



# Arrays: Classical Problems

Course on Basic Data Structures (C++)

Given an array of integers. (Size is  $N$ )

Print  $N$  numbers where  $ans[i] = \text{product}(a[j])$   
 $j \neq i$

Input  $\rightarrow$  5 1 2 3

Output  $\rightarrow$  3 5 5 5

$$\text{pre-max}[i] = \max(a[0], a[1] \dots a[i])$$

$$\text{suf-max}[i] = \max(a[i], a[i+1] \dots a[n-1]);$$

$O(N)$

for ( $i = 0; i < n; ++i$ )

pre-max[i] = ( $i == 0$ ) ? a[i] : max(pre-max[i-1], a[i])

[i-1]

i

$$\text{pre-sum}[i] = \text{pre-sum}[i-1] + a[i]$$

$$\text{pre\_min}(i) = \min(\text{pre\_min}[i-1], a[i]),$$

$$\text{pre\_fun}(i) = \text{fun}(\text{pre\_fun}[i-1], a[i])$$

gcd

lcm

product

mod

or

xor

$$\gcd(a, b, c) \Rightarrow \underline{\gcd(a, b)}, c$$

$$f(a[0], a[1] \dots a[i]) = f(f(a[0] \dots a[i-1]), a[i])$$



Given an array with integers.

Queries:  $m$

$l \quad r \quad x$

$\hookrightarrow a[l] + z \cdot x, a[l+1] + z \cdot x \dots a[r] + z \cdot x$

We want to retrieve the final array.

1 2 1 5 6 -3

1 2 3 1 5 1

u v x

→ 5 5 1

→ 1 6 -2

→ 3 4 2

⇒ 1 2 1 5 7 -3

2) -1 0 -1 3 5 -5

2) -1 0 1 5 5 -5

m<sub>2</sub> 3

Time -  $O(N \times M)$

Required -  $O(N + M)$



2D - vector (1-based indexing)

↳ queries.

$i_1, j_1$

↓

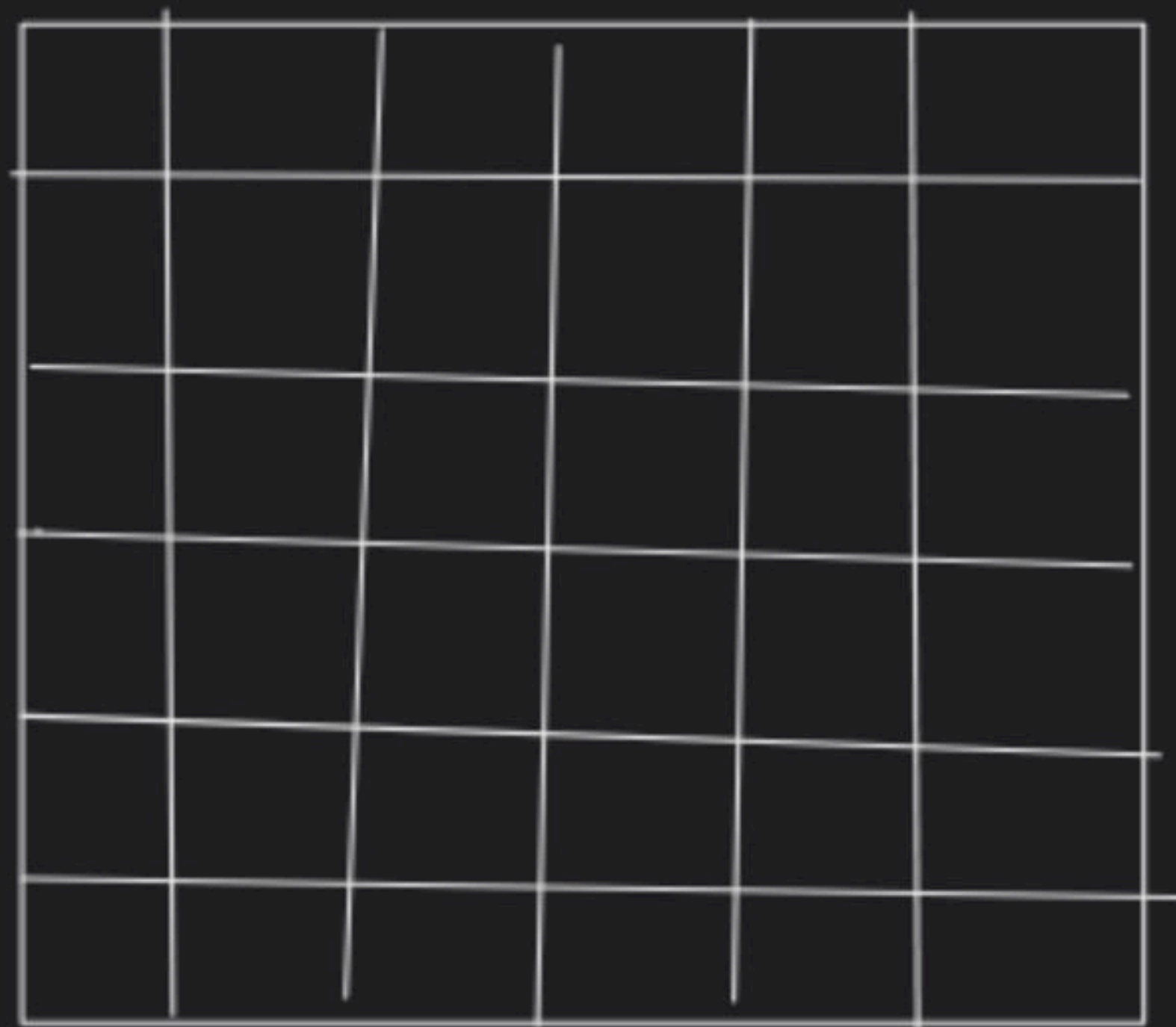
bottom-left

$i_2, j_2$

↓

bottom-right

$O(i_2 - i_1)$  or  $O(j_2 - j_1) \Rightarrow O(1)$





	0	1	2	3
b[2]	0	3	1	4

$$b[1] + 1 \cdot b[0] \Rightarrow 3$$

$$b[2] + 2 \cdot b[1] \Rightarrow 4$$

$$b[3] + 2 \cdot b[2] \Rightarrow 8$$

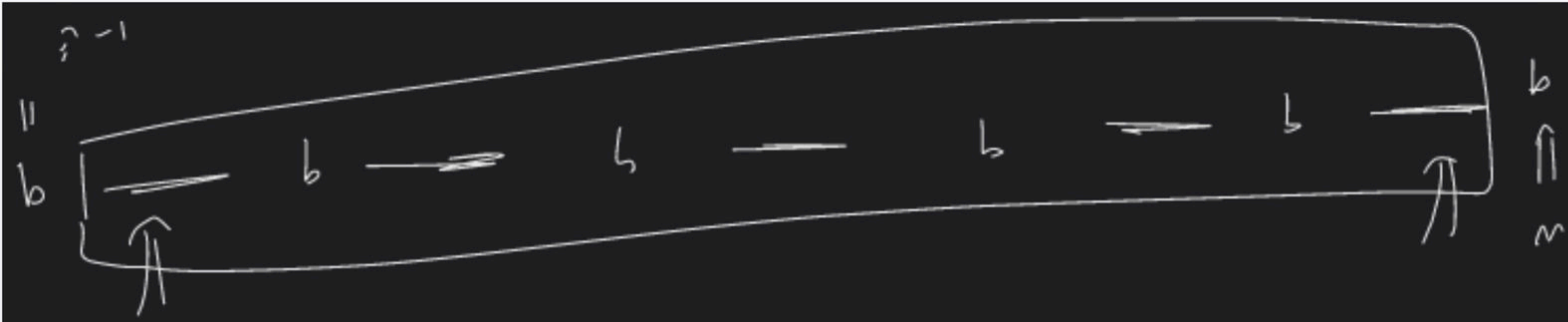
$$\begin{array}{ccccccc}
 \underline{w w w w} & b & \underline{w w} & b & \underline{w w w} & b & \underline{w w w w} \\
 \frac{4 \times (4+1)}{2} & - & \frac{2 \times (2+1)}{2} & + & \frac{3 \times (3+1)}{2} & + & \frac{4 \times (4+1)}{2}
 \end{array}$$

Vi v

b w w w w b w w b b w w w b

0, 5, 8, 9, 13  
└─┘ └─┘ └─┘ └─┘  
↓ ↓ ↓ ↓  
4 2 0 3





Given an array, divide the array  
into maximum no. of contiguous parts,  
such that if we sort the  
individual parts, the array also  
gets sorted.

3   2   1   5   4   6

↘ ↘ 3

1, 2, 3, 4, 5, 6

1 2 3 4 5 6 1



C = 2

3 1 2

$\Rightarrow$  ans = 2



ans = 3

5

5

5

5

5

Ans = 5



3, 2, 1, | 4, | 5, | 6

3 2 5 1 0 6  $\Rightarrow$  3

1 2 3 4 5 6

$$\text{pre-max}[i] \leq \text{sf-min}[i+1]$$

1.) Max parts

2.)  $i_1 j_1$  to  $i_2 j_2 \leftarrow \text{sub-matrix}$   
 $O(1)$

3.) Check if it is possible to  
rearrange the array s.t.  
sum of no 2 consec. elements is  
divisible by 3.

# 4) Maximum Subarray Sum

3 2 1  
↘ ↗

1 2 3  
2 1 3

3 4 5  
||  
4 5 6  
5 6 7

$m_2$  5

7 0 1