

Parallel Programming

Parallelism

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Why Parallelism?

- > High performance computing
 - Solve large problems faster (or be able to solve them in time at all)
 - Solve larger problems in the same time
- Enterprise computing
 - Increase throughput
- Computing in general
 - Satisfy an increasing performance demand while avoiding physical limitations (frequency, power, ...)
 - > Save money (lots of small chips are cheaper than a complex one)
 - Save energy (small or specialized chips are more energy efficient for certain tasks)

Perilin

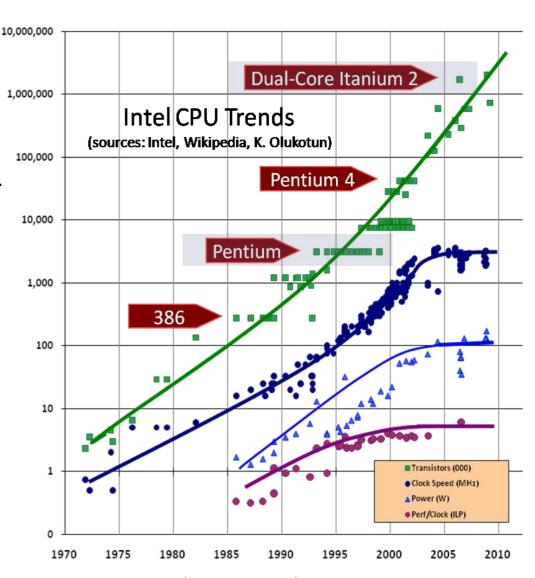
Limits in Processors

Power Consumption

- Limited to about 150W for air cooling
- Higher frequency needs higher core voltages -> more power

> Memory

- Upon a cache miss, execution might stall (despite Instruction Level Parallelism)
- Instruction Level Parallelism
 - It is hard to extract more ILP out of (general) programs



Source: The Free Lunch Is Over, H. Sutter



Power Wall

> Power consumption of a processor

$$P_{dynamic} \sim \alpha C \cdot f \cdot V^{2}$$

$$P_{static} \sim V \cdot I_{leak}$$

- > Frequency and voltage are not independent
 - > Transistor switching speed depends on voltage
 - > Critical path must settle during a clock cycle
- > Lower frequencies are more energy efficient

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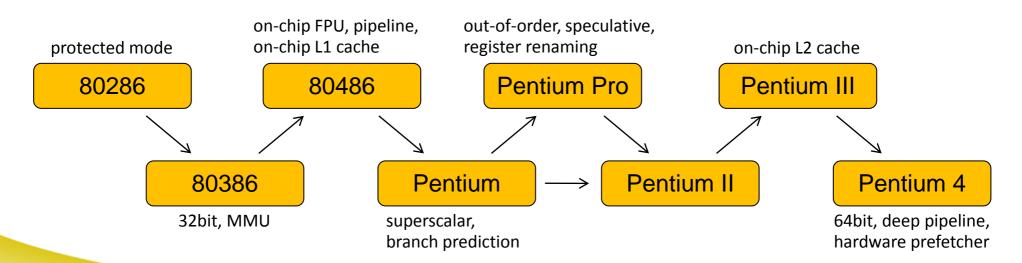
Memory Wall

- > Memory is slow
 - E. g. a read latency of up to 190 processor cycles on an Intel Xeon X5570
- Various efforts to hide that
 - > Multi-level cache hierarchies
 - > Lower latencies (e. g. 4 cycles for L1, 10 for L2, 38 for L3 cache)
 - > Memory prefetching
 - > Out-of-order execution
 - > Execute something else while waiting
- Still, the performance of many tasks does not scale linearly with the processor frequency



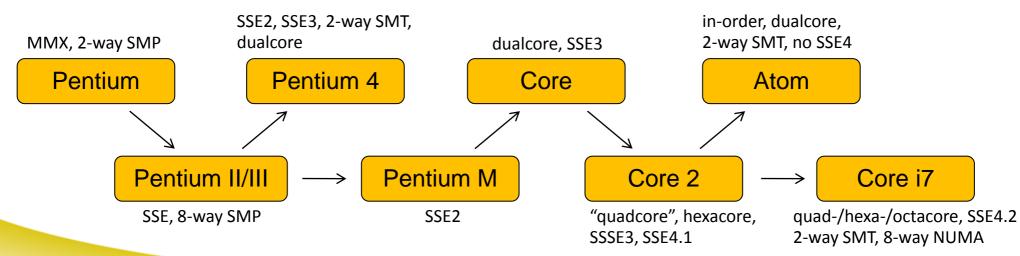
Historical Processor Development (i)

- > Starting with (relatively) simple processors
- Increase clock frequency
- Increase IPC (Instructions Per Cycle)
 - > By adding on-chip memory
 - > By increasing ILP (Instruction Level Parallelism)
- Increase word size (Bit Level Parallelism)



Historical Processor Development (ii)

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- Parallelism was introduced at various locations in various combinations
 - > Multisocket systems (first SMP, later NUMA architectures)
 - > Multicore processors
 - > Simultaneous Multithreading (SMT)
 - > SIMD units



Typical Hardware

Network of Computers

(Cluster, Grid, Internet, ...)

Computer

(Parallel, Server, Desktop, Mobile, Embedded, ...)

Processor

(CPU, GPU, ...)

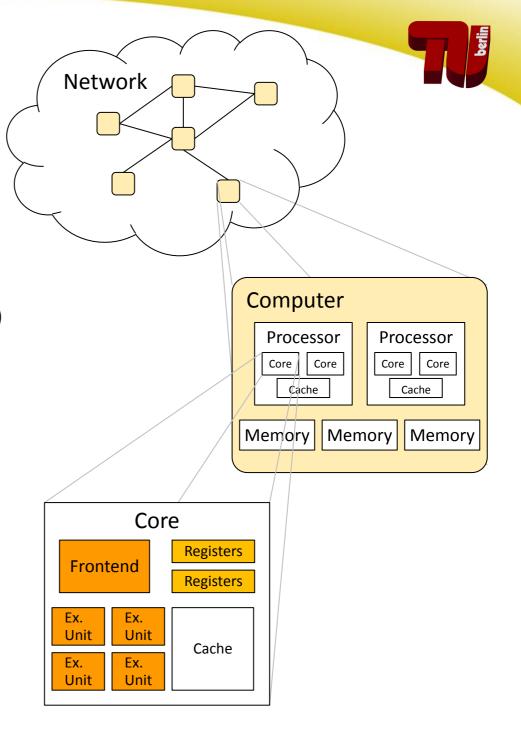
Core

(single threaded/SMT, in-order/out-of-order, ...)

Execution Unit

(Integer, Floating Point, SIMD, ...)

Bits



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Levels of Parallelism

- > Bit level parallelism
 - > Parallel processing of bits of a word
- > Instruction level parallelism
 - > Parallel processing of instructions of the same task
- > Data level parallelism
 - > Identical operations on sets of data
- > Task level parallelism
 - > Independent blocks of operations



Typical "Desktop" Software

- Mostly single threaded software
 - > Threads only to provide responsiveness

> Seldom optimized for target architecture

- Only special purpose software is aware of the available parallelism, e. g.
 - > Multimedia applications
 - > Number crunching applications



Parallelism in Sequential Software

- > Sequential software is already (a little bit) parallel
 - > Realized by compiler and/or hardware
- > Bit Level Parallelism
- Instruction Level Parallelism (ILP)
 - > To increase instruction throughput
 - > Many different techniques
 - > Pipelining, superscalar processors, out-of-order execution, speculative execution, ...
- Data Level Parallelism
 - > SIMD instructions to speed up some applications



Parallelism in Parallel Programs

Parallel programs exploit more parallelism than sequential programs!

- Data Level Parallelism
 - Identical/similar operations on data structures are captured on a larger scale and carried out in parallel.

- > Task Level Parallelism
 - > Independent blocks of operations are carried out in parallel.



New Problems with Parallelism

- You have to care about
 - > Problem decomposition
 - Coordination of interacting processes/threads
 - > Correctness, deadlocks, race conditions
 - > Via shared memory or message passing
 - > Scalability
 - > Environment
 - Interconnection network (latency, bandwidth, topology)
 - > Heterogeneity
 - > Exclusiveness



Means to create Parallel Programs

- Parallelizing compilers
 - Pro: compile sequential code without changes into a parallel program
 - > Con: very hard to do efficiently, limited applicability
- Parallel programming languages
 - > Pro: easy to exploit parallelism
 - > Con: low acceptance, compiler support needed
- > Augmented sequential programming languages
 - > Pro: higher acceptance
 - > Con: also needs compiler support
- Libraries enabling parallelism
 - > Pro: high acceptance, no compiler support needed
 - > Con: tends to be a little complex

Parallel Programming Environments from the 1990s



"C* in C ABCPL **ACE** ACT++ **ADDAP** Adl Adsmith **AFAPI ALWAN** AM **AMDC** Amoeba **AppLeS** ARTS Athapascan-0b Aurora **Automap** bb threads Blaze BlockComm **BSP C*** Č** C4 CarlOS Cashmere

COOL **CORRELATE** CparPar CPS **CRL CSP Cthreads CUMULVS DAGGER DAPPLE** Data Parallel C DC++ DCE++ DDD DICE DIPC Distributed Smalltalk JAVAR **DOLIB DOME DOSMOS** DRL **DSM-Threads** Ease **ECO** Eilean **Emerald EPL** Excalibur **Express** Falcon **Filaments FLASH** FM Fork Fortran-M FX

GA

GAMMA Glenda GLU **GUARD HAsL HORUS HPC** HPC++ **HPF IMPACT** ISETL-Linda ISIS **JADA JADE** Java RMI iavaPG **JavaSpaces** JIDL **Joyce** Karma Khoros **KOAN/Fortran-S** LAM Legion Lilac Linda LiPS Locust Lparx Lucid Maisie Manifold Mentat Meta Chaos

Midway

Millipede

Mirage Modula-2* Modula-P **MOSIX** DαM MPC++ MPI Multipol Munin Nano-Threads NESL NetClasses++ Nexus Nimrod NOW **Objective Linda** Occam Omega **OOF90** Orca P++ P-RIO P3L P4-Linda Pablo **PADE PADRE** Panda **Papers** Para++ Paradigm Parafrase2 **Paralation Parallaxis** Parallel Haskell

ParLin **Parlog Parmacs** Parti рC pC++ PCN PCP: PCU **PEACE PENNY** PET **PETSc** PH **Phosphorus** POET **Polaris POOL-T POOMA POSYBL PRESTO Prospero Proteus PSDM** PSI **PVM** QPC++ Quake Quark **Quick Threads** Sage++ SAM SCANDAL **SCHEDULE**

SciTL **SDDA**

ParLib++

SHMEM SIMPLE Sina **SISAL** SMI SONIC Split-C SR **Sthreads** Strand **SUIF** SuperPascal Synergy TCGMSG **Telegraphos** The FORCE Threads.h++ **TRAPPER TreadMarks** UC uC++ UNITY V ViC* Visifold V-NUS **VPF** Win32 threads

Parallel-C++

ParC

CC++

Chu

Cid

Cilk

Code

Charlotte

Charm++

CM-Fortran

Converse

Concurrent ML

Charm

WinPar

Zounds

XPC

ZPL

WWWinda

XENOOPS

Today's Parallel Programming Environments – a subjective list



- > Clusters
 - > Established: MPI
- > Multi-core systems
 - > Established: explicit threading
 - > Java, C#, pthreads, Win32 threads, ...
 - > Upcoming: OpenMP
- > GPUs
 - > Established: CUDA
 - > Upcoming: OpenCL