

The Maximum Facility Location Problem

Mitchell Jones

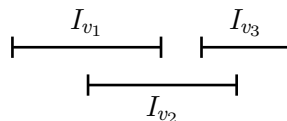
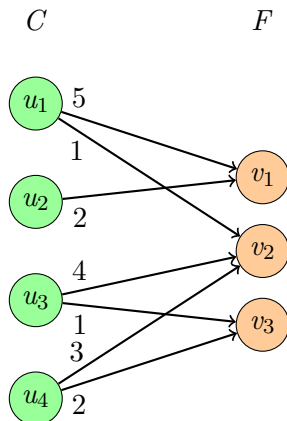
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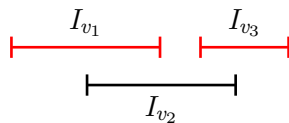
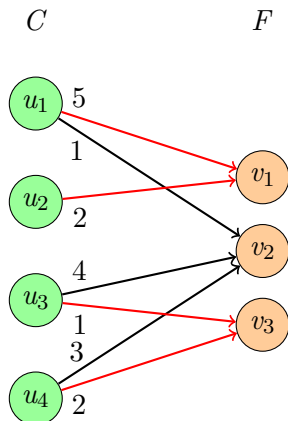


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Example



Example (cont.)



Value = 10

Maximum Facility Location

- ▶ C is the set of clients
- ▶ F is the set of facilities
- ▶ Weight function $w : C \times F \rightarrow \mathbb{R}$
- ▶ Each facility $v \in F$ has an interval I_v on the real line

Goal

Select an independent set $S \subseteq F$ maximising cost of assigning clients to facilities.

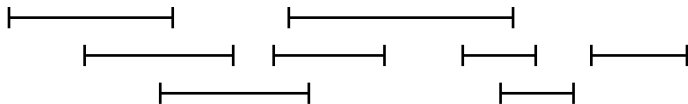
Linear Program

$$\begin{aligned}
 &\text{maximise} && \sum_{u \in C, v \in F} w_{uv} x_{uv} \\
 &\text{subject to} && \sum_{v \in F} x_{uv} \leq 1 \quad \forall u \in C \\
 & && x_{uv} \leq y_v \quad \forall u \in C, v \in F \\
 & && \sum_{v \in F: p \in I_v} y_v \leq 1 \quad \forall p \in P \\
 & && x_{uv} \geq 0 \quad \forall u \in C, v \in F \\
 & && y_v \geq 0 \quad \forall v \in F
 \end{aligned} \tag{LP}$$

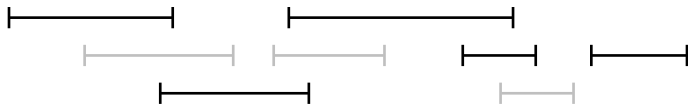
Overview of Results

- ▶ Motivated by genome mapping
 - ▶ Algorithms perform up to $1000\times$ faster
 - ▶ Comparable precision/recall
- ▶ Two new algorithms
 - ▶ Independent rounding: 0.19-approximation
 - ▶ Dependent rounding: Non-constant approximation
- ▶ Establish the hardness of the problem

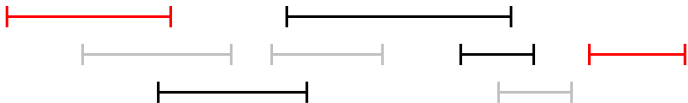
Independent Rounding: Setup



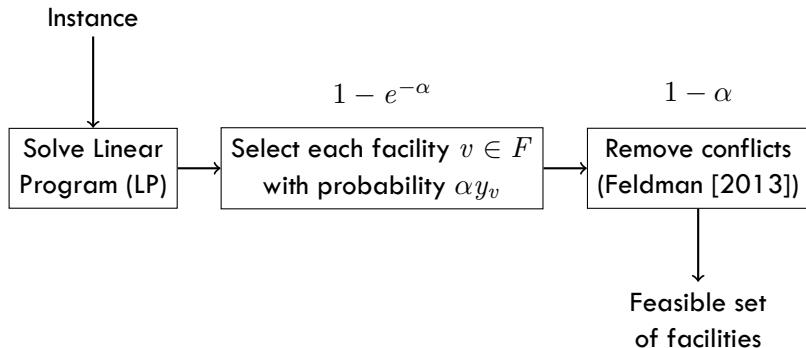
Independent Rounding: Random Selection



Independent Rounding: Filtering



Independent Rounding: Key Result



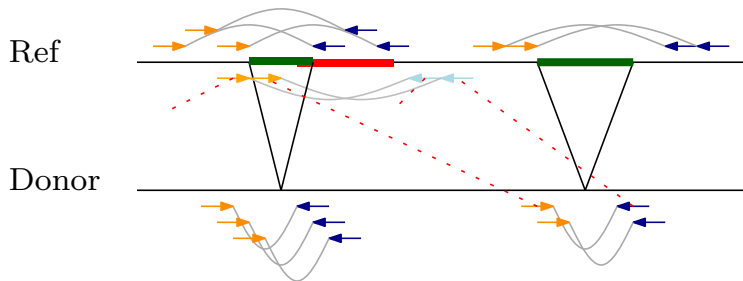
- ▶ $(1 - e^{-\alpha})(1 - \alpha)$ -approximation
- ▶ Choosing $\alpha = 0.44$ leads to 0.19-approximation

Hardness

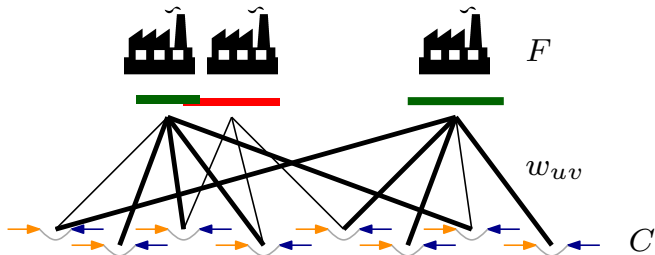
Theorem

For any $\varepsilon > 0$, there is no $(1 - 1/e + \varepsilon)$ -approximation algorithm for maximum facility location, unless $P = NP$.

Problem



Modelling the Problem



Experiments

- ▶ Two data sets
 - ▶ Craig Venter's genome (simulated)
 - ▶ NA12878 genome (real data, no ground truth)
- ▶ Compare precision/recall of predicted vs. actual deletions
- ▶ Two previous greedy algorithms by Hormozdiari et al. [2010]

Craig Venter

	Recall	Precision	Runtime (mins)
IP	65.80	40.26	10.4
Independent	65.02	40.63	15.8
Dependent	65.59	40.49	4.5
Greedy	65.44	41.47	1840.0
wGreedy	65.29	41.42	4095.0

NA12878

	Recall	Precision	Runtime (mins)
1000G (Mills et al. [2011])			
IP	60.3	69.3	10.5
Independent	60.2	69.11	7.0
Dependent	60.1	69.03	3.1
Greedy	60.2	69.79	1737.5
wGreedy	60.3	69.99	4947.2
long-read (Layer et al. [2014])			
IP	67.1	78.9	10.5
Independent	67.4	79.5	7.0
Dependent	67.4	79.4	3.1
Greedy	67.1	80.1	1737.5
wGreedy	66.8	79.8	4947.2

Concluding Statements

- ▶ Modelled the problem as an integer/linear program
- ▶ Developed two new approximation algorithms
- ▶ Established hardness of problem
- ▶ Biological application is faster than existing methods, without sacrificing solution quality

Future Work

- ▶ An algorithm with better approximation ratio
 - ▶ Gap between best known result and ≈ 0.63
 - ▶ Or a tighter hardness bound
- ▶ Analyse guarantee of non-constant approximation algorithm
- ▶ New conflict resolution heuristics

Questions?

Thanks to Stefan Canzar & Khaled Elbassioni

References

Moran Feldman. Maximization Problems with Submodular Objective Functions. PhD thesis, Israel Institute of Technology, 2013.

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