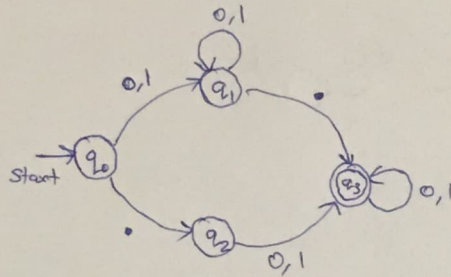


1. Consider the following DFA:

$$D = \langle \{q_0, q_1, q_2, q_3\}, \underbrace{\{0, 1\}}_Q, \underbrace{\{q_0\}}_S, \underbrace{\{q_1\}}_F, \underbrace{\{0, 1, -, \delta\}}_A \rangle$$



with start state q_0 , end state q_3 . This generates the language of unsigned real numbers in binary given by the regex $(011)^* \cdot (01)^+ \mid (011)^+ \cdot (011)^*$.

2. Consider the CFG $G = \langle \{A, B, C, S\}, \{0, 1, \cdot\}, P, S \rangle$ given by the production rules,

$$S \rightarrow 0A \mid 1A \mid \cdot B$$

$$A \rightarrow 0A \mid 1A \mid \cdot C$$

$$B \rightarrow 0C \mid 1C$$

$$C \rightarrow 0C \mid 1C \mid \epsilon.$$

This is a right linear grammar which generates the language of unsigned real numbers in binary.