16

Indian Institute of Technology Delhi Department of Computer Science and Engineering

CSL302 April 22, 2014

Programming Languages 10 minutes

Quin 5
Maximum Marks: 89

Assume we have a functional language with call by name i.e., lazy semantics for function call. The language has at least the following constructs:

$$e \in Exp := x \mid \dots \mid \lambda x.e_1 \mid e_1(e_2) \mid \dots \mid \text{let } x \stackrel{\text{def}}{=} e_1 \text{ in } e_2 \text{ end } \mid (e_1,e_2) \mid \text{proj}_i^{(2)} \in Exp := x \mid \dots \mid \lambda x.e_1 \mid e_1(e_2) \mid \dots \mid \text{let } x \stackrel{\text{def}}{=} e_1 \text{ in } e_2 \text{ end } \mid (e_1,e_2) \mid \text{proj}_i^{(2)} \in Exp := x \mid \dots \mid \lambda x.e_1 \mid e_1(e_2) \mid \dots \mid \text{let } x \stackrel{\text{def}}{=} e_1 \text{ in } e_2 \text{ end } \mid (e_1,e_2) \mid \text{proj}_i^{(2)} \in Exp := x \mid \dots \mid \lambda x.e_1 \mid e_1(e_2) \mid \dots \mid \text{let } x \stackrel{\text{def}}{=} e_1 \text{ in } e_2 \text{ end } \mid (e_1,e_2) \mid \text{proj}_i^{(2)} \in Exp := x \mid \dots \mid \lambda x.e_1 \mid e_1(e_2) \mid \dots \mid \text{let } x \stackrel{\text{def}}{=} e_1 \text{ in } e_2 \text{ end } \mid (e_1,e_2) \mid \text{proj}_i^{(2)} \in Exp := x \mid \dots \mid \lambda x.e_1 \mid e_1(e_2) \mid \dots \mid \text{let } x \stackrel{\text{def}}{=} e_1 \text{ in } e_2 \text{ end } \mid (e_1,e_2) \mid \text{proj}_i^{(2)} \in Exp := x \mid \dots \mid \lambda x.e_1 \mid e_1(e_2) \mid \dots \mid \text{let } x \stackrel{\text{def}}{=} e_1 \text{ in } e_2 \text{ end } \mid (e_1,e_2) \mid \text{proj}_i^{(2)} \in Exp := x \mid \dots \mid \lambda x.e_1 \mid e_1(e_2) \mid \dots \mid \text{let } x \stackrel{\text{def}}{=} e_1 \text{ in } e_2 \text{ end } \mid (e_1,e_2) \mid \text{proj}_i^{(2)} \in Exp := x \mid \dots \mid \lambda x.e_1 \mid e_1(e_2) \mid \dots \mid e_1(e_1,e_2) \mid$$

Q1. (10 marks) Big-step call-by-name semantics. Recall the Principle of Correspondence, and accommingly provide big-step/Natural semantics rules for the call-by-name semantics in a "closure transition style" for only the following constructs:

let $x \stackrel{\text{def}}{=} e_1$ in e_2 end $| (e_1, e_2) | \text{proj}_i^{(2)} e$ a ∈ Answers ::= ... | Wal closures / 0.5 7 = Tables = (X -> Ansux 18) = (variables to closures mapping) cl & Closures = (Emp X Jable) of the form Ke, 77 ~ 05 Big-step rules $(e_1, \sqrt[3]{x \mapsto (e_1, \sqrt[3]{y})} \Rightarrow vel$ $(e_2, \sqrt[3]{x \mapsto (e_1, \sqrt[3]{y})} \Rightarrow vel$

-c. (da A.)