

## Lecture 2

## Operations Counting

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## Estimating Time by Counting Operations

- We will start by assuming that one command, one programming statement, takes a fixed amount of time.
  - If we have  $m$  statement, it will take  $km$  milliseconds to follow them.
  - Linearly proportional to each other.
- For simplicity, we will assume that any statement takes the same amount time to compute.
  - For example,  $x=3$ ; or `Math.abs(-178)` are consider to take the same amount of time to compute.
  - This might not be true, but good enough for our purpose.
  - Note that `int m=n/2`; is count as 3 operations: declaration, assignment, and division.
- Let start counting

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## isPrime0()

	code	operation count
1	<code>static boolean isPrime0(int n) {</code>	
2	<code>    if (n==1) return false;</code>	1
3	<code>    if (n&lt;=3) return true;</code>	1
4	<code>    int m=n/2;</code>	3
5	<code>    for(int i=2; i&lt;=m; i++) {</code>	$2+m+(m-1) = 2m+1$
6	<code>        if (n%i==0) return false;</code>	$2(m-1)$
7	<code>    }</code>	
8	<code>    return true;</code>	1
9	<code>}</code>	
total =		$4m+5$
		$= 2n+5$

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Let Count isPrime1()

code	operation count
1 static boolean isPrime1(int n) {	
2   if (n==1) return false;	1
3   if (n<=3) return true;	1
4   int m = (int)Math.sqrt(n);	4
5   for(int i=2; i<=m; i++) {	$2+m+(m-1) = 2m+1$
6     if (n%i==0) return false;	$2(m-1)$
7   }	
8   return true;	1
9 }	
total	$= 4m + 6$
	$= 4\sqrt{n} + 6$

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isPrime2()

code	operation count
1 static boolean isPrime2(int n) {	
2   if (n==1) return false;	1
3   if (n<=3) return true;	1
4   if ((n%2==0)    (n%3==0)) return false;	5
5   int m = (int)Math.sqrt(n);	4
6   for(int i=5; i<=m; i+=6) {	$2+(\frac{m}{6}+1)+\frac{m}{6} = 2\frac{m}{6}+3$
7     if (n%i==0) return false;	$2\frac{m}{6}$
8     if (n%(i+2)==0) return false;	$3\frac{m}{6}$
9   }	
10   return true;	1
11 }	
total	$= 7\frac{m}{6} + 15$
	$= 7\frac{\sqrt{n}}{6} + 15$

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Operation Function

- These are functions of number operations where n is the size of the input

$$f_0(n) = 2n + 5$$

$$f_1(n) = 4\sqrt{n} + 6$$

$$f_2(n) = 7\frac{\sqrt{n}}{6} + 15$$

Let call them Operation Function

- Let plot some graph

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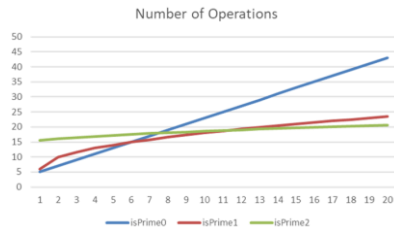
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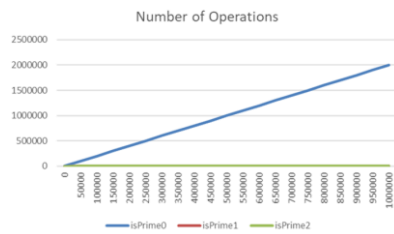
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Small input ( $n \leq 20$ )

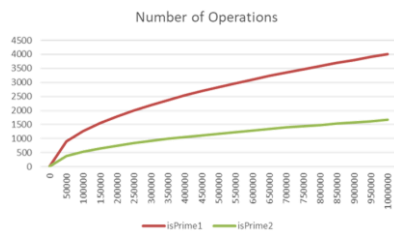
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## Large Input



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## Large Input w/o isPrime0



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## Some Discussion

$$f_0(n) = 2n + 5 \quad f_1(n) = 4\sqrt{n} + 6 \quad f_2(n) = 7\frac{\sqrt{n}}{6} + 15$$

- The result of statement counting is similar to our benchmark results.
  - isPrime0 grows at the same rate as the input.
  - isPrime1 and isPrime2 grows proportion to the square root of the input.
  - isPrime1 and isPrime2 are comparable, while isPrime0 is by far the slowest.
- It is typically possible to count number of operations of algorithms.
- It is also applicable to counting space.
- While the actual time depends on machines, number of operations does not.
  - Make it good for directly compare algorithms
- Rather than remember the functions of all known algorithms, we will compare them to well-known functions using **Asymptotic Analysis**.

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## Summary

- Rather than measuring the time, we can count number of operations in algorithms
- Number of operations are just an estimate number, which is good enough for our purpose.
- If we know the function of operations of algorithms, we can compare them easily
- However, we will not have to memorize the exact function of operations of each algorithm, we will use **Asymptotic Analysis** to compare them to well-known function.

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