

# Cutland Computability Exercises

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May 15th, 2024

## 1 Chapter 1 Exercises

**Exercise 1.1** (2.2).

8, 4, 2, 0, 0, 0,  $\dots$

$I_1$

8, 4, 2, 0, 0, 0,  $\dots$

$I_2$

8, 5, 2, 0, 0, 0,  $\dots$

$I_3$

8, 5, 3, 0, 0, 0,  $\dots$

$I_4$ , don't jump

8, 5, 3, 0, 0, 0,  $\dots$

$I_5$ , jump to 2

8, 5, 3, 0, 0, 0,  $\dots$

$I_2$

8, 6, 3, 0, 0, 0,  $\dots$

$I_3$

8, 6, 4, 0, 0, 0,  $\dots$

$I_4$ , don't jump

8, 6, 4, 0, 0, 0,  $\dots$

$I_5$ , jump to 2

8, 6, 4, 0, 0, 0,  $\dots$

$I_2$   
 $8, 7, 4, 0, 0, 0, \dots$   
 $I_3$   
 $8, 7, 5, 0, 0, 0, \dots$   
 $I_4$ , don't jump  
 $8, 7, 5, 0, 0, 0, \dots$   
 $I_5$ , jump to 2  
 $8, 7, 5, 0, 0, 0, \dots$   
 $I_2$   
 $8, 8, 5, 0, 0, 0, \dots$   
 $I_3$   
 $8, 7, 6, 0, 0, 0, \dots$   
 $I_4$ , jump to 6  
 $8, 7, 6, 0, 0, 0, \dots$   
 $I_6$   
 $6, 7, 6, 0, 0, 0, \dots$

**Exercise 1.2 (2.3).** We proceed with a proof by induction. After executing each instruction, we want to show that  $r_1 < r_2$  and the program does not terminate. Consider the first executed instruction,  $J(1, 2, 6)$ . Since  $2 < 3$ , we do not jump and  $r_1 < r_2$ . Suppose the program has not terminated after instruction  $n$ , and  $r_1 < r_2$ . Consider instruction  $n + 1$ . The instruction must be  $S(2)$ ,  $S(3)$ ,  $J(1, 2, 6)$ ,  $J(1, 1, 2)$ ,  $T(3, 1)$ .

Case 1:  $S(2)$ . Then we have that  $r_1 < r_2 < r_2 + 1$ . There is a next instruction, so this case is solved.

Case 2:  $S(3)$ . Since  $r_1, r_2$  are not affected, it remains that  $r_1 < r_2$ . The next line is instruction 4, so the program does not terminate.

Case 3:  $J(1, 2, 6)$ . Since  $r_1 < r_2$ , by the induction hypothesis, we do not jump, and instead run instruction 5. It remains that  $r_1 < r_2$ .

Case 4:  $J(1, 1, 2)$ . Since jumps do not affect our registers,  $r_1 < r_2$  remains. In addition, we jump to instruction 2, so the program does not terminate.

Case 5:  $T(3, 1)$ . For the instruction to have fired, we must be on line 6. Since line 5 always jumps to line 2, it is the case that we must have executed  $J(1, 2, 6)$  previously. However,  $r_1 < r_2$  for the previous step, so this case is invalid.

Thus after executing  $n + 1$  instructions, the program does not terminate. Thus the program will never terminate.

**Exercise 1.3** (3.3.1a).

$$I_1 : J(1, 2, 5)$$

Checks if  $x = 0$

$$I_2 : S(2)$$

If  $x \neq 0$ , add one

$$I_3 : T(2, 1)$$

Move 1 or 0 into the output register

**Exercise 1.4** (3.3.1b).

$$I_1 : S(2)$$

$$I_2 : S(2)$$

$$I_3 : S(2)$$

$$I_4 : S(2)$$

$$I_5 : S(2)$$

$$I_6 : T(2, 1)$$

This one is self-explanatory

**Exercise 1.5** (3.3.1c).

$$I_1 : J(1, 2, 3)$$

Check if  $x = y$

$$I_2 : S(3)$$

If  $x \neq y$ , add one to register 3

$$I_3 : T(3, 1)$$

Move 0 or 1 to the output register

**Exercise 1.6** (3.3.1d).

$$I_1 : J(1, 2, 9)$$

Check to see if  $x = y$

$$I_2 : T(1, 3)$$

Move  $x$  to register 3

$$I_3 : T(2, 4)$$

Move  $y$  to register 4

$$I_4 : J(1, 2, 9)$$

$$I_5 : J(3, 4, 12)$$

Check to see if  $y + k = x$  or  $x + k = y$

$$I_6 : S(1)$$

$$I_7 : S(4)$$

Add one to  $x$  and  $y$

$$I_8 : J(1, 1, 4)$$

Jump back to the checks

$$I_9 : T(5, 1)$$

$$I_{10} : J(1, 1, 14)$$

Put 0 in the output register

$$I_{12} : S(5)$$

$$I_{13} : T(5, 1)$$

Put 1 in the output register

**Exercise 1.7** (3.3.1e).

$$I_1 : J(1, 3, 6)$$

$$I_2 : S(2)$$

$$I_3 : S(3)$$

$$I_4 : S(3)$$

$$I_5 : S(3)$$

$$I_6 : T(2, 1)$$

**Exercise 1.8** (3.3.1f).

$$I_1 : J(1, 2, 6)$$

$$I_2 : S(2)$$

$$I_3 : S(3)$$

$$I_4 : S(3)$$

$$I_5 : J(1, 1, 1)$$

$$I_6 : T(3, 1)$$

Getting  $2X$  into the output register

$$I_7 : Z(2)$$

$$I_8 : Z(3)$$

$$I_9 : S(2)$$

$$I_{10} : S(3)$$

$$I_{11} : S(3)$$

Setting up for  $3k + 1, 3k + 2, 3k$

$$I_{12} : J(1, 4, 25)$$

$$I_{13} : S(2)$$

$$I_{14} : S(2)$$

$$I_{15} : S(2)$$

$$I_{16} : S(3)$$

$$I_{17} : S(3)$$

$$I_{18} : S(3)$$

$$I_{19} : S(4)$$

$$I_{20} : S(4)$$

$$I_{21} : S(4)$$

$$I_{22} : J(1, 2, 25)$$

$$I_{23} : J(1, 2, 25)$$

$$I_{23} : J(1, 2, 25)$$

If any are hit, it means the floor is found in register 4

$$I_{24} : J(1, 1, 13)$$

$$I_{25} : T(4, 1)$$

Transfer register 4 to output

**Exercise 1.9** (3.3.2).

$$f_P^{(2)}(x, y) = \begin{cases} x - y & x > y \\ \text{undefined} & \text{o/w} \end{cases}$$

**Exercise 1.10** (3.3.3). We proceed via proof by induction. We want to prove that after executing  $n$  instructions, for all registers  $R_k$ , we have that either  $r_k = m$  or  $r_k = x + a$ . Since there are no jump instructions, we execute the same number instructions no matter what the input is. Consider  $n = 0$ . We have that  $r_1 = x = x + 0$  and  $r_b = 0$ , for  $b \neq 1$ . Suppose after executing  $n$  instructions, for all  $k \in \mathbb{N}$ , we have that  $r_k = m_k$  or  $r_k = x + a_k$ , for some  $m_k$  and  $a_k$ . Consider the  $n + 1$ th instruction.

Case 1: It is  $Z(s)$  for some  $s \in \mathbb{N}$ . Then  $r_s = 0$  and the rest of the registers remain the same. Thus this case is closed.

Case 2: It is  $S(l)$  for some  $l \in \mathbb{N}$ . Then by the induction hypothesis, either  $r_l = m_l$  or  $r_l = x + a_l$ . Then after applying the successor operation,  $r_1 = m_l + 1$ , still a constant or  $r_l = x + (a_l + 1)$ , still in the same form. The other registers remain the same.

Case 3:  $T(p, o)$ . Now  $r_o = r_p$ , where  $r_p = m_p$  or  $r_p = x + a_p$ . Thus we have proven the induction step for all three cases.

Thus when the program terminates, the value at the output register is either  $m$  or  $x + a$  for some  $m, a \in \mathbb{N}$ . Since the value of  $x$  was arbitrary, it must be the case that  $f(x) = m$  for all  $x \in \mathbb{N}$  or  $f(x) = x + a$  for all  $x \in \mathbb{N}$ .

**Exercise 1.11** (3.3.4).

$$I_1 : Z(n)$$

$$I_2 : J(n, m, 5)$$

$$I_3 : S(n)$$

$$I_4 : J(1, 1, 2)$$

**Exercise 1.12** (4.3b).

$$I_1 : J(1, 2, 10)$$

$$I_2 : T(1, 3)$$

$$I_3 : T(2, 4)$$

$$I_4 : J(1, 2, 9)$$

$$I_5 : J(3, 4, 10)$$

$$I_6 : S(1)$$

$$I_7 : S(4)$$

$$I_8 : J(1, 1, 4)$$

$$I_9 : S(5)$$

$$I_{10} : T(5, 1)$$

**Exercise 1.13** (4.3b).

$$I_1 : S(2)$$

$$I_2 : S(2)$$

$$I_3 : S(2)$$

$$I_4 : J(1, 2, 6)$$

$$I_5 : S(3)$$

$$I_6 : T(3, 1)$$

**Exercise 1.14** (4.3c).

$$\begin{aligned}
 I_1 &: S(2) \\
 I_2 &: J(1, 2, 10) \\
 I_3 &: J(1, 3, 9) \\
 I_4 &: S(2) \\
 I_5 &: S(2) \\
 I_6 &: S(3) \\
 I_7 &: S(3) \\
 I_8 &: J(1, 1, 2) \\
 I_9 &: S(4) \\
 I_{10} &: T(4, 1)
 \end{aligned}$$

**Exercise 1.15** (5.2.1).

$$f \circ \alpha^{-1}(m) = \begin{cases} m & \text{if } m \text{ is even,} \\ -(m+1) & \text{if } m \text{ is odd.} \end{cases}$$

$$\alpha \circ f \circ \alpha^{-1}(m) = \begin{cases} 2m & \text{if } m \text{ is even,} \\ 2m+1 & \text{if } m \text{ is odd.} \end{cases}$$

The set of instructions is as follows:

$$\begin{aligned}
 I_1 &: J(1, 4, 6) \\
 I_2 &: S(4) \\
 I_3 &: S(5) \\
 I_4 &: S(5) \\
 I_5 &: J(1, 1, 1) \\
 I_6 &: S(2) \\
 I_7 &: J(1, 2, 14) \\
 I_8 &: J(1, 3, 15) \\
 I_9 &: S(2) \\
 I_{10} &: S(2) \\
 I_{11} &: S(3)
 \end{aligned}$$

$$I_{12} : S(3)$$

$$I_{13} : J(1, 1, 7)$$

$$I_{14} : S(5)$$

$$I_{15} : T(5, 1)$$

**Exercise 1.16.** (5.2.2)

$$f \circ \alpha^{-1}(m) = \begin{cases} 1 & \text{if } m \text{ is even,} \\ 0 & \text{if } m \text{ is odd.} \end{cases}$$

$$\alpha \circ f \circ \alpha^{-1}(m) = \begin{cases} 2 & \text{if } m \text{ is even,} \\ 1 & \text{if } m \text{ is odd.} \end{cases}$$

The set of instructions is as follows:

$$I_1 : S(2)$$

$$I_2 : J(1, 2, 10)$$

$$I_3 : J(1, 3, 9)$$

$$I_4 : S(2)$$

$$I_5 : S(2)$$

$$I_6 : S(3)$$

$$I_7 : S(3)$$

$$I_8 : J(1, 1, 2)$$

$$I_9 : S(4)$$

$$I_{10} : S(4)$$

$$I_{11} : T(4, 1)$$