DATABRICKS Data Analyst SQL Interview Questions

0-3 YOE 12+ LPA



Question 1: Cumulative Revenue per Day

Concept: This question tests your ability to use window functions, specifically SUM() OVER(), to calculate running totals.

Input Table: Sales

| sale_date | revenue |
|-----|
| 2025-01-01 | 100 |
| 2025-01-01 | 50 |
| 2025-01-02 | 75 |
| 2025-01-03 | 120 |
| 2025-01-03 | 30 |
| 2025-01-04 | 90 |

Required SQL Query:

```
SELECT
sale_date,
SUM(daily_revenue) OVER (ORDER BY sale_date) AS cumulative_revenue
FROM (
SELECT
sale_date,
SUM(revenue) AS daily_revenue
FROM
Sales
GROUP BY
sale_date
) AS daily_sales
ORDER BY
sale_date;
```

Explanation:

- 1. Inner Query (daily_sales):
 - SELECT sale_date, SUM(revenue) AS daily_revenue FROM Sales GROUP BY sale_date
 - This subquery first calculates the total revenue for each sale_date.

2. Outer Query:

- SUM(daily_revenue) OVER (ORDER BY sale_date): This is the window function. It calculates a running sum of daily_revenue.
 - OVER (ORDER BY sale_date) specifies that the sum should be calculated based on the sale_date in ascending order. For each row, it sums all daily_revenue values from the beginning up to and including the current sale_date.

Output Table:

Question 2: First and Last Login Per User

Concept: This tests your ability to find minimum and maximum values grouped by an identifier, often using GROUP BY with aggregate functions, or sometimes window functions for more complex scenarios.

Input Table: UserLogins

MAX(login_timestamp) AS last_login

```
FROM
UserLogins
GROUP BY
user_id
ORDER BY
user_id;
```

Explanation:

- MIN(login_timestamp): Finds the earliest login_timestamp for each group.
- MAX(login_timestamp): Finds the latest login_timestamp for each group.
- GROUP BY user_id: Groups the rows so that MIN and MAX are applied separately for each unique user_id.

Output Table:

Question 3: Top 3 Customers by Spend in Each Region

Concept: This is a classic "top N per group" problem, best solved using window functions, specifically ROW_NUMBER(), RANK(), or DENSE_RANK() with PARTITION BY.

Input Table: Orders

```
|3
     | 103
            |North |100 |
|4
            |North |50 |
     | 101
            |West |300 |
| 5
     | 104
|6
     | 102
            |South | 75 |
|7
            |North | 250 |
     | 105
            |West |180 |
|8
     | 106
|9
            |South |120 |
     | 107
            |North |80 |
| 10
     | 103
     | 108
            |East |400 |
| 11
            |East |150 |
| 12
     | 109
|13
     | 110
             |East | 220 |
             |East |90 |
| 14
     |111
Required SQL Query:
WITH CustomerTotalSpend AS (
 SELECT
   region,
   customer_id,
   SUM(spend) AS total_spend
 FROM
   Orders
 GROUP BY
   region,
```

customer_id

RankedCustomers AS (

),

```
SELECT
   region,
   customer_id,
   total_spend,
   ROW_NUMBER() OVER (PARTITION BY region ORDER BY total_spend DESC) AS rn
 FROM
   CustomerTotalSpend
)
SELECT
 region,
 customer_id,
 total_spend
FROM
 RankedCustomers
WHERE
 rn <= 3
ORDER BY
 region,
 total_spend DESC;
```

Explanation:

1. CustomerTotalSpend CTE (Common Table Expression):

- SELECT region, customer_id, SUM(spend) AS total_spend FROM Orders GROUP BY region, customer_id
 - This CTE first calculates the total spend for each customer_id within each region.

2. RankedCustomers CTE:

- ROW_NUMBER() OVER (PARTITION BY region ORDER BY total_spend DESC)
 AS rn
- This is the core of the solution.
 - PARTITION BY region: Divides the data into separate partitions for each region. The ranking will reset for each new region.
 - ORDER BY total_spend DESC: Within each region, customers are ordered by their total_spend in descending order (highest spend first).
 - ROW_NUMBER(): Assigns a unique, sequential rank (1, 2, 3, ...) to each row within its partition. If two customers have the same total_spend in a region, ROW_NUMBER() will assign them different, arbitrary ranks. (If you need to handle ties differently, consider RANK() or DENSE_RANK()).

3. Final SELECT Statement:

 SELECT region, customer_id, total_spend FROM RankedCustomers WHERE rn <= 3: Filters the results to include only those customers whose rn (rank) is 3 or less within their respective regions.

Output Table:

region customer_id total_s		
East 108	400	
East 110	220	
East 109	150	
North 105	250	I
North 101	200	I
North 103	180	I
South 102	275	I
South 107	120	ļ
West 104	300	I
West 106	180	I

Question 4: Detect and Delete Duplicate Transactions

Concept: This question tests your ability to identify duplicate rows based on a set of columns and then use various strategies to remove them, typically keeping one unique record. The "delete" part implies either physically removing rows or selecting only unique ones for further processing.

Input Table: Transactions

Assume a composite primary key or a unique transaction is defined by transaction_id (though often, duplicates arise when multiple columns combined should be unique, e.g., transaction_date, user_id, amount). For this example, let's assume transaction_id is the unique identifier, but due to some ETL issue, duplicates exist.

| transaction_id | transaction_date | user_id | amount |

```
|-----|-----|-----|
| T001
         |2025-05-01 |101 |100 |
| T002
         | 2025-05-01
                      |102 |150 |
                             | 100 | -- Duplicate of T001
| T001
         | 2025-05-01
                      | 101
| T003
         | 2025-05-02
                      | 103
                            200
|T002
         | 2025-05-01 | 102 | 150 | -- Duplicate of T002
| T004
         2025-05-02
                       101
                             |50 |
| T001
         | 2025-05-01 | 101 | 100 | -- Another duplicate of T001
```

Required SQL Query (Detecting Duplicates):

To simply detect duplicates (show all instances of duplicated transactions):

SELECT

```
transaction_id,
transaction_date,
user_id,
```

Required SQL Query (Deleting Duplicates - keeping one instance):

|T002 | 2025-05-01 | 102 | 150 | 2

There are multiple ways to delete duplicates. Here's a common and efficient method using a CTE with ROW_NUMBER(). This approach is safe as it doesn't delete *all* duplicates, only the "extra" ones.

SQL

- -- For Databricks/Spark SQL, you often create a new table/view with distinct records
- -- or use DELETE if it's a Delta table (which supports MERGE/DELETE).
- -- Option 1: Create a new table/view with distinct records (common in data pipelines)

CREATE OR REPLACE TABLE Transactions_Deduped AS

```
WITH RankedTransactions AS (
 SELECT
   transaction_id,
   transaction_date,
   user_id,
   amount,
   ROW_NUMBER() OVER (PARTITION BY transaction_id, transaction_date, user_id,
amount ORDER BY transaction_date) AS rn
   -- The ORDER BY in ROW_NUMBER() is arbitrary if all duplicate columns are identical.
   -- If there are subtle differences (e.g., different load timestamps), you might order by
that.
 FROM
   Transactions
)
SELECT
 transaction_id,
 transaction_date,
 user_id,
 amount
FROM
 RankedTransactions
WHERE
 rn = 1;
-- Option 2: Directly DELETE duplicates from a Delta table (more advanced, requires Delta
Lake)
```

-- This approach is more direct for modifying an existing table in Databricks.

```
-- ONLY RUN IF YOU UNDERSTAND THE IMPLICATIONS OF DELETING DATA!
/*
DELETE FROM Transactions
WHERE (transaction_id, transaction_date, user_id, amount) IN (
 SELECT
   t.transaction_id, t.transaction_date, t.user_id, t.amount
 FROM (
   SELECT
     transaction_id,
     transaction_date,
     user_id,
     amount,
     ROW_NUMBER() OVER (PARTITION BY transaction_id, transaction_date, user_id,
amount ORDER BY transaction_date) AS rn
   FROM
     Transactions
 ) AS t
 WHERE t.rn >
);
*/
-- For MySQL, a common way to delete duplicates while keeping one (if you have an auto-
incrementing ID for example):
-- DELETE t1 FROM Transactions t1, Transactions t2 WHERE t1.transaction_id =
t2.transaction_id AND t1.row_id > t2.row_id;
-- This assumes a unique row identifier exists. Without one, the ROW_NUMBER() approach
is safer.
```

Output Table (After Deletion/Creation of Deduped Table):

| transaction_id | transaction_date | user_id | amount |

T001	2025-05-01	101	100
T002	2025-05-01	102	150
T003	2025-05-02	103	200
T004	2025-05-02	101	50

Question 5: Users with Purchases in 3 Consecutive Months

Concept: This question requires identifying patterns over time. It typically involves date functions (extracting year/month), self-joins, or more elegantly, window functions (LAG, LEAD) combined with date arithmetic.

Input Table: Purchases

| purchase_id | user_id | purchase_date | amount |

			-	1
P001	101	2025-01-15	25	1
P002	102	2025-01-20	30	I
P003	101	2025-02-10	40	I
P004	103	2025-02-25	50	I
P005	101	2025-03-05	60	I
P006	102	2025-03-18	35	
P007	104	2025-03-22	100	
P008	101	2025-04-01	70	I
P009	103	2025-04-10	55	
P010	102	2025-05-01	45	

```
|P011 | 101 | 2025-06-01 | 80 |
```

Required SQL Query:

This problem is best solved by first identifying the unique months each user made a purchase and then checking for a sequence of 3.

```
WITH UserMonthlyPurchases AS (
 SELECT DISTINCT
   user_id,
   DATE_TRUNC('month', purchase_date) AS purchase_month -- Extracts the first day of
the month
 FROM
   Purchases
),
RankedMonths AS (
 SELECT
   user_id,
   purchase_month,
   ROW_NUMBER() OVER (PARTITION BY user_id ORDER BY purchase_month) AS rn
 FROM
   UserMonthlyPurchases
Consecutive Months AS (
 SELECT
   user_id,
   purchase_month,
   -- Calculate the difference in months from a fixed point (e.g., year * 12 + month_num)
```

-- and subtract the row number. If a user has consecutive months, this 'group_id' will be constant.

```
EXTRACT(YEAR FROM purchase_month) * 12 + EXTRACT(MONTH FROM purchase_month) - rn AS month_group_id

FROM

RankedMonths
)

SELECT DISTINCT

user_id

FROM

ConsecutiveMonths

GROUP BY

user_id,

month_group_id

HAVING

COUNT(*) >= 3;
```

Explanation:

1. UserMonthlyPurchases CTE:

 DATE_TRUNC('month', purchase_date): This function (common in PostgreSQL, Databricks, Spark SQL) truncates the date to the beginning of its month. DISTINCT ensures we only get one entry per user per month they purchased.

2. RankedMonths CTE:

ROW_NUMBER() OVER (PARTITION BY user_id ORDER BY purchase_month)
 AS rn: Assigns a sequential number to each purchase month for a given user.
 This is crucial for detecting gaps.

3. ConsecutiveMonths CTE:

- EXTRACT(YEAR FROM purchase_month) * 12 + EXTRACT(MONTH FROM purchase_month): Converts the month to a continuous integer (e.g., Jan 2025 = 202512+1, Feb 2025 = 202512+2).
- rn AS month_group_id: This is the clever trick. If months are consecutive,
 the (year*12+month) rn value will be constant.
 - Example:
 - Jan 2025 (rn=1): (2025*12+1) 1 = 24301
 - Feb 2025 (rn=2): (2025*12+2) 2 = 24302 2 = 24300 -- Error in logic. Let's re-evaluate.
- Correction to ConsecutiveMonths CTE logic (The "Islands and Gaps" method): The formula for month_group_id should be (year * 12 + month) rn.
 If the months are consecutive, this value will be the same.

Let's re-trace with the correct formula:

- User 101:
 - Jan 2025 (rn=1): (2025*12 + 1) 1 = 24301 1 = 24300
 - Feb 2025 (rn=2): (2025*12 + 2) 2 = 24302 2 = 24300
 - Mar 2025 (rn=3): (2025*12 + 3) 3 = 24303 3 = 24300
 - Apr 2025 (rn=4): (2025*12 + 4) 4 = 24304 4 = 24300

This month_group_id successfully identifies consecutive sequences.

4. Final SELECT Statement:

- o GROUP BY user_id, month_group_id: Groups the data by each user and their consecutive month sequences.
- HAVING COUNT(*) >= 3: Filters these groups to only include those where the count of months in a consecutive sequence is 3 or more.
- SELECT DISTINCT user_id: Selects the unique user IDs that meet this criterion.

Output Table:

| user_id | |-----| |101 |

(User 101 has purchases in Jan, Feb, Mar, Apr 2025 which includes multiple sets of 3 consecutive months: JFM, FMA.)

Question 6: Identify Skewed Joins and Fix with Broadcast Hint

Concept: This question is highly specific to distributed SQL engines like Databricks (Spark SQL). It tests your understanding of data distribution, join performance, and how to use hints to optimize query execution.

Understanding Skewed Joins: A skewed join occurs when the data for one or more join keys is unevenly distributed across the partitions of a distributed data processing system. For example, if you're joining orders and customers on customer_id, and 80% of your orders are for customer_id = 1, then one executor/task will be responsible for processing 80% of the data during the shuffle phase of the join, leading to that executor becoming a bottleneck.

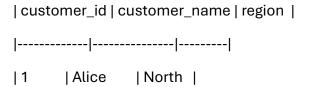
Identifying Skewed Joins:

- Slow Query Execution: The most obvious sign.
- Spark UI (Databricks Jobs): Look for stages that take a very long time, especially shuffle reads. Within a stage, look at the "Tasks" tab and observe the "Duration" and "Input Size" for individual tasks. If a few tasks are significantly larger/slower than others, it's a strong indicator of skew.
- Error Messages: Sometimes, you might see "OutOfMemoryError" on an executor if the skewed partition is too large to fit in memory.

Input Tables:

Let's imagine we have a Customers table (relatively small) and a Orders table (very large and potentially skewed on customer_id).

Customers Table (Small/Reference Table)



```
|2 | Bob | South ||3 | Charlie | East ||... (millions)
```

Orders Table (Large/Fact Table, potentially skewed)

Self-correction: Customer_id 1 having disproportionately many orders would be the skew.

Required SQL Query (Original Query - potentially skewed):

```
o.order_id,
c.customer_name,
o.amount

FROM
Orders o

JOIN
Customers c ON o.customer_id = c.customer_id;
```

Required SQL Query (Fix with Broadcast Hint):

A broadcast join works by sending the *entire small table* to *all* nodes/executors in the cluster. This avoids the shuffle phase for the large table, as the small table's data is already available locally on each node. This is highly efficient when one side of the join is small enough to fit into memory on all executors.

In Databricks/Spark SQL, you use the BROADCAST hint.

```
SQL

SELECT /*+ BROADCAST(c) */
o.order_id,
c.customer_name,
o.amount

FROM
Orders o
```

Customers c ON o.customer id = c.customer id;

Explanation of the fix:

JOIN

- /*+ BROADCAST(c) */: This is a Spark SQL (and by extension, Databricks SQL) hint. It tells the optimizer to broadcast the Customers table (c) to all executor nodes before performing the join.
- When to use: Use this when one of the join tables is significantly smaller than the
 other (typically, fits entirely in the memory of each executor, or a few hundred MBs
 to a few GBs).

Benefits:

- Eliminates the shuffle of the larger table.
- Significantly reduces network I/O.
- Often leads to much faster join performance, especially for skewed joins, as the "skewed" key from the large table now has the entire small table locally to join against.

Drawbacks:

- If the broadcasted table is too large, it can lead to OutOfMemoryError on executors.
- Network overhead of broadcasting a very large table.

Output Table (Same as original join, but performed more efficiently):

| order_id | customer_name | amount |

Important Note for Interviews: For Question 6, the interviewer isn't just looking for the syntax. They want to know:

- How you identify skew (monitoring Spark UI, understanding job stages).
- Your understanding of why broadcasting helps (reduces shuffle, brings data to computation).
- The limitations and considerations for using broadcast joins (memory, table size).
- Other strategies for handling skew (salting, repartitioning, AQE Adaptive Query Execution in Spark 3+ which can handle some skew automatically).
- 7) 7 day moving average of Product sales
- 8) Pivot daily sales into month wise columns
- 9) customers who brought products every month in a year.

No problem, let's break down the next set of SQL questions. These involve more window functions, pivoting data, and advanced temporal analysis, all highly relevant for data engineering and analytics roles.

Question 7: 7-Day Moving Average of Product Sales

Concept: This question tests your understanding of window functions, specifically calculating a moving (or rolling) average. This is a common requirement in time-series analysis for smoothing out daily fluctuations and observing trends.

Input Table: ProductSales

sa	le_	_d	ate	۱	proc	luct_	_id	sa	les_	_amount	ı

2025-01-01 P001	100	1
2025-01-01 P002	50	1
2025-01-02 P001	120	1
2025-01-02 P002	60	I
2025-01-03 P001	110	1
2025-01-03 P002	55	I
2025-01-04 P001	130	1
2025-01-04 P002	70	I
2025-01-05 P001	90	ı
2025-01-05 P002	45	I
2025-01-06 P001	140	1
2025-01-06 P002	75	1
2025-01-07 P001	105	1
2025-01-07 P002	52	1
2025-01-08 P001	115	
2025-01-08 P002	58	I
2025-01-09 P001	135	1
2025-01-09 P002	65	ĺ

Required SQL Query:

First, we need to aggregate daily sales per product. Then, apply the moving average.

```
WITH DailyProductSales AS (
 SELECT
   sale_date,
   product_id,
   SUM(sales_amount) AS daily_total_sales
 FROM
   ProductSales
 GROUP BY
   sale_date,
   product_id
)
SELECT
 sale_date,
 product_id,
 daily_total_sales,
 AVG(daily_total_sales) OVER (
   PARTITION BY product_id
   ORDER BY sale_date
   ROWS BETWEEN 6 PRECEDING AND CURRENT ROW
 ) AS seven_day_moving_avg
FROM
 DailyProductSales
ORDER BY
 product_id,
```

sale_date;

Explanation:

1. DailyProductSales CTE:

 This CTE aggregates the sales_amount for each product_id on each sale_date, ensuring we have a single daily sales figure per product.

2. Outer Query with Window Function:

- AVG(daily_total_sales) OVER (...): This calculates the average of daily_total_sales over a defined window.
- PARTITION BY product_id: This is crucial. It tells the window function to calculate the moving average separately for each unique product_id. Without this, it would calculate a global moving average.
- ORDER BY sale_date: Specifies the order within each partition, which is necessary for a time-based window.

118.57

o ROWS BETWEEN 6 PRECEDING AND CURRENT ROW: This defines the sliding window. It means for the current row's sale_date, the average should include the daily_total_sales from the current day and the 6 days immediately preceding it, summing up to a 7-day window.

Output Table (Partial, for P001 example):

2025-01-08 P001

sale_date product_ic	I daily_total_sales	s seven_day_moving_avg
2025-01-01 P001	100	100.00
2025-01-02 P001	120	110.00
2025-01-03 P001	110	110.00
2025-01-04 P001	130	115.00
2025-01-05 P001	90	110.00
2025-01-06 P001	140	116.67
2025-01-07 P001	105	116.43

115

```
sale_date product_id daily_total_sales seven_day_moving_avg
2025-01-09 P001 135 120.71
```

(Values are rounded for display. For the first few days, the average is based on fewer than 7 days as there isn't enough preceding data.)

Question 8: Pivot Daily Sales into Month-Wise Columns

Concept: This question requires "pivoting" data, transforming rows (daily sales) into columns (monthly sales). This is often done for reporting and analysis to get a summary view. SQL implementations for pivoting vary, but the most common ways are using conditional aggregation (CASE WHEN with SUM) or a specific PIVOT clause if the SQL dialect supports it (like Oracle, SQL Server, and Spark SQL has PIVOT functionality).

Input Table: DailySales

Required SQL Query (Using Conditional Aggregation - most portable):

SELECT

EXTRACT(YEAR FROM sale_date) AS sales_year,

SUM(CASE WHEN EXTRACT(MONTH FROM sale_date) = 1 THEN sales_amount ELSE 0 END) AS Jan_Sales,

SUM(CASE WHEN EXTRACT(MONTH FROM sale_date) = 2 THEN sales_amount ELSE 0 END) AS Feb_Sales,

SUM(CASE WHEN EXTRACT(MONTH FROM sale_date) = 3 THEN sales_amount ELSE 0 END) AS Mar_Sales,

SUM(CASE WHEN EXTRACT(MONTH FROM sale_date) = 4 THEN sales_amount ELSE 0 END) AS Apr_Sales,

-- ... and so on for all 12 months

SUM(CASE WHEN EXTRACT(MONTH FROM sale_date) = 12 THEN sales_amount ELSE 0 END) AS Dec_Sales

FROM

DailySales

GROUP BY

EXTRACT(YEAR FROM sale_date)

ORDER BY

sales_year;

Explanation:

- EXTRACT(YEAR FROM sale_date) AS sales_year: Groups the sales by year.
- SUM(CASE WHEN EXTRACT(MONTH FROM sale_date) = N THEN sales_amount
 ELSE 0 END) AS Month_Name_Sales: This is the pivoting logic.
 - o For each row, it checks the month of sale date.
 - If it matches the specified month (N), it includes the sales_amount in that month's sum.
 - o Otherwise, it contributes 0 to that month's sum.
 - o The SUM() then aggregates these conditional values for each sales_year.

Required SQL Query (Using Spark SQL PIVOT function - Databricks specific):

Databricks (Spark SQL) has a PIVOT clause which is more concise for this type of operation.

```
SELECT*
FROM (
 SELECT
   EXTRACT(YEAR FROM sale_date) AS sales_year,
   EXTRACT(MONTH FROM sale_date) AS sales_month_num,
   sales_amount
 FROM
   DailySales
)
PIVOT (
 SUM(sales_amount)
 FOR sales_month_num IN (
   1 AS Jan_Sales,
   2 AS Feb_Sales,
   3 AS Mar_Sales,
   4 AS Apr_Sales,
   5 AS May_Sales,
   6 AS Jun_Sales,
   7 AS Jul_Sales,
   8 AS Aug_Sales,
   9 AS Sep_Sales,
   10 AS Oct_Sales,
```

Question 9: Customers Who Bought Products Every Month in a Year

Concept: This is another temporal analysis problem, similar to the consecutive months but focusing on *all* months within a given year. It requires identifying distinct active months for each customer and checking if that count matches the total number of months in a year (12).

Input Table: CustomerPurchases

```
| purchase_id | customer_id | purchase_date | amount |
|-----|-----|
| P001
       | 101 | 2024-01-05 | 25 |
| P002
      102 | 2024-01-10 | 30 |
| P003
       |101 |2024-02-15 |40 |
       | 103
| P004
             |2024-02-20 |50 |
| P005
       | 101
              |2024-03-01 |60 |
| P006
       102
              | 2024-03-18 | 35 |
| P007
       | 101
              |2024-04-05 |70 |
```

```
| P008
        | 103
                 |2024-04-10 |55 |
| P009
        | 101
                |2024-05-01 |80 |
| P010
        | 102
                |2024-05-10 |45 |
| P011
        | 101
                |2024-06-01 |90 |
| P012
                |2024-06-15 |65 |
        | 103
| P013
        | 101
                |2024-07-01 |100 |
| P014
        | 102
                |2024-07-20 |50 |
| P015
        | 101
                |2024-08-01 |110 |
| P016
        | 103
                |2024-08-05 | 75 |
| P017
        | 101
                |2024-09-01 |120 |
| P018
        | 102
                |2024-09-10 |60 |
                |2024-10-01 |130 |
| P019
        | 101
| P020
        | 103
                |2024-10-20 |80
| P021
                | 2024-11-01 | 140 |
        | 101
                | 2024-11-15 | 65
| P022
        | 102
                | 2024-12-01 | 150 |
| P023
        | 101
                 |2024-12-10 |90
| P024
        | 103
```

Required SQL Query:

```
WITH CustomerMonthlyActivity AS (
```

SELECT DISTINCT

EXTRACT(YEAR FROM purchase_date) AS purchase_year,

EXTRACT(MONTH FROM purchase_date) AS purchase_month,

customer_id

FROM

CustomerPurchases

```
SELECT
customer_id,
purchase_year
FROM
CustomerMonthlyActivity
GROUP BY
customer_id,
purchase_year
HAVING
```

COUNT(DISTINCT purchase_month) = 12; -- Check if they had distinct purchases in all 12 months

Explanation:

1. CustomerMonthlyActivity CTE:

- EXTRACT(YEAR FROM purchase_date) AS purchase_year: Extracts the year of the purchase.
- EXTRACT(MONTH FROM purchase_date) AS purchase_month: Extracts the month of the purchase.
- DISTINCT is crucial here. It ensures that if a customer makes multiple purchases in the same month, that month is only counted once for their activity.

2. Outer Query:

- GROUP BY customer_id, purchase_year: Groups the distinct monthly activities by each customer and the year.
- HAVING COUNT(DISTINCT purchase_month) = 12: This is the filtering condition. It checks if the count of *distinct* months a customer was active in a given year is exactly 12. If it is, they purchased in every month of that year.

Output Table (for 2024):

```
| customer_id | purchase_year |
|-----|
| 101 | 2024 |
```

Question 10: Rank Products by Sales Per Year

Concept: This question tests your ability to use window functions, specifically ranking functions (RANK(), DENSE_RANK(), ROW_NUMBER()), to assign ranks to products based on their sales within each year.

Input Table: ProductSales

```
| sale_date | product_id | sales_amount |
|-----|
| 2024-01-01 | P001 | 100
| 2024-01-05 | P002
                    | 150
| 2024-02-10 | P001
                     | 200
| 2024-03-01 | P003
                     | 120
| 2024-04-15 | P002
                     | 180
| 2024-05-20 | P0<mark>01</mark>
                    250
| 2024-06-01 | P003
                     90
| 2025-01-01 | P001 | 300
| 2025-01-10 | P002
                    | 220
| 2025-02-15 | P003
                     | 180
| 2025-03-01 | P001
                     | 350
| 2025-04-05 | P002
                     | 280
| 2025-05-10 | P003
                     | 200
```

Required SQL Query:

We need to first calculate the total sales for each product per year, then apply the ranking. Let's use RANK() here, which assigns the same rank to ties and leaves gaps in the sequence. DENSE_RANK() would assign sequential ranks without gaps for ties, and ROW_NUMBER() assigns unique ranks even for ties. Choose based on specific requirements for tie-breaking.

```
WITH YearlyProductSales AS (
 SELECT
   EXTRACT(YEAR FROM sale_date) AS sales_year,
   product_id,
   SUM(sales_amount) AS total_sales
 FROM
   ProductSales
 GROUP BY
   EXTRACT(YEAR FROM sale_date),
   product_id
)
SELECT
 sales_year,
 product_id,
 total_sales,
 RANK() OVER (PARTITION BY sales_year ORDER BY total_sales DESC) AS sales_rank
FROM
 YearlyProductSales
ORDER BY
 sales_year,
 sales_rank;
```

Explanation:

1. YearlyProductSales CTE:

 This CTE aggregates sales_amount by sales_year and product_id to get the total sales for each product in a given year.

2. Outer Query with Window Function:

- RANK() OVER (PARTITION BY sales_year ORDER BY total_sales DESC);
 - PARTITION BY sales_year: This is crucial. It tells the ranking function to restart the rank calculation for each new year.
 - ORDER BY total_sales DESC: Within each year, products are ordered by their total_sales in descending order (highest sales get rank 1).
 - RANK(): Assigns a rank. If two products have the same total sales, they get the same rank, and the next rank skips a number (e.g., 1, 2, 2, 4).

Output Table:

| sales_year | product_id | total_sales | sales_rank | |-----2024 | P001 | 550 12 2024 | P002 330 | 2024 | P003 | 210 | P001 | 650 | 2025 2025 | P002 | 500 | 2 | P003 380 2025 |3

Question 11: Employees Earning More Than Department Average

Concept: This is a common pattern for comparing individual records to an aggregate value within their group. It can be solved using subqueries or window functions. Window functions are generally preferred for performance and readability in modern SQL.

```
Input Table: Employees
```

```
| employee_id | employee_name | department | salary |
|-----|
| 101
       | Alice
                 |Sales | 60000 |
| 102
       | Bob
                 |Sales | 75000 |
| 103
       | Charlie
                  | Sales
                         | 50000 |
| 104
       | David
                 | Marketing | 80000 |
| 105
       | Eve
                | Marketing | 70000 |
| 106
       | Frank
                 | Marketing | 90000 |
| 107
       | Grace
                | HR
                         |65000 |
| 108
       | Heidi
                 | HR
                        |60000 |
Required SQL Query (Using Window Function - recommended):
SQL
WITH DepartmentAvgSalary AS (
 SELECT
   employee_id,
   employee_name,
   department,
   salary,
   AVG(salary) OVER (PARTITION BY department) AS dept_avg_salary
 FROM
   Employees
)
SELECT
 employee_id,
 employee_name,
```

```
department,
salary,
dept_avg_salary
FROM
DepartmentAvgSalary
WHERE
salary > dept_avg_salary
ORDER BY
department,
employee_id;
```

Explanation:

1. DepartmentAvgSalary CTE:

 AVG(salary) OVER (PARTITION BY department): This calculates the average salary for each department. The PARTITION BY department ensures that the average is computed separately for each department, and this average is then "attached" to every employee in that department in the result set.

2. Outer Query:

 WHERE salary > dept_avg_salary: Filters the results to include only those employees whose individual salary is greater than their department's average salary.

Output Table:

Alternative (Using Subquery):

```
SELECT
 e.employee_id,
 e.employee_name,
 e.department,
 e.salary,
 da.avg_salary AS dept_avg_salary
FROM
 Employees e
JOIN (
 SELECT
   department,
   AVG(salary) AS avg_salary
 FROM
   Employees
 GROUP BY
   department
) da ON e.department = da.department
WHERE
 e.salary > da.avg_salary
ORDER BY
 e.department,
 e.employee_id;
```

This approach first calculates the average salary per department in a subquery, then joins it back to the main Employees table to perform the comparison.

Question 12: Find Median Transaction Amount

Concept: Finding the median is tricky in SQL because it's a positional measure, not a straightforward aggregate like AVG or SUM.

- Built-in: Some SQL databases (like SQL Server 2012+, Oracle, PostgreSQL 9.4+)
 have PERCENTILE_CONT or MEDIAN aggregate functions. Spark SQL/Databricks
 also provides PERCENTILE_CONT and PERCENTILE_APPROX.
- Manual Method: For databases without built-in median, you typically use window functions (ROW_NUMBER()) to assign ranks and then select the middle value(s).

Input Table: Transactions

| transaction_id | transaction_amount |

T001	100	I
T002	250	I
T003	50	I
T004	300	I
T005	120	I
T006	180	1
T007	200	1
T008	70	
T009	160	

Sorted amounts: 50, 70, 100, 120, **160**, 180, 200, 250, 300 (Median is 160)

Required SQL Query (Using Built-in Function - Spark SQL/Databricks):

SELECT

PERCENTILE_CONT(0.5) WITHIN GROUP (ORDER BY transaction_amount) AS median_transaction_amount

FROM

Transactions;

- -- OR (for approximate median, which can be faster for large datasets)
- -- SELECT PERCENTILE_APPROX(transaction_amount, 0.5) AS approximate_median_transaction_amount
- -- FROM Transactions;

Explanation (Built-in):

- PERCENTILE_CONT(0.5) WITHIN GROUP (ORDER BY transaction_amount): This
 function calculates the 50th percentile (which is the median) of
 transaction_amount.
 - PERCENTILE_CONT (continuous) interpolates values if the median falls between two numbers (e.g., for an even number of rows).
 - WITHIN GROUP (ORDER BY transaction_amount) specifies the column to order by for the percentile calculation.

Output Table (Built-in):

```
| median_transaction_amount |
|-----|
| 160.0 |
```

WITH RankedTransactions AS (

Required SQL Query (Manual Method - more portable across SQL dialects):

This method handles both odd and even numbers of rows.

SELECT

transaction_amount,

ROW_NUMBER() OVER (ORDER BY transaction_amount) AS rn,

COUNT(*) OVER () AS total_rows

```
FROM
```

Transactions

)

SELECT

AVG(transaction_amount) AS median_transaction_amount

FROM

RankedTransactions

WHERE

```
rn = CEIL(total_rows / 2.0) OR rn = FLOOR(total_rows / 2.0) + 1;
```

- -- For odd total_rows, both conditions point to the same middle row.
- -- For even total_rows, it selects the two middle rows to average.

Explanation (Manual Method):

1. RankedTransactions CTE:

- o ROW_NUMBER() OVER (ORDER BY transaction_amount) AS rn: Assigns a sequential rank to each transaction amount after ordering them.
- COUNT(*) OVER () AS total_rows: This is a window function that counts the total number of rows in the entire result set (the entire table in this case), and this count is added to every row.

2. Outer Query:

- WHERE rn = CEIL(total_rows / 2.0) OR rn = FLOOR(total_rows / 2.0) + 1: This is the core logic to select the middle row(s).
 - If total_rows is odd (e.g., 9): CEIL(9/2.0) = 5, FLOOR(9/2.0)+1 = 4+1=5.
 It selects the 5th row.
 - If total_rows is even (e.g., 8): CEIL(8/2.0) = 4, FLOOR(8/2.0)+1 = 4+1=5.
 It selects the 4th and 5th rows.
- AVG(transaction_amount): Averages the selected transaction amount(s). If only one row is selected (odd count), it's just that value. If two rows are selected (even count), it averages them.

Output Table (Manual Method):

```
| median_transaction_amount |
|-----|
| 160.0 |
```

Question 13: Users Who Placed Their First Order in the Last 30 Days

Concept: This question involves identifying a user's *first* activity and then filtering based on a time window. It combines aggregation (MIN), date functions, and filtering.

Input Table: Orders

Given today's date is 2025-06-29 (as per your prompt's current time). Last 30 days would be from 2025-05-30 to 2025-06-29.

Required SQL Query:

```
SELECT

user_id,

MIN(order_date) AS first_order_date

FROM

Orders

GROUP BY

user_id

HAVING

MIN(order_date) >= DATE('2025-06-29') - INTERVAL '30' DAY AND MIN(order_date) <= DATE('2025-06-29');

-- In Databricks/Spark SQL, you might use:

-- MIN(order_date) >= date_sub(current_date(), 30) AND MIN(order_date) <= current_date();

-- Or for literal date:

-- MIN(order_date) >= date_sub('2025-06-29', 30) AND MIN(order_date) <= '2025-06-29';
```

Explanation:

- SELECT user_id, MIN(order_date) AS first_order_date FROM Orders GROUP BY user_id: This part aggregates the data to find the earliest order_date for each user_id. This effectively identifies their first order.
- HAVING MIN(order_date) >= DATE('2025-06-29') INTERVAL '30' DAY AND
 MIN(order_date) <= DATE('2025-06-29'): This HAVING clause filters the grouped
 results. It keeps only those users whose first_order_date falls within the last 30 days
 from the current date (2025-06-29).
 - DATE('YYYY-MM-DD') INTERVAL 'X' DAY is common syntax (e.g., in MySQL, PostgreSQL).
 - For Databricks/Spark SQL, date_sub(current_date(), 30) is the typical way to get a date 30 days ago. current_date() would be used if the query runs on the current day, otherwise, a specific date literal like '2025-06-29' is used.

Output Table:

```
| user_id | first_order_date |
|------|
| 103 | 2025-06-01 |
| 104 | 2025-06-05 |
| 105 | 2025-06-15 |
| 106 | 2025-06-25 |
| 107 | 2025-06-28 |
```

Question 14: Compare Price Change Between Two Dates Per Product

Concept: This requires comparing a product's price at a start date versus an end date. This can be done by joining the table to itself or, more elegantly, using window functions like LAG or by filtering and joining. Let's assume price is recorded daily for simplicity.

Input Table: ProductPrices

```
| product_id | price_date | price |
|-----|----|
| P001 | 2025-01-01 | 10.00 |
| P002
       | 2025-01-01 | 20.00 |
| P001
       | 2025-01-15 | 10.50 |
| P002 | 2025-01-15 | 21.00 |
| P001
       | 2025-01-31 | 11.00 |
| P002
       | 2025-01-31 | 20.50 |
| P003
        | 2025-01-01 | 5.00 |
| P003
        | 2025-01-31 | 5.50 |
```

Required SQL Query (Let's compare between '2025-01-01' and '2025-01-31'):

```
WITH StartPrices AS (
 SELECT
   product_id,
   price AS start_price
 FROM
   ProductPrices
 WHERE
   price_date = '2025-01-01'
),
EndPrices AS (
 SELECT
   product_id,
   price AS end_price
 FROM
   ProductPrices
 WHERE
   price_date = '2025-01-31'
SELECT
 s.product_id,
 s.start_price,
 e.end_price,
 e.end_price - s.start_price AS price_change,
 ROUND((e.end_price - s.start_price) / s.start_price * 100, 2) AS percentage_change
FROM
```

StartPrices s

JOIN

EndPrices e ON s.product_id = e.product_id

ORDER BY

s.product_id;

Explanation:

- 1. StartPrices CTE: Filters ProductPrices to get prices on the specified start_date.
- 2. **EndPrices CTE:** Filters ProductPrices to get prices on the specified end_date.
- 3. Outer Query:
 - JOIN StartPrices s JOIN EndPrices e ON s.product_id = e.product_id: Joins the two CTEs on product_id to bring the start and end prices onto the same row.
 - e.end_price s.start_price AS price_change: Calculates the absolute change.
 - ROUND((e.end_price s.start_price) / s.start_price * 100, 2) AS
 percentage_change: Calculates the percentage change. ROUND is used for
 cleaner output.

Output Table:

| product_id | start_price | end_price | price_change | percentage_change |

Question 15: Customers Whose First and Last Transaction on Same Day

Concept: This requires finding the minimum and maximum transaction dates for each customer and then comparing them.

Input Table: CustomerTransactions

```
| transaction_id | customer_id | transaction_date | amount |
|-----|-----|
| T001
        | 101
               |2025-04-01 |100 |
|T002
        | 102 | 2025-04-05
                            |150 |
| T003
        | 101
              | 2025-04-01 | 50 | -- Same day as T001
|T004
        | 103
               |2025-04-10 |200 |
               | 2025-04-05 | 75 | -- Same day as T002
| T005
       | 102
                            | 300 |
|T006
        | 104
                | 2025-04-12
| T007
        | 103
              | 2025-04-11
                            | 90 |
```

|180 |

Required SQL Query:

| 105

| T008

```
SELECT

customer_id,

MIN(transaction_date) AS first_transaction_date,

MAX(transaction_date) AS last_transaction_date

FROM

CustomerTransactions

GROUP BY

customer_id

HAVING

MIN(transaction_date) = MAX(transaction_date)

ORDER BY

customer_id;
```

| 2025-04-15

Explanation:

- SELECT customer_id, MIN(transaction_date) AS first_transaction_date, MAX(transaction_date) AS last_transaction_date FROM CustomerTransactions GROUP BY customer_id: This aggregates the data to find the earliest (MIN) and latest (MAX) transaction_date for each customer_id.
- 2. HAVING MIN(transaction_date) = MAX(transaction_date): This HAVING clause filters the grouped results, keeping only those customers where their first and last transaction dates are identical.

Output Table:

Question 16: Percentage of Returning Users Each Month

Concept: This is a classic cohort analysis or user retention problem. It requires identifying first-time users, then identifying returning users in subsequent months, and finally calculating the percentage.

Input Table: UserLogins (or UserPurchases, assuming any activity counts as "usage")

```
| 101 | 2025-02-01 | -- Returning in Feb

| 103 | 2025-02-05 | -- Returning in Feb

| 105 | 2025-02-10 | -- First activity (Feb 2025)

| 102 | 2025-03-01 | -- Returning in Mar

| 101 | 2025-03-05 | -- Returning in Mar

| 104 | 2025-03-10 | -- Returning in Mar
```

Required SQL Query:

This is a multi-step process:

- 1. Find each user's first activity month.
- 2. Find distinct active users per month.
- 3. For each month, count users who are *not* first-time users in that month (returning users).
- 4. Calculate the percentage.

```
WITH UserFirstActivity AS (

SELECT

user_id,

DATE_TRUNC('month', MIN(activity_date)) AS first_activity_month

FROM

UserLogins

GROUP BY

user_id

),

MonthlyActiveUsers AS (

SELECT DISTINCT

DATE_TRUNC('month', ul.activity_date) AS activity_month,
```

```
ul.user_id,
   ufa.first_activity_month
 FROM
   UserLogins ul
 JOIN
   UserFirstActivity ufa ON ul.user_id = ufa.user_id
),
MonthlyUserCounts AS (
 SELECT
   activity_month,
   COUNT(user_id) AS total_active_users,
   COUNT(CASE WHEN activity_month > first_activity_month THEN user_id ELSE NULL
END) AS returning_users,
   COUNT(CASE WHEN activity_month = first_activity_month THEN user_id ELSE NULL
END) AS new_users
 FROM
   MonthlyActiveUsers
 GROUP BY
   activity_month
)
SELECT
 activity_month,
 total_active_users,
 new_users,
 returning_users,
 CASE
```

WHEN total_active_users > 0 THEN ROUND((CAST(returning_users AS DECIMAL) / total_active_users) * 100, 2)

ELSE 0

END AS percentage_returning_users

FROM

MonthlyUserCounts

ORDER BY

activity_month;

Explanation:

- 1. **UserFirstActivity CTE:** Determines the first_activity_month for each user.
- 2. **MonthlyActiveUsers CTE:** Gets all distinct user_id and their activity_month (current month of activity) along with their first_activity_month.

3. MonthlyUserCounts CTE:

- COUNT(user_id) AS total_active_users: Counts all unique users active in a given activity_month.
- COUNT(CASE WHEN activity_month > first_activity_month THEN user_id ELSE NULL END) AS returning_users: This is the key for returning users. A user is "returning" in a given activity_month if their activity_month is later than their first_activity_month. COUNT(CASE WHEN ... THEN ... ELSE NULL END) effectively counts only those where the condition is true.
- COUNT(CASE WHEN activity_month = first_activity_month THEN user_id
 ELSE NULL END) AS new_users: Counts users whose activity_month is their first_activity_month (i.e., new users for that month).

4. Final SELECT:

- Calculates the percentage_returning_users by dividing returning_users by total_active_users and multiplying by 100.
- o CAST(... AS DECIMAL) ensures floating-point division.
- CASE WHEN total_active_users > 0 THEN ... ELSE 0 END handles division by zero.

Output Table:

| activity_month | total_active_users | new_users | returning_users | percentage_returning_users |

		-	-			
2024-12-01	2	2	0	0.00	1	
2025-01-01	4	2	2	50.00	1	
2025-02-01	3	1	2	66.67	1	
2025-03-01	3	0	3	100.00	I	

These are excellent, practical SQL problems often encountered in data analysis and engineering