

angel 2

## HashMap &amp; Heap Questions

1. Count of Subarrays with equal 0's & 1's. → gfg
2. Contiguous subarray (Longest subarray with equal 0's & 1's). → gfg & Leetcode

3. Count of Subarray with equal 0, 1, 2. → gfg.
4. Longest subarray with equal 0, 1, 2.  
↳ pepcoding

5. Number of subarrays sum exactly k.  
↳ gfg

6. Longest sub array with Sum K. → gfg.

7. Subarray sum divisible by k. → gfg.

8. Longest subarray with sum divisible by k.

↳ leetcode, gfg.

# HASHMAP AND HEAP STARTED

Date → 11 March.

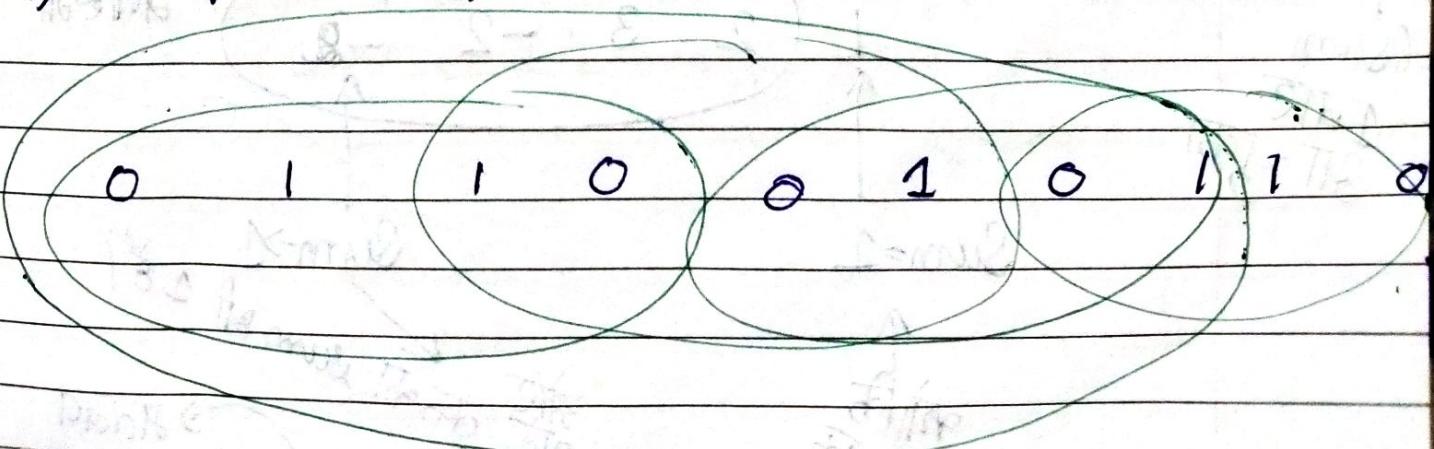
→ subarrays are continuous

## 1. COUNT OF SUBARRAYS WITH EQUAL 0'S AND 1'S. → Ifg

एक array given है। इसे बताना होगा कि इसे किसे किसी भी subarray में equal no. of 0's & 1's होते हैं।

किसी भी array के  $\frac{n(n+1)}{2}$  subarrays होते हैं।

→ O(n) में ये problem solve करनी है।



कैसे करें  
 0 → -1 add करें  
 1 → 1 add करें

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0 1 1 1 1 0 0 1 0 1 1 0

Sum. 0 -1 0 1 2 1 0 1 0 1 2 1

→ यह sum variable को update करते चले।  
 और हर बार जो sum की value आयगी, उसका  
 Hashmap में add करे as a key और उसके value होगी  
 frequency of this sum.

Sum

0 1 1 1 0 0 1 0 1 1 0  
 ↓  
 0 -1 0 1 2 1 0 1 0 1 2 1  
 ↑ ↑ ↑ ↑ ↑ ↑ ↑ ↑ ↑ ↑ ↑ ↑  
 0 के लिए +1 +1 +1 +1 -1 +1 add +1 +1 +1 -1  
 -1 add  
 किया किया

0 → 1 0 → 2 1 → 1 2 → 1 0 → 3 1 → 3 0 → 4 1 → 2 3 → 3 1 → 5

→ जो भी Subarray  
of 0 sum बनाएगा

↑  
 0sum  
 1वाले  
 आये हुए

1 2 3 -2 -3

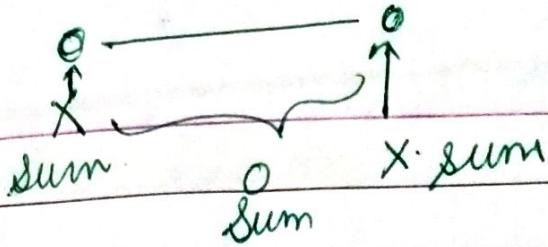
Sum=1

पहली  
बार तक  
की sum.  
ही है

Sum>1  
 और 2वाले तक की sum of 1 है

मतलब  
1 sum

बाजाने  
लाले 1 Subarray



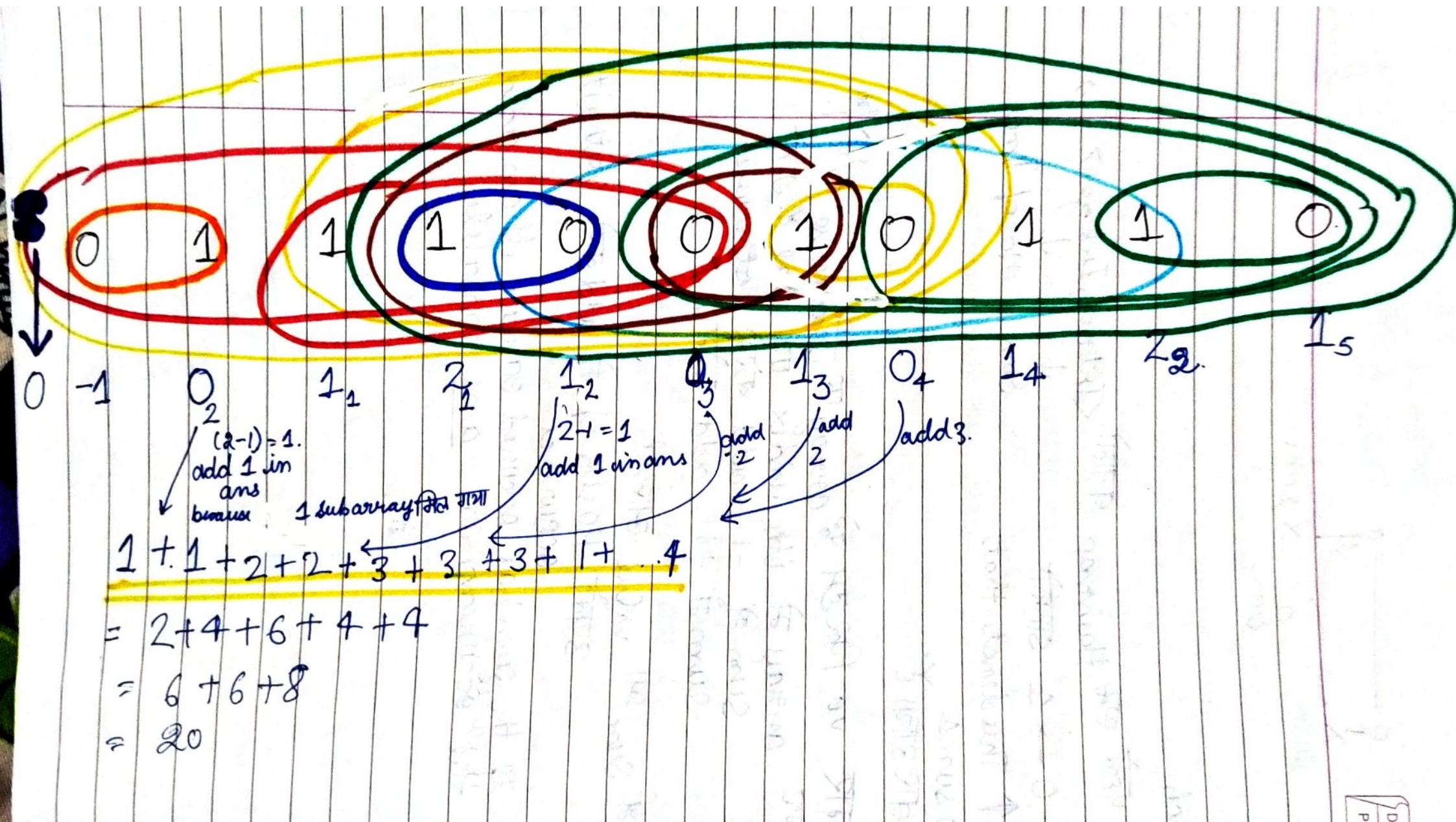
### Approach :

- ① सबसे पहले एक Hashmap बनाते हैं,  $\langle \text{Integer}, \text{Integer} \rangle$  की,
- ② उसमें  $0 \rightarrow 1$  लिखते हैं।  
This denotes that

$0 \text{ sum} 1$   
हार आया है।

- ③ फिर हम उसे array पर traverse करेंगे।  
अगर array में ith index पर val  $\neq 0$  हो  
sum में -1 add करेंगे otherwise  
sum में +1 add करेंगे।
- ④ जो भी sum भी value आएगी।  
उसको Hashmap में add करो निचे its occurrence.

If the sum has occurred earlier than add the occurrence in ans variable.



## public class Solution

```
static int CountSubArrWithEqualZeroAndOne  
(int arr[], int n).
```

```
{ int ans = 0;
```

```
HashMap<Integer, Integer> map = new HashMap<>();  
map.put(0, 1);
```

Initially sum & 0 add

0 के कोटि परिवर्तन

मैंने 0 के लिए sum 0 रखा

```
int sum = 0;
```

```
for (int val : arr)
```

```
{ if (val == 0)
```

Sum += -1;

val 0 से -1 add

-1 add करता है

otherwise +1

add 1 करता है

```
else
```

Sum += 1;

Hashmap में sum already

exists value get करता है

ans +1 add करता है

```
if (map.containsKey(sum))
```

```
{ int x = map.get(sum);
```

ans += x;

```
map.put(sum, x + 1);
```

Hashmap में sum की value +1 करता है

```
else
```

```
{ map.put(sum, 1);
```

otherwise Hashmap में sum की value 1 से start करता है

```
} return ans;
```

## 2. Longest Subarray with Equal 0 & 1.

0	0	1	0	1	0	1	1	0	1	1	1	
0	1	2	3	4	5	6	7	8	9	10	11	12

Sum -1	-2	-1	-2	-1	-2	-10	+1	-2	-1	0	1
--------	----	----	----	----	----	-----	----	----	----	---	---

Hashmap में sum और <sup>जोकि</sup> index add करो

Then when again the same sum occur again,  
then check the difference in their indices.

→ पहले सबसे पहले कहीं sum आए हो उसे add करदो

Hashmap में, with its current index  
& when the same sum occur again,  
just check the differences in their index

that will be the length of subarray

→ Keep on updating the length of max  
subarray

class Solution

{  
public int findMaxLength (int [ ] arr)

int sum = 0;

int MaxLength = 0;

HashMap < Integer, Integer > map = new HashMap <> ();  
map.put (0, -1);

for (int i = 0; i < arr.length; i++)

int val = arr[i];

if (val == 0)

sum += -1;

else

sum += 1;

if (map.containsKey (sum))

int x = map.get (sum);

int currentLength = i - x;

if (currentLength > MaxLength)

maxLength = currentLength;

Account  
Index  
Previous  
Index  
Subarray  
Length

{ map.put (sum, i);

return maxLength;

if the sum occurred  
for first time,  
put it in  
the Hashmap  
& its index

}

{

int sum = 0;

int MaxLength = 0;

HashMap < Integer, Integer > map = new HashMap <> ();  
map.put (0, -1);

for (int i = 0; i < arr.length; i++)

int val = arr[i];

if (val == 0)

sum += -1;

else

sum += 1;

if (map.containsKey (sum))

int x = map.get (sum);

int currentLength = i - x;

if (currentLength > MaxLength)

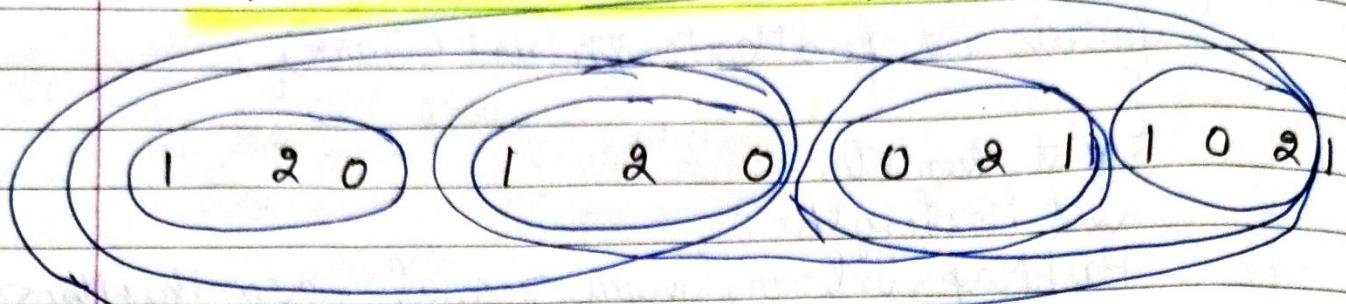
maxLength = currentLength;

{ map.put (sum, i);

return maxLength;

if the sum occurred  
for first time,  
put it in  
the Hashmap  
& its index

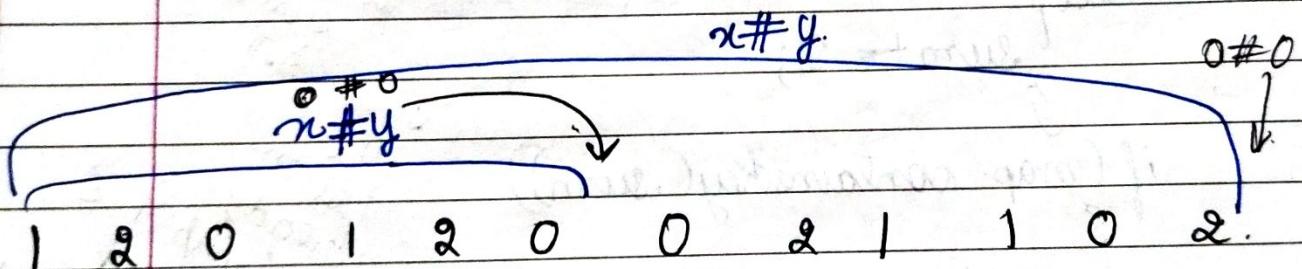
### 3. COUNT OF SUBSTRING WITH EQUAL 0, 1, 2.



gap 1,0    ~~#~~ gap 2,1

2) ~~दो~~ दोरी द्वारा keys in the Hashmap.

& the <sup>no. of</sup> occurrence of a particular key will be its value.



If ~~मात्र~~ same key occur again,  
it means the no. of 0, 1, 2 in  
between are equal

So we ~~would~~ keep a count of 0, 1, 2  
and get the key  $\rightarrow$  gap 10 # gap 21

~~जितना~~ gap in the 1st  
gap in no. of 0 & no. of 1  
~~जितना~~ gap in no. of 2 & no. of 1

$$\begin{aligned}
 \text{ans} &= 1 + 1 + 1 + 2 + 3 + 2 + 4 + 3 \\
 &= 3 + 2 + 3 + 2 + 4 + 3 \\
 &= 5 + 5 + 7 \\
 &= 17.
 \end{aligned}$$

<code>

```
public static long NumberOfSubArrayWithEqualSum  
(int arr[])
```

```
{
```

```
int count0=0;
```

```
int count1=0;
```

```
int count2=0;
```

```
String key = "0#0";
```

} Initially  
count of 0, 1, 2 is 0

↑  
gap1 ↑  
gap2 ↑

} Initially  
the difference (gap)  
of 1 & 0 and  
gap of 2 & 1 is

```
HashMap<String, Integer> map = new HashMap<>();  
map.put(key, 1);
```

↓  
0#0 key put करता है

HashMap में

Because initially it  
count of 0, 1, 2 are all zero.

```
int ans=0;
```

for (int i=0; i < arr.length; i++)

{  
    uint val = arr[i];  
    if (val == 0)

{  
    count0++;

?  
    close if (val == 1)

{  
    count1++;

?  
else

{  
    count2++;

key = (count1 - count0)  $\neq$  (count2 - count1);  
if (map.contains(key))

{  
    int x = map.get(key).  
    ans += x;  
    map.put(key, x+1);

Add  
the earlier  
value in  
answer

↓  
add 1 to  
the already  
present  
value

else { map.put(key, 1);

↓  
if this key  
is not already  
present, then  
add it with  
value 1.

return ans;

↓  
return  
the ans

#### 4. Longest Subarray with equal 0/1/2.

0	1	0	2	0	1	0
Count0	0	1	2	2	3	4
Count1	0	0	1	1	1	2
Count2	0	0	0	1	1	1
0#0 -1#0 0#-1 -1#-1 -1#0 -2#0 -1#-1 -2#-1						
Index -1	0	1	2	3	4	5

Put the key & value in hash map

Subtract 0 from 3 = 3  
Length = 3  
Max Length = 3

Subtract 2 from 5 = 3  
Length = 3  
Max Length = 3.

दूसरी तरफ length का लिया जाएगा जिसमें count of 0, 1, 2 equal होंगे और length of this subarray should be maximum of all subarrays with equal 0, 1, 2.

#### → Approach

- दूसरे point पर 0 की count, 1 की count, 2 की count calculate करते चलेगा और फिर check करेगा कि  $x_1 - x_0$  और  $x_2 - x_1$  की value क्या हैं।

$$x_1 - x_0 \neq x_2 - x_1$$

This Ray  
is a  
Strong

If the same gap occurs again at some other index, it means the count of 0, 1, 2 in b/w 2 gaps are equal.

2. If the same key occurs again at some other index, then calculate the difference b/w both index, this will be the length of subarray.
3. Find the maximum length of all lengths of subarrays with equal 0, 1, 2.

count	1	0	1	2	0	1	0	1	2	1	2	1	2	2	0	1	1
key	0#0	1#-1	2#-2	2#-1	1#-1	2#-2	1#-2	2#-3	2#-2	3#-3	3#-2	3#-1	2#-1	3#-1	3#-2		
index	-1	0	1	2	3	4	5	6	7	8	9	10	11	12			

Initially  
0#0 is  
the key &  
index is -1  
Put 0#0 → -1  
in the Hashmap.

This key  
is there in  
map at  
index 0

$$\text{Length} = 3 - 0 = 3$$

This key  
is present  
at index 1

$$\text{Length} = 4 - 1 = 3$$

This key  
is present  
at index 1

$$7 - 1 = 6$$

This key  
is present  
at index 2

$$11 - 2 = 9$$

This key is  
already present  
at index 3

$$12 - 9 = 3$$

$$\begin{aligned} \text{map} = & 0\#0 \rightarrow -1 \\ & 1\#-1 \rightarrow 0 \\ & 2\#-2 \rightarrow 1 \\ & 2\#-1 \rightarrow 2 \\ & 1\#-2 \rightarrow 5 \\ & 2\#-3 = 6 \end{aligned}$$

$$\begin{aligned} & 3\#-3 = 8 \\ & 3\#-2 = 9 \\ & 3\#-1 = 10 \end{aligned}$$

↓  
Maximum  
length = 9

public static int solution (int arr[])

{  
    int count0 = 0;  
    int count1 = 0;  
    int count2 = 0;

} Initially  
count of 0,  
count 1, 2 is  
count of 0.

String key = "0#0"; The is  
the key initially

HashMap<String, Integer> map = new HashMap<>();

map.put(key, -1);

This is  
the hashmap.

int length = 0;

int MaxLength = 0;

Put the  
key sets under  
as ->

for (int i=0; i < arr.length; i++)

{ int val = arr[i];

if (val == 0)

We will travel  
on all the  
values in  
array

    count0++;

    else if (val == 1)

If the  
val  
is 0  
count0

        count1++;

    If val  
    is 1  
    count1  
    increase  
    by 1

    else {

        count2++;

If val is 2,  
increase count2 by 1

}

key = (count1 - count0) + "#" + (count2 - count1);

(The key  
will be gap of 1 & 0 # gap of 2 & 1

If this key is already present in hashmap then find the difference between current index & key value.

if (map::containsKey(key))

{ int length = i - x;

maxLength = Math::max(length, maxLength);

}

else

{ map::put(key, value);

The maximum length will be the max of length & maxLength.

If the key is current index  
not already present)

then put the key value  
with current  
as current index

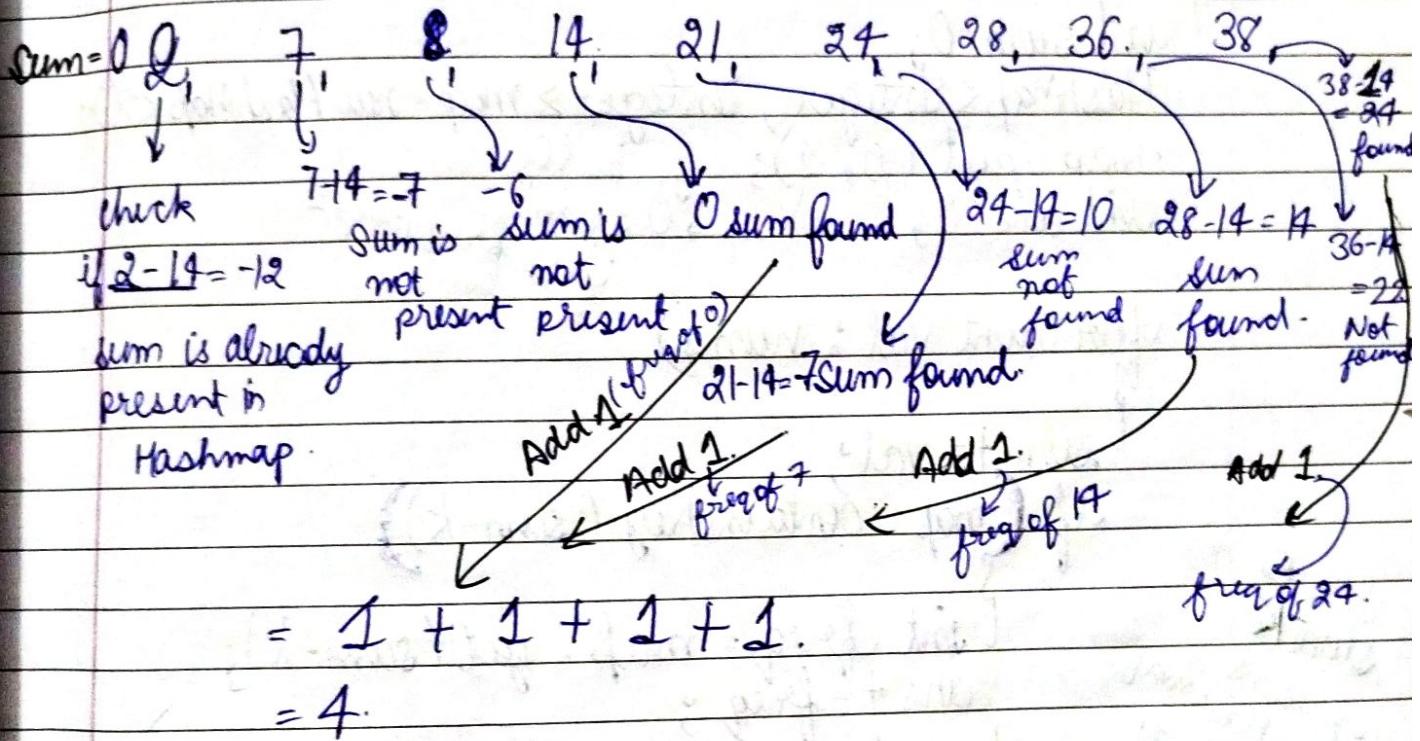
return maxLength;

}

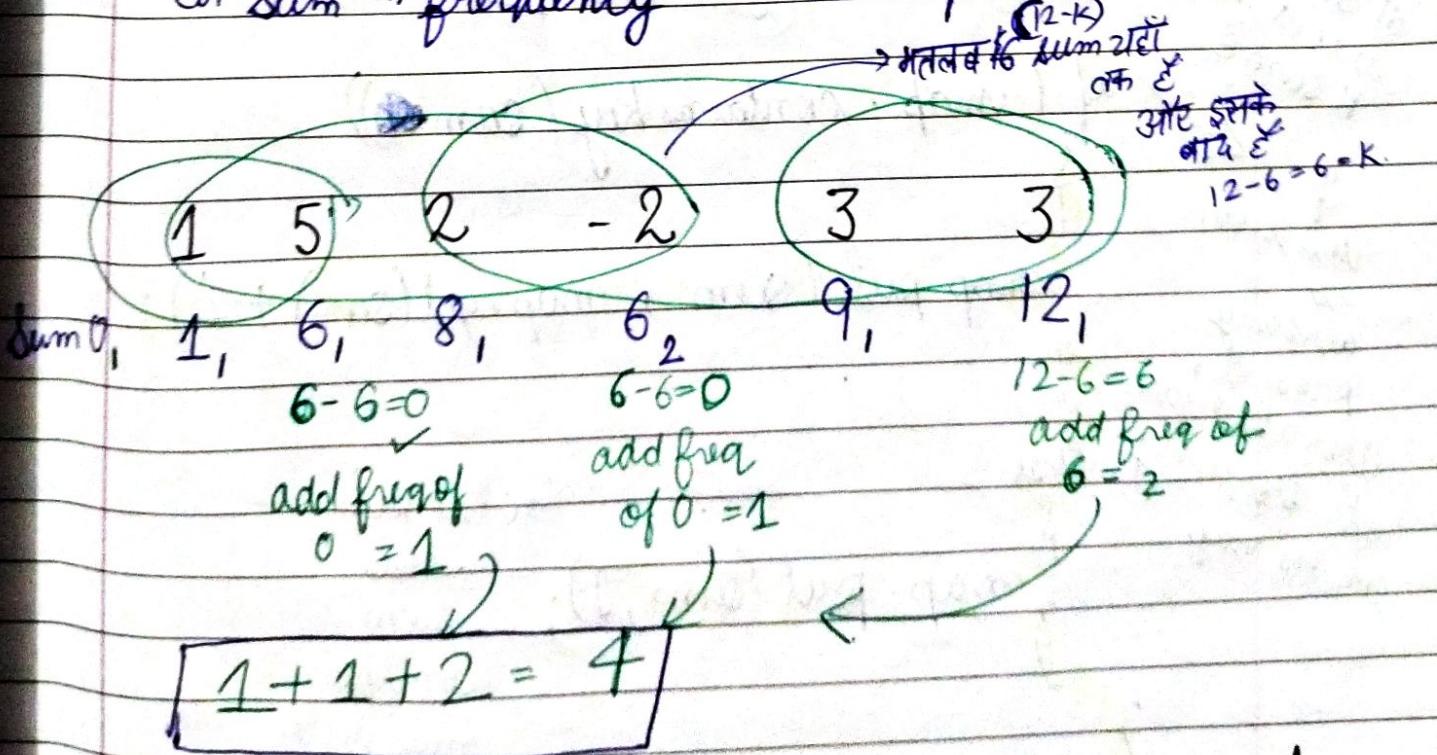
## 5. Count of Subarray with Sum equal to K.

Sum = 14.

2 5 1 6 7 3 4 8 2



EH Sum की frequency का hashmap बनाओ



Check if  $\sum - k$  is present in Hashmap or not.

If yes, then add its frequency in answer.

&lt;Code&gt;

public int subarraysum (int [] nums, int k)

{  
    int sum = 0;

    HashMap<Integer, Integer> map = new HashMap<X>;  
    map.put(0, 1);     Initially  
    int ans = 0;     → 0 sum  
                      has frequency 1

    for (int val : nums)

        sum += val;

        if (map.containsKey(sum - k))

            int freq = map.get(sum - k);

            ans += freq;

            ↳ add the freq of  
(sum - k) in ans

        if (map.containsKey(sum))

            if this sum

                is already

                present,

                then increase

                its frequency

                in the map

            map.put(sum, map.get(sum) + 1);

        else

            map.put(sum, 1);

    else put

    this

    sum

    in the

    map with

    its freq = 1

    return ans;

## 6 Longest Sub Array with Sum K.

$$K=15$$

0 1 2 3 4 5  
10 5 2 7 1 9

Sum 0

10 15 17 24 25 34

$$15-15=0 \checkmark \quad 17-15=2 \quad 24-15=9X \quad 25-15=10 \checkmark \quad 34-15=19X.$$

subtract  
1 from 1 = 2      4 - 0 = 4

Put 0 with  
+ index in  
Map

Check is

Sum-K present?

If yes then subtract  
its index from  
current index

This will be length  
of current subarray.

We need to compare  
all the lengths  
& return  
Maxlength.

map → length

10	0	→ 1 - (-1) = 2
15	1	
24	2	
25	3	
34	4	→ 4 - 0 = 4
	5	

So maxlen = 4

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Approach ① add(0, -1) in map



- ② Create array at  $\text{arr}[\text{index}]$  value of sum variable  
& add arr[0] = 0.
- ③ if  $(\text{sum} - k)$  key is present in HashMap,  
then

the length of the corresponding subarray  
will be  $\text{currentIndex} - (\text{sum} - k) \text{Index}$

- ④ If the current sum is not present in the  
HashMap, Put it with the current index as  
its value.

$(\text{sum}, \text{Index})$

$\curvearrowright$  current index

- ⑤ Compare length of the subarrays & return  
the maximum length.

<code>

```
public static int lenOfLongSubarr
    (int nums[], int k)
```

{

```
    int sum = 0;
```

```
    HashMap<Integer, Integer> map = new HashMap<>();
    map.put(0, -1);
```

```
    int length = 0;
```

```
    int MaxLength = 0;
```

```
    for (int i = 0; i < nums.length; i++)
```

{

```
        int val = nums[i];
```

```
        sum += val;
```

```
        if (map.containsKey(sum - k))
```

```
            int index = map.get(sum - k);
```

```
            length = i - index;
```

```
            MaxLength = Math.max(length, MaxLength);
```

}

```
    if (map.get(sum) == null)
```

{

```
        map.put(sum, i);
```

}

}

```
return MaxLength;
```

}

7.

## Subarray sum divisible by k.

$K=4$ .

2 5 7 8 3 5 7

Hashmap में इसे store करें  $\frac{\text{sum} \% K}{\text{remainder}}$  with its frequency

2 5 7 8 3 5 7.

Sum / K = 2

अपर्याप्त तक  
के sum का  
remainder  
 $2 \not| \Sigma$

Sum / K = 2.

अपर्याप्त तक के  
sum का  
remainder  
 $2 \not| \Sigma$

तो  $2 \not| \text{sum}$

जोमाला

$n(K)$

multiple of K  
(divisible  
by K).

$k=7$

a b c d e f g h i j

$$S_1 = 51.$$

$$S_1 \% K = 51 \% 7 = 2.$$

$$S_1 = km + x.$$

$$S_2 - S_1 = k(m - n)$$

This is a subarray  
divisible by  $K$ .

$S_2$

$$\begin{aligned} S_2 \% K &= x \\ 79 \% 7 &= 2 \\ S_2 &= kn + x \end{aligned}$$

$k=7$

11 6 -32 4 12 34 6 14 9 quiet array

0 11 17 -15 -11 1 35 41 55 64 ← Sum

Initially  
sum 0

$O_1 \quad 4_1 \quad 3_1 \quad -1_1 \quad 1_1 \quad 0_2 \quad 6_2 \quad 6_3 \quad 1_2 \quad K$

$$= 1 + 1 + 1 + 2 + 1 = 6$$

$\leftarrow \overbrace{1 \quad -6 \quad -5 \quad -4 \quad -3 \quad -2 \quad -1 \quad 0 \quad 1 \quad 2 \quad 3 \quad 4 \quad 5 \quad 6 \quad 7}^{\text{Sum}} = 7 \cdot 5 - 3$

$$\begin{aligned} \text{Sum} &= Kx - \alpha + K - K \\ &= K(x-1) + (K-\alpha) \\ &= K(x-1) + (K-\alpha) \end{aligned}$$

$$\begin{aligned} &= (7 \cdot 5) - (3 + 7) - 7 \\ &= 7(5-1) + 4 \\ &= 7 \cdot 4 + 4. \end{aligned}$$

public static int solution (int[] arr, int k)

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{   
 HashMap<Long, Long> map = new HashMap<>();

long sum = 0;

long ans = 0;

Initially sum  
& answer is  
0

map.put(sum, 1L); This is however put 1 (Integer)  
as long  
for (int i=0; i<arr.length; i++). (1L means Integer value)

{ long val = arr[i];

sum = sum + val;

long remainder = val % k;

if (remainder < 0).

remainder  
of value  
of arr[i]

{ remainder += k;

3H R remainder

-ve वे तो 3H R करा

if (map.containskey(remainder))

remainder  
+ve ए जाय

{ Long freq = map.get(remainder);

remainder map.put(remainder, freq + 1);

map.get(remainder) ans += freq;

already  
0  
else

frequency  
ans & add 1

frequency  
को  
increase

{ map.put(remainder, 1L);

3H R  
remainder  
already present

ans & add  
with frequency

return ans;

ans &  
value  
return 1.