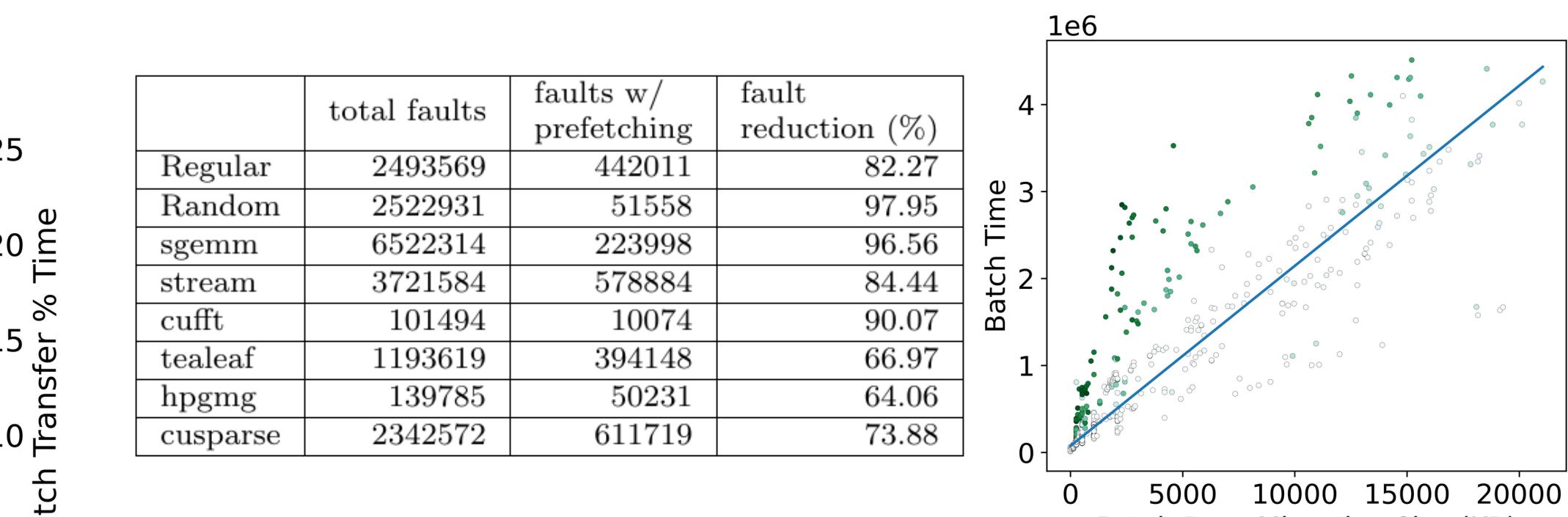


Prefetching

- Page-level access pattern of several applications show the workload of this prefetcher:

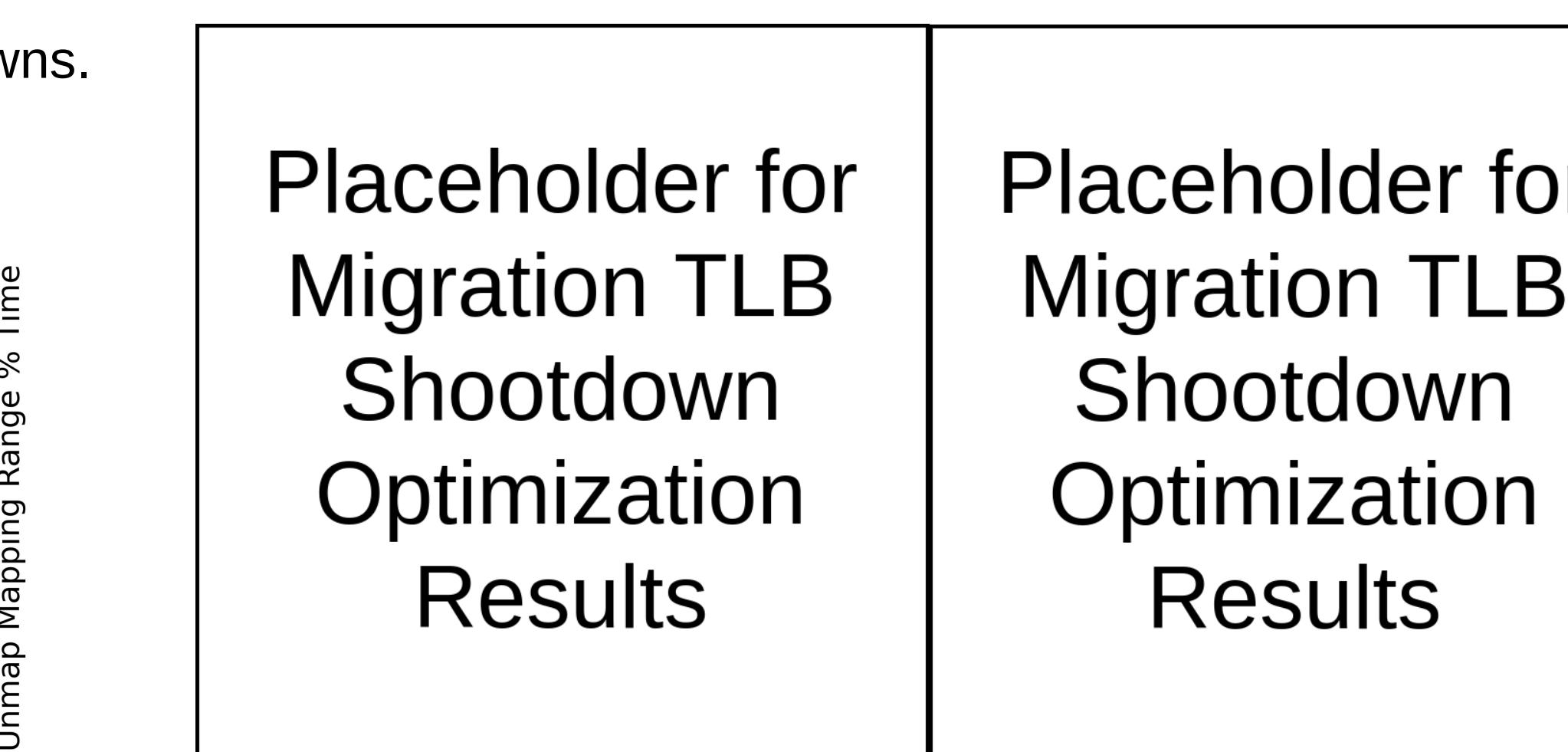


- Prefetching reduces the total page faults by up to 96% (left).
- Key Insight:** Prefetching is very effective but cannot remove overhead for lower-level system functionality.
- Costs like **page unmapping** and DMA setup (right) **cannot be avoided through the existing prefetcher** and are a larger percentage of overhead with prefetching.
- Prefetching notably decreases the number of *batches* seen, reducing the baseline overhead despite bulk costs.
- The trade-offs for prefetching increase when oversubscription is enabled, potentially increasing the amount of evicted data.



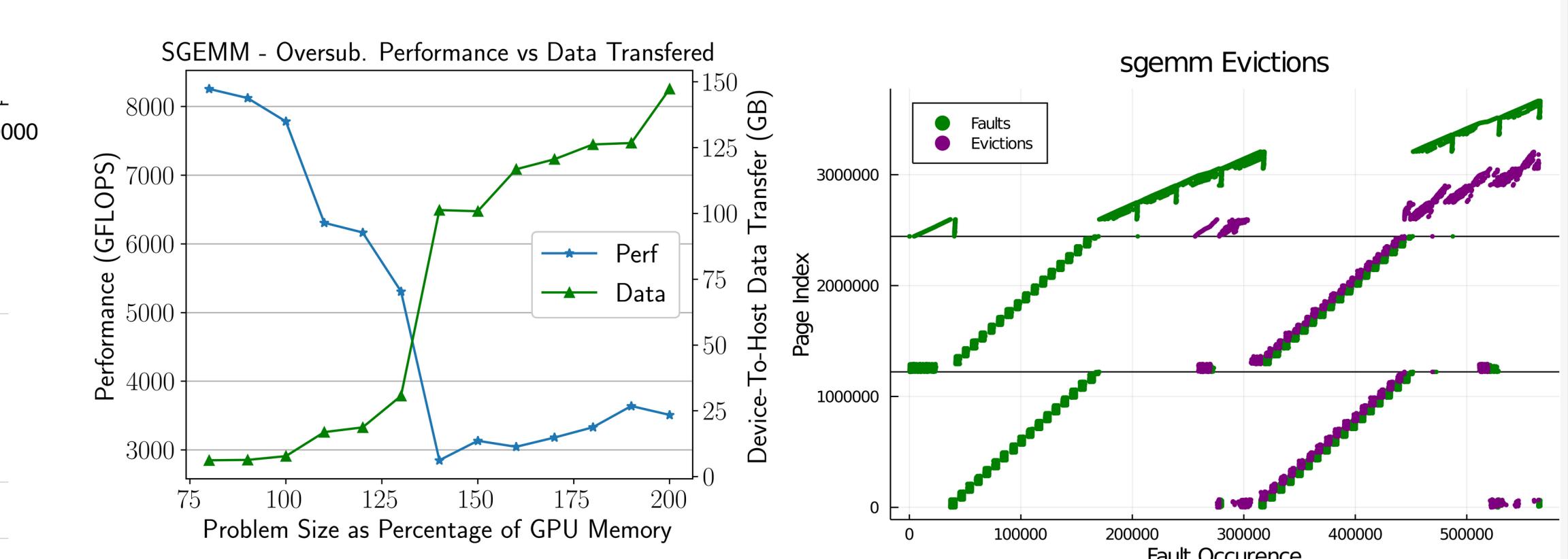
Optimization

- Optimization focusing on page unmapping and system-level improvements is in-progress and to be added in September.
- Optimizations focus on TLB shootdowns specifically for data migration, impacting heterogeneous shared memory systems but also other migration scenarios, such as swap storage.

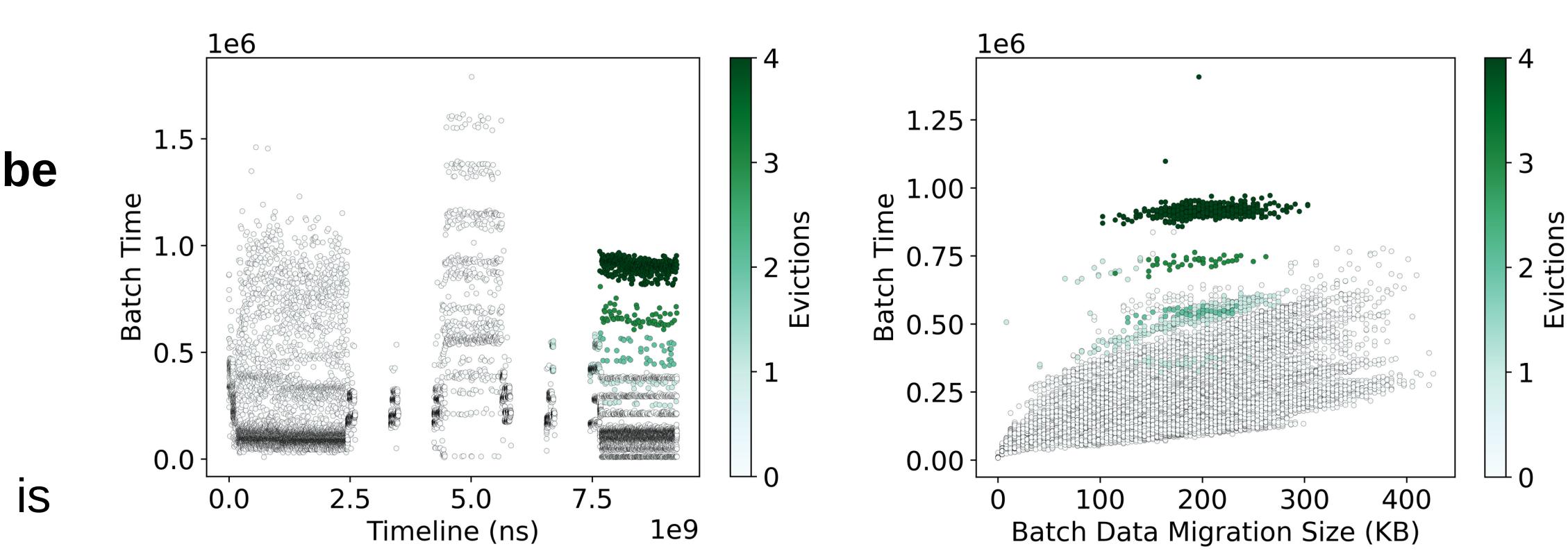


Oversubscription

- Eviction uses LRU from host driver perspective.
 - Page fault raises virtual address block (2MB) in LRU list.
 - Highly-used regions can be evicted early.
 - Fully-resident regions cannot move up in LRU list.
 - Evicted pages must be paged back (left), and quantity of evicted data is highly correlated with performance (right).



- Eviction occurs once memory is fully allocated. (left)
- Key Insight:** Eviction has flat cost based on number of virtual address blocks evicted (right).



Conclusions/Discussion

- Data transfer is not the primary cost of UVM; higher bandwidth can improve performance but is not the main issue.
- System costs like building DMA connections and TLB shootdowns are a primary source of overhead.
- Prefetching is effective but cannot eliminate these overheads.
- Oversubscription/eviction has flat cost and does not antagonize system costs, but is subject to them.
- Provided unmapping optimizations improve baseline application performance without source code changes.
- Results applicable to other related systems, such as Heterogeneous Memory Management (HMM).

Publications

- Tyler Allen, Rong Ge, “In-Depth Analyses of Unified Virtual Memory System for GPU Accelerated Computing”, Supercomputing 2021
- Tyler Allen, Rong Ge, “Demystifying GPU UVM Cost with Deep Runtime and Workload Analysis,” 2021 IEEE International Parallel and Distributed Processing Symposium (IPDPS), 2021, pp. 141-150, 10.1109/IPDPS49936.2021.00023

