

## Lecture E1: Review

### CS 1555: Database Management Systems

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<http://db.cs.pitt.edu/courses/cs1555/current.term/>

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Lectures based: P. Chrysanthos & N. Farnan Lectures

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## ACID

- Atomicity
  - Either all the operations associated with a transaction happen or none of them happens
- Consistency
  - A transaction is a correct program segment. It satisfies the database's integrity constraints at its boundaries
- Isolation
  - Transactions are independent, the result of the execution of concurrent transactions is the same as if transactions were executed serially, one after the other
- Durability (a.k.a. Permanency)
  - The effects of completed transactions become permanent surviving any subsequent failure(s)



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## Useful Terms

- Cardinality of a relation  $r(R)$ : # of tuples in  $r(R)$  (denoted by  $|r(R)|$ )
- Arity or degree of  $r(R)$ : # of attributes in  $R$  (denoted by  $|R|$ )

$$|r(R)| = 3$$

$$|R| = 4$$

SID	Degree	Major	Year
123	BS	Math	1992
064	BA	History	1991
445	PhD	CS	1999

- ◆  $|r(R)| \geq 0$  And  $|R| > 0$
- ◆ Cardinality is property of a relation
- ◆ Arity is property of relation schema or a relation



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## Relational Database Schema

- A **database schema** is a set of relation schemas and a set of **integrity constraints**



### Integrity Constraints

- **Structural** Integrity Constraints
  - **key** constraints: uniqueness of keys
  - **entity integrity** constraint: no primary key value can be **NULL**
  - **referential integrity** constraint
- **Semantic** Integrity Constraints
  - E.g., ??



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## Relational Algebra

- Operations on entire relations  $- + \times \div$ 
  - Operands are (constant or variable) relations
  - Result is a relation
- Set theory operations:
  - Union, Intersection, Difference and Cartesian Product (product for short)
- Specific relational operations:
  - Selection, Projection, Join and Division
- Complete set of relational algebra operations:
  - Select, project, product, union and difference
- SQL is based on concepts from relational algebra



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## Relational Algebra

- Suppose relation Student has 20 tuples. What is the minimum and maximum number of tuples in the result of this expression:

$$\rho_{s1(i1,n1,g,h)} Student \bowtie \rho_{s2(i2,n2,g,h)} Student$$

- minimum = 0, maximum = 400
- minimum = 20, maximum = 20
- minimum = 20, maximum = 400 ✓
- minimum = 40, maximum = 40



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## Relational Algebra

- Which of the following English sentences describes the result of this expression:

$$\pi_{cName} College - \pi_{cName} (Apply \bowtie (\pi_{sID} (\sigma_{GPA > 3.5} Student) \cap \pi_{sID} (\sigma_{major = CS} Apply)))$$

- All colleges with no GPA>3.5 applicants who applied for a CS major at that college
- All colleges with no GPA>3.5 applicants who applied for a CS major at any college ✓
- All colleges where all applicants either have GPA>3.5 or applied for a CS major at that college
- All colleges where no applicants have GPA>3.5 or no applicants applied for a CS major at that college



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## Outer Join Examples

- STUDENT(SID, Name, Class, Major)
- ENROLLS(CID, SID, Term, Grade)

- Q1:

```
SELECT *
FROM (STUDENT S LEFT OUTER JOIN ENROLLS E
      ON S.SID=E.SID)
ORDER BY S.SID;
```

- Q2:

```
SELECT SID, S.Name, S. Major
FROM STUDENT S NATURAL LEFT OUTER JOIN ENROLLS E
WHERE E.Term IS NULL;
```



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## Outer Join Q1 Execution

### Students

SID	Name	Class	Major
123	John	3	CS
124	Mary	3	CS
999	Newman	1	CS

### Enroll

SID	CID	Term	Grade
123	CS1520	Fall 10	3.75
124	CS1520	Fall 10	4
123	CS1555	Fall 10	4
124	CS1555	Fall 10	NULL

### Q1 RESULT

S.SID	S.Name	S.Class	S.Major	E.SID	E.CID	E.Term	E.Grade
123	John	3	CS	123	CS1520	Fall 10	3.75
123	John	3	CS	123	CS1555	Fall 10	4
124	Mary	3	CS	124	CS1520	Fall 10	4
124	Mary	3	CS	124	CS1555	Fall 10	NULL
999	Newman	1	CS	NULL	NULL	NULL	NULL

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## Outer Join Q2 Execution

□ **SELECT** SID, S.Name, S. Major  
**FROM** STUDENT S NATURAL LEFT OUTER JOIN ENROLLS E  
**WHERE** E.Term IS NULL;

### Students

SID	Name	Class	Major
123	John	3	CS
124	Mary	3	CS
999	Newman	1	CS

### Enroll

SID	CID	Term	Grade
123	CS1520	Fall 10	3.75
124	CS1520	Fall 10	4
123	CS1555	Fall 10	4
124	CS1555	Fall 10	NULL

### Q2 RESULT

S.SID	S.Name	S. Major
999	Newman	CS

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## Pattern Matching...

- Retrieve all students with *local* phone numbers (any area code) which start with 6 and whose third digit is 3.

```
SELECT Name
FROM STUDENT
WHERE Phone LIKE '___6_3%';
```

- Escape defines the escape character that causes SQL to interpret a wildcard char (%) as itself in a string:

```
SELECT VideoName
FROM RENTALS
WHERE Discount LIKE '10&%' ESCAPE '&';
```

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## Natural JOIN

- Suppose relation R(A,C) has the following tuples and relation S(B,C,D) has the following tuples:
- Compute the natural join of R and S. Which of the following tuples is in the result? Assume each tuple has schema (A,B,C,D).

A	C
3	3
6	4
2	3
3	5
7	1

B	C	D
5	1	6
1	5	8
4	3	9

- a) (6, 4, 3, 9)
- b) (2, 4, 3, 9) ✓
- c) (2, 3, 1, 6)
- d) (5, 1, 6, 4)

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## Theta JOIN

- Suppose relation R(A,C) has the following tuples and relation S(B,C,D) has the following tuples:
- Compute the theta-join of R and S with the condition  $R.B = S.B \text{ AND } R.A < S.C$ . Which of the following tuples is in the result? Assume each tuple has schema (A, R.B, S.B, C, D).

A	B
1	a
7	t
2	g
4	c
9	t

B	C	D
c	5	6
a	7	8
t	8	9

- a) (2, g, c, 5, 6)
- b) (4, c, c, 7, 8)
- c) (1, a, c, 5, 6)
- d) (4, c, c, 5, 6) ✓



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## Projection

- Suppose relation R(A,B,C) has the following tuples
- Compute the projection  $\pi_{C,B}(R)$ . Which of the following tuples is in the result?

A	B	C
1	2	3
4	2	3
4	5	6
2	5	3
1	2	6

- a) (6,2) ✓
- b) (5,6)
- c) (5,3)
- d) (2,6)



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## Difference

- Suppose relation R(A,B,C) has the following tuples and relation S(A,B,C) has the following tuples:
- Compute  $(R - S) \cup (S - R)$ , often called the "symmetric difference" of R and S. Which of the following tuples is in the result?

A	B	C
1	2	3
4	2	3
4	5	6
2	5	3
1	2	6

A	B	C
2	5	3
2	5	4
4	5	6
1	2	3

- a) (1,2,3)
- b) (2,5,4) ✓
- c) (4,5,6)
- d) (1,5,6)



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## Intersection

- Suppose relation R(A,B,C) has the following tuples and relation S(A,B,C) has the following tuples:
- Compute the intersection of the relations R and S. Which of the following tuples is in the result?

A	B	C
1	2	3
4	2	3
4	5	6
2	5	3
1	2	6

A	B	C
2	5	3
2	5	4
4	5	6
1	2	3

- a) (2,2,6)
- b) (2,5,4)
- c) (1,2,3) ✓
- d) (2,4,3)



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## Query Rewriting

### View

```
CREATE VIEW CS_Students AS
SELECT name, age
FROM Student
WHERE Major = 'CS';
```



### Original Query (user)

```
SELECT name
FROM CS_Students
where age > 19;
```



### Modified Query (DBMS)

```
SELECT name
FROM Student
WHERE Major = 'CS'
AND age > 19;
```



## Views

- Consider tables  $R(A,B)$  and  $S(B,C)$  and a view  $V = \text{select } A,C \text{ from } R,S \text{ where } R.B=S.B$ . Suppose  $R=\{(1,5),(2,5)\}$  and  $S=\{(5,10)\}$ , so  $V=\{(1,10),(2,10)\}$ . The user wants to delete tuple  $(2,10)$  from  $V$ . Which of the following modifications to  $R$  and/or  $S$  does NOT correctly reflect this modification?

- delete  $(2,5)$  from  $R$
- update  $(2,5)$  to  $(2,6)$  in  $R$
- update  $(2,5)$  to  $(1,6)$  in  $R$
- delete  $(5,10)$  from  $S$  ✓



## Views

- Consider tables  $R(A,B)$  and  $S(B,C)$  and a query  $Q = \text{select } A,C \text{ from } R,S \text{ where } R.B=S.B \text{ and } A < 10 \text{ and } C > 20$ . Which of the following materialized views can NOT be used to help evaluate  $Q$ ?

- $V1 = \text{select } A,C \text{ from } R,S \text{ where } R.B=S.B$
- $V2 = \text{select } A,C \text{ from } R,S \text{ where } A < 10 \text{ and } C > 20$  ✓
- $V3 = \text{select } A,R.B,S.B,C \text{ from } R,S \text{ where } A < 10 \text{ and } C > 20$
- $V4 = \text{select } * \text{ from } R \text{ where } A < 10$

