



# Hash Map: Classical Problems

Course on Basic Data Structures (C++)

Q. Given 2 arrays  $x$  &  $y$ , find no. of pairs of distinct integers  $(i, j)$  s.t.  $x[i] - x[j] = y[i] - y[j]$ .

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eg.  $x =$ 

0	1	2
3	1	5

 $y =$ 

9	2	11
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$(0, 2)$   
 $(2, 0)$

Ans = 2

Sol:-

$$x[i] - x[j] = y[j] - y[i]$$

$$\Rightarrow x[i] - y[i] = x[j] - y[j]$$

um  $f$ ;

for ( $i=0$ ;  $i < n$ ;  $++i$ )

$f[x[i] - y[i]]++$ ;

ans = 0;

for (auto  $f$  :  $f$ )

{ if ( $n \geq f.second$ ;

ans +=  $n * (n-1)$ ;

}

2 distinct  
& order matters

$$n * (n-1)$$

$$5 \rightarrow x_i - y_i = 5$$

for  $i = 1, 2, 5, 8, 10, 11$

$$n = 6$$

$$\underline{n * (n-1)}$$



We have a list of strings. Find the number of pairs  $(i, j)$  s.t.  $0 \leq i < j \leq n-1$

$\Delta$   $w[i] + w[j]$  can be shuffled to form a palindrome.

$w = [ab, bcde, cba]$   $(1, 2) \rightarrow bcde + cba \Rightarrow$   
 $d b c e d e c b a$

$(0, 1) \rightarrow ab + bcde \Rightarrow$  No

$(0, 2) \rightarrow ab + cba \Rightarrow$  No

Ans = 1

- 1) Order doesn't matter
- 2) Parities of frequencies of diff. characters matter
- 3) In  $w[u] + w[v]$ , number of characters with odd freq. should be less than 2.

$$\left\{ \begin{array}{cccccccc} 'a', & 'b', & 'c', & 'd', & - & - & - & - & 'z' \\ 0, & 1, & 1, & 0, & 1 & - & - & - & - \end{array} \right\}$$

Example  
of how  
to hash

$$0 \times 2^0 + 1 \times 2^1 + 1 \times 2^2 + 0 \times 2^3 + 1 \times 2^4 - \dots$$

$$\Rightarrow \text{hash}(s_1 + s_2) = \text{hash}(s_1) \oplus \text{hash}(s_2)$$

$s_1$ ,  $s_2$

$\text{hash}(s_1)$

$\text{hash}(s_2)$

$b_1$	$b_2$	$b_1 \oplus b_2$
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0	0	0
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0	1	1
---	---	---

1	0	1
---	---	---

1	1	0
---	---	---



0, 2, 4, 8, 16 — — — — —  $2^{25}$

↑  
possible values of xor

⇒ Min-time vels remove

2) Value of a key

2) Update time of a particular.  $\rightarrow$  most recent time

2) Remove a particular key

key  $\rightarrow$  time

key  $\rightarrow$  value

time  $\rightarrow$  key

Sol. 1:-

Get  $\rightarrow O(1)$

Put  $\rightarrow O(N)$

Sol. 2  $\rightarrow$

Get  $\rightarrow O(\log N)$

Put  $\rightarrow O(\log N)$

Required:

Get  $\rightarrow O(1)$

Put  $\rightarrow O(1)$



