

ECE 0402 - Pattern Recognition

Lecture 8 on 2/9/2022

Review:

- Challenge: Number of hypothesis is $|\mathcal{H}|$ potentially infinite
- Better: Narrow the scope to the finite training set in order to replace easily infinite m . Dichotomies allow us that.
- $h : \{x_1, \dots, x_n\} \mapsto \{-1, 1\} \implies 2^n$ different way of labeling, max! so dichotomy is the way of labeling THAT particular data set
- Hence, in general, $|\mathcal{H}| > |\mathcal{H}(x_1, \dots, x_n)|$. In English number of hypothesis \uparrow number of dichotomies
- So maybe a dichotomies are a better measure of “richness” of the set.
- And then we introduced the idea of “growth function that gets rid of the dependence of dichotomy to a particular of our training set x_1, \dots, x_n .”
- Growth function: $m_{\mathcal{H}}(n) = \max_{x_1, \dots, x_n \in \mathcal{X}} |\mathcal{H}(x_1, \dots, x_n)|$.

We went through some examples of Growth functions:

- Positive rays: $m_{\mathcal{H}}(n) = n + 1$
- Positive intervals: $m_{\mathcal{H}}(n) = \frac{1}{2}n^2 + \frac{1}{2}n + 1$
- Convex sets: $m_{\mathcal{H}}(n) = 2^n$
- Linear classifiers in \mathbb{R}^2 :

$$m_{\mathcal{H}}(1) = 2$$

$$m_{\mathcal{H}}(2) = 4$$

$$m_{\mathcal{H}}(3) = 8$$

$$m_{\mathcal{H}}(4) = 14$$

$$m_{\mathcal{H}}(n) = ?$$

So in the previous lecture we left with linear classifiers example, and we didn't actually calculate the growth function, we just worked out for first four values of n ...

Recall

$$\mathbb{P}[|\hat{R}_n(h^*) - R(h^*)| > \epsilon] \leq 2me^{-2\epsilon^2 n}$$

Another way to write this, by setting $2me^{-2\epsilon^2 n} = \delta$

$$R(h^*) \leq \hat{R}_n(h^*) + \sqrt{\frac{1}{2n} \log \frac{2m}{\delta}}$$

If m is "infinite," it's already exponential

If $m \propto e^n$, we have a problem...

No matter how big n gets $\sqrt{\frac{1}{2n} \log \frac{2m}{\delta}}$ will never be smaller...

What if we replace with m with $m_{\mathcal{H}}(n)$? Suppose that for any $\delta \in (0, 1)$, we can guarantee at least $1 - \delta$

$$R(h^*) \leq \hat{R}_n(h^*) + \sqrt{\frac{1}{2n} \log \frac{2m_{\mathcal{H}}(n)}{\delta}}$$

- If $m_{\mathcal{H}}(n) = 2^n$ then $\sqrt{\frac{1}{2n} \log \frac{2m_{\mathcal{H}}(n)}{\delta}}$ is constant worst case - shattering
- If $m_{\mathcal{H}}(n)$ is a polynomial in n , $\sqrt{\frac{1}{2n} \log \frac{2m_{\mathcal{H}}(n)}{\delta}}$ decays like $\sqrt{\frac{\log n}{n}}$.

When is learning feasible? Instead of just memorizing based on sheer quantity of hypotheses

Assuming that we are indeed allowed to substitute $m_{\mathcal{H}}(n)$ for m , we can argue that for a given set of hypothesis \mathcal{H} learning is possible provided that $m_{\mathcal{H}}(n)$ is a polynomial.

Key idea: Break points

def'n: If no data set of size k can be shattered by \mathcal{H} , then k is a **break point** for \mathcal{H} .

$$m_{\mathcal{H}(k)} < 2^k$$

This also implies that if k is a break point, then so is any $k' > k$.

Examples of Break points

- Positive rays: $m_{\mathcal{H}}(n) = n + 1$
 - break point: $k = 2$
- Positive intervals: $m_{\mathcal{H}}(n) = \frac{1}{2}n^2 + \frac{1}{2}n + 1$ not n^2
 - break point: $k = 3$
- Convex sets: $m_{\mathcal{H}}(n) = 2^n$
 - break point: $k = \infty$
- Linear classifiers in \mathbb{R}^2 :
 - break point: $k = 4$



If there exists any break point, then $m_{\mathcal{H}}(n)$ is polynomial in n



If no break points, then $m_{\mathcal{H}}(n) = 2^n$

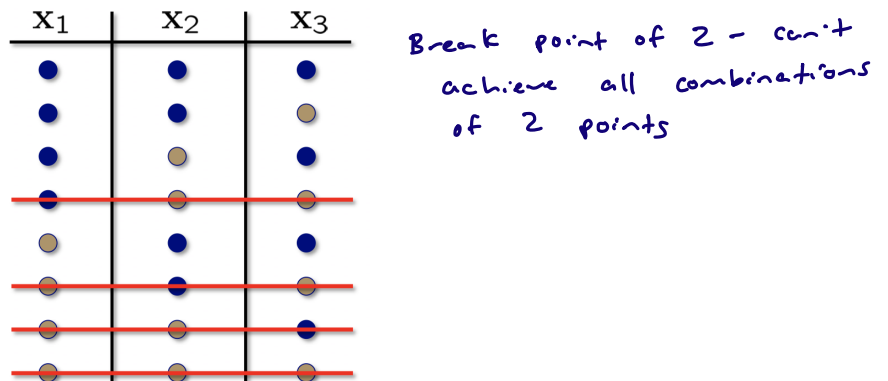
As soon as we have a single break point, this starts eliminating tons of dichotomies.

- We can show that $m_{\mathcal{H}}(n)$ is polynomial in n .
- To show that we worry too much and show that $m_{\mathcal{H}}(n) \leq \text{some polynomial}$
- Main approach will center around:
 - $B(n, k) :=$ maximum number of dichotomies on n points such that no subset of size k can be shattered by these dichotomies
 - Notice that this is a purely combinatorial quantity
 - By definition, $m_{\mathcal{H}}(n) \leq B(n, k)$

Example: how many dichotomies?

You are given a hypothesis set which has a break point of 2.

How many dichotomies can you get on 3 data points?



Summary: $B(n, k)$ is the combinatorial quantity that's an upper bound on the growth function for any possible set of classifiers.

You can bound $B(n, k)$ recursively which is an algorithmic proof. We will just skip the analytical proof as it is pages and pages math. There is also a “proof by picture” for this which I like, you may review that from “Learning from Data” if you are interested.

$$B(n, k) \leq B(n - 1, k) + B(n - 1, k - 1)$$

Analytical solution: $B(n, k) \leq \sum_{i=0}^{k-1} \binom{n}{i}$ You can prove that it is actually equal,

$$B(n, k) = B(n-1, k) + B(n-1, k-1)$$

but all we really need is an upper bound, so that is all we will prove here.

Proof by induction:

$$B(n, k) \leq B(n-1, k) + B(n-1, k-1)$$

- Base case

$$B(n, 1) = 1$$

$$B(1, k) = \begin{cases} 1 & \text{if } k = 1 \\ 2 & \text{otherwise} \end{cases}$$

- Inductive step

– suppose the inequality is true for $B(n-1, k)$ and $B(n-1, k-1)$

$$\begin{aligned} B(n, k) &\leq \sum_{i=0}^{k-1} \binom{n-1}{i} + \sum_{i=0}^{k-2} \binom{n-1}{i} \\ &= 1 + \sum_{i=0}^{k-1} \binom{n-1}{i} + \sum_{i=1}^{k-1} \binom{n-1}{i-1} \\ &= 1 + \sum_{i=1}^{k-1} \left(\binom{n-1}{i} + \binom{n-1}{i-1} \right) \\ &= 1 + \sum_{i=1}^{k-1} \binom{n}{i} = \sum_{i=0}^{k-1} \binom{n}{i} \end{aligned}$$