



# RTES theory and analysis:

## Rate Monotonic Theory - Optimal Priority Policy

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## ■ Segment Outline

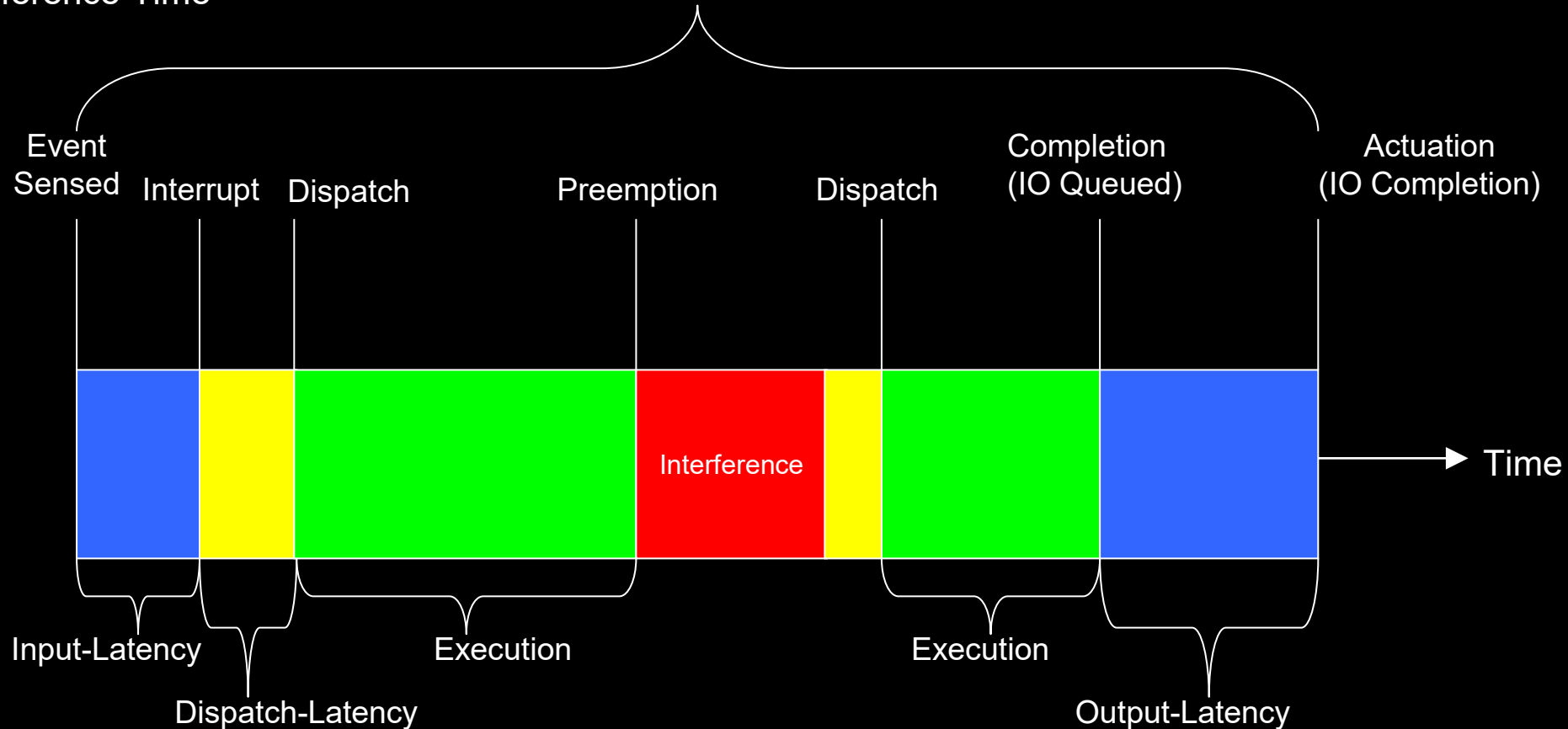
- **The Real-Time Service Response Timeline**
- **Argument for Optimality of RM Priority Assignment Policy**
- **Policy - Highest Frequency (Shortest Period) Services is assigned highest priority**
- **Same Frequency given same priority, dispatched in FIFO order from ready queue**
- **Adjust priority up by assuming higher frequency than actual (period transform)**

# A Service Release and Response

- $C_i$ , WCET
- Input/Output Latency
- Interference Time

$$\text{Response Time} = \text{Time}_{\text{Actuation}} - \text{Time}_{\text{Sensed}}$$

(From Release to Response)



## ■ RM Assumptions and Constraints

- **A1: Periodic Service Requests, the Period is Constant**
- **A2: Completion-Time < Period**
- **A3: Service Requests are Independent (No Known Phasing)**
- **A4: Run-time is Known and Deterministic (WCET may be Used)**
- **C1: Deadline = Period by Definition**
- **C2: Fixed Priority, Preemptive, Run-to-Completion Scheduling**
- **Critical Instant: longest response time for a service occurs when all system services are requested simultaneously (maximum interference case for lowest priority service)**
- **No Other Shared Resources – Not in Paper, but key assumption – e.g. shared memory**

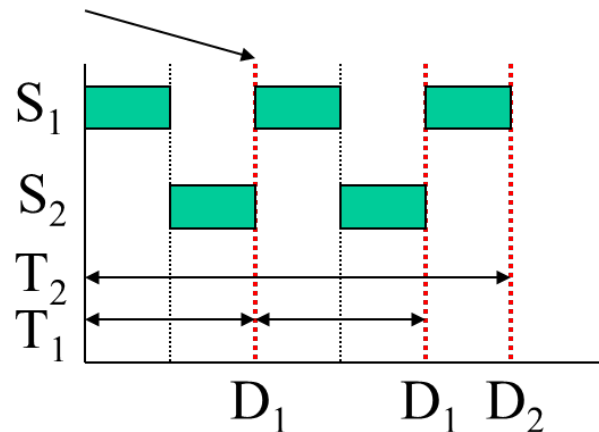
# RM Priority Assignment Policy

## RM Policy:

Given: services  $S_1, S_2$  with periods  $T_1$  and  $T_2$  and  $C_1$  and  $C_2$  with  $T_2 > T_1$

E.g.  $T_1=2, T_2=5, C_1=1, C_2=2$ , then if  $\text{prio}(S_1) > \text{prio}(S_2)$ , note that:

$S_1$  Makes Deadline if  $\text{prio}(S_1) > \text{prio}(S_2)$

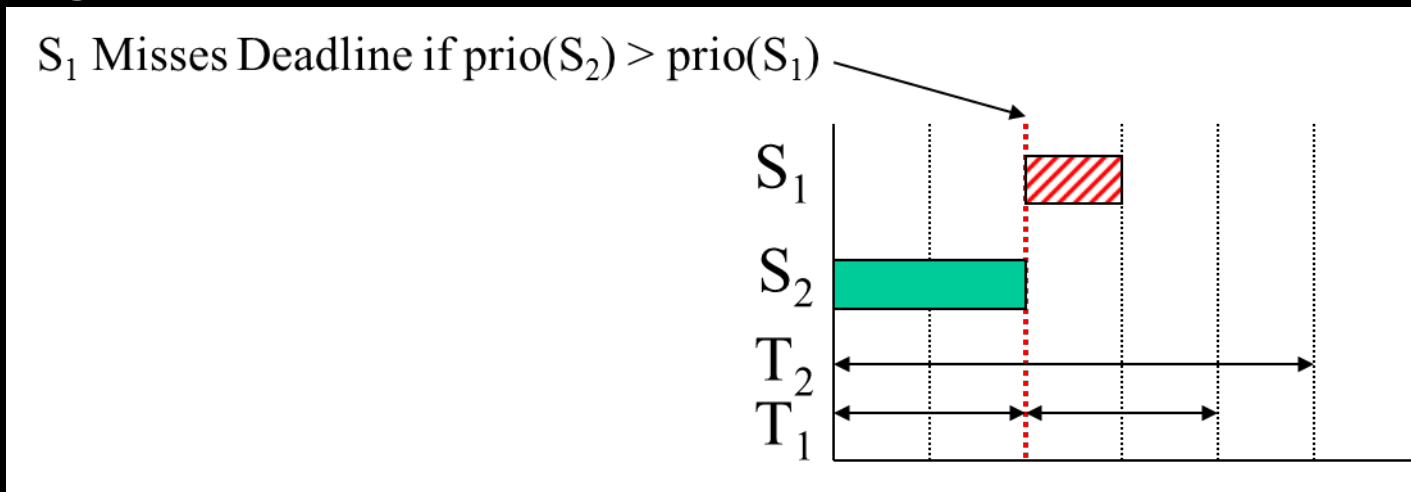


# RM Priority Assignment Policy

## Consider Alternative Policy:

Given: services  $S_1$ ,  $S_2$  with periods  $T_1$  and  $T_2$  and  $C_1$  and  $C_2$  with  $T_2 > T_1$

E.g.  $T_1=2$ ,  $T_2=5$ ,  $C_1=1$ ,  $C_2=2$ , then if  $\text{prio}(S_2) > \text{prio}(S_1)$ , note that:



## Conclusion:

If  $\{S_n\}$  feasible with  $\text{prio}(S_2) > \text{prio}(S_1)$ , then  $\{S_n\}$  is also always feasible given  $\text{prio}(S_1) > \text{prio}(S_2)$ , but converse is not necessarily TRUE!!

Therefore,  $\text{prio}(S_1) > \text{prio}(S_2)$  is OPTIMAL

## ■ Summary

- The Real-Time Service Response Timeline
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