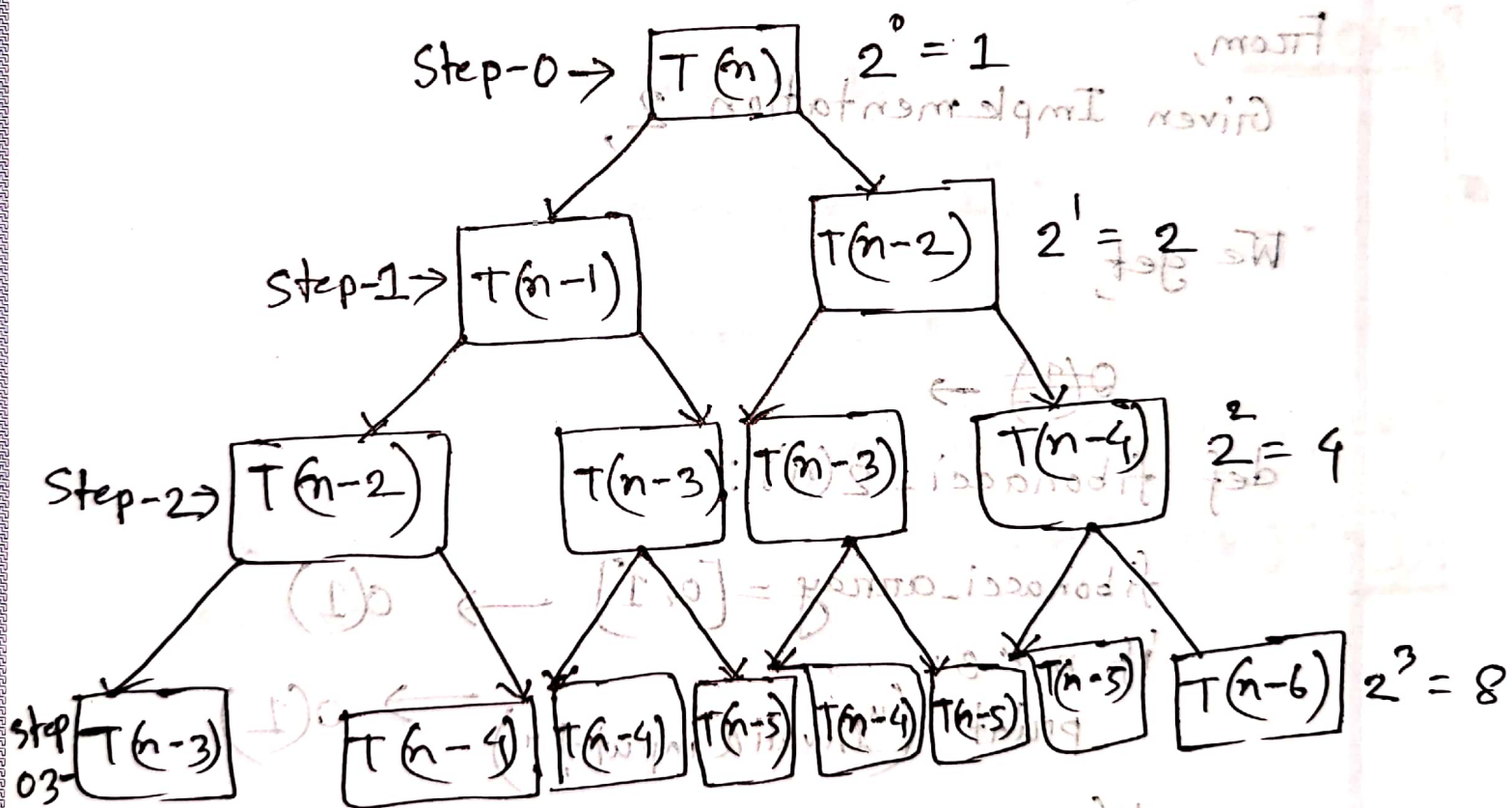


Ans. to the ques. no. 2

ID: 19101465

From

Given Implementation 1,



Step-k  $\rightarrow 2^k$

Suppose,  $k = n$ , Then  $T(n-k) = T(n-n) = T(0)$

As  $T(0)$  is a base case,  $T(0)$  gives 1 operation

$\therefore$  Total operation  $= 2^0 + 2^1 + 2^2 + 2^3 + \dots + 2^n = 2^{n+1} - 1$

So, Time complexity for Implementation 1:

$$O(2^{n+1} - 1) = O(2^{n+1}) = O(2^n)$$

From,

Given Implementation 2,

We get,

~~def~~ fibonacci\_2(n):  
 fibonacci\_array = [0, 1] → O(1)  
 if n < 0:  
 print("Invalid input!") → O(1)

elif n <= 2:

return fibonacci\_array[n-1] → O(1)

~~else:~~

else:

for i in range(2, n): → O(n+2)

fibonacci\_array.append(fibonacci\_array

[i-1] + fibonacci\_array[i-2])

return fibonacci\_array[-1]

(\*) `n = int(input("Enter a number: "))`

`nth_fib = fibonacci-2(n)`

`print("The %d-th fibonacci number  
is %d" % (n, nth_fib))`

$$\therefore \text{Time complexity} = O(1) + O(n-2)$$

$$= O(1) + O(n-2)$$

$$= O(n-2)$$

$$= O(n)$$



Ans. to the ques. no. - 4

```
def Multiply_matrix (A, B):
```

```
    global n
```

```
    C = [[0] * n for i in range (n)]
```

```
    for i in range (0, n, 1):
```

```
        for j in range (0, n, 1):
```

```
            for k in range (0, n, 1):
```

```
                C[i][j] += int(A[i][k]) *  
                             int(B[k][j])
```

```
    return C
```

$O(n)^3$

$\therefore$  Time complexity =  $O(n)^3$

=  $O(n)^3$