

## Tutorial 5

Q.1) What's difference between DFS & BFS. write applications of both algorithm.

A.1.)

DFS

BFS

- |   |   |
|---|---|
| <ul style="list-style-type: none"> <li>- Use stack to find shortest path</li> <li>- Stands for Depth first Search</li> <li>- Suitable for searching vertices closer to given source</li> <li>- Recursive algorithm - like backtracking</li> <li>- Requires less memory</li> <li>- Considers <sup>not</sup> all neighbours first i.e. visit then explore all path through this decision</li> </ul> <p>eg:- acyclic graph, topological order.</p> | <ul style="list-style-type: none"> <li>- Use Queue to find shortest path.</li> <li>- stands for Breadth first Search</li> <li>- suitable where solutions are closer to source</li> <li>- NO concept of backtracking</li> <li>- Requires more Memory.</li> <li>- Considers all neighbours first i.e. not suitable for decision making.</li> </ul> <p>eg:- Bipartite graph, shortest st path etc.</p> |
|---|---|

Q.2.) Which Data structure is used to implement DFS & BFS & why?

A.2.) DFS Algorithm traverses a graph in depthwise motion & uses stack to remember to get next vertex to search, when a dead end occurs in any iteration. So basically in order to reach to the deepest node first stack - used.

- BFS uses queue to avoid recursion. Queue keeps track of locations remaining to be searched as opposed to needing a potentially large amount of processor stack space for recursive solution.

Q.3) what do you mean by sparse & dense graphs? which representation of graph is better for both.

A.3))

Dense Graph

Sparse Graph

- A graph in which number of edges is close to maximal no. of edges.
- Graphs are densely connected.
- $G = (V, E)$  in which  $|E| = O(|V|^2)$
- If a graph is dense we should store it as adjacency matrix.
- A graph in which no. of edges is minimal to no. of edges.
- Graphs are sparsely connected (e.g. Trees).
- $G = (V, E)$  in which  $|E| = O(|V|)$ .
- If graph - sparse - store it as a list of edges.

Q.4.) How to detect cyclic graph using BFS & DFS?

A.4.) using BFS:-

- compute in-degree (no. of incoming edges) for each vertex present in graph & initialize count of visited nodes at 0.
  - Pick all vertices with in-degree as 0's, add to queue. (Enqueue operation).
  - Remove vertex from queue.
- Increment count of visited nodes by 1

- Decrease in-degree by 1 for all neighbouring nodes.
  - If in-degree of a neighbouring node is reduced to zero,  
add to queue.
- 4.) Repeat step 3 until queue is empty.
- Q.)

Using DFS:-

- 1.) Create graph using given number of edges & vertices.
- 2.) Create recursive (f) that have current index or vertices,  
visited array & parent node.
- 3.) Mark unvisited node as visited.
- 4.) find unvisited vertices that are adjacent to current node.
- 5.) If adjacent node not parent - already visited - return true

Q.5) What do you mean by disjoint set data structure?

Explain 3 operations along with examples, which  
can be performed on disjoint sets.

A.5.) Disjoint-set data structure:- also called union find  
data structure or merge-find set is data structure  
that stores a collection of disjoint (non-overlapping) sets.  
These operations performed by Disjoint-set data struc-  
ture :-

- 1.) Making new set containing a new element.
  - 2.) Finding representative of set containing given element.
  - 3.) Merging two sets.
- 1.) Making New Set:- Makeset operation adds a new element  
into a new set containing only the new element, &

new set is added to data structure.

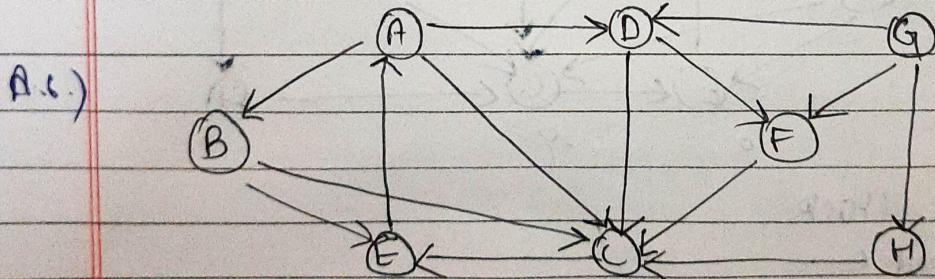
### ii) Finding Set Representatives

Find operation follows chain of parent pointers from a specified query node  $x$  until it reaches a root element. Root element represents set to which  $x$  belongs and may be  $x$  itself.

### iii) Merging Two Sets :-

Operation union( $x, y$ ) replaces set containing  $x$  & set containing  $y$  with their union.

Q.5) Run DFS & BFS on graph shown on R.H.S.

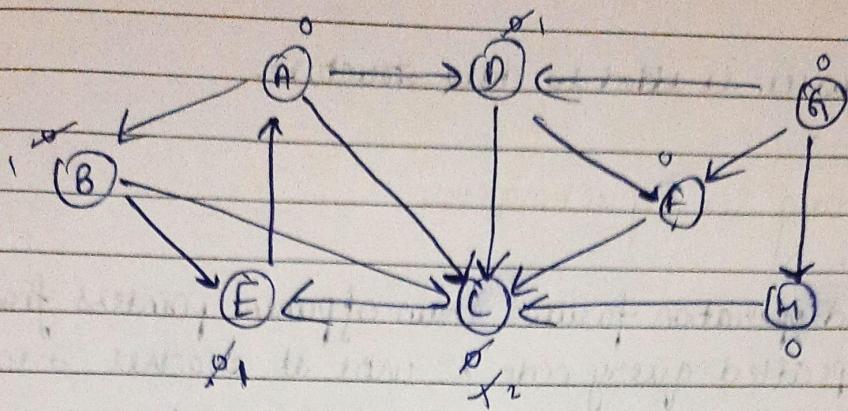


B.5) BFS :-  $\text{A} \rightarrow \text{parent node}$

$B \leftarrow C \rightarrow D \rightarrow \text{child node}$

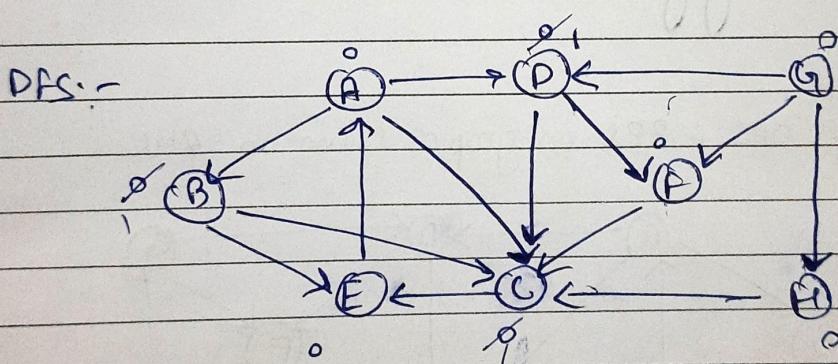
- taking any one child node, say 'B'

$E \leftarrow C \rightarrow (A, B) \rightarrow \text{furthermore} \dots$



Parent	A	A	A	B	B	D	D	F
Node	B	C	D	F	C	C	F	C

path :- A → B → C.



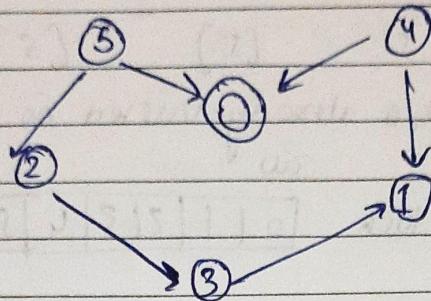
nodes STACK  
processed ←

A	A
BCE	ABED
B	BE
E	BEC
C	EC

path A → B → C

Q.7) Apply topological sorting & DFS on graph having vertices from 0 to 5

(Q.7)



Adj cent list (G)

$0 \rightarrow 1$

$1 \rightarrow$

$2 \rightarrow 3$

$3 \rightarrow 1$

$4 \rightarrow 0, 1$

$5 \rightarrow 0, 2$

0 1 2 3 4 5

visited - - - - -

stack (empty)

step 1: Topological sort(0), visited[0] = true



list-empty (no recursion call).

stack [0]

Step 2: Topological sort[2], visited[2] = true  
 " " [3] " [3] = true

- 1 - already visited, no more recursion

CALL

Stack  $\boxed{0 \mid 1 \mid 3 \mid 2}$

Step 5:-       $+^i [5]$     "  $[5] = \text{true}$

$2, 0$  already visited no recursion  
can

Stack  $\boxed{0 \mid 1 \mid 2 \mid 3 \mid 4 \mid 5}$

Print all elements of stack from top  $\rightarrow$  bottom.

Q.8) Can heap be used to implement priority queue?

Name few algorithm where you need to use priority queue?

A.8) Heaps can be used to implement priority queue. It takes  $O(\log N)$  time to insert & delete each element in priority queue.

Algorithm where priority queue can be applied are Dijkstra, Prim's algorithm

Q.9) Difference between Max & Min heap?

It's used to maximum element in heap

It's used to access minimum element in heap