# Operating System 2 Programming Assignment 6

## Tanmay Garg CS20BTECH11063

- There are 2 files in the folder
  - o xv6modified.tar.gz
  - ProgAsm6\_CS20BTECH11063.pdf
- To compile and run the entire xv6 operating system

```
$ make
$ make qemu
```

#### To run pgtprint function

\$ pgtprint

#### To run demandpaging

- \$ demandpaging
- The entire repository of xv6 has been cloned from the following github link
  - o xv6-public
- To add a new system call which prints current date and time
  - In syscall.h the following changes are made
    - #define SYS pgtprint 23 is added
  - In syscall.c the following changes are made
    - extern int sys\_pgtprint(void); is added which is an external system call defined elsewhere
    - [SYS pgtprint] sys pgtprint is added inside static int (\*syscalls[])(void)
    - Inside sysproc.c a system call sys\_pgtprint() is added to print page table entires
    - Inside user.h a function which would be called by user to print date and time is added, int pgtprint()
    - Inside usys.S, SYSCALL(pgtprint) is added, so that the user can now call this function to print the date and time
    - Added pgtprint to the UPROGS definition in Makefile
    - A file pgtprint.c has been created which calls pgtprint() to print the page tables
- The working and design of the program:

- To create a system call, we first should assign a system call number to our new system call
- When the user inputs the name of the system call in xv6, the console reads the name of the function and find the relevant number which corresponds to the function
- A system call is executed in kernel mode as the program needs access low level functions
- The program generates a trap
- During a system call, the processor switches from user mode to kernel mode, then it completes the process and returns to user mode
- Traps are handled in trapasm.S
- The file makes a call trap(struct trapframe \*) in trap.c
- A system call can also take arguments to kernel mode or take return value back to user mode
- o syscall.c stores an array of function pointers that return int
- There is a sys\_ prefix so that assembly instructions for each syscall are not written
- A particular syscall takes eax register value that we placed on usys.S with its corresponding syscall number and then calls sys\_<name of syscall> and puts the return value into eax register
- When a user function is passed with any parameters, they are stored in the user stack and are accessed using argint argstr and argptr given in syscall.c
- Pointers to the memory location are passed to avoid unnecessary overhead
- Important Points to note
  - To use pgtprint call, enter pgtprint in the console and press Enter
  - The entire page table entry will be printed as per requirements mentioned in the assignment

Here are some of the sample outputs for various situations as mentioned in the assignment

• Normal pgtprint without any arrays

```
(kaliltanmay@DESKTOP-54NA06G) - [~/.../OS 2/ProgAsm6/modifed/xv6-public]

$ make qemu
qemu-system-i386 -serial mon:stdio -drive file=fs.img,index=1,media=disk,format=raw -drive file=xv6.img,index=0,media=disk,format=raw -smp 2 -m 512
xv6...
cpu0: starting 0
sb: size 1000 nblocks 941 ninodes 200 nlog 30 logstart 2 inodestart 32 bmap start 58
init: starting sh
$ pgtprint
Entry number: 0, Virtual address: 8dee2027, Physical address: dee2027
Entry number: 1, Virtual address: 8dedf067, Physical address: dedf067
$ pgtprint
Entry number: 0, Virtual address: 8df2c027, Physical address: df2c027
Entry number: 1, Virtual address: 8df74067, Physical address: df74067
$ pgtprint
Entry number: 0, Virtual address: 8dfbc027, Physical address: df74067
$ pgtprint
Entry number: 0, Virtual address: 8dfbc027, Physical address: dfbc027
Entry number: 1, Virtual address: 8dfbc027, Physical address: dfbc027
Entry number: 1, Virtual address: 8def067, Physical address: dedf067
$ publication of the file of the file
```

• pgtprint with int arrGlobal[100000];

```
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```

• pgtprint with int arrLocal[100000];

```
(kaliltanmay® DESKTOP-54NA06G) -[~/.../OS 2/ProgAsm6/modifed/xv6-public]

$ make qemu

qemu-system-i386 -serial mon:stdio -drive file=fs.img, index=1, media=disk, format=raw -drive file=xv6.img, index=0, media=disk, format=raw -smp 2 -m 512 xv6...

qu0: starting 0

sb: size 1000 nblocks 941 ninodes 200 nlog 30 logstart 2 inodestart 32 bmap start 58

init: starting sh

$ pgtprint

0

Entry number: 0, Virtual address: 8dee2027, Physical address: dee2027

Entry number: 1, Virtual address: 8df2c027, Physical address: dedf067

$ pgtprint

0

Entry number: 0, Virtual address: 8df74067, Physical address: df74067

$ pgtprint

0

Entry number: 0, Virtual address: 8df74067, Physical address: df74067

$ pgtprint

0

Entry number: 0, Virtual address: 8df6c027, Physical address: df6c027

Entry number: 1, Virtual address: 8df6c027, Physical address: df6c027

Entry number: 1, Virtual address: 8dedf067, Physical address: df6c027

Entry number: 1, Virtual address: 8dedf067, Physical address: dedf067

Entry number: 1, Virtual address: 8dedf067, Physical address: dedf067
```

In the case of using local array the pagetable printing does not change while in the case of using global array we can see that many pagetable entries are being printed

### Assignment part 2

The following changes have been made for part 2 of the assignment

- The above changes have been made at line 56 and 58
- trap.c
  - switch case for T\_PGFLT has been added at line 80
- demandpaging.c
  - The following file has been added to give user program functionality
- defs.h
  - int map\_to\_page(pde\_t \*pgdir, void \*va, uint size, uint pa, int perm);
  - The declaration of function has been made in the file to map memory to pages
- vm.c

- o int map\_to\_page(pde\_t \*pgdir, void \*va, uint size, uint pa, int perm)
- This function is same as mapppages
- o Both the functions are exactly same

Here are some of the sample outputs for demandpaging in the operating system

```
kaliltanmay⊕ DESKTOP-54NA06G) -[~/.../OS 2/ProgAsm6/modifed/xv6-public]

$ make qemu
qemu-system-i336 - serial mon:stdio -drive file=fs.img,index=1,media=disk,format=raw -drive file=xv6.img,index=0,media=disk,format=raw -smp 2 -m 512
xv6...

cpu0: starting 0
sb: size 1000 mblocks 941 ninodes 200 nlog 30 logstart 2 inodestart 32 bmap start 58
init: starting sh
$ demandpaging
global addr from user space: B00
page fault occurred, doing demand paging for address: 0x1000
Pgdir number 0 Entry number: 0, Virtual address: 8dee1000, Physical address: dee1000
Pgdir number 0 Entry number: 1, Virtual address: 8dee1014, Physical address: dee1014
page fault occurred, doing demand paging for address: 0x2000
Pgdir number 0 Entry number: 2, Virtual address: 8dee1000, Physical address: dee1004
Pgdir number 0 Entry number: 1, Virtual address: 8dee1000, Physical address: dee1004
Pgdir number 0 Entry number: 1, Virtual address: 8dee1004, Physical address: dee1004
Pgdir number 0 Entry number: 1, Virtual address: 8dee1004, Physical address: dee1004
Pgdir number 0 Entry number: 3, Virtual address: 8dee1004, Physical address: dee1004
Pgdir number 0 Entry number: 3, Virtual address: 8dee1004, Physical address: dee1004
Pgdir number 0 Entry number: 1, Virtual address: 8dee1004, Physical address: dee1004
Pgdir number 0 Entry number: 0, Virtual address: 8dee1004, Physical address: dee1004
Pgdir number 0 Entry number: 2, Virtual address: 8dee1004, Physical address: dee1004
Pgdir number 0 Entry number: 2, Virtual address: 8dee1004, Physical address: dee1008
Pgdir number 0 Entry number: 3, Virtual address: 8dee1006, Physical address: dee1008
Pgdir number 0 Entry number: 3, Virtual address: 8dee1004, Physical address: dee1004
Pgdir number 0 Entry number: 3, Virtual address: 8dee1004, Physical address: dee1004
Pgdir number 0 Entry number: 4, Virtual address: 8dee1004, Physical address: dee1004
Pgdir number 0 Entry number: 3, Virtual address: 8dee1004, Physical address: dee1004
Pgdir number 0 Entry number: 4, Virtual address: 8dee1004, Physical
```

The above assignment has been done in collaboration with Dr. Rajesh Kedia and students for page fault creation in exec. c Their roll numbers and names are:

- Aayush Patel CS20BTECH11001
- Vikhyath CS20BTECH11056

The following are some of the references used to solve the assignment

- https://pdos.csail.mit.edu/6.828/2017/homework/xv6-zero-fill.html
- https://oslab.kaist.ac.kr/wpcontent/uploads/esos\_files/courseware/undergraduate/UKD/homework/homework04.pdf
- https://pdos.csail.mit.edu/6.828/2008/lec/l5.html
- https://www.cs.virginia.edu/~cr4bd/4414/F2018/paging-and-protection.html#tocAnchor-1-3
- https://www.cs.columbia.edu/~junfeng/11sp-w4118/lectures/exec.pdf
- https://www.cse.iitd.ac.in/~sbansal/os/previous\_years/2011/xv6-book/mem.pdf