

236EX5

siweitang

August 2020

## 1 Proof of Correctness

Define  $P(n)$ : Let  $A$  be a list of integers, let  $p, q$  be integers, let  $n = q - p$ ,  $0 \leq p < q \leq \text{len}(A)$ ,  $\text{IndexMin}(A, p, q)$  terminates and returns the index of a minimum element in  $A[p: q]$ .

**Base case:** Let  $n = 1$ . So  $q - p = 1$ , that means  $A[p: q]$  contains only one element, that is the element with index  $p$ . So the smallest element in  $A[p: q]$  is the only element with index  $p$ . Hence  $P(1)$  holds.

**Inductive Step:** Let  $n \geq 1$ . So  $q - p \geq 1$ , it means that  $A[p: q]$  contains more than one elements. Let  $j \in N, 1 \leq j < n$ ,  $P(j)$  holds. WTP: for all  $n \in N, n \geq 1$ ,  $P(n)$  holds.

Since  $n \geq 1$ , so  $q - p \geq 1$ , the if statement on Line 1 is false. Goes to Line 4 else statement part. It runs Line 4-10.

On line 4,  $m = \lfloor \frac{p+q}{2} \rfloor \Rightarrow p < m < q$  (by hint)

So  $0 \leq p < m \leq \text{len}(A)$ .

Also, since  $m > p$ , so  $1 \leq m - p < n$ .

By line 5 and inductive hypothesis,  $j$  is the index of the smallest element in  $A[p: m]$

Since  $m < q$ , so  $0 \leq m < q \leq \text{len}(A)$ . Also,  $1 \leq q - m < n$

By line 6 and inductive hypothesis,  $k$  is the index of the smallest element in  $A[m: q]$

So there are two cases.

Case 1: if  $A[j] \leq A[k]$ , so  $j$  is the smallest element in  $A[p: q]$ , then by line 8,  $j$  is returned.

Case 2: if  $A[j] > A[k]$ , so  $k$  is the smallest element in  $A[p: q]$ , then by line 10,  $k$  is returned.

## 2 Proof

### 2.1 a

loop invariant LI(k):

- i) The elements of  $A[0: k]$  are rearranged in sorted(nondecreasing) order
- ii)  $0 \leq k \leq \text{len}(A) - 1$
- iii) All the elements of  $A[0: k]$  is less than or equal to all the elements of  $A[k: \text{len}(A) - 1]$

### 2.2 b

$m = \text{len}(A) - 1 - k$

### 2.3 c

Suppose the loop terminates. By the second condition of loop invariant,  $0 \leq k \leq \text{len}(A) - 1$

when the while loop ends,  $k \geq \text{len}(A) - 1$

$\Rightarrow k = \text{len}(A) - 1$

By the first condition of loop invariant, all the elements of  $A[0: k]$  are rearranged so they are sorted(nondecreasing) order.

Since  $k = \text{len}(A) - 1$ , so  $k$  is the last index of  $A \Rightarrow$  all the elements of  $A$  are rearranged in sorted(nondecreasing) order.