

COL216

Computer Architecture

Pipelined processor : Handling hazards

24th February, 2022

Executing branch instructions

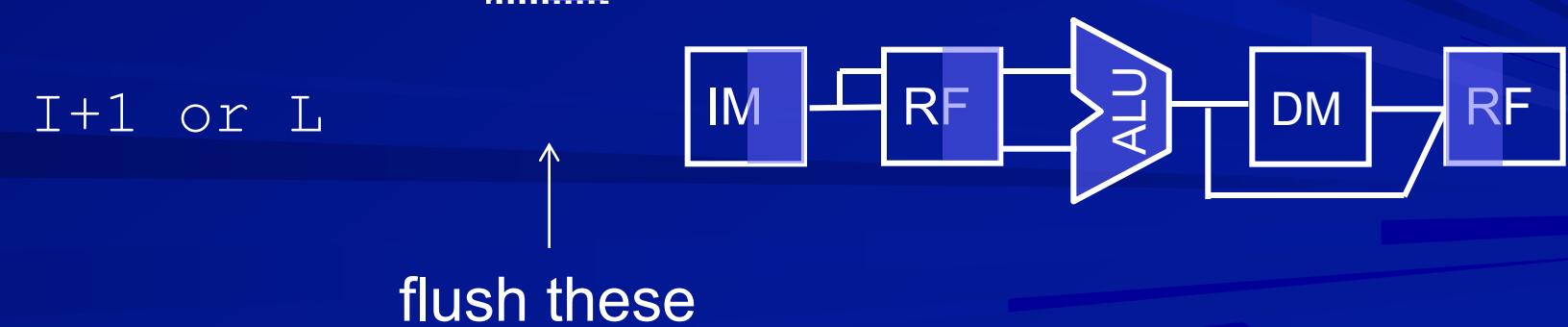
- In which cycle the instruction is found to be a branch instruction? 2
- In which cycle the branch decision is known? 2 or 3
- In which cycle the target address is computed? 3

example

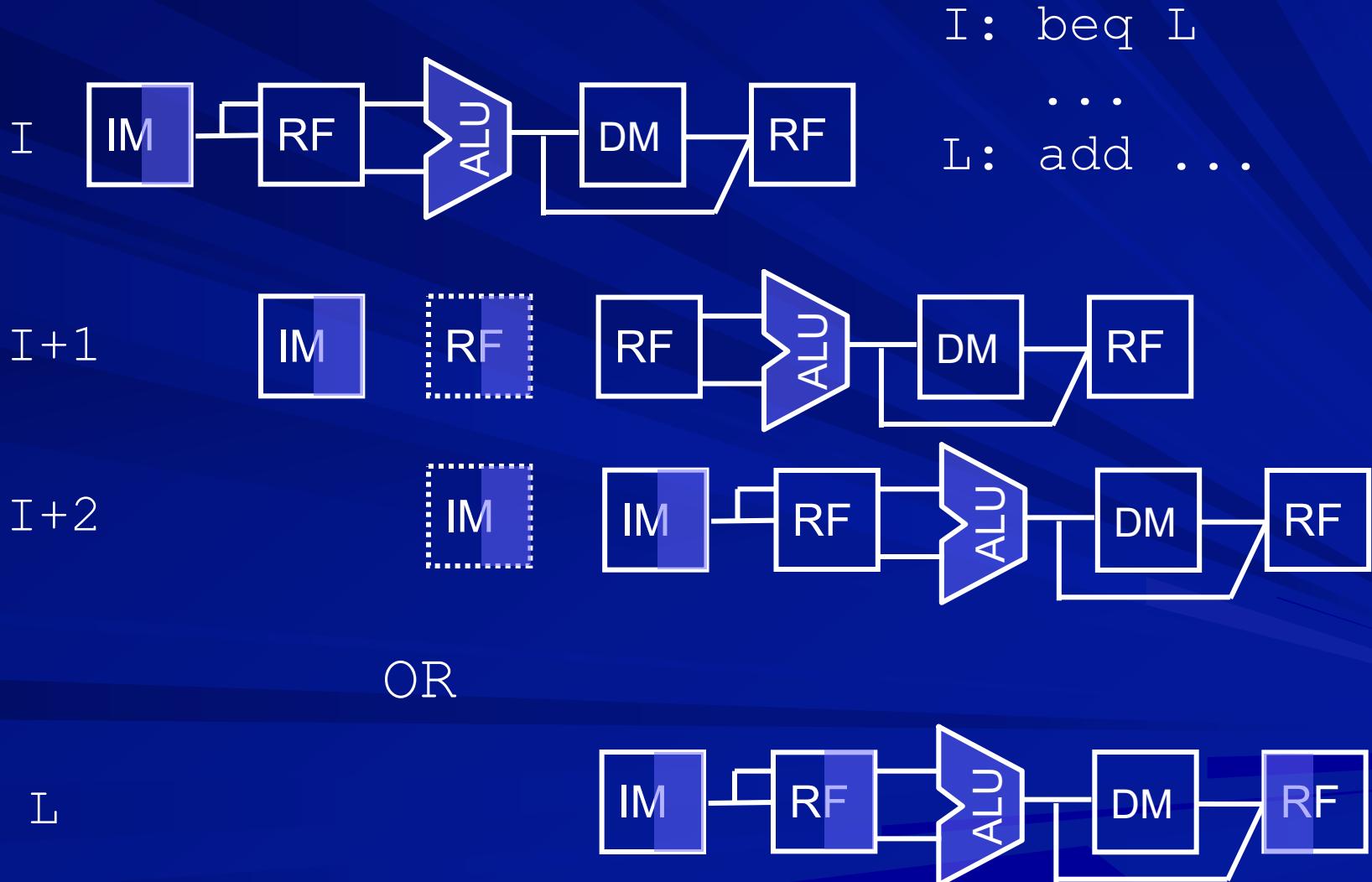
Handling control hazards

- Flush the inline instructions
- Freeze (stall) the inline instructions
- Allow the inline instructions to continue
- Delayed branch
- Predict the branch decision
- Predict the target address

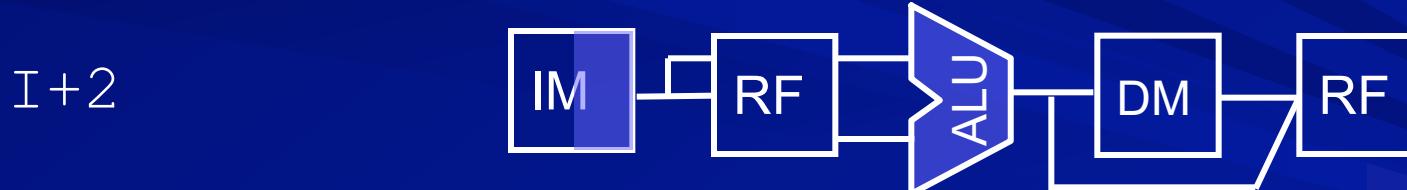
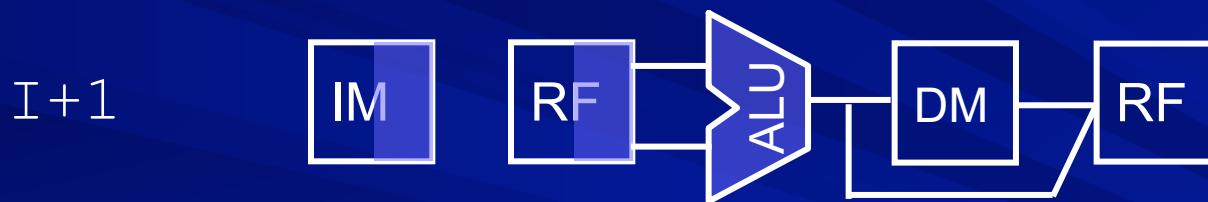
Flush inline instructions



Freeze inline instructions



Allow inline instructions

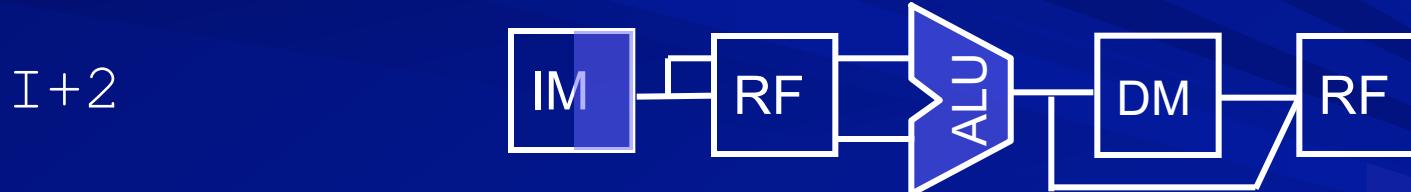


OR

L



Delayed branch



OR

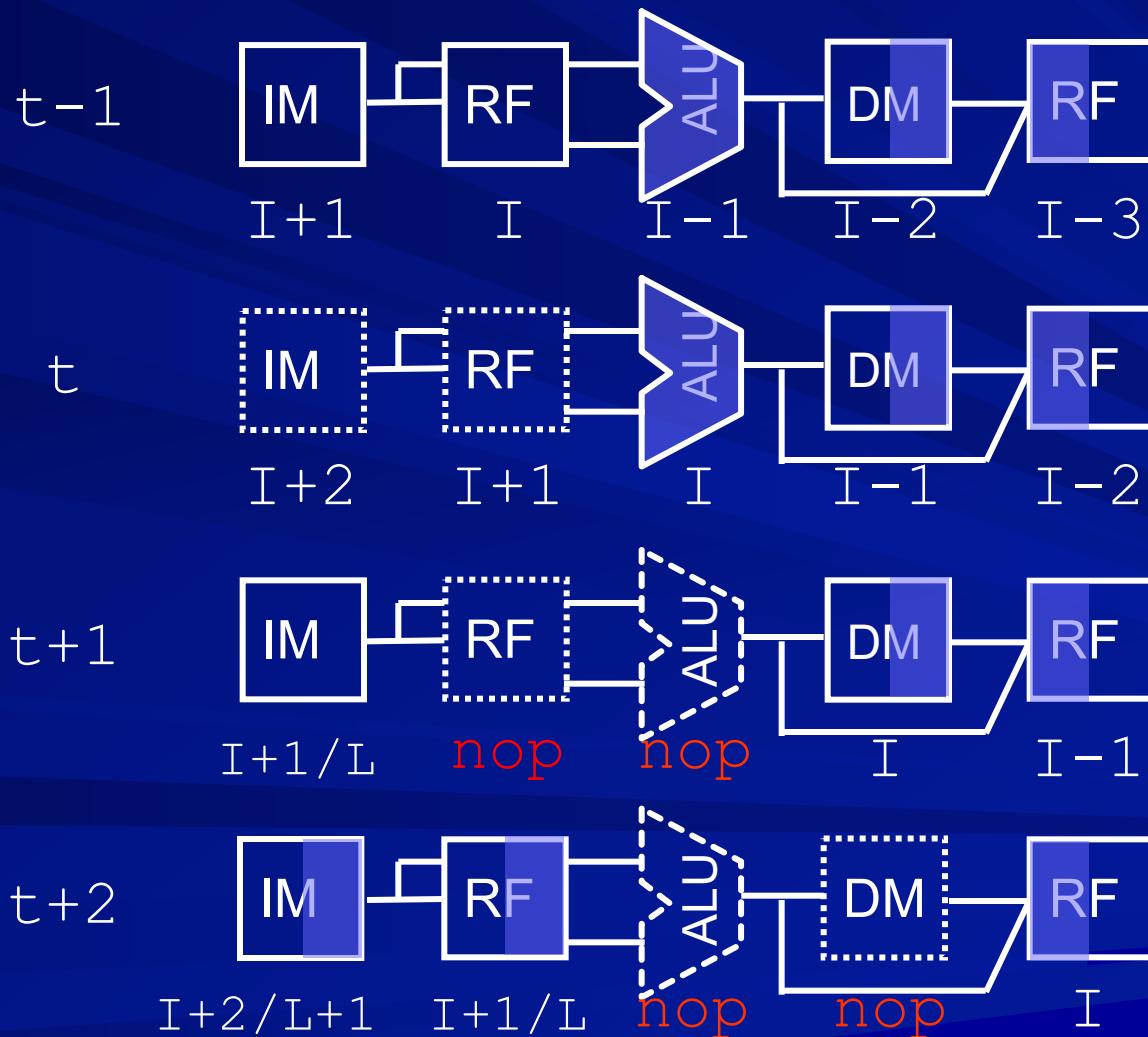
L



Stalls due to control hazards

flush: stage-wise view

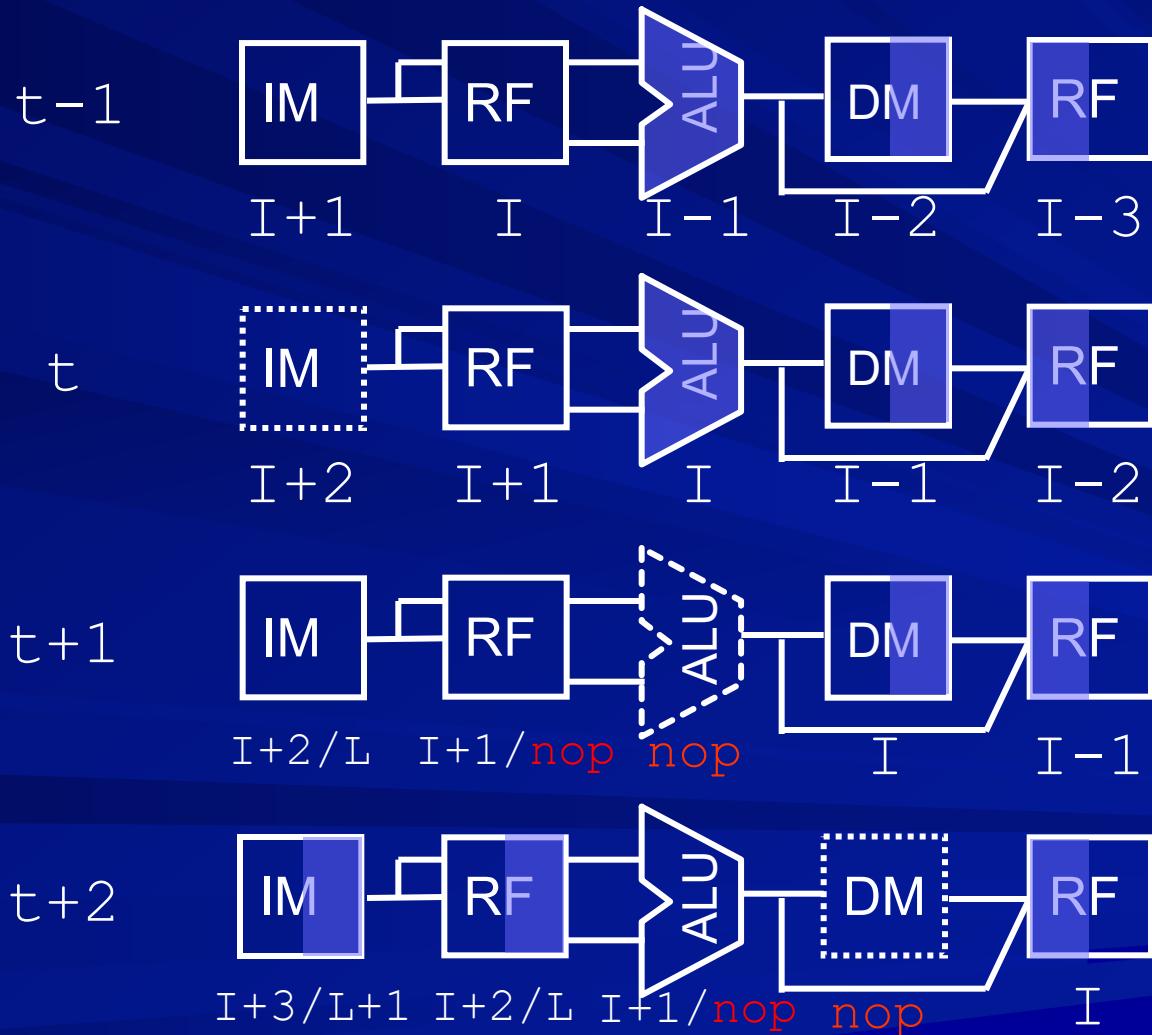
I: beq L



Stalls due to control hazards

freeze: stage-wise view

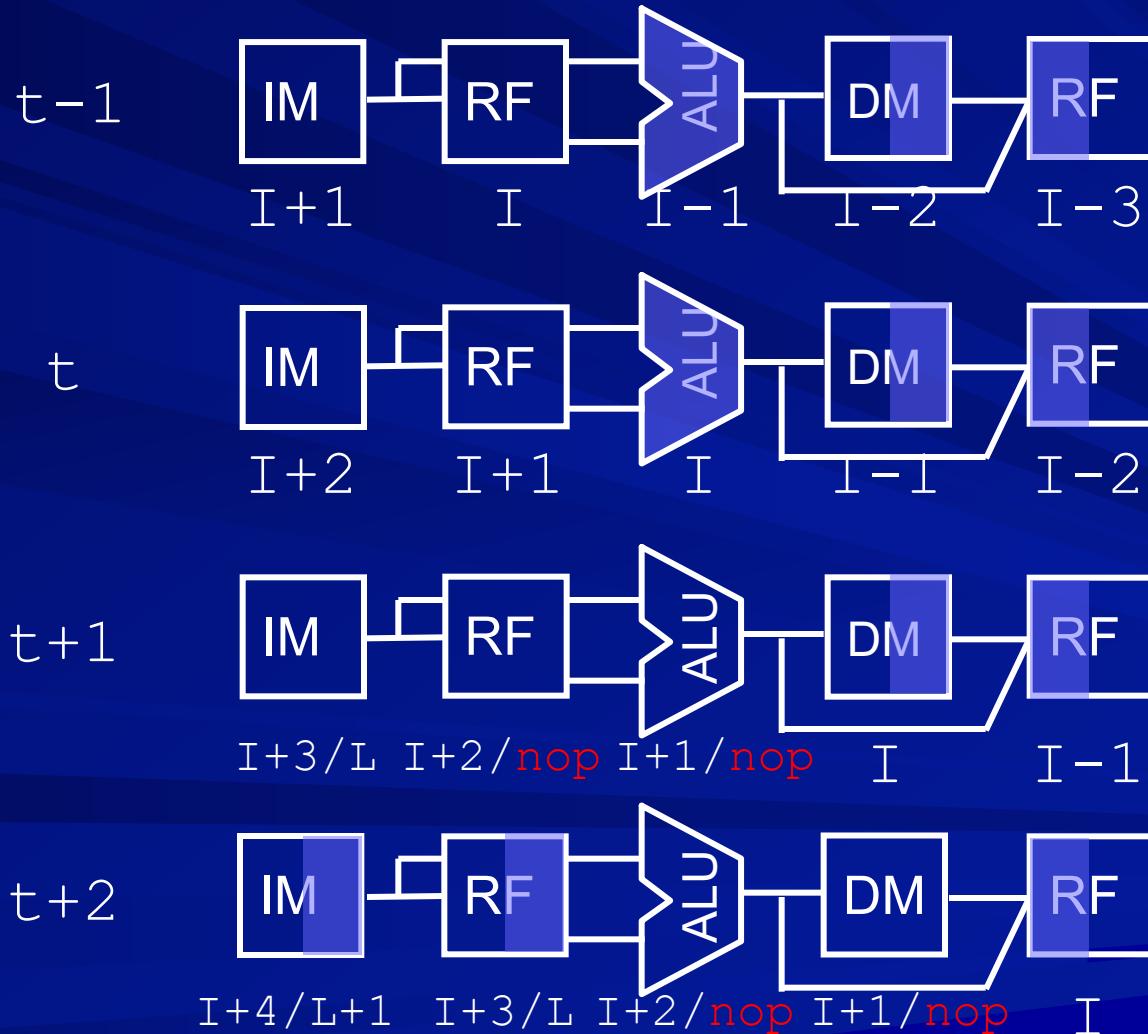
I: beq L



Stalls due to control hazards

continue: stage-wise view

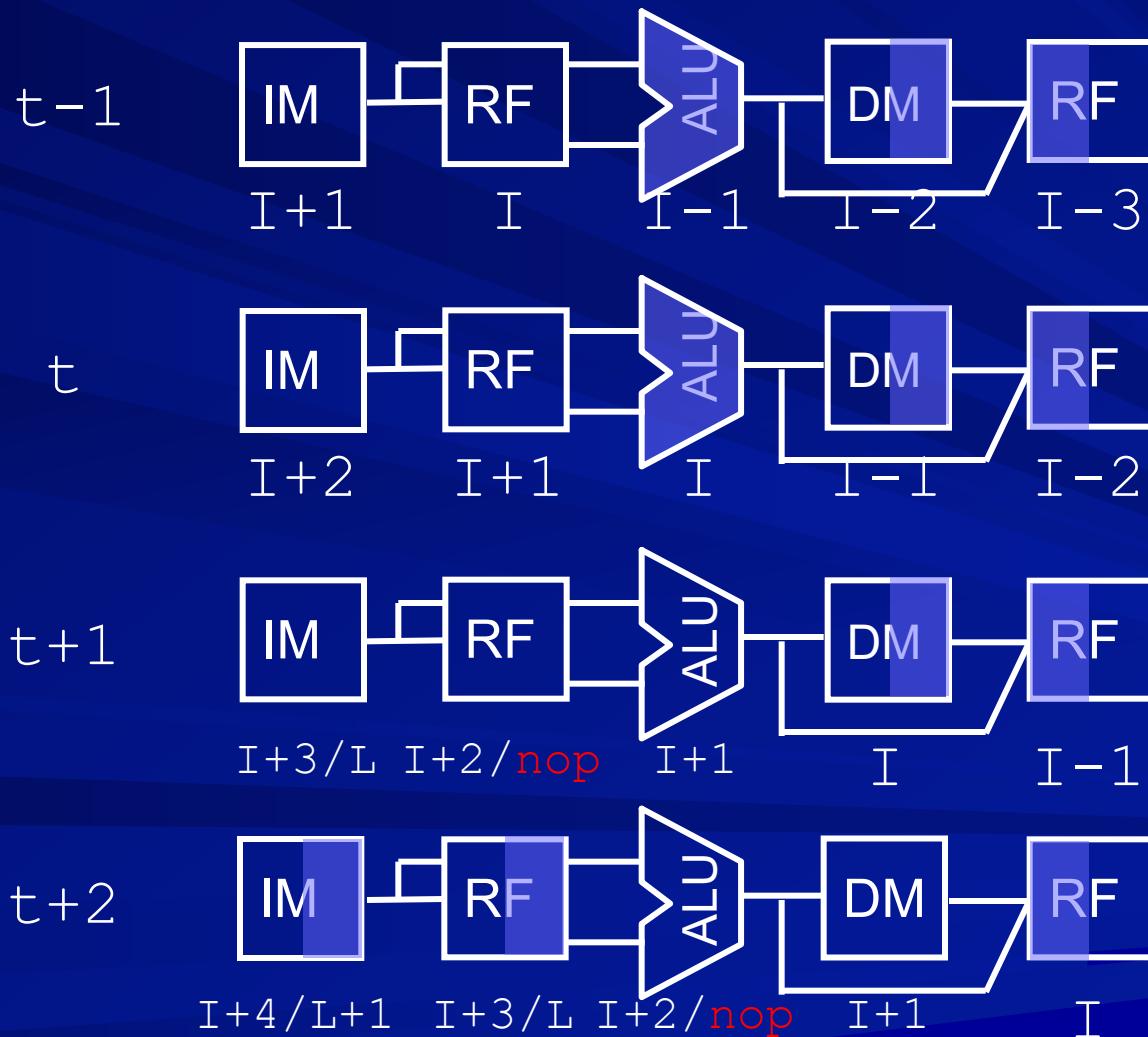
I: beq L



Stalls due to control hazards

delayed branch: stage-wise view

I: beq L



Branch Prediction

- Treat conditional branches as unconditional branches / NOP
- Undo if necessary

Strategies:

- Static
- Dynamic

Branch Prediction

- Fixed
 - *always predict inline*
- Static
 - *predict on the basis of instruction type, target address or profiling information*
- Dynamic
 - *predict based on recent history*

Dynamic Branch Prediction - basic idea

Predict based on the history of previous
branch

loop: xxx

xxx

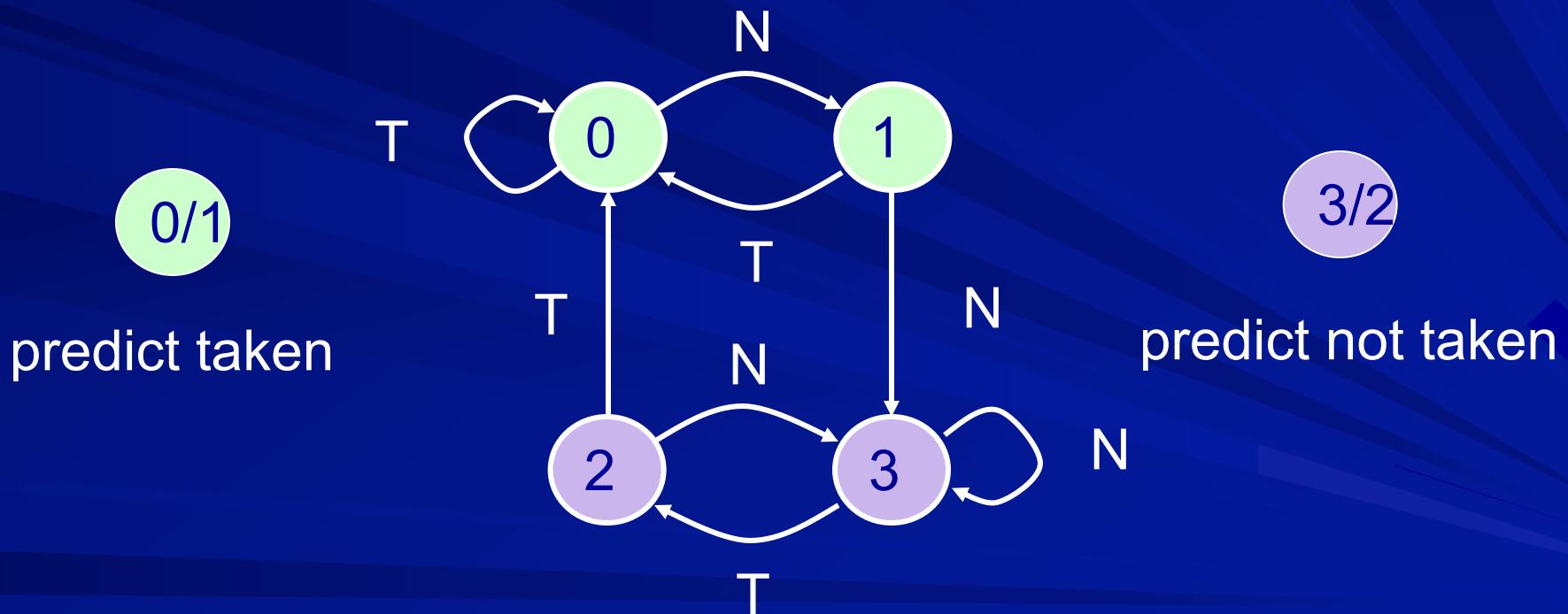
xxx

xxx

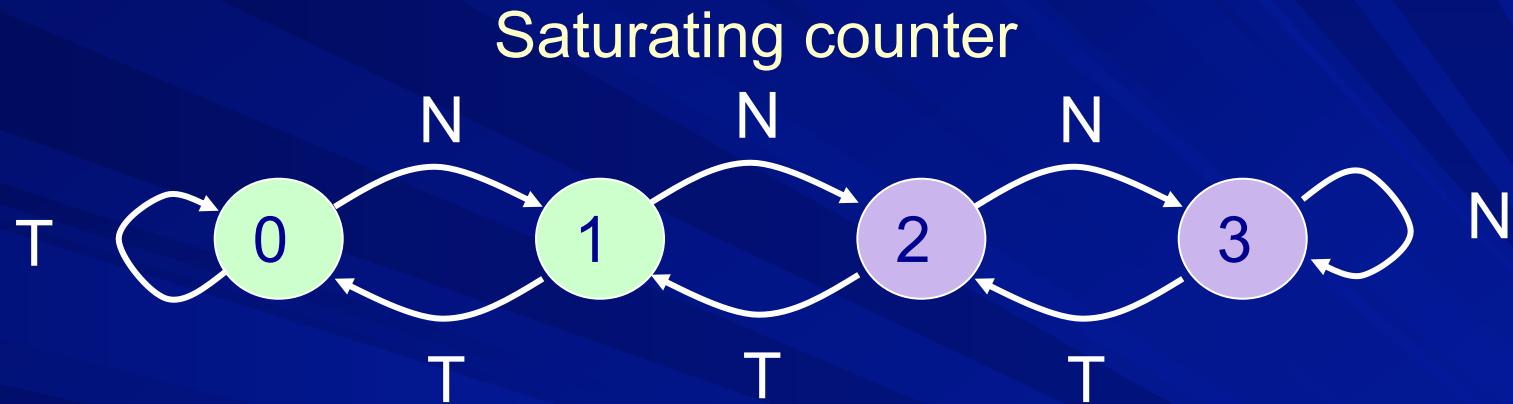
b<cond> loop

2 mispredictions
for every
occurrence of
the loop

Dynamic Branch Prediction - A 2-bit prediction scheme



Another 2-bit prediction scheme



0/1

predict taken

3/2

predict not taken

Dynamic information about branch

- Previous branch decisions
- Previous target address or target instruction
- Stored in
 - Cache
 - Separate buffer

Branch History Table (BHT)

Branch Target Buffer (BTB)

Target Instruction Buffer (TIB)

Branch Target Buffer



- hit \Rightarrow explicit prediction using prediction bits
 - prediction \rightarrow go target (use target info)
 - \rightarrow go inline (ignore target info)
- miss \Rightarrow go inline

Accessing BTB

- In which cycle do you access BTB?
- Before checking condition?
- Before address computation?
- Just after decoding?
- Along with instruction fetch?

Correlation between branches

B1: if (x)

...

B2: if (y)

...

$z = x \&& y$

B3: if (z)

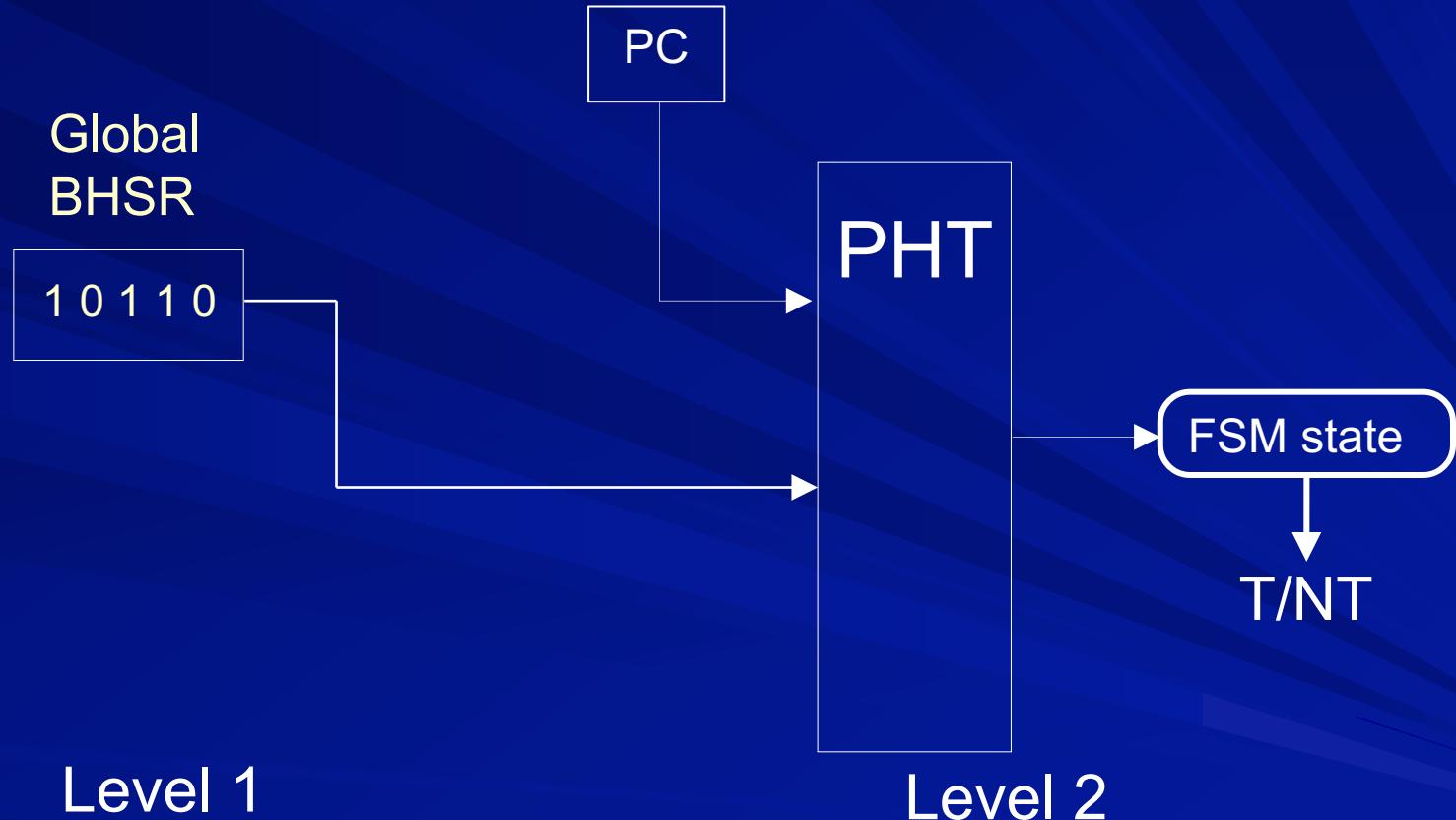
...

- B3 can be predicted with 100% accuracy based on the outcomes of B1 and B2

Two-Level Branch Predictors

- Level 1
 - Branch History Shift Register (BHSR) - last n occurrences
 - Captures patterned behavior of groups of branches
- Level 2
 - Pattern History Table (PHT) - states of predictor FSMs
 - Captures behavior of individual branches

A two-level branch predictor



Bits from PC and BHSR are combined to index PHT

Optimizing programs for pipeline

- Instruction reordering
- Moving instructions across branches
- Delayed branches
- Predication

Running program on pipeline

```
    mov r1, #1
    mov r2, #100
L1:   str r1, [r2, #0]
        add r2, r2, #4
        add r1, r1, #1
        cmp r1, #11
        bne L1
        mov r3, #0
        mov r2, #100
L2:   sub r1, r1, #1
        cmp r1, #0
        beq Over
        ldr r4, [r2, #0]
        add r3, r3, r4
        add r2, r2, #4
        b L2
```

Over:

Running program on pipeline

Without stalls

	mov	r1, #1	1		
	mov	r2, #100	2		
L1:	str	r1, [r2, #0]	3	8	.. 48
	add	r2, r2, #4	4		
	add	r1, r1, #1	5		
	cmp	r1, #11	6		
	bne	L1	7	12	.. 52
	mov	r3, #0	53		
	mov	r2, #100	54		
L2:	sub	r1, r1, #1	55	62	.. 118
	cmp	r1, #0	56		
	beq	Over	57		
	ldr	r4, [r2, #0]	58		
	add	r3, r3, r4	59		
	add	r2, r2, #4	60		
	b	L2	61	68	.. 124
Over:			128		

Running program on pipeline

L1:	mov r1, #1	1	
	mov r2, #100	2	
	str r1, [r2, #0]	5	14 .. 86
	add r2, r2, #4	6	
	add r1, r1, #1	7	
	cmp r1, #11	10	
	bne L1	11	20 .. 92
	mov r3, #0	95	
	mov r2, #100	96	
L2:	sub r1, r1, #1	97	112.. 232
	cmp r1, #0	100	
	beq Over	101	
	ldr r4, [r2, #0]	104	
	add r3, r3, r4	107	
	add r2, r2, #4	108	
	b L2	109	124..244
Over:		254	

With stalls

Over:

Execute inline instructions after branch

L1:	mov r1, #1	1		
	mov r2, #100	2		
	str r1, [r2, #0]	5	14 .. 86	
	add r2, r2, #4	6		
	add r1, r1, #1	7		
	cmp r1, #11	10		
	bne L1	11	20 .. 92	
	mov r3, #0	95		
	mov r2, #100	96		
L2:	sub r1, r1, #1	97	110.. 214	
	cmp r1, #0	100		
	beq Over	101		
	ldr r4, [r2, #0]	102		
	add r3, r3, r4	105		
	add r2, r2, #4	106		
	b L2	107	120..224	
		234		

Note: this benefit is applicable to instruction “bne L1” as well, but not considered in the calculations shown here.

Over:

Reorder instructions

	mov r1, #1	1	
	mov r2, #100	2	
L1:	str r1, [r2, #0]	5	14 .. 86
	add r2, r2, #4	6	
	add r1, r1, #1	7	
	cmp r1, #11	10	
	bne L1	11	20 .. 92
	mov r3, #0	95	
	mov r2, #100	96	
L2:	sub r1, r1, #1	97	110.. 214
	cmp r1, #0	100	
	beq Over	101	
	ldr r4, [r2, #0]	102	
	add r3, r3, r4	105	
	add r2, r2, #4	106	
	b L2	107	120..224
Over:		234	

After
reordering

	mov r1, #1	1	
	mov r2, #100	2	
L1:	str r1, [r2, #0]	5	13 .. 77
	add r1, r1, #1	6	
	add r2, r2, #4	7	
	cmp r1, #11	9	
	bne L1	10	18 .. 82
	mov r3, #0	85	
	mov r2, #100	86	
L2:	sub r1, r1, #1	87	99 .. 195
	cmp r1, #0	90	
	beq Over	91	
	ldr r4, [r2, #0]	92	
	add r2, r2, #4	93	
	add r3, r3, r4	95	
	b L2	96	108..204
Over:		214	

Further reordering

	mov	r1, #1	1		
	mov	r2, #100	2		
L1:	str	r1, [r2, #0]	5	13 .. 77	
	add	r1, r1, #1	6		
	add	r2, r2, #4	7		
	cmp	r1, #11	9		
	bne	L1	10	18 .. 82	
	mov	r3, #0	85		
	mov	r2, #100	86		
L2:	sub	r1, r1, #1	87	99 .. 195	
	cmp	r1, #0	90		
	beq	Over	91		
	ldr	r4, [r2, #0]	92		
	add	r2, r2, #4	93		
	add	r3, r3, r4	95		
	b	L2	96	108..204	
Over:			214		

After
reordering

	mov	r1, #1	1	
	mov	r2, #100	2	
L1:	str	r1, [r2, #0]	5	13 .. 77
	add	r1, r1, #1	6	
	add	r2, r2, #4	7	
	cmp	r1, #11	9	
	bne	L1	10	18 .. 82
	sub	r1, r1, #1	85	
	mov	r3, #0	86	
	mov	r2, #100	87	
L2:	cmp	r1, #0	88	97 .. 169
	beq	Over	89	
	ldr	r4, [r2, #0]	90	
	add	r2, r2, #4	91	
	sub	r1, r1, #1	92	
	add	r3, r3, r4	93	
	b	L2	94	103..175
Over:			182	

Delayed branch

	mov	r1, #1	1	
	mov	r2, #100	2	
L1:	str	r1, [r2, #0]	5	13 .. 77
	add	r1, r1, #1	6	
	add	r2, r2, #4	7	
	cmp	r1, #11	9	
	bne	L1	10	18 .. 82
	sub	r1, r1, #1	85	
	mov	r3, #0	86	
	mov	r2, #100	87	
L2:	cmp	r1, #0	88	97 .. 169
	beq	Over	89	
	ldr	r4, [r2, #0]	90	
	add	r2, r2, #4	91	
	sub	r1, r1, #1	92	
	add	r3, r3, r4	93	
	b	L2	94	103..175
Over:			182	

Delayed branch

Instructions
in delay
slots

L1:	mov	r1, #1	1	
	mov	r2, #100	2	
	str	r1, [r2, #0]	5	
	add	r1, r1, #1	6	13 .. 69
	add	r2, r2, #4	7	
	cmp	r1, #11	9	
	bne	L1	10	
	str	r1, [r2, #0]	11	18 .. 74
	sub	r1, r1, #1	75	
	mov	r3, #0	76	
	mov	r2, #100	77	
	cmp	r1, #0	78	
L2:	beq	Over	79	87 .. 151
	ldr	r4, [r2, #0]	80	
	add	r2, r2, #4	81	
	sub	r1, r1, #1	82	
	add	r3, r3, r4	83	
	b	L2	84	
	cmp	r1, #0	85	93 .. 157
			162	

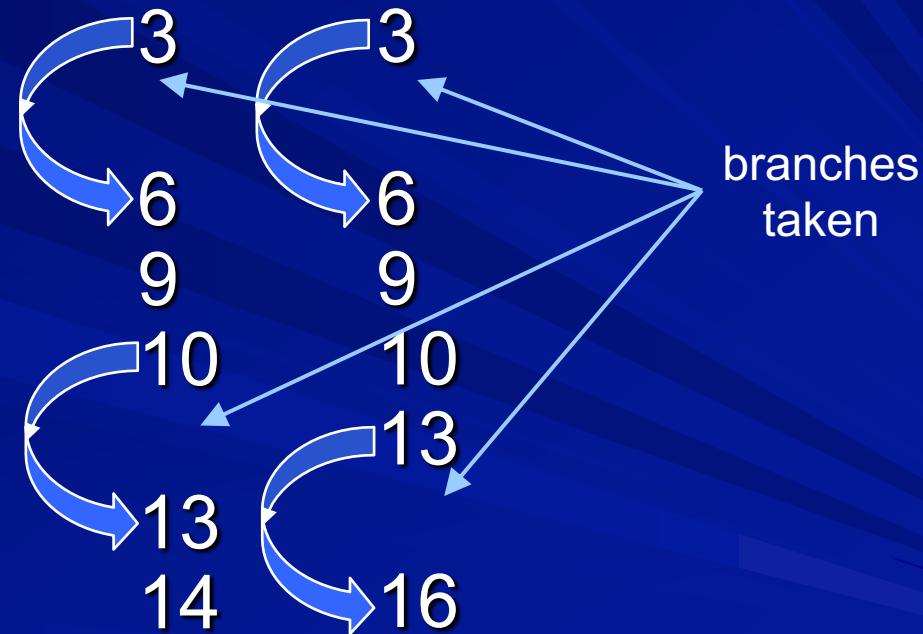
Summary

■ Original code without stalls	128
■ Original code with stalls	254
■ Execute inline instr after branch	234
■ After first re-ordering	214
■ After second reordering	182
■ Using delayed branch	162

Branch elimination example

(find max element in array)

func:	str	lr, [sp, #-4]!	1	1
	cmp	r2, #0	2	2
	bne	L1	3	3
	b	Ret	6	6
L1:	ldr	r3, [r1]	9	9
	cmp	r0, r3	10	10
	blo	L2	13	13
	b	L3	14	16
L2:	mov	r0, r3	15	17
L3:	add	r1, r1, #4	16	18
	sub	r2, r2, #1		
	bl	func		
Ret:	ldr	pc, [sp], #4		



Branch elimination example

func: str lr, [sp, #-4]!

cmp r2, #0

bne L1

b Ret



beq Ret

L1: ldr r3, [r1]

cmp r0, r3

blo L2

b L3

L2: mov r0, r3



L2: movlo r0, r3

L3: add r1, r1, #4

sub r2, r2, #1

bl func

Ret: ldr pc, [sp], #4

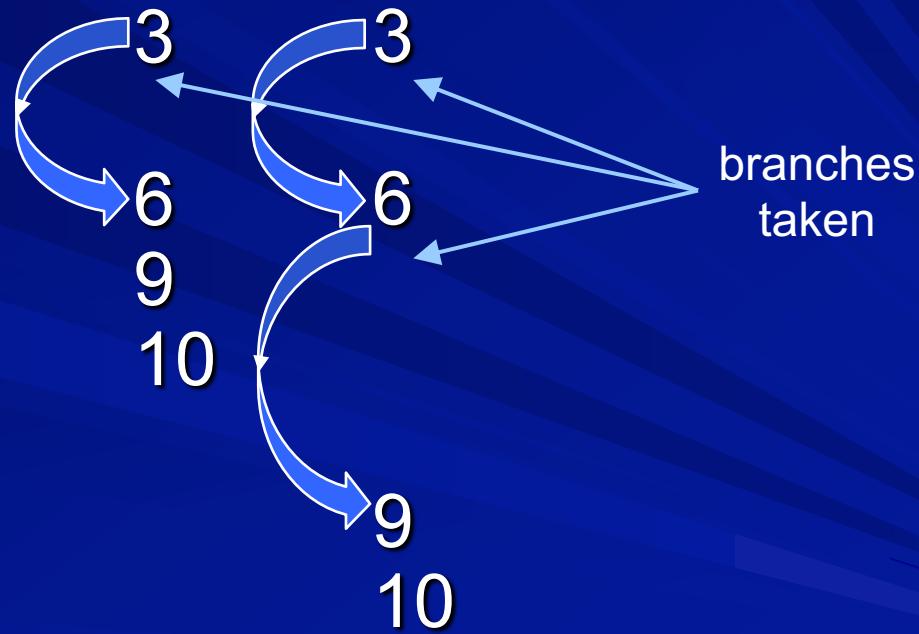
Reduced branches

func:	str	lr, [sp, #-4]!	1
	cmp	r2, #0	2
	beq	Ret	3
L1:	ldr	r3, [r1]	6
	cmp	r0, r3	9
L2:	movlo	r0, r3	10
L3:	add	r1, r1, #4	11
	sub	r2, r2, #1	12
	bl	func	13
Ret:	ldr	pc, [sp], #4	

Another example

(GCD)

```
func: str  lr, [sp, #-4]!  1   1
      cmp r0, r1    2   2
      bne L1
      b    Ret
L1:  bhs L2
      sub r1, r1, r0  9
      bl  func
      b    Ret
L2:  sub r0, r0, r1  10
      bl  func
Ret: ldr pc, [sp], #4
```



Another example

func: str lr, [sp, #-4]!

cmp r0, r1

bne L1

b Ret

L1: bhs L2

sub r1, r1, r0

bl func

b Ret

L2: sub r0, r0, r1

bl func

Ret: ldr pc, [sp], #4



beq Ret



sublo r1, r1, r0
subhs r0, r0, r1
bl func

Reduced branches

func:	str	lr, [sp, #-4]!	1
	cmp	r0, r1	2
	beq	Ret	3
L2:	sublo	r1, r1, r0	6
	subhs	r0, r0, r1	9
	bl	func	10
Ret:	ldr	pc, [sp], #4	

data
dependence

Moving “sublo r1,r1,ro” up

func:	str	lr, [sp, #-4]!	1
	cmp	r0, r1	2
	sublo	r1, r1, r0	3
	beq	Ret	4
L2:	subhs	r0, r0, r1	7
	bl	func	8
Ret:	ldr	pc, [sp], #4	

Put “sublo r1,r1,ro” in delay slot

func:	str	lr, [sp, #-4]!	1
	cmp	r0, r1	2
	beq	Ret	3
	sublo	r1, r1, r0	4
L2:	subhs	r0, r0, r1	7
	bl	func	8
Ret:	ldr	pc, [sp], #4	

THANKS