# **HPC ASSIGNMENT 2**

# **Group 14**

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## 1)Merge sort technique (by considering bottom up approach)

Suppose we wish to redesign merge sort to run on a parallel computing platform. Just as it it useful for us to abstract away the details of a particular programming language and use pseudocode to describe an algorithm, it is going to simplify our design of a parallel merge sort algorithm to first consider its implementation on an abstract PRAM machine. However, we will also consider the realities of practical platforms and discuss a likely version that would get implemented in practice.

Suppose we have a PRAM machine with  $\mathbf{n}$  processors. A theoretically possible case, but unlikely in practice, would be that that we would have plenty of processors, so that  $\mathbf{n} >= \mathbf{N}$ , the size of our data to be sorted. This relatively impractical case is referred to as fine-grained parallelism. Fine-grained case provides a useful starting point for eventually arriving at course-grained solutions.

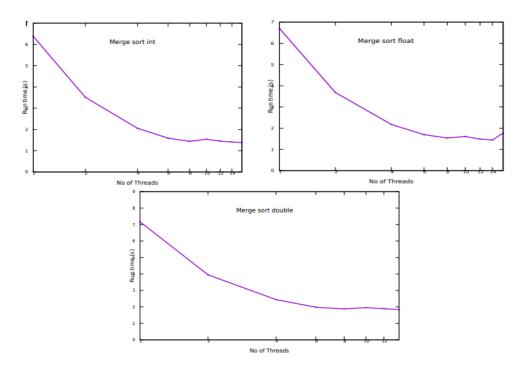
When thinking about the sequential merge sort algorithm, a way to visualize what is happening and reason about its complexity is to look at its recursion tree. For example, if there were 8 items to be sorted, The number of items each recursive call would be working on. The merge function is called on every node but the leaves of this tree, where the input m is a single element.

The sequential steps of this algorithm are taking place by following a depth-first traversal through this tree following the left children first. Take a moment to visualize, starting from the top node, which node begins executing the next merge sort function, and the next, and so on.

Now let's examine which of these operations that were running one after the other in the sequential version could be run simultaneously on separate processors. A natural way to split the work that can be run 'in parallel' is to do the work required at each level of the tree. Note that when considering a parallel solution, we use an **iterative approach** and envision starting our work at the bottom of the tree, moving up one level at each iteration. Each individual process is simply executing a merge on a range of values in the array (a single array could be used to sort in place, or a result could be used if you wish to preserve the original input).

We have 8 elements to be sorted and what would result at each level of the tree. At the level of the leaves of the tree, there is no real work, and we could begin by envisioning that level as our N data

items, all shared in memory by our processors. At the next level, each of 4 processors can work on disjoint sets of 2 separate data items and merge them together. Once those are done, 2 processors can each work on merging to create a list of 4 sorted items, and finally, one last processor can merge those 2 lists of 4 to create the final list of 8 sorted items. This is parallel merge sort algorithm.



#### 2) N-Queens on the N $\times$ N chessboard in non-attackable positions, N $\in$ {10, 12, 14, 16}.

#### N-queens Problem sequential Solution

The N-queens problem is a generalization of the well-known 8-queens problem, which consists in arranging 8 queens on a chessboard so that none can take another. A queen attacks another if they are located on the same diagonal, row or column. In the case of the N-queens, N queens are placed on a NxN board. There exists a known number of solutions; for instance, there are 92 solutions for placing 8 queens on a 8x8 board . Other problems are derivable from this. Among them, it is worth mentioning the problem which, given a NxN board, looks for the lesser quantity of queens that can be placed so that all the board squares are attacked by some queen . An initial solution to the N-queens problem, by way of a sequential algorithm, consists in testing all the possible placement combinations of queens on the board and choosing the valid ones. The combination in which no queen of the board is attacked by another is considered as valid. This solution can be upgraded by discarding, during the search, those ways by which a solution to the problem cannot be found . We used 03 optimization to get values faster.

### Parallelization of the N-queens problem

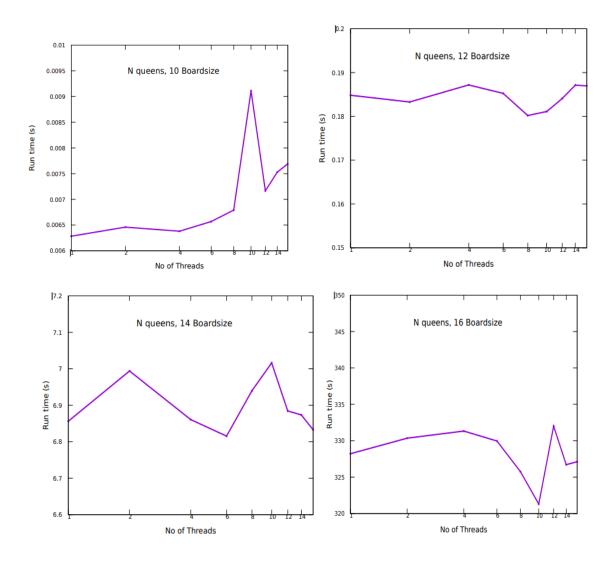
All the algorithms perform N stages where, in each stage I, they try to place queens on all the valid position of row i.

1- Solution with processor pipe: Let P1..PN be processors and Ti the set of boards with queens placed on valid positions on the first i rows. A processor Pi is in charge of placing the queens on row

i. The work begins with processor P1, which places a queen on each possible position of row 1, thus obtaining a set of initial boards (T1). Each of these obtained boards passes to processor P2, which places a queen on all the valid positions of row 2 in each of them, passing the obtained boards (T2) to the next processor. The process is repeated until N processor is reached. This receives from processor PN-1 the boards with queens located accurately on the first N-1 rows (TN-1), and its task is to place the queens on row N, in order to obtain the set of solutions to the problem (TN).

#### 2.-Loosely coupled processor N Solution

Let P1..PN be processors and Ti,j the set of boards where i represents the column in which the queen was placed on the first row, and j the row up to which the board has queens. For instance: T2,4 is the set of all the boards with queens placed accurately on the first four rows, being the queen of the first row placed on the second column. A processor Pi is in charge of finding the set of boards (Ti,N) with all the possible solutions having placed the queen on the first row column i. In this case there are no communication between the processors during the search of the solutions. Each processor works independently.



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necque-NP-Pro-SFF-409-G9-Besktop-PC: /DesktopS gcc -o3 n_queens.c -o queens.o -fopennp
mccque-NP-Pro-SFF-409-G9-Besktop-PC: /DesktopS ./queens.o
humber of threads: 1
Itte taken: 6.00020271
humber of solutions: 724
Humber of solutions: 724
Humber of solutions: 727
Humber of threads: 8
humber of solutions: 724
Humber of solutions: 725
Humber of solutions: 724
Humber of solutions: 726
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