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8.8

- a) It is $O(\log n)$
- b) The worst case scenario for non-balanced tree is $O(n)$

8.9

- a) Asymptotic time complexity is $O(n^2)$
- b) Asymptotic space complexity is $O(1)$

8.10

- a) Asymptotic time complexity is $O(n)$. Making the assumption that the CPU is optimized to do the `std::swap` function in-place, since there will be no extra allocation of memory, the space complexity is $O(1)$;
- b) The time complexity is $O(n)$ due to the for-loop. Since the elements are copied into a vector, there is a space overhead of n , making the space complexity $O(n)$;
- c) Looking at both the time complexity and the space complexity, they both have the same time complexity but `reverse_array_1` has a better space complexity. Hence, `reverse_array_1` is a better function to use.

8.12

Overall speedup for each part:

A:

$$1/((1-0.05)+(0.05/10)) = 1.047$$

B:

$$1/((1-0.5)+(0.5/1.05)) = 1.024$$

C :

$$1/((1-0.1)+(0.1/3)) = 1.071$$

The overall speedup is the most in case of C, hence C should be optimized.

8.13

- a) Here, input is the number of bits in the integer. Hence, the time complexity is $O(n)$. There aren't any extra space being used, so the space complexity is $O(1)$.
- b)

Unsigned int `hamming_2(unsigned int n)`

```
{  
    return( arr([n & 0xFFFF] + arr[n >>16] ));  
}
```

Here, arr is an array storing the bitcounts of all the unsigned integers, from min to max inclusive. Since the array lookup is constant, the time complexity here is $O(1)$.

This is a significant increase in the time complexity, which is an advantage. However, the space complexity will be unreasonably large due to storing the bit counts in an array lookup table; this is disadvantage.

c) The input parameter is an unsigned int, which gives a large number as an input value. Although the numbers are not going to infinity, I believe the n is sufficiently large enough to consider asymptotic analysis in situation like this.