



# CHANDIGARH UNIVERSITY

Discover. Learn. Empower.

## PROJECT REPORT

### Subarray Summit – Maximum Subarray Problem

**Subject Name:** Design and Analysis of Algorithms **Subject**  
**Subject Code :** 23CSH-301

**Submitted to:**

Er. Mohammad Shaqlain  
(E17211)

**Submitted by:**

Name: Tanisha Kumari  
UID: 23BCS12542  
Section: KRG\_2B



## Aim:

To design and implement algorithms to solve the Maximum Subarray Problem, which involves finding the contiguous subarray within a one-dimensional array that has the largest sum.



## Objectives:

1. To understand and implement different algorithmic approaches for solving the maximum subarray problem.
2. To compare the efficiency of Brute Force, Divide and Conquer, and **Kadane's Algorithm**.
3. To analyze the time complexity and identify the most optimal solution.
4. To demonstrate the use of algorithm design paradigms like **Divide & Conquer** and **Dynamic Programming**.



## Problem Definition

Given an array of integers (both positive and negative), find the contiguous subarray which has the maximum possible sum.

### **Example:**

Input Array = [-2, -3, 4, -1, -2, 1, 5, -3]

Output = 7

(Contiguous subarray = [4, -1, -2, 1, 5])



## Algorithms Used

We implement three algorithms to show efficiency improvement:

## 1. Brute Force Approach

Compute the sum of all possible subarrays and track the maximum sum.

**Time Complexity:**  $O(n^2)$

**Space Complexity:**  $O(1)$

---

## 2. Divide and Conquer Approach

Divide the array into two halves.

Find:

Maximum subarray sum in the left half

Maximum subarray sum in the right half

Maximum subarray sum that crosses the midpoint

Return the maximum of these three.

**Recurrence Relation:**

$$T(n) = 2T(n/2) + O(n)$$

**Solving gives:**  $T(n) = O(n \log n)$

**Space Complexity:**  $O(\log n)$  (due to recursion stack)

---

## 3. Kadane's Algorithm (Dynamic Programming)

Traverse array, maintaining a running sum and reset it to 0 if it becomes negative.

Keep track of the maximum sum seen so far.

**Time Complexity:**  $O(n)$

**Space Complexity:**  $O(1)$



## C++ Implementation

```
#include <iostream>
```

```
#include <climits>
```

```
using namespace std;
```

```
// ----- BRUTE FORCE APPROACH -----
```

```
int maxSubarrayBruteForce(int arr[], int n) {  
    int maxSum = INT_MIN;  
    for (int i = 0; i < n; i++) {  
        int currentSum = 0;  
        for (int j = i; j < n; j++) {  
            currentSum += arr[j];  
            maxSum = max(maxSum, currentSum);  
        }  
    }  
    return maxSum;  
}
```

```
// ----- DIVIDE AND CONQUER APPROACH -----
```

```
int maxCrossingSum(int arr[], int l, int m, int h) {  
    int sum = 0;  
    int left_sum = INT_MIN;  
    for (int i = m; i >= l; i--) {  
        sum += arr[i];  
        left_sum = max(left_sum, sum);  
    }  
  
    sum = 0;  
    int right_sum = INT_MIN;
```

```

        for (int i = m + 1; i <= h; i++) {
            sum += arr[i];
            right_sum = max(right_sum, sum);
        }

        return left_sum + right_sum;
    }

int maxSubarrayDivideConquer(int arr[], int l, int h) {
    if (l == h)
        return arr[l];

    int m = (l + h) / 2;

    int left_max = maxSubarrayDivideConquer(arr, l, m);
    int right_max = maxSubarrayDivideConquer(arr, m + 1, h);
    int cross_max = maxCrossingSum(arr, l, m, h);

    return max({left_max, right_max, cross_max});
}

// ----- KADANE'S ALGORITHM -----
int maxSubarrayKadane(int arr[], int n) {
    int maxSoFar = INT_MIN, maxEndingHere = 0;

    for (int i = 0; i < n; i++) {

```

```

        maxEndingHere += arr[i];
        maxSoFar = max(maxSoFar, maxEndingHere);

        if (maxEndingHere < 0)
            maxEndingHere = 0;
    }

    return maxSoFar;
}

// ----- MAIN FUNCTION -----
int main() {
    int arr[] = {-2, -3, 4, -1, -2, 1, 5, -3};
    int n = sizeof(arr) / sizeof(arr[0]);

    cout << "Array: ";
    for (int x : arr) cout << x << " ";
    cout << "\n\n";

    cout << "Brute Force Max Sum: " << maxSubarrayBruteForce(arr, n) << endl;
    cout << "Divide & Conquer Max Sum: " << maxSubarrayDivideConquer(arr, 0,
n - 1) << endl;

    cout << "Kadane's Algorithm Max Sum: " << maxSubarrayKadane(arr, n) <<
endl;

    return 0;
}

```

```
Array: -2 -3 4 -1 -2 1 5 -3
```

```
Brute Force Max Sum: 7  
Divide & Conquer Max Sum: 7  
Kadane's Algorithm Max Sum: 7
```

Figure 1: Output



## Complexity Analysis

Approach	Time Complexity	Space Complexity	Remarks
Brute Force	$O(n^2)$	$O(1)$	Simple but inefficient
Divide & Conquer	$O(n \log n)$	$O(\log n)$	Uses recursion
Kadane's Algorithm	$O(n)$	$O(1)$	Most efficient

### ✓ Recurrence Relation (for Divide & Conquer)

$$T(n) = 2T(n/2) + O(n)$$

By applying the **Master Theorem**,

$$T(n) = O(n \log n)$$



## Applications

1. **Financial Analysis:** To identify the most profitable period in stock price changes.
2. **Data Analytics:** Detect the most positive trend in sensor or signal data.

3. **Machine Learning:** Feature extraction for time-series data.
  4. **Genomics:** Identifying regions with maximum activity in DNA sequences.
  5. **Weather Forecasting:** Finding longest warm or cold streaks.
- 



## Conclusion

The **Maximum Subarray Problem** efficiently demonstrates the evolution of algorithmic design:

- From **Brute Force ( $O(n^2)$ )**
- To **Divide & Conquer ( $O(n \log n)$ )**
- To **Kadane's Algorithm ( $O(n)$ )**

Kadane's Algorithm provides the most optimal solution, making it ideal for large datasets and real-world applications where quick computation is essential.