

MATH 178 Homework #10

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My color: Red, My partner's color: Blue

RSA

1.

$p = 7, q = 97, e = 257, n = pq = 679, \varphi(679) = \varphi(7)\varphi(97) = 96 * 6 = 576.$

$$d = e^{-1}(\text{mod } 576) \equiv 1 \quad (1)$$

$$576 = 2 * 257 + 62 \quad (2)$$

$$257 = 4 * 62 + 9 \quad (3)$$

$$62 = 6 * 9 + 8 \quad (4)$$

$$9 = 1 * 8 + 1 \quad (5)$$

$$1 = 9 - 1 * 8 \quad (6)$$

$$1 = 9 - 1(62 - 6 * 9) \quad (7)$$

$$1 = 7 * 9 - 62 \quad (8)$$

$$1 = 7(257 - 4 * 62) - 62 \quad (9)$$

$$1 = 7 * 257 - 29 * 62 \quad (10)$$

$$1 = 7 * 257 - 29 * (576 - 2 * 257) \quad (11)$$

$$1 = 65 * 257 - 29 * 576 \quad (12)$$

$$1 = 65 * 257 \quad (13)$$

$$d = 65 \quad (14)$$

To decrypt 146, use $146^{65} \text{ mod } (679)$

Reduce $146^{65}(\text{mod } 679)$

$b = 65, n = 146, m = 679, S[] = \{1, 0, 0, 0, 0, 0, 1\}, k = 6$

| $n(\text{mod } 679)$ | s | a |
|----------------------|----------|-----|
| | | 1 |
| 146 | $s[0]=1$ | 146 |
| $146^2=267$ | $s[1]=0$ | 146 |
| $267^2=673$ | $s[2]=0$ | 146 |
| $673^2=36$ | $s[3]=0$ | 146 |
| $36^2=617$ | $s[4]=0$ | 146 |
| $617^2=449$ | $s[5]=0$ | 146 |
| $449^2=617$ | $s[6]=1$ | 454 |

To decode 454 to a key pair a, b , use $146/26 = 17 = a, 146 \text{ mod } 26 = 12 = b.$

Now we have the encrypting pair a, b we need to find the decrypting pair a', b'

$$26 = 1 * 17 + 9 \quad (1)$$

$$17 = 1 * 9 + 8 \quad (2)$$

$$9 = 1 * 8 + 1 \quad (3)$$

$$1 = 9 - 1 * 8 \quad (4)$$

$$1 = 9 - (17 - 1 * 9) \quad (5)$$

$$1 = 2 * 9 - 17 \quad (6)$$

$$1 = 2(26 - 1 * 17) - 17 \quad (7)$$

$$1 = 2 * 26 - 3 * 17 \quad (8)$$

$$1 = -3 * 17 \quad (9)$$

$$1 = 23 * 17 \quad (10)$$

We can now decrypt the message. Using the encoding $a = 0, b = 1...$

$$P \equiv a'(C - b)(\text{mod}26) \quad (1)$$

$$\equiv 23(C - 12)(\text{mod}26) \quad (2)$$

$$P \rightarrow 23(15 - 12)(\text{mod}26) = 17 = R \quad (3)$$

$$B \rightarrow 23(1 - 12)(\text{mod}26) = 7 = H \quad (4)$$

$$E \rightarrow 23(4 - 12)(\text{mod}26) = 24 = Y \quad (5)$$

$$X \rightarrow 23(23 - 12)(\text{mod}26) = 19 = T \quad (6)$$

$$B \rightarrow 23(1 - 12)(\text{mod}26) = 7 = H \quad (7)$$

$$I \rightarrow 23(8 - 12)(\text{mod}26) = 12 = M \quad (8)$$

The original plaintext was RHYTHM.

Using a new encrypting pair $a, b = \{11, 21\}$. $11 * 26 + 21 = 307$. To encrypt this, reduce $307^{19} \text{ mod } 681$.

Reduce $307^{19}(\text{mod } 681)$

$b = 19, n = 307, m = 681, S[] = \{1, 0, 0, 1, 1\}, k = 4$

| $n(\text{mod}681)$ | s | a |
|--------------------|----------|-----|
| | | 1 |
| 307 | $s[0]=1$ | 307 |
| $307^2=271$ | $s[1]=1$ | 307 |
| $271^2=574$ | $s[2]=0$ | 115 |
| $574^2=553$ | $s[3]=0$ | 115 |
| $553^2=40$ | $s[4]=1$ | 514 |

The reduced number that I would send to (you) is 514.

2.

$p = 50453156304734882643265661554998829784271380864981$
 $n = 1859486070465862481250084737025244499643330 \dots 2729$
 $q = 36855693610814891619906843543364439878667990665909$
 $e = 5746381426879165805597951012558256671104182 \dots 7277$
 $\varphi(n) = 1859486070465862481250084737025244499643 \dots 1840$
 $d = 1781807991792385566284160318478558467393146 \dots 9173$

3.

The message is:

53285491632384029814060608150978435305953163003130
0070853873325328959028610977865082224016588557854.

This was the encryption $m^e \pmod n$. Therefore to decrypt we need $m^d \pmod n$.
 $\text{Mod}(c, n)^d = 9260$. This is our key.

The ciphertext was 7FCF82CF65F8F607EB4B9D7CFA....

Therefore, our plaintext was:

fresh freaken crawdad. whats my middle name.

4.

Since my partner's color is *blue*, he also has a different n and e .

The key is 51913.

$\text{Mod}(m, n)^e = 1663 \dots 7914$

The key is 51914.

$\text{Mod}(m, n)^e = 9656 \dots 1560$

4.5.

Person 1 $\rightarrow n_1 = r_1 r_2$

Person 2 $\rightarrow n_2 = r_1 r_2$

Person 3 $\rightarrow n_3 = r_3 r_4$

Person 4 $\rightarrow n_4 = r_3 r_5$

- i) Person 2 is a danger to Person 1, because they share the same n, p, q . This means that Person 2 can decode any of Person 1's private messages.
- ii) If all of the n_i are the same, and one prime was being shared, r_3, r_4, r_5 specifically, then $r_5 = r_4$ if they are sharing the same n and r_3 . This means that Person 4 is a danger to Person 3.