

MATH 51 Homework #2

Tamir Enkhjargal

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Section 1.3

16.

(a) For the first example, I will also give the name of the identity used per step.

$$\begin{array}{lll}
 [\neg p \wedge (p \vee q)] \rightarrow q & \text{Given} & (1) \\
 \neg([\neg p \wedge (p \vee q)]) \vee q & \text{Implication} & (2) \\
 [p \vee \neg(p \vee q)] \vee q & \text{DeMorgan's \& Double Negation} & (3) \\
 [p \vee (\neg p \wedge \neg q)] \vee q & \text{DeMorgan's} & (4) \\
 [(p \vee \neg p) \wedge (p \vee \neg q)] \vee q & \text{Distributive} & (5) \\
 [T \wedge (p \vee \neg q)] \vee q & \text{Negation} & (6) \\
 (T \vee q) \wedge ((p \vee \neg q) \vee q) & \text{Distributive} & (7) \\
 T \wedge (p \vee (\neg q \vee q)) & \text{Domination \& Associative} & (8) \\
 T \wedge (p \vee T) & \text{Negation} & (9) \\
 T \wedge T & \text{Domination} & (10) \\
 T & \text{Q.E.D.} & (11)
 \end{array}$$

(b)

$$\begin{array}{lll}
 [(p \rightarrow q) \wedge (q \rightarrow r)] \rightarrow (p \rightarrow r) & (1) \\
 [(\neg p \vee q) \wedge (\neg q \vee r)] \rightarrow (\neg p \vee r) & (2) \\
 \neg[(\neg p \vee q) \wedge (\neg q \vee r)] \vee (\neg p \vee r) & (3) \\
 [(p \wedge \neg q) \vee (q \wedge \neg r)] \vee (\neg p \vee r) & (4) \\
 [(p \vee q) \wedge (p \vee \neg r)] \wedge ((\neg q \vee q) \wedge (\neg q \vee \neg r)) \vee (\neg p \vee r) & (5) \\
 [((p \vee q) \wedge (p \vee \neg r)) \wedge (T \wedge (\neg q \vee \neg r))] \vee (\neg p \vee r) & (6) \\
 [((p \vee q) \wedge (p \vee \neg r)) \wedge (\neg q \vee \neg r)] \vee (\neg p \vee r) & (7) \\
 [((p \vee q) \wedge (p \vee \neg r)) \vee (\neg p \vee r)] \wedge [(\neg q \vee \neg r) \vee (\neg p \vee r)] & (8) \\
 [((p \vee q) \wedge (p \vee \neg r)) \vee (\neg p \vee r)] \wedge [(\neg q \vee \neg p) \vee (\neg r \vee r)] & (9) \\
 [((p \vee q) \wedge (p \vee \neg r)) \vee (\neg p \vee r)] \wedge [(\neg q \vee \neg p) \vee T] & (10) \\
 [((p \vee q) \wedge (p \vee \neg r)) \vee (\neg p \vee r)] \wedge T & (11) \\
 [((p \vee q) \wedge (p \vee \neg r)) \vee (\neg p \vee r)] & (12) \\
 [(p \vee q) \vee (\neg p \vee r)] \wedge [(p \vee \neg r) \vee (\neg p \vee r)] & (13) \\
 [(p \vee \neg p) \vee (q \vee r)] \wedge [(p \vee \neg p) \vee (\neg r \vee r)] & (14) \\
 [T \vee (q \vee r)] \wedge [T \vee T] & (15) \\
 T \wedge T & (16) \\
 T & (17)
 \end{array}$$

(c)

$$\begin{aligned}
& [p \wedge (p \rightarrow q)] \rightarrow q & (1) \\
& [p \wedge (\neg p \vee q)] \rightarrow q & (2) \\
& \neg[p \wedge (\neg p \vee q)] \vee q & (3) \\
& [\neg p \vee (p \wedge \neg q)] \vee q & (4) \\
& [(\neg p \vee p) \wedge (\neg p \vee \neg q)] \vee q & (5) \\
& [T \wedge (\neg p \vee \neg q)] \vee q & (6) \\
& (\neg p \vee \neg q) \vee q & (7) \\
& \neg p \vee (\neg q \vee q) & (8) \\
& \neg p \vee T & (9) \\
& T & (10)
\end{aligned}$$

(d)

$$\begin{aligned}
& [(p \vee q) \wedge (p \rightarrow r) \wedge (q \rightarrow r)] \rightarrow r & (1) \\
& [(p \vee q) \wedge (\neg p \vee r) \wedge (\neg q \vee r)] \rightarrow r & (2) \\
& \neg[(p \vee q) \wedge (\neg p \vee r) \wedge (\neg q \vee r)] \vee r & (3) \\
& [(\neg p \wedge \neg q) \vee (p \wedge \neg r) \vee (q \wedge \neg r)] \vee r & (4) \\
& [((\neg p \vee p) \wedge (\neg p \vee \neg r)) \vee (q \wedge \neg r)] \vee r & (5) \\
& [((T \wedge (\neg p \vee \neg r)) \vee (q \wedge \neg r)] \vee r & (6) \\
& [\neg p \vee \neg r \vee (q \wedge \neg r)] \vee r & (7) \\
& \neg p \vee (q \wedge \neg r) \vee (\neg r \vee r) & (8) \\
& \neg p \vee (q \wedge \neg r) \vee T & (9) \\
& T & (10)
\end{aligned}$$

24.

$p \leftrightarrow q$ is a true statement only when they are the same truth value. On the other hand, $p \oplus q$ is only true when they are the opposite truth value. Negating the XOR would mean that these two are logically equivalent.

30.

$$\neg p \rightarrow (q \rightarrow r) \equiv p \vee \neg q \vee r \quad (1)$$

$$\equiv \neg q \vee p \vee r \quad (2)$$

$$q \rightarrow (p \vee r) \equiv \neg q \vee p \vee r \quad (3)$$

$$\therefore \neg p \rightarrow (q \rightarrow r) \equiv q \rightarrow (p \vee r) \quad (4)$$

32.

$p \leftrightarrow q$ is a true statement only when they are the same truth value. Even if you negate both p and q , the statement $\neg p \leftrightarrow \neg q$ is true when they are the same truth value.

p	q	$\neg p$	$\neg q$	$p \leftrightarrow q$	$\neg p \leftrightarrow \neg q$
F	F	T	T	T	T
F	T	T	F	F	F
T	F	F	T	F	F
T	T	F	F	T	T