

Lecture 6. Relations

Ref.: K H Rosen Chapter 7

Binary Relation

- A **binary relation** R from a set A to a set B is a subset $R \subseteq A \times B$.

E.g. Let $A = \{a, b, c\}$

$B = \{1, 2, 3, 4\}$

R is defined by the ordered pairs
 $\{(a, 1), (a, 2), (c, 4)\}$

aRb denotes that $(a, b) \in R$ while $a \not R b$ means
 $(a, b) \notin R$

When aRb , a is said to be related to b by R .

Relation on a Set

- A binary relation R on a set A is a relation from A to A .

E.g. Let $A = \{a, b, c\}$

$R = \{(a, a), (a, b), (c, c)\}$

Question: How many binary relations are there on a set A with n elements?

Representations of Relations

Connection Matrices

Let R be a relation from set $A = \{a_1, a_2, \dots, a_m\}$ to $B = \{b_1, b_2, \dots, b_n\}$.

A *$m \times n$ connection matrix M* for R is defined by

$$M_{ij} = \begin{cases} 1 & \text{if } \langle a_i, b_j \rangle \text{ is in } R; \\ 0 & \text{otherwise.} \end{cases}$$

Example

- assume the rows are labeled with the elements of A and the columns are labeled with the elements of B.

Let $A = \{a, b, c\}$, $B = \{e, f, g, h\}$ and
 $R = \{ \langle a, e \rangle, \langle c, g \rangle \}$.

Then the connection matrix **M** for R is

$$\begin{bmatrix} 1 & 0 & 0 & 0 \\ 0 & 0 & 0 & 0 \\ 0 & 0 & 1 & 0 \end{bmatrix}$$

Representations of Relations

Graphical representation

Definition: A *Directed graph* or a *Digraph* D from A to B is a collection of *vertices* and a collection of *edges*. If there is an ordered pair $e = \langle x, y \rangle$ in R then there is an *arc* or *edge* from x to y in D . The elements x and y are called the *initial* and *terminal vertices* of the edge e .

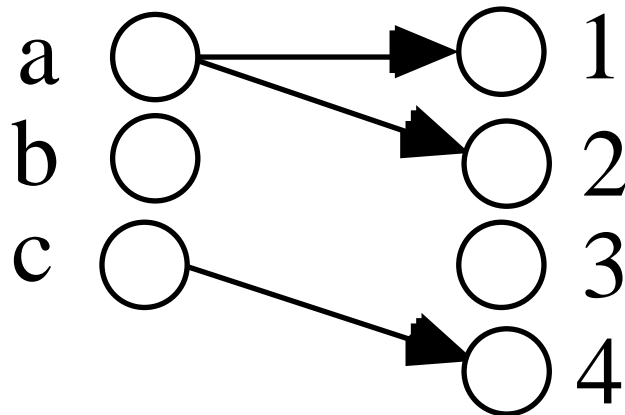
Example

Let $A = \{a, b, c\}$ and $B = \{1, 2, 3, 4\}$

R is defined by the ordered pairs or edges

$\{ \langle a, 1 \rangle, \langle a, 2 \rangle, \langle c, 4 \rangle \}$

R can be represented by the digraph D :

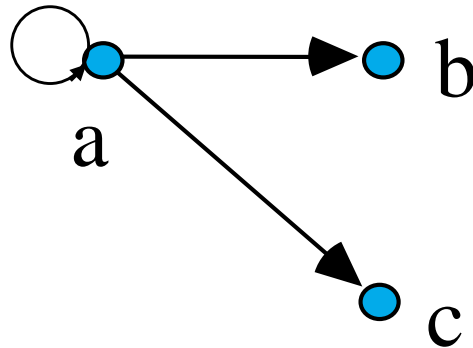


Example

For a relation R on a set $A = \{a, b, c\}$

$$R = \{ \langle a, a \rangle, \langle a, b \rangle, \langle a, c \rangle \}.$$

Then a digraph representation of R is:



An arc of the form $\langle x, x \rangle$ on a digraph is called a *loop*.

Special Properties of Binary Relations

Given: a Universe **U** and a binary relation **R** on a subset **A** of **U**

- **R** is *reflexive* iff $(a,a) \in R$ for every element $a \in A$

A relation **R** on a set **A** is reflexive if every element of **A** is related to itself.

E.g. Is the “Divides” relation on the set of positive integers reflexive?

Yes, since $a|a$ for all positive integer a .

Special Properties of Binary Relations

- R is **symmetric** iff $(a,b) \in R$ whenever $(b,a) \in R$ for every $a, b \in A$.
- R is called **antisymmetric** such that $(a,b) \in R$ and $(b,a) \in R$ only if $a=b$ for every $a, b \in A$.

Symmetric and antisymmetric are not opposite since a relation can have both or lack both of them.

(Can you think of some examples?)

Special Properties of Binary Relations

- R is *transitive* iff whenever $(a,b) \in R$ and $(b,c) \in R$, then $(a,c) \in R$, for every $a, b, c \in A$.

This is the most difficult one to check.
Algorithms have been developed to check this.

Relation properties in Matrix

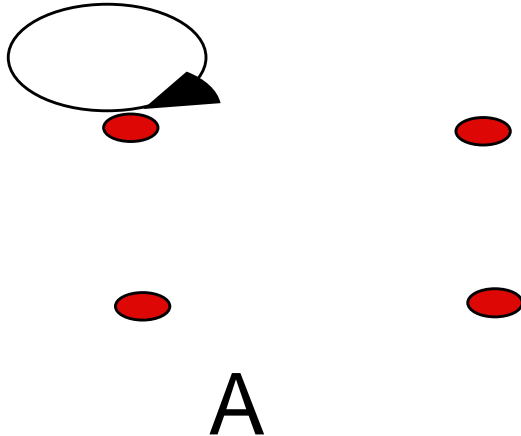
Let R be a binary relation on a set A and let M be its connection matrix. Then

- R is **reflexive** iff $M_{ii} = 1$ for all i .
- R is **symmetric** iff M is a symmetric matrix: $M = M^T$
- R is **antisymmetric** iff $M_{ij} = 0$ or $M_{ji} = 0$ for all $i \neq j$.

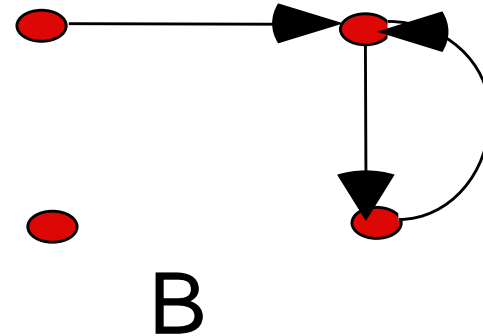
Relation properties in Digraph

- A relation is **reflexive** iff there is a loop at **every** vertex of its digraph.
 - A relation is **symmetric** iff for every edge between distinct vertices in its digraph there is an edge in the opposite direction.
 - A relation is **antisymmetric** iff there are never two edges in the opposite directions between distinct vertices.
 - A relation is **transitive** iff whenever there is a edge from vertex x to y and from y to z , there is a edge from x to z .
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Example

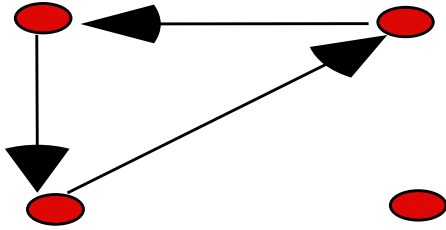


A: not reflexive
 symmetric
 antisymmetric
 transitive



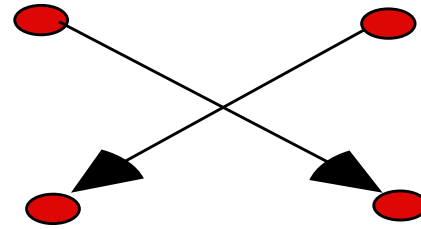
B: not reflexive
 not symmetric
 not antisymmetric
 not transitive

Example



C

C: not reflexive
 not symmetric
 antisymmetric
 not transitive



D

D: not reflexive
 not symmetric
 antisymmetric
 transitive

Combining Relations

Since Relations from A to B are subset of $A \times B$, two relations from A to B can be combined in any way two sets can be combined.

- Union $R_1 \cup R_2$
- Intersection $R_1 \cap R_2$
- Difference $R_1 - R_2$
- Compliment $\overline{R_1}$

Combining Relations: Example

Let $A = \{1, 2, 3\}$ and $B = \{1, 2, 3, 4\}$

The Relations $R_1 = \{(1,1), (2,2), (3,3)\}$ and $R_2 = \{(1,1), (1,2), (1,3), (1,4)\}$ can be combined to obtained as

- $R_1 \cup R_2 = \{(1,1), (1,2), (1,3), (1,4), (2,2), (3,3)\}$
- $R_1 \cap R_2 = \{(1,1)\}$
- $R_1 - R_2 = \{(2,2), (3,3)\}$
- $R_2 - R_1 = \{(1,2), (1,3), (1,4)\}$

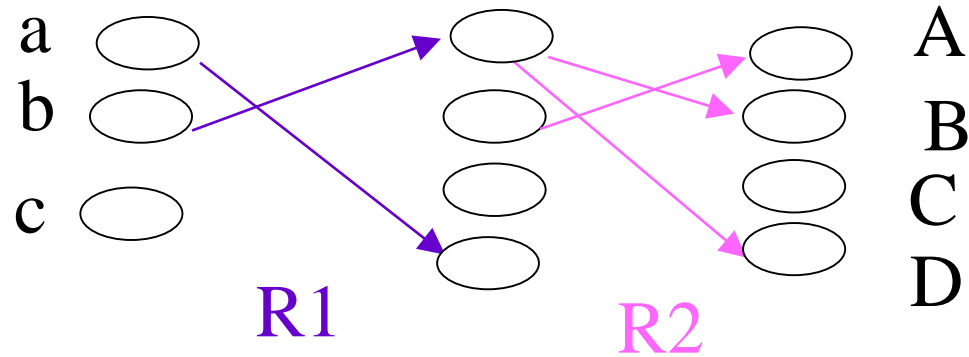
Composition of Relations

Suppose $R1$ is a relation from set A to B , $R2$ is a relation from set B to C . Then the **composition of $R2$ with $R1$** , denoted $R2 \circ R1$ is the relation from A to C :

If (a,b) is a member of $R1$ and (b,c) is a member of $R2$ then (a,c) is a member of $R2 \circ R1$.

- For (a, c) to be in the composite relation $R2 \circ R1$ there must exist a b in B
- We read them right to left as in functions.

Example



$$R2 \circ R1 = \{(b, D), (b, B)\}$$

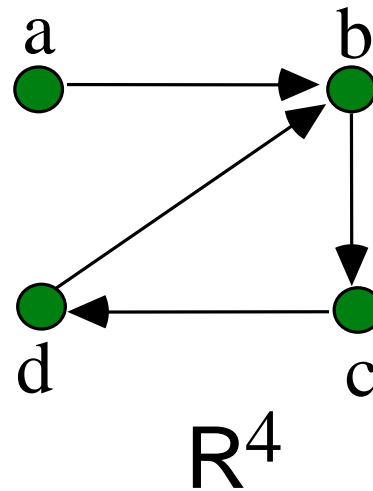
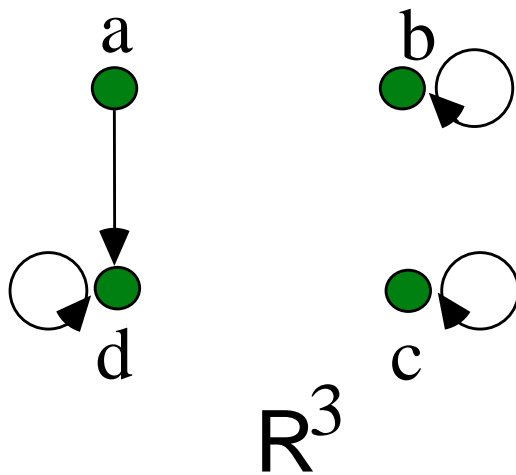
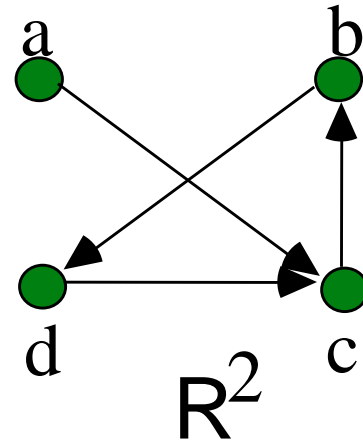
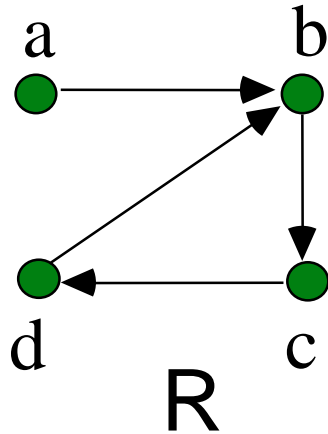
Powers of a Relation

Let R be a binary relation on A . Then R^n :
the power of R are defined inductively by

Basis: $R^1 = R$

Induction: $R^{n+1} = R^n \circ R$

Example



Inverse of R

Let R be a relation on A . Then R^{-1} or the **inverse of R** is the relation

$$R^{-1} = \{ (y, x) \mid (x, y) \in R \}$$

Example:

Let $A = \{1, 2, 3, 4\}$

The Relation $R = \{(1, 1), (1, 2), (1, 3), (1, 4)\}$

$$R^{-1} = \{ (1, 1), (2, 1), (3, 1), (4, 1) \}$$

This relation is sometimes denoted as R^T or R^c and called the *converse* of R .

n-ary Relation

Let A_1, A_2, \dots, A_n be sets. An *n*-ary relation on these sets is a subset of $A_1 \times A_2 \times \dots \times A_n$. The sets A_1, A_2, \dots, A_n are called the **domains** of the relation, and n is called its **degree**.

E.g.

Let R be the relation consisting of a 4-tuples (Student, ID, School, Age) representing student particulars in Curtin. If 21-year-old John with ID 1234567 is a student in CS, (John, 1234567, CS, 21) belongs to R .

Equivalence Relations

Definition: A relation R on a set A is an *equivalence relation* iff R is reflexive, symmetric and transitive.

Examples

- Let R be a relation on the strings of English letters such that aRb iff $l(a)=l(b)$ where $l(x)$ is the length of the string x .

R is an equivalence relation since:

- aRa Reflexive
- $aRb \rightarrow bRa$ Symmetric
- $aRb \wedge bRc \rightarrow aRc$ Transitive

Examples

- Let R be a relation on the set of real numbers such that aRb iff $a-b$ is an integer.

R is an equivalence relation since:

- $a-a=0$ Reflexive
- if $a-b=k$, $b-a=-k$ (integer) Symmetric
- if $a-b=i$ and $b-c=j$, $a-c=(a-b)+(b-c)=i+j$ Transitive

Examples

- Let m be a positive integer greater than 1. Show that $R = \{(a,b) \mid a \equiv b \pmod{m}\}$ is an equivalence relation on the integer set.

R is called the **congruence modulo m** .
(This is the most widely used equivalence relation.)

Examples

$a \equiv b \pmod{m}$ iff m divides $a-b$

Hence,

$a-a=0$ divisible by m . $\therefore a=a \pmod{m}$. Reflexive.

If $a=b \pmod{m}$, $a-b=km$ for certain k .

$b-a=(-k)m$. $\therefore b=a \pmod{m}$. Symmetric.

If $a=b \pmod{m}$ and $b=c \pmod{m}$

$a-b=im$ and $b-c=jm$ for certain i, j .

$a-c=(a-b)+(b-c)=im+jm=(i+j)m$

$\therefore a=c \pmod{m}$. Transitive.

Equivalence class

Definition: Let R be an equivalence relation on A . The set of all elements that are related to an element a in A is called the **equivalence class of a** .

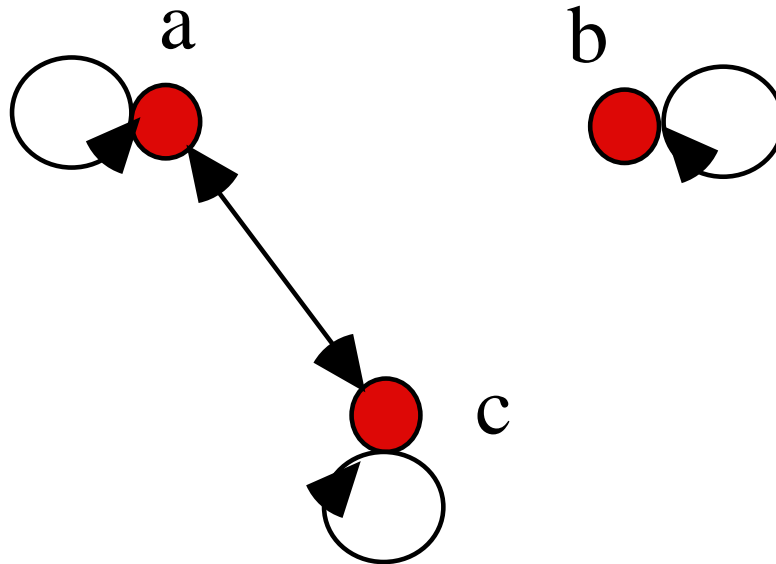
The equivalence class of a with respect to R is denoted as $[a]_R$.

When only one relation is under consideration, the subscript R can be deleted and write it as $[a]$.

$$[a] = \{x \mid (a, x) \in R\}$$

a is called a **representative** of this class.

Example



$$[a] = \{a, c\}, [c] = \{a, c\}, [b] = \{b\}.$$

Partition

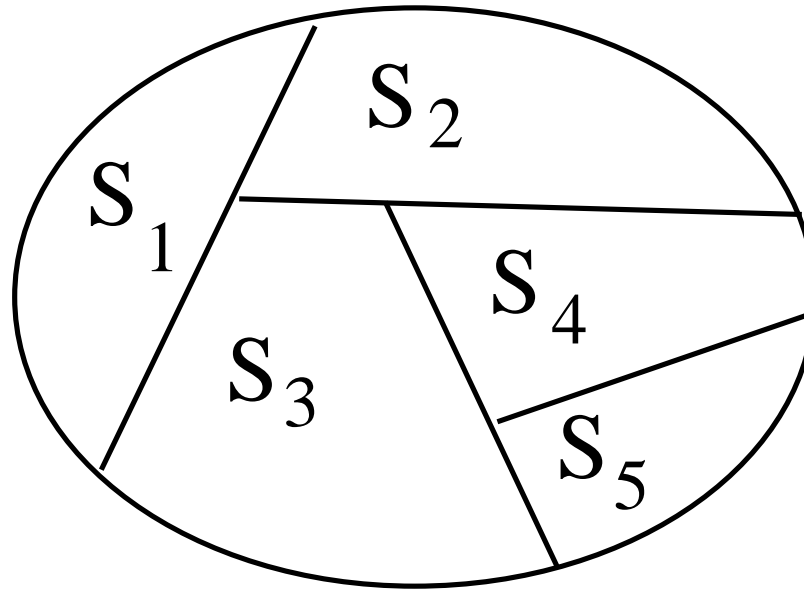
Definition: Let S_1, S_2, \dots, S_n be a collection of subsets of A . Then the collection forms a **partition of A** if the subsets are **nonempty, disjoint and exhaust A** :

$$S_i \neq \emptyset$$

$$S_i \cap S_j = \emptyset \text{ if } i \neq j$$

$$\bigcup S_i = A$$

Partition: Example

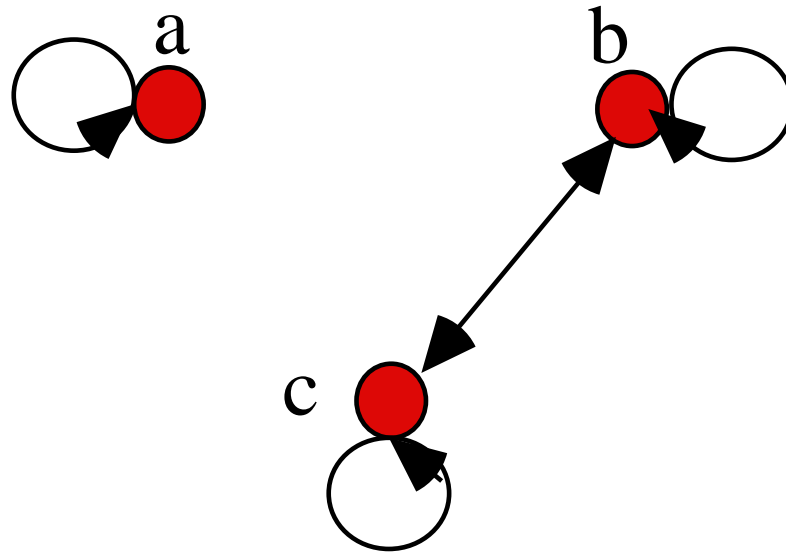


A

Partition

Theorem: Let R be an equivalence relation on a set A . Then the equivalence classes of R form a partition of A . Conversely, given a partition $\{S_i \mid i \in I\}$ of the set A , there is an equivalence relation R that has the sets S_i as its equivalence classes.

Example



$$A = [a] \cup [b] = [a] \cup [c] = \{a\} \cup \{b, c\}$$

Example

What are the sets in the partition of integers arising from congruence modulo 4?

The four congruence classes, $[0]_4$, $[1]_4$, $[2]_4$, $[3]_4$ form a partition of integers.

They are the sets:

$$[0]_4 = \{\dots, -8, -4, 0, 4, 8, \dots\}$$

$$[1]_4 = \{\dots, -7, -3, 1, 5, 9, \dots\}$$

$$[2]_4 = \{\dots, -6, -2, 2, 6, 10, \dots\}$$

$$[3]_4 = \{\dots, -5, -1, 3, 7, 11, \dots\}$$

Partial Order

Definition: Let R be a relation on A . Then R is a ***partial order*** iff R is **reflexive**, **antisymmetric** and **transitive**.

(A, R) is called a **partially ordered set** or a ***poset***.

Note: It is not required that every two elements be related under a partial order. That's the ***partial*** part of it.

Example

Show that the divisibility relation “|” is a partial ordering on the set of positive integers.

“|” is **reflexive** since for any $a \in \mathbb{Z}^+$, $a|a$;

“|” is **antisymmetric** since for any $a, b \in \mathbb{Z}^+$,

$$a|b \wedge b|a \rightarrow a=b;$$

“|” is **transitive** since for any $a, b, c \in \mathbb{Z}^+$,

$$a|b \wedge b|c \rightarrow a|c;$$

Hence $(\mathbb{Z}^+, |)$ is a poset.

Example

Show that relation “ \subseteq ” is a partial ordering on the power set of a set S .

“ \subseteq ” is reflexive since for any $A \in P(S)$, $A \subseteq A$;
“ \subseteq ” is antisymmetric since for any $A, B \in P(S)$,
 $A \subseteq B \wedge B \subseteq A \rightarrow A = B$;
“ \subseteq ” is transitive since for any $A, B, C \in P(S)$,
 $A \subseteq B \wedge B \subseteq C \rightarrow A \subseteq C$;
Hence $(P(S), \subseteq)$ is a poset.

Notation \leq

In a poset (A, R) , the notation $a \leq b$ is often used to denote that aRb . Hence poset is often written as (A, \leq) .

Note that the symbol “ \leq ” is used to denote the relation in any poset, not just the “less than or equals” relation.

The notation $a < b$ denotes that $a \leq b$ but $a \neq b$. It's also read as “ a is less than b ” or “ b is greater than a ”.

Notation \leq

Definition: The elements a and b of a poset (A, \leq) are called **comparable** if **either $a \leq b$ or $b \leq a$** . Otherwise a and b are called **incomparable**.

E.g.

In the poset $(P(Z), \subseteq)$, $\{1,2\}$ and $\{1,3\}$ are not comparable.

In the poset $(Z^+, |)$, 3 and 9 are comparable, while 5 and 7 are not comparable.

Total Order

Definition: If **every two** elements in a poset (A, \leq) are **comparable**, A is called a ***total ordered*** or ***linearly ordered*** or ***simply ordered set***, and \leq is called a total order or linear order or simple order.

A totally ordered set is also called a ***chain***.

Example

- (\mathbb{Z}, \leq) is a poset. In this case either $a \leq b$ or $b \leq a$ so two elements are always related. Hence, \leq is a total order and (\mathbb{Z}, \leq) is a chain.
- If S is a set then $(P(S), \subseteq)$ is a poset. For any two set $A, B \in P(S)$, it may not be the case that $A \subseteq B$ or $B \subseteq A$. Hence, $(P(S), \subseteq)$ is not totally ordered.
- $(\mathbb{Z}^+, |)$ is a poset which is not a chain.

Summary

- **Binary relation**
- **Relation on a set**
- **Combining and composition of relations**
- **Power and Inverse of relation**
- **Special properties of relation:**
 - **reflexive**
 - **symmetric**
 - **antisymmetric**
 - **transitive**

Summary

- **Equivalence relations and classes**
- **Partitions**
- **Partial order and Total order**