

Task 03

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prob-1 (only Dijkstra part)

import queue

import math

def Dijkstra (Graph, source):

dist

dist = [null, 0]

visited = []

for i in range (0, len (Graph) + 1):
visited.append (False) } $O(N)$

Q = queue.PriorityQueue() $\rightarrow O(1)$

for i in range (2, len (Graph) + 1):
dist.append (float ('inf')) } $O(N)$

prev = []

for i in range (0, len (Graph) + 1):
prev.append (float ('inf')) (null') } $O(N)$

j = 1

for i in Graph
Q.put ((dist [i], i)) } $O(N \log N)$

j = j + 1

while not Q.empty():
u = Q.get() [2] $\rightarrow O(\log N)$
if visited [int (u)]:
continue

visited[int(u)] = True

for x in Graph[u]: $\rightarrow O(M)$

v = x[0]

- alt = dist[int(u)] + int(u[1])

If alt < dist[int(v)]:

dist[int(v)] = alt

prev[int(v)] = u

Q.put((dist[int(v)], v))

return dist

$\rightarrow O(\log N)$

Time complexity = $O(N) + O(N \log N) + O(N \log N)$
 $+ O(N \log N) + O(N) + O(M \log N)$

$= O(N) + O(N \log N) + O(M \log N)$

$= O(N \log N) + O(M \log N)$

$= O(N \log N + M \log N)$

Problem 2 (Only Dijkstra part)

'import queue

'import math

def Dijkstra(Graph, source):

dist = ['null', 0]

visited = []

for i in range(0, len(Graph)+1):

visited.append(False)

Q = queue.PriorityQueue()

for i in range(2, len(Graph)+1):

dist.append(float('inf'))

prev = []

for i in range(0, len(Graph)+1):

prev.append('null')

j = 1

$O(N \log N)$

for i in Graph: $\rightarrow O(N)$

Q.put((dist[j], i)) $O(N \log N)$

j = j + 1

while not Q.empty(): $\rightarrow O(N)$

u = Q.get()[1] $\rightarrow O(\log N)$

if visited[int(u)]:

continue

visited[int(u)] = true

for x in Graph[u]: $\rightarrow O(M)$

v = x[0]

alt = dist[int(u)] + int(x[1])

if alt < dist[int(v)]:

dist[int(v)] = alt

prev[int(v)] = u

Q.put((dist[int(v)], v))
 $\rightarrow O(\log N)$

m = []

j = len(Graph)

m.append(str(j))

while j != 1:

y = prev[j]

m.append(y)

j = int(y)

return m

$O(N \log N) + O(N) + O(M \log N)$

$$\begin{aligned}
 \text{Time complexity} &= O(N) + O(N \log N) + O(N) + \\
 &\quad O(M \log N) \\
 &= O(N \log N) + O(M \log N) \\
 &= O(N \log N + M \log N)
 \end{aligned}$$

If the number of titans in each road is exactly 1, there is an $O(N+M)$ algorithm to solve this problem. That is BFS (Breadth First Search).

In the first line of input there will be the number of vertex and then the number in the second line will be the same weight for each edge.

Then the rest number of lines will show the connections of the graph. So, the input will be for the graph in problem 1 \Rightarrow

5				
1				
1	2	4		
2	1	3	5	
3	2	4	5	
4	1	3		
5	2	3		