# Lecture 1. Propositional Logics

Ref.: K. H. Rosen, Section 1.1 & 1.2

### **Definition**

#### **Definition:**

A Proposition is a declarative statement that is either true or false but not both.

True



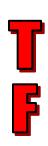
False

### **Examples**

#### Propositions:

"Qilin is a CS lecturer."

"
$$1 + 1 = 3$$
"



#### Non-Propositions:

"Is he a CS student?"

"
$$X + 1 = 3$$
"

a paradox

"This sentence is false."

Propositions cannot contain undefined variables!

### **Denotations**

Letters (p, q, r, s...) are used to denote propositions

E.g., p is defined as "Canberra is the capital of Australia"

"Canberra is the capital of Australia"

1

 $\mathcal{L}$ 

How to denote "Canberra is the capital of Australia and Canberra is located at the southeast of Australia"?

### **Compound Propositions**

New propositions can be generated by combining existing propositions using logical operators, and they are called compound propositions.

# Negation

"Today is Tuesday"

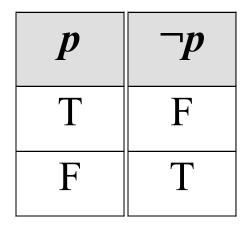
"Today is not Tuesday"

#### **Definition:**

Let *p* be a proposition. The statement "It is not the case that *p*" is another proposition, called the negation of *p*.

The negation of p is denoted by  $\neg p$ . The proposition  $\neg p$  is read "not p"

## **Negation: Truth table**



• A truth table displays the truth value of complex propositions ( $\neg p$ ) corresponding to the truth value of elementary propositions (p)

# Conjunction

"Today is Tuesday" (p) "It is raining today" (q) "Today is Tuesday and it is raining today" p and q

• We also use the symbol  $\wedge$  to represent *and*.

$$p \wedge q$$

## Conjunction

#### **Definition:**

Let p and q be propositions. The propositions "p and q" denoted by  $p \wedge q$ , is true when p and q are both true and is false otherwise.

The proposition  $p \wedge q$  is called the Conjunction of p and q.

# Conjunction

p	q	p ^ q
T	T	T
T	F	F
F	T	F
F	F	F

Given n elementary propositions, the number of rows equals to 2<sup>n</sup>!

## Disjunction

"Today is Tuesday" (p) "It is raining today" (q)
"Today is Tuesday or it is raining today"
(both might be true)

p or q

• When using logical or it is represented by the symbol  $\lor$ 

$$p \vee q$$

## Disjunction

#### **Definition:**

Let p and q be propositions. The proposition "p or q" denoted by  $p \vee q$ , is the proposition that is false when p and q are both false and true otherwise.

The proposition  $p \lor q$  is called the Disjunction of p and q.

# Disjunction

p	q	p v q
T	T	T
T	F	T
F	T	T
F	F	F

### Exclusive or

"You may have ice cream" (p) "You may have cake" (q)

"As desert you may have either ice cream *or* cake" (but not both!)

• We use the symbol  $\oplus$  to represent the *exclusive or* 

$$p \oplus q$$

### Exclusive or

#### **Definition:**

Let p and q be propositions. The exclusive or of p and q, denoted by  $p \oplus q$ , is the proposition that is true when exactly one of p and q is true and is false otherwise.

### Exclusive or

p	q	$p \oplus q$
T	T	F
T	F	T
F	T	T
F	F	F

"It is sunny today" (p) "We go to the beach" (q) If p, then q

"If it is sunny today then we go to the beach"

• When using logical *implication* it is represented by the symbol →

$$p \rightarrow q$$

#### **Definition:**

Let p and q be propositions. The implication  $p \rightarrow q$  is the proposition that is false when p is true and q is false and true otherwise.

In this implication *p* is called the hypothesis and *q* is called the conclusion.

p	q	$p \rightarrow q$
T	T	T
T	F	F
F	T	T
F	F	T

• Think of implication as a rule

"If Tom is a cat (p) then Tom has four legs (q)"

- Tom is a cat and Tom has four legs
- Tom is a cat and Tom doesn't have four legs
- Tom is not a cat and Tom has four legs
- Tom is not a cat and Tom doesn't have four legs

Sounds good

Something is wrong!

No problem

Think of implication as a promise

"If you score 80% or above (p) then I will give you a dollar (q)"

- If you did score above 80%, and I gave you a dollar, the promise is kept (p is True, q is True,  $p \rightarrow q$  is True)
- If you did score above 80%, and I didn't give you a dollar, the promise is broken (p is True, q is False,  $p \rightarrow q$  is False)
- If you did **not** score above 80%, no matter I give you a dollar or not, you cannot complain the promise is broken
   (p is False, p → q is True, no matter q)

A false statement implies anything.

### Meaning of $p \rightarrow q$ :

```
If p then q
p implies q
p only if q
p is sufficient for q
q if p
q whenever p
q is necessary for p
```

```
Meaning of p \rightarrow q: Example
      if x=1+3, then x=4
                     If x=1+3 then x=4
If p then q:
                     x=1+3 implies x=4
p implies q:
                     x=1+3 only if x=4
p only if q:
                     x=1+3 is sufficient for x=4
p is sufficient for q:
                     x = 4 if x = 1 + 3
q if p:
                     x=4 whenever x=1+3
q whenever p:
```

q is necessary for p: x=4 is necessary for x=1+3

```
Meaning of p \rightarrow q: Example
 if you finish all exercises in the textbook,
 then you'll pass the exam.
If p then q:
p implies q:
p only if q:
p is sufficient for q:
q if p:
q whenever p:
q is necessary for p:
```

### **Biconditional**

```
"The polygon has exactly 3 sides" (p) a triangle" (q)
```

"The polygon has exactly 3 sides

if and only if

the polygon is a triangle"

p if and only if q

• When using logical *biconditional* it is represented by the symbol ↔

$$p \leftrightarrow q$$

### **Biconditional**

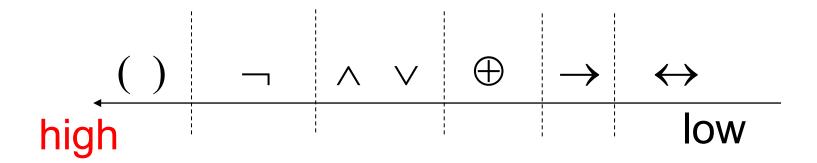
#### **Definition:**

Let p and q be propositions. The biconditional  $p \leftrightarrow q$  is the proposition that is true when p and q have the same truth values and is false otherwise.

### **Biconditional**

p	q	$p \leftrightarrow q$
T	T	T
T	F	F
F	T	F
F	F	T

### Precedence of connectives



^ and ∨ have co-equal priority, hence it is necessary to include parentheses to avoid ambiguity in some expression;

e.g. 
$$(p \land q) \lor r$$
 vs.  $p \land (q \lor r)$ 

### Precedence of connectives (cont.)

• when more than 2 instances of binary connectives of equal priority are not separated by (), the leftmost one has the highest priority;

E.g. 
$$p \rightarrow q \rightarrow r \equiv (p \rightarrow q) \rightarrow r$$

• when more than 2 instances of ¬ are not separated by (), the right most one has precedence;

E.g., 
$$\neg \neg \neg \neg p \equiv \neg(\neg(\neg(\neg p)))$$

### **A Few Terminologies**

- A compound proposition that is always true, no matter what the truth values of the propositions that occur in it, is called a tautology.
- A compound proposition that is always false is called a contradiction.
- Finally, a proposition that is neither a tautology nor a contradiction is called a contingency.

## Examples

 $p \lor \neg p$  is a tautology

 $p \land \neg p$  is a contradiction

p	$\neg p$	$p \lor \neg p$	$p \land \neg p$
T	F	T	F
F	T	T	F

## Logical equivalence

 Two statement forms are logically equivalent iff they have identical truth values for all possible combinations of truth values of their propositional symbols

 $p \Leftrightarrow q$  denotes that p and q are logically equivalent. Sometimes  $p \equiv q$ .

### Remember biconditional?

#### **Definition:**

Let p and q be propositions. The biconditional  $p \leftrightarrow q$  is the proposition that is true when p and q have the same truth values and is false otherwise.

Hence, the propositions p and q are called logically equivalent if  $p \leftrightarrow q$  is a tautology.

## Example

pq	$p \vee q$	$\neg (p \lor q)$	¬р	$\neg q$	$\neg p \wedge \neg q$
TT	T	F	F	F	F
TF	T	$\mathbf{F}$	F	T	$\mathbf{F}$
FT	T	$\mathbf{F}$	T	F	$\mathbf{F}$
F F	F	T	T	T	T

¬(p  $\vee$  q) is equivalent to ¬p  $\wedge$  ¬q ¬(p  $\vee$  q)  $\leftrightarrow$  ¬p  $\wedge$  ¬q is a tautology

# Example

pq	$p \wedge q$	$\neg (p \land q)$	¬р	$\neg q$	$\neg p \lor \neg q$
TT	T	F	F	F	F
TF	F	$\mathbf{T}$	F	T	$\mathbf{T}$
FT	F	T	T	F	T
F F	F	T	T	T	T

¬(p ∧ q) is equivalent to ¬p ∨ ¬q

## Example

- Negation of "it is sunny <u>but</u> it is not hot"
   it is not sunny or 'it is hot' or 'both'
- $p \equiv$  "John is clever and he is rich"  $\equiv a \land r$
- $p' \equiv$  "John is clever and rich"  $\equiv a \land r$
- $\neg p' \equiv$  "John is not 'clever and rich'"

$$\neg(a \land r) \equiv \neg a \lor \neg r$$
? vs.  $\neg a \land \neg r$ ?



## Equivalent form for implication

p	q	$p \rightarrow q$	$\neg p \lor q$
T	T	T	T
T	F	F	F
F	T	T	T
F	F	T	T

"If it is sunny then we go to the beach"



"It is not sunny or we go to the beach"

## **Negation of implication**

$$\neg (p \rightarrow q) \equiv \neg(\neg p \lor q)$$

$$\equiv \neg (\neg p) \land \neg q$$

$$\equiv p \land \neg q$$

E.g. Negate "If it is sunny then we go to the beach"



"It is sunny and we do not go to the beach"

Show that  $(p \land q \rightarrow p)$  and  $(p \rightarrow q \lor p)$  are tautologies.

Method 1: use truth table:

p	q	p∧q	q∨p	$p \land q \rightarrow p$	$p \rightarrow q \lor p$
T	T	T	T	T	T
T	F	F	T	T	T
F	T	F	T	T	T
F	F	F	F	T	T

#### Method 2:

```
(p \land q \rightarrow p) \equiv \neg (p \land q) \lor p
                        \equiv (\neg p \lor \neg q) \lor p
                         \equiv (\neg p \lor p) \lor \neg q
                        \equiv t \vee \neg q \equiv t
(p \rightarrow q \lor p) \equiv \neg p \lor (q \lor p)
                          \equiv (\neg p \lor p) \lor q
                          \equiv t \vee q \equiv t
```

```
Show: p \lor q \rightarrow r \equiv (p \rightarrow r) \land (q \rightarrow r)
    Method 1: use truth table;
     Method 2:
            p \lor q \rightarrow r \equiv (p \lor q) \rightarrow r
                                \equiv \neg (p \lor q) \lor r
                                \equiv (\neg p \land \neg q) \lor r
                                \equiv (\neg p \lor r) \land (\neg q \lor r)
                                \equiv (p \rightarrow r) \land (q \rightarrow r)
```

## Three terms about $p \rightarrow q$

• The contrapositive of  $p \rightarrow q$  is  $\neg q \rightarrow \neg p$ 

In fact, 
$$\neg p \rightarrow \neg q \equiv q \rightarrow p$$

• The converse of  $p \rightarrow q$  is  $q \rightarrow p$ 

$$q \rightarrow p - p \rightarrow q$$
?

• The inverse of  $p \rightarrow q$  is  $\neg p \rightarrow \neg q$ 

$$\neg p \rightarrow \neg q = p \rightarrow q?$$

• Clearly,  $\neg p \rightarrow \neg q \equiv q \rightarrow p$ 

E.g. "If it is sunny then we go to the beach."

#### Contrapositive:

"If we do not go to the beach then it is not sunny."

#### Converse:



"If we go to the beach then it is sunny."

#### Inverse:



"If it is not sunny then we do not go to the beach."

### **Biconditional**

$$p \leftrightarrow q \equiv p \text{ if, and only if, } q$$
  
 $\equiv p \text{ if } q, \text{ and, } p \text{ only if, } q$ 

pq	$p \rightarrow q$	$q \rightarrow p$	$p \leftrightarrow q$
TT	T	T	T
TF	F	T	$\mathbf{F}$
FT	T	F	$\mathbf{F}$
FF	T	T	T

$$p \leftrightarrow q \equiv (p \rightarrow q) \land (\neg p \rightarrow \neg q) \equiv (p \rightarrow q) \land (q \rightarrow p)$$

De Morgan's laws

$$\neg(p \land q) \Leftrightarrow \neg p \lor \neg q$$
$$\neg(p \lor q) \Leftrightarrow \neg p \land \neg q$$

Commutative

$$p \wedge q \Leftrightarrow q \wedge p$$
  
 $p \vee q \Leftrightarrow q \vee p$ 

Associative

$$(p \lor q) \lor r \Leftrightarrow p \lor (q \lor r)$$
  
 $(p \land q) \land r \Leftrightarrow p \land (q \land r)$ 

Distributive

$$p \land (q \lor r) \Leftrightarrow (p \land q) \lor (p \land r)$$
  
 $p \lor (q \land r) \Leftrightarrow (p \lor q) \land (p \lor r)$   
[Remember?  $a(b+c)=ab+ac$ ]

Identity (t is tautology, c is contradiction)

$$p \wedge t \Leftrightarrow p \qquad p \vee c \Leftrightarrow p$$

Universal bound

$$p \lor t \Leftrightarrow t$$

$$\boldsymbol{p} \wedge \boldsymbol{c} \Leftrightarrow \boldsymbol{c}$$

Negation

$$p \land \neg p \Leftrightarrow c$$

$$p \lor \neg p \Leftrightarrow t$$

Double Negation

$$\neg(\neg p) \Leftrightarrow p$$

Idempotent

$$p \land p \Leftrightarrow p \qquad p \lor p \Leftrightarrow p$$

Absorption

$$p \lor (p \land q) \Leftrightarrow p \quad p \land (p \lor q) \Leftrightarrow p$$

## **Example 1: Proof of Distributive Law**

$$p \lor (q \land r) \Leftrightarrow (p \lor q) \land (p \lor r)$$

p q r	$q \wedge r$	$p \vee (q \wedge r)$	$p \vee q$	$p \vee r$	$(p \lor q) \land (p \lor r)$
T T T T F T F T T F T T F T	T F F F T	T T T T	T T T T	T T T T T	T T T T
F T F F F F	F F F	F F F	T F F	F T F	F F F

## Example 2:

Show that  $\neg(p \lor (\neg p \land q))$  and  $\neg p \land \neg q$  are logically equivalent  $\neg (p \lor (\neg p \land q)) \Leftrightarrow \neg p \land \neg (\neg p \land q)$ De Morgan  $\Leftrightarrow \neg p \land (\neg (\neg p) \lor \neg q)$  De Morgan  $\Leftrightarrow \neg p \wedge (p \vee \neg q)$ Double negation  $\Leftrightarrow (\neg p \not \neg p) \lor (\neg p \land \neg q)$  Distributive  $\Leftrightarrow$  F  $\checkmark$  ( $\neg p \land \neg q$ ) known contradiction  $\Leftrightarrow (\neg p \land \neg q) \checkmark F$ Commutative Identity  $\Leftrightarrow \neg p \land \neg q$ 

#### Proof of Absorption

$$p \lor (p \land q) \Leftrightarrow p \quad p \land (p \lor q) \Leftrightarrow p$$

pq	$p \wedge q$	$p\lor(p\land q)$	$p \vee q$	$p \land (p \lor q)$
TT	T	T	T	T
TF	F	T	T	$\mathbf{T}$
FT	F	F	T	$\mathbf{F}$
F F	F	F	F	F

## **Translating English Sentences**

Statement 1: John is short or Marry is pretty, and, John is not short or Marry is not pretty

Statement 2: It is not the case that both John is short and Marry is pretty, or, both John is not short and Marry is not pretty

#### Logical Forms:

1. 
$$(P \lor Q) \land (\neg P \lor \neg Q)$$

2. 
$$\neg ((P \land Q) \lor (\neg P \land \neg Q))$$

## **Translating English Sentences**

They are equivalent:

$$\neg((P \land Q) \lor (\neg P \land \neg Q))$$

$$\Leftrightarrow \neg(P \land Q) \land \neg(\neg P \land \neg Q)$$

$$\Leftrightarrow (\neg P \lor \neg Q) \land (P \lor Q)$$

$$\Leftrightarrow (P \lor Q) \land (\neg P \lor \neg Q)$$

Use the Truth Table:

## **Translating English Sentences**

рq	$\neg p \neg q$	$p \vee q$	$\neg p \vee \neg q$	$(p \lor q) \land (\neg p \lor \neg q)$	$p \oplus q$
T T T F F T F F	F F T T T	T T T F	F T T	F T T F	F T T F

Hence the two confusing sentences are logically equivalent to the more legible form:

Either John is short or Marry is pretty, but not both.

# Only If

"Emma eats dinner only if she is hungry"

- "If she is not hungry then Emma does not eat dinner"
- = "If Emma eats dinner then she is hungry"

p only if q  $\equiv (\neg q \rightarrow \neg p)$   $\equiv p \rightarrow q$ 

### Sufficient condition

- "Vandalizing others' property is a sufficient condition for Michael to be fined"
- = "If Michael vandalizes others' property then he will be fined"

*p is a sufficient condition for*  $q \equiv p \rightarrow q$ 

## **Necessary condition**

- "Being over 16 is a necessary condition for a person to get a driver's licence"
- = "If a person has not turned 16 then he/she cannot get a driver's licence"
- = "If a person get his/her driver's licence then he/she is over 16"

$$p$$
 is a necessary condition for  $q \equiv \neg p \rightarrow \neg q$   
$$\equiv q \rightarrow p$$

Are the following 3 statements logically equivalent?

- "You fix my ceiling or I won't pay my rent"  $x: P \lor \neg Q$
- "If you do not fix my ceiling then I won't pay my rent"

$$y: \neg P \rightarrow \neg Q$$

• "I will pay my rent only if you fix my ceiling"  $z: Q \rightarrow P$ 

Rewrite 'if n is prime, then n is odd or n is 2' in a few other (logically equivalent) ways

A prestige company wrote: 'a person will be hired only if he majors in mathematics or computer science, get a *B* average or better, and take accounting'; an applicant with such qualifications was turned down; did the company lie?

No. They stated the necessary conditions, not sufficient condition. You will be considered only with the qualification, but not guaranteed to be hired.

## **Summary**

- Definition of Proposition
- Truth table of propositions
- Logical Connectives of proposition:
  - $not \neg$ ,  $and \land$ ,  $or \lor$ , exclusive or  $\oplus$ , implication  $\rightarrow$ , and biconditional  $\leftrightarrow$
- Translation to and from English
- Tautology and Contradiction

## **Summary**

- Logical Equivalence
- Common Logical Laws
- Negation of  $p \rightarrow q$
- Contrapositive, Converse and inverse of  $p \rightarrow q$
- Sufficient condition, necessary condition, bicondition.