# Operating System 2 Programming Assignment 6

## Tanmay Garg CS20BTECH11063

- There are 2 files in the folder
  - o xv6modified.tar.gz
  - ProgAsm6\_CS20BTECH11063.pdf
- To compile and run the entire xv6 operating system

```
$ make
$ make qemu
```

#### To run pgtprint function

\$ pgtprint

#### To run demandpaging

- \$ demandpaging
- The entire repository of xv6 has been cloned from the following github link
  - o xv6-public
- To add a new system call which prints current date and time
  - In syscall.h the following changes are made
    - #define SYS pgtprint 23 is added
  - In syscall.c the following changes are made
    - extern int sys\_pgtprint(void); is added which is an external system call defined elsewhere
    - [SYS pgtprint] sys pgtprint is added inside static int (\*syscalls[])(void)
    - Inside sysproc.c a system call sys\_pgtprint() is added to print page table entires
    - Inside user.h a function which would be called by user to print date and time is added, int pgtprint()
    - Inside usys.S, SYSCALL(pgtprint) is added, so that the user can now call this function to print the date and time
    - Added pgtprint to the UPROGS definition in Makefile
    - A file pgtprint.c has been created which calls pgtprint() to print the page tables
- The working and design of the program:

- To create a system call, we first should assign a system call number to our new system call
- When the user inputs the name of the system call in xv6, the console reads the name of the function and find the relevant number which corresponds to the function
- A system call is executed in kernel mode as the program needs access low level functions
- The program generates a trap
- During a system call, the processor switches from user mode to kernel mode, then it completes the process and returns to user mode
- Traps are handled in trapasm.S
- The file makes a call trap(struct trapframe \*) in trap.c
- A system call can also take arguments to kernel mode or take return value back to user mode
- o syscall.c stores an array of function pointers that return int
- There is a sys\_ prefix so that assembly instructions for each syscall are not written
- A particular syscall takes eax register value that we placed on usys.S with its corresponding syscall number and then calls sys\_<name of syscall> and puts the return value into eax register
- When a user function is passed with any parameters, they are stored in the user stack and are accessed using argint argstr and argptr given in syscall.c
- Pointers to the memory location are passed to avoid unnecessary overhead
- Important Points to note
  - To use pgtprint call, enter pgtprint in the console and press Enter
  - The entire page table entry will be printed as per requirements mentioned in the assignment

Here are some of the sample outputs for various situations as mentioned in the assignment

• Normal pgtprint without any arrays

```
(kaliltanmay@DESKTOP-54NA06G) - [~/.../OS 2/ProgAsm6/modifed/xv6-public]

$ make qemu
qemu-system-i386 -serial mon:stdio -drive file=fs.img,index=1,media=disk,format=raw -drive file=xv6.img,index=0,media=disk,format=raw -smp 2 -m 512
xv6...
cpu0: starting 0
sb: size 1000 nblocks 941 ninodes 200 nlog 30 logstart 2 inodestart 32 bmap start 58
init: starting sh
$ pgtprint
Entry number: 0, Virtual address: 8dee2027, Physical address: dee2027
Entry number: 1, Virtual address: 8dedf067, Physical address: dedf067
$ pgtprint
Entry number: 0, Virtual address: 8df2c027, Physical address: df2c027
Entry number: 1, Virtual address: 8df74067, Physical address: df74067
$ pgtprint
Entry number: 0, Virtual address: 8dfbc027, Physical address: df74067
$ pgtprint
Entry number: 0, Virtual address: 8dfbc027, Physical address: dfbc027
Entry number: 1, Virtual address: 8dfbc027, Physical address: dfbc027
Entry number: 1, Virtual address: 8def067, Physical address: dedf067
$ publication of the file of the file
```

• pgtprint with int arrGlobal[100000];

```
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```

• pgtprint with int arrLocal[100000];

```
(kaliltanmay® DESKTOP-54NA06G) -[~/.../OS 2/ProgAsm6/modifed/xv6-public]

$ make qemu
qemu-system-i386 -serial mon:stdio -drive file=fs.img, index=1, media=disk, format=raw -drive file=xv6.img, index=0, media=disk, format=raw -smp 2 -m 512
xv6...
cpu0: starting 0
sb: size 1000 nblocks 941 ninodes 200 nlog 30 logstart 2 inodestart 32 bmap start 58
init: starting sh
$ pgtprint
0
Entry number: 0, Virtual address: 8dee2027, Physical address: dee2027
Entry number: 1, Virtual address: 8dedf067, Physical address: dedf067
$ pgtprint
0
Entry number: 0, Virtual address: 8df2c027, Physical address: df2c027
Entry number: 1, Virtual address: 8df74067, Physical address: df74067
$ pgtprint
0
Entry number: 0, Virtual address: 8df6c027, Physical address: df6c027
Entry number: 1, Virtual address: 8df6c027, Physical address: df6c027
Entry number: 1, Virtual address: 8dedf067, Physical address: df6c027
Entry number: 1, Virtual address: 8dedf067, Physical address: dedf067

$ \textstyre \textstyre
```

In the case of using local array the pagetable printing does not change while in the case of using global array we can see that many pagetable entries are being printed

In the case of global variables values were assigned statically on compile time and in the case of local variables it is assigned on the heap dynamically

### Assignment part 2

The following changes have been made for part 2 of the assignment

```
    exec.c
    if((sz = allocuvm(pgdir, sz, ph.vaddr + ph.filesz)) == 0)
    sz += ph.memsz - ph.filesz;
    The above changes have been made at line 54-70
```

- trap.c
  - switch case for T\_PGFLT has been added at line 80
- demandpaging.c
  - The following file has been added to give user program functionality
- defs.h
  - o int map\_to\_page(pde\_t \*pgdir, void \*va, uint size, uint pa, int perm);

- The declaration of function has been made in the file to map memory to pages
- vm.c
  - int map\_to\_page(pde\_t \*pgdir, void \*va, uint size, uint pa, int perm)
  - This function is same as mapppages
  - Both the functions are exactly same
- Added \_demandpaging to the UPROGS definition in Makefile

Here are some of the sample outputs for demandpaging in the operating system

```
(kaliltammay@DESKTOP-54NA06G) -[~/.../OS 2/ProgAsm6/modifed/xv6-public]

qemm-system-i386 -serial mon:stdio -drive file=fs.img,index=1,media=disk,format=raw -drive file=xv6.img,index=0,media=disk,format=raw -smp 2 -m 512 xv6...

cpu0: starting 0

sb: size 1000 nblocks 941 ninodes 200 nlog 30 logstart 2 inodestart 32 bmap start 58

init: starting sh

5 demandpaging

global addr from user space: B00

page fault occurred, doing demand paging for address: 0x1000

Pgdir number 0 Entry number: 0, Virtual address: 8dee1000, Physical address: dee1004

Pgdir number 0 Entry number: 1, Virtual address: 8dee1014, Physical address: dee1014

page fault occurred, doing demand paging for address: 0x2000

Pgdir number 0 Entry number: 0, Virtual address: 8dee1004, Physical address: dee1004

Pgdir number 0 Entry number: 0, Virtual address: 8dee1004, Physical address: dee1004

Pgdir number 0 Entry number: 1, Virtual address: 8dee1004, Physical address: dee1004

Pgdir number 0 Entry number: 2, Virtual address: 8dee1004, Physical address: dee1004

Pgdir number 0 Entry number: 3, Virtual address: 8dee1004, Physical address: dee1004

Pgdir number 0 Entry number: 3, Virtual address: 8dee1004, Physical address: dee1004

Pgdir number 0 Entry number: 3, Virtual address: 8dee1004, Physical address: dee1004

Pgdir number 0 Entry number: 1, Virtual address: 8dee1004, Physical address: dee1004

Pgdir number 0 Entry number: 1, Virtual address: 8dee1000, Physical address: dee1004

Pgdir number 0 Entry number: 2, Virtual address: 8dee1000, Physical address: dee1004

Pgdir number 0 Entry number: 3, Virtual address: 8dee1000, Physical address: dee1004

Pgdir number 0 Entry number: 4, Virtual address: 8dee1000, Physical address: dee1004

Pgdir number 0 Entry number: 4, Virtual address: 8dee1004, Physical address: dee1004

Pgdir number 0 Entry number: 4, Virtual address: 8dee1004, Physical address: dee1004

Pgdir number 0 Entry number: 4, Virtual address: 8dee1004, Physical address: dee1004

Pgdir number 0 Entry number: 4, Virtual address: 8dee1004,
```

The above assignment has been done in collaboration with Dr. Rajesh Kedia and students for page fault creation in exec.c Their roll numbers and names are:

- Aayush Patel CS20BTECH11001
- Vikhyath CS20BTECH11056

The following are some of the references used to solve the assignment

- https://pdos.csail.mit.edu/6.828/2017/homework/xv6-zero-fill.html
- https://oslab.kaist.ac.kr/wp-content/uploads/esos\_files/courseware/undergraduate/UKD/homework/homework04.pdf
- https://pdos.csail.mit.edu/6.828/2008/lec/l5.html
- https://www.cs.virginia.edu/~cr4bd/4414/F2018/paging-and-protection.html#tocAnchor-1-3
- https://www.cs.columbia.edu/~junfeng/11sp-w4118/lectures/exec.pdf
- https://www.cse.iitd.ac.in/~sbansal/os/previous\_years/2011/xv6-book/mem.pdf

#### Note

- The above results may differ based on the opprating system xv6 is run on
- It may also depend on the version of xv6 using and the architecture on which it is running
- The address printed where page fault occured is the address got from rcr2()