

Design and Analysis of Algorithms (COMP3001)

Tutorial 3

Divide and Conquer

Question 1.

- a) Using Figure 2.4 (textbook 3rd edition) as a model, illustrate the operation of MERGESORT on the array $A = \langle 13, 19, 9, 5, 12, 8, 7, 4, 11, 2, 6, 21 \rangle$.
- b) Using Figure 7.1 (textbook 3rd edition) as a model, illustrate the operation of PARTITION on the array
 - $A = \langle 13, 19, 9, 5, 12, 8, 7, 4, 11, 2, 6, 21 \rangle$
 - $A = \langle 13, 19, 9, 5, 12, 8, 7, 4, 11, 2, 6 \rangle$
- c) Using Figure 2.2 (textbook new edition) as a model, illustrate the operation of INSERTION-SORT on the array $A = \langle 13, 19, 9, 5, 12, 8, 7, 4, 11, 2, 6, 21 \rangle$. Read the analysis of Insertion Sort.

Question 2.

Exercise 7.2-2 (textbook, Cormen, et al). What is the running time of QUICKSORT when all elements of array A have the same value? Is this the best case for QUICKSORT? Why?

Question 3.

- a) Show that Quicksort's best case running time is $\Omega(n \log n)$.
- b) Show that the running time of QUICKSORT is $\Theta(n^2)$ when the array A contains distinct elements and is sorted in decreasing order.

Question 4.

- a) Design an algorithm that will take as input a set of real numbers S and a target real number x , and output TRUE if there are two numbers in S that sum to x , and FALSE otherwise.
- b) How does the running time of your algorithm change if S is sorted or unsorted?

Question 5.

Chapter 8 (Textbook) provides the following Counting Sort algorithm. The algorithm considers an input array $A[1..n]$, and assumes each element in array A is an integer in the range 0 to k . Note that $A.length$ is size of array A , and thus $A.length = n$. Further, the algorithm uses array $B[1..n]$ to keep the sorted output, and array $C[0..k]$ for temporary storage.

Counting-Sort (A, B, k)

```
1   let  $C[0..k]$  be a new array
2   for  $i = 0$  to  $k$ 
3        $C[i] = 0$ 
4   for  $j = 1$  to  $A.length$ 
5        $C[A[j]] = C[A[j]] + 1$ 
6   //  $C[i]$  now contains the number of elements equal to  $i$ 
7   for  $i = 1$  to  $k$ 
8        $C[i] = C[i] + C[i - 1]$ 
9   //  $C[i]$  now contains the number of elements less than or equal to  $i$ 
10  for  $j = A.length$  downto 1
11       $B[C[A[j]]] = A[j]$ 
12       $C[A[j]] = C[A[j]] - 1$ 
```

- a) Show that Counting-Sort has a running time complexity of $O(n)$.
- b) (Exercise 8.2.1 – textbook). Using Figure 8.2 as a model, illustrate the operation of Counting-Sort on the array $A = [6, 0, 2, 0, 1, 3, 4, 6, 1, 3, 2]$, i.e.,
 - Show the contents of array A and C after line 5.
 - Show the content of array C after line 8.
 - Show the content of array B and C after one, two, and three iterations of the loop in lines 10-12
 - Show the final sorted output array B .

Question 6.

Consider an array $A[1..n]$ of integers. Design a divide and conquer algorithm to find the minimum element in A . For example, if $A = [3, 6, 1, 5, 7, 2, 1]$, your algorithm should obtain 1. **Your algorithm must be computationally as efficient as possible.**

- (i) Write the pseudocode of your algorithm. **Hint.** Similar to MERGE_SORT.
- (ii) Show how your algorithm works on the input $A = [3, 6, 1, 5, 7, 2, 1]$.

- (iii) Write the recurrence of the time complexity of your algorithm and solve the recurrence to find its time complexity in Big Oh.

Question 7.

- a) Strassen's algorithm (discussed in the lecture) generates the following equations to produce:

$$\begin{pmatrix} C_{11} & C_{12} \\ C_{21} & C_{22} \end{pmatrix} = \begin{pmatrix} A_{11} & A_{12} \\ A_{21} & A_{22} \end{pmatrix} \times \begin{pmatrix} B_{11} & B_{12} \\ B_{21} & B_{22} \end{pmatrix}$$

$$\begin{aligned} S_1 &= B_{12} - B_{22} & S_6 &= B_{11} + B_{22} \\ S_2 &= A_{11} + A_{12} & S_7 &= A_{12} - A_{22} \\ S_3 &= A_{21} + A_{22} & S_8 &= B_{21} + B_{22} \\ S_4 &= B_{21} - B_{11} & S_9 &= A_{11} - A_{21} \\ S_5 &= A_{11} + A_{22} & S_{10} &= B_{11} + B_{12} \end{aligned}$$

$$\begin{aligned} P_1 &= A_{11} \times S_1 \\ P_2 &= S_2 \times B_{22} \\ P_3 &= S_3 \times B_{11} \\ P_4 &= A_{22} \times S_4 \\ P_5 &= S_5 \times S_6 \\ P_6 &= S_7 \times S_8 \\ P_7 &= S_9 \times S_{10} \end{aligned}$$

$$\begin{aligned} C_{11} &= P_5 + P_4 - P_2 + P_6 \\ C_{12} &= P_1 + P_2 \\ C_{21} &= P_3 + P_4 \\ C_{22} &= P_5 + P_1 - P_3 - P_7 \end{aligned}$$

Verify that the following equations are correct, and hence Strassen's algorithm is correct.

$$\begin{aligned} C_{11} &= A_{11} \times B_{11} + A_{12} \times B_{21} \\ C_{12} &= A_{11} \times B_{12} + A_{12} \times B_{22} \\ C_{21} &= A_{21} \times B_{11} + A_{22} \times B_{21} \\ C_{22} &= A_{21} \times B_{12} + A_{22} \times B_{22} \end{aligned}$$

- b) Exercise 4.2-1. Use Strassen's algorithm to compute the matrix product. Show your work.

$$\begin{pmatrix} 1 & 3 \\ 7 & 5 \end{pmatrix} \begin{pmatrix} 6 & 8 \\ 4 & 2 \end{pmatrix}$$