

# Hash Map: Intuition and Implementation

Course on Basic Data Structures (C++)

## Hashing

Mapping elements to an integer. in a way.

that if  $e_1 \neq e_2$ , then  $\text{hash}(e_1) \neq \text{hash}(e_2)$

---

## strings

```
int hash(string s)
{
    return 5;
}
```

An example  
of a bad / useless  
hash function.

We have a string. You need to store the frequency of each character that is there in the string. ( $s[i] \rightarrow 'a' \text{ to } 'z'$ )

```
int cnt[26] = {0};
```

```
for (char ch : s)  
    cnt[ch - 'a']++;
```

```
int hash(char ch)  
{  
    return ch - 'a';  
}
```



Frequency of numbers in an array.

$a[i] \in [-100, 100]$

```
{ int freq[201] = {0};  
  for (int i : a)  
    freq[i + 100]++;  
}
```

```
int hash(int x)  
{  
  return 100 + x;  
}
```

An array of strings. ( $\text{len}(s[i]) \leq 3$ ).

Store the freq. of each string

a h a

b a c

z x y

a b

int freq[27][27][27]

↑        ↑        ↑

Look at these 2 scenarios:

$$1) \quad a[i] \in [-10^9, 10^9]$$

$$2) \quad 1 \leq \text{len}(s(i)) \leq 10^5$$

1.) Have an array of size  $n$ ;

$$2.) \text{hash}(i) = i \% n \quad [0, n-1]$$

# Hash Maps

stores { key, value } pairs.

<int, int>

<string, int>

<string, string>

<double, string>

<int, vector<int>>

<string, vector<string>>

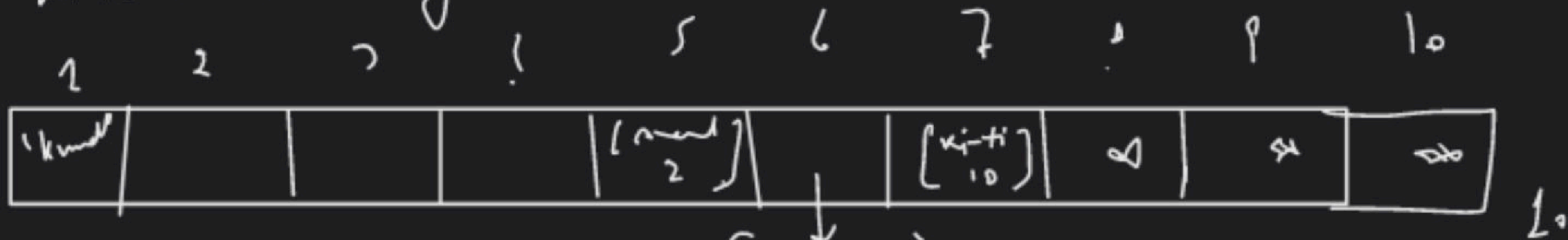


## hash\_function

- 1) Maps the key to some integer  $(i)$   
[0, size-1]  
↑
- 2) The {key, value} pair is stored  
at  $i$ .

# Taking Care of Collisions :

## 1) Linear Probing



hash(s)  $\rightarrow$   $(\text{length}(s) - 1) \% 10$

$\left[ \begin{array}{c} \text{"akash"} \\ 2 \end{array} \right]$   $\downarrow$   $m[\text{"akash"}] = 1$

$m[\text{"Anand"}] = 2$

$m[\text{"Kirti"}] = 10$

$m[\text{"kundi"}]$

## 2) Chaining / Bucketing



vector<list> v();

vector

LL



{  
2  
}

{  
1  
}

↓  
{  
10  
}

↓  
{  
33  
}

1) Update/Insert  $\leftarrow O(\text{len})$

2) Delete  $\leftarrow O(\text{len})$

3) Search  $\leftarrow O(\text{len})$

$\Rightarrow$  Any Insert/Delete/Search  $\Rightarrow O(1)$



1) Insert / Update

if we insert the key to its  
correct position if list was sorted.

$O(N)$  — )

2) Delete

BS  $\Rightarrow$  search  $\Rightarrow$  To delete  $O(N)$

## Terminologies

size  $\rightarrow$  no. of keys in hm.

bucket\_count  $\leftarrow$  no. of bucket in

load\_factor  $\leftarrow$  size / bucket\_count  $\rightarrow$  (Default 1.0)

max\_load\_factor  $\leftarrow$  bucket\_count is increased to

the smallest prime  $\geq 2 \times$  cur bucket\_count if  $lf > mx$  if

hash(string s)

```
{ int ans = 0;
  for (char ch : s)
    ans = (int) ch;
  return ans % buck-size;
}
```

Base  $\leftarrow p$  (prime  $> 256$ )

Polynomial Rolling  
Hash Function.

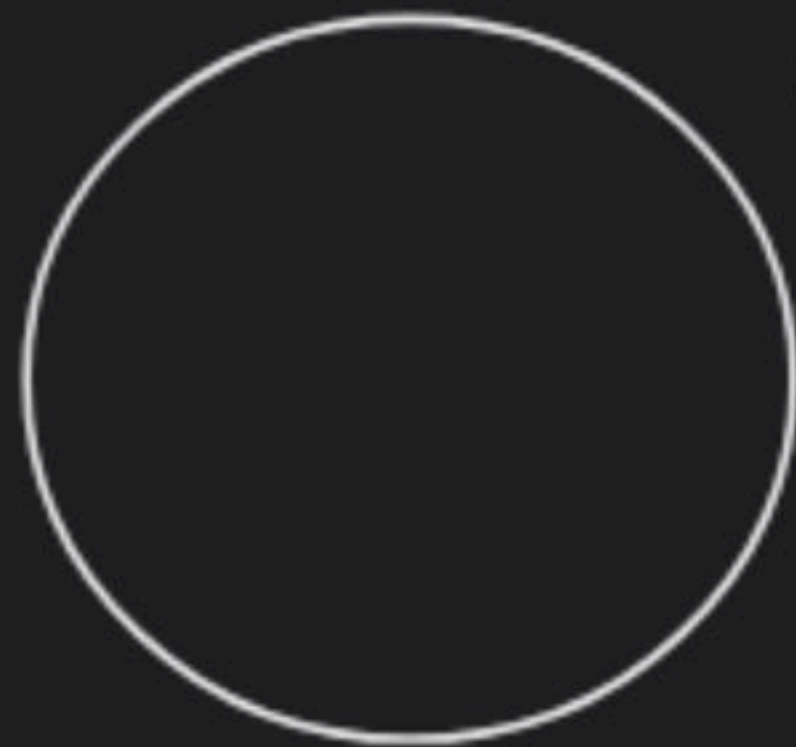
$$\text{hash}(s) = \left( s[0] * p^0 + s[1] * p^1 + s[2] * p^2 + \dots + s[n-1] * p^{n-1} \right)$$

.  $\leq$  buck-size.



unordered set

Insert  
a value



Check if  
a value is  
present or not



Delete a  
value



No. of subarrays with sum = k.


$$l \text{ to } r \Rightarrow \text{pref}[r] - \text{pref}[l-1] = k$$

For  $r$ , calculate no. of indices  $i$   
s.t.  $i < r$  and  $\text{pref}[i] = \text{pref}[r] - k$ .

↳ Sum of these = Ans  $[i+1, r]$

Find no. of sub Arrays  $S$  s.t.

$$\text{sum}(S) = \text{length}(S)$$

  $[1 \ 2 \ 0 \ 5]$

Ans: 3