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Indian Institute of Technology Delhi
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Programming Languages

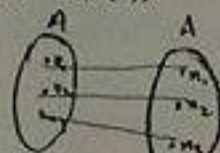
10 minutes

Quiz 1

Maximum Marks: 10

Q1. (2+4 marks) Total functions. A function $f: A \rightarrow B$ is called *total* if $\text{preDom}(f) = \text{dom}(f)$. Show that (i) the identity maps $\text{id}_A: A \rightarrow A$ are total functions; (ii) if $f: A \rightarrow B$ and $g: B \rightarrow C$ are total functions, then so is $f \circ g: A \rightarrow C$, where $(f \circ g)(x) = g(f(x))$.

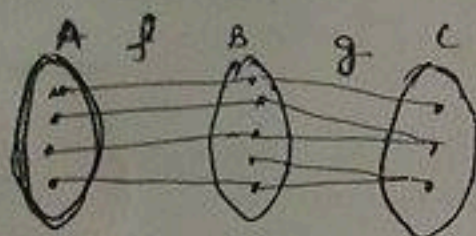
(i) $\text{id}_A: A \rightarrow A$



for every x_i in A , there is ~~some~~ $\text{id}(x_i) = x_i$ (only one value)
There is no diversion from one to many. \therefore It is a valid fn.

for every x_i in A , there is a corresponding $\text{id}(x_i) = x_i$ in A .
 $\Rightarrow \text{preDom}(\text{id}_A) = \text{dom}(\text{id}_A)$.
 \therefore total fn.

(ii)



$f \circ g: A \rightarrow C$

$f(x)$ takes x from A , for every $x \in A$, there is $f(x)$ in B (given)

also given, for every $b \in B$, there is $g(b) \in C$

\therefore for composite $g(f(x))$, take x from A , there is $b \in B$ and for that $b \in B$, there is $c \in C$.
 $\therefore f \circ g: A \rightarrow C$ is total.

Q2. (4 marks) Recursive data types. Suppose we have a simple language of expressions consisting of (i) Integers; (ii) Addition of two expressions; (iii) Multiplication of two expressions. Write a data type definition (in OCaml) to represent expressions. Please provide appropriate constructors for each case.

type exp = int | S of exp;; Negative Int?

(i) Let rec integer function (if n=0) then Z, else

X