Algorithm Code Book

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Chapter 1

Data Structure

1.1 Trie

1.1.1 Static Trie

```
1 #define Max 10005
2 int getId(char c)
3 {
       return c>='a'?c-'a':c-'A'+26;
4
5 }
6 struct Trie
7
8
       struct Tree
9
           int Next[52];
11
           bool word;
12
           void clear()
           {
13
                word=false;
14
                memset(Next, -1, sizeof(Next));
15
16
       } T[Max];
17
       int ptr;
18
       void clear()
19
21
           ptr=1;
           T[0].clear();
22
           memset(T[0].Next,0,sizeof(T[0].Next));
23
24
       void Insert(const char *str)
25
26
           int p=0;
27
           for (int i=0; str[i]; i++)
28
29
                int id=getId(str[i]);
                if(T[p].Next[id] <= 0)
31
32
                    T[p].Next[id]=ptr;
33
                    T[ptr++].clear();
34
35
                p=T[p].Next[id];
36
37
           T[p]. word=true;
38
39
       bool Search (const char *str)
```

```
41
            int p=0;
42
            for(int i=0; str[i]; i++)
43
44
                int id=getId(str[i]);
45
                if (T[p]. Next[id]>0)
                {
                     p=T[p]. Next[id];
48
49
                else return false;
50
51
           return T[p].word;
52
53
54 };
55 Trie A;
```

1.2 RMQ

1.2.1 Bit

1D Bit

```
1 #define MaxVal 100000
2 int Bit [MaxVal];
3 /**find sum from 1 to idx**/
4 int read(int idx)
5 {
6
       int sum = 0;
       while (idx > 0)
8
           sum += Bit[idx];
9
           i\,\mathrm{d} x \ -\!\!= \ (\,i\,\mathrm{d} x \ \& \ -i\,\mathrm{d} \, x\,)\;;
10
11
       return sum;
12
13 }
  /**update value ind to MaxVal**/
14
void update(int idx ,int val)
16
       while (idx \le MaxVal)
17
18
19
            Bit[idx] += val;
20
           idx += (idx \& -idx);
21
22
23
  /**Find the value of idx**/
24
  int readSingle(int idx)
25
26
       int sum = Bit[idx]; /// sum will be decreased
27
       if (idx > 0) /// special case
28
29
            int z = idx - (idx & -idx); /// make z first
30
            idx--; /// idx is no important any more, so instead y, you can use
31
      idx
            while (idx != z) /// at some iteration idx (y) will become z
32
33
                sum -= Bit[idx];/// substruct Bit frequency which is between y
34
      and "the same path"
                idx = (idx \& -idx);
35
            }
36
```

```
return sum;
}
```

2D Bit

```
void updatey(int x , int y , int val)
2 {
       while (y \le \max_{y})
3
4
           tree[x][y] += val;
5
6
           y += (y \& -y);
7
8 }
9
  void update(int x , int y , int val)
10 {
       while (x \le \max_{x} x)
11
12
           updatey(x , y , val); // this function should update array tree[x]
13
           x += (x \& -x);
14
15
16 }
```

1.2.2 Square Root Decomposition

```
1 #include < bits / stdc++.h>
2 using namespace std;
3 \text{ const int } sz = 100005;
4 const int inf=(1 < < 28);
5 template < typename t> t MIN3(t a, t b, t c)
6 {
7
       return min(a, min(b, c));
8 }
9 int BLOCK[400];
10 int arr[sz];
  int getId(int indx,int blockSZ)
11
12
13
       return indx/blockSZ;
14 }
void init (int sz)
16
       for (int i=0; i \le sz; i++)BLOCK[i]=inf;
17
18
  void update(int val, int indx, int blockSZ)
19
20 {
       int id=getId(indx, blockSZ);
21
       BLOCK[id]=min(BLOCK[id], val);
22
23 }
24 int query (int L, int R, int blockSZ)
25 {
       int lid=getId(L, blockSZ);
26
       int rid=getId(R, blockSZ);
27
       if(lid=rid)
28
29
            int ret=inf;
30
            for(int i=L; i<=R; i++)ret=min(ret, arr[i]);</pre>
31
           return ret;
32
33
       int ml=inf,m2=inf,m3=inf;
       for (int i=L; i<(lid+1)*blockSZ; i++)ml=min(m1, arr[i]);
35
       for (int i=lid+1; i<rid; i++)m2=min(m2,BLOCK[i]);</pre>
36
       for (int i=rid*blockSZ; i \le R; i++)m3=min(m3, arr[i]);
```

```
return MIN3(m1,m2,m3);
38
39 }
40 int main()
41 {
       int N,Q;
42
       scanf ("%d %d",&N,&Q);
43
       int blockSZ=sqrt(N);
44
       init (blockSZ);
45
       for (int i=0; i < N; i++)
46
47
            int x;
48
            scanf("%d",&x);
49
            arr[i]=x;
50
51
            update(x, i, blockSZ);
       }
53
       while (Q--)
55
            int x, y;
            scanf("%d %d",&x,&y);
56
            printf("%d\n", query(x,y, blockSZ));
57
58
       return 0;
59
60
```

1.2.3 MO's Algorithm

```
2
      MO's Algorithm
      \verb|problem:| | \texttt{http://www.spoj.com/problems/DQUERY}|
3
4
       MOs algorithm is just an order in which we process the queries.
5
      We were given M queries, we will re-order the queries in a particular
6
      order and then process them.
       Clearly, this is an off-line algorithm. Each query has L and R, we will
      call them opening and closing.
       Let us divide the given input array into Sqrt(N) blocks.
8
       Each block will be N / Sqrt(N) = Sqrt(N) size.
       Each opening has to fall in one of these blocks.
10
       Each closing has to fall in one of these blocks.
11
12
       All the queries are first ordered in ascending order of their block
13
      number (block number is the block in which its opening falls).
       Ties are ordered in ascending order of their R value.
14
15
16 **/
17 #include < bits / stdc++.h>
18 using namespace std;
19 #define Mx 30005
20 #define MxNum 1000005
21 int BlockSize;
22 int Answer;
int Freq[MxNum],Num[Mx];
24 struct info
25 {
26
       int L,R,qno;
       info(int L=0,int R=0,int qno=0):L(L),R(R),qno(qno) {};
27
28
       bool operator < (const info &a) const
29
           if (L/BlockSize!=a.L/BlockSize) return L/BlockSize <a.L/BlockSize;</pre>
           return R<a.R;
31
```

```
33 } Query [200005];
34 int StoreAnswer [200005];
35 void Add(int indx)
36 {
37
       Freq[Num[indx]]++;
38
       if (Freq [Num[indx]]==1) Answer++;
39
  void Remove(int indx)
40
41 {
       Freq[Num[indx]] - -;
42
       if(Freq[Num[indx]]==0)Answer--;
43
44 }
45 int main()
46
47
       int N;
       scanf("%d",&N);
48
49
       BlockSize=sqrt(N);
       for (int i=0; i< N; i++)
50
51
            scanf("%d",&Num[i]);
53
       int Q;
54
       scanf("%d",&Q);
55
       for(int i=0; i<Q; i++)
56
57
       {
            int x, y;
58
59
            scanf("%d %d",&x,&y);
            Query [i] = info(x-1,y-1,i);
60
61
       sort(Query,Query+Q);
62
       int currentL=0,currentR=0;
63
       Answer=0;
64
       for (int i=0; i<Q; i++)
65
66
            int L=Query[i].L;
67
            int R=Query[i].R;
68
            while (currentL<L)
69
            {
                 Remove(currentL);
                 currentL++;
72
73
            while (currentL>L)
74
75
                 Add(currentL -1);
76
                 currentL --;
77
            }
78
            while (currentR<=R)
79
            {
                 Add(currentR);
81
82
                 currentR++;
            }
83
            while (currentR>R+1)
84
            {
85
                 Remove (\operatorname{current} R - 1);
86
                 currentR --;
87
88
            StoreAnswer [Query [i].qno]=Answer;
89
90
       for (int i=0; i < Q; i++)
91
92
       {
            printf("%d\n", StoreAnswer[i]);
93
```

```
94 }
95 return 0;
96 }
```

1.2.4 Segment Tree

Lazy Propagration1

```
1 /**
2 **You are given an array of N elements, which are initially all 0. After **
      that you will be given C commands. They are
  **0 p q v - you have to add v to all numbers in the range **of p to q (
      inclusive), where p and q are two indexes of the array.
  **1 p q - output a line containing a single integer which is the sum of all
      **the array elements between p and q (inclusive)
5 */
6 #include < bits / stdc++.h>
7 using namespace std;
8 typedef long long LLD;
9 LLD tree [3*100005];
10 LLD lazy [3*100005];
  void update(int left,int right,int index,int x,int y,int value)
11
12
  {
13
       if (x<=left&&y>=right)
14
       {
            tree[index] += (LLD)(right - left + 1)*value;
           lazy[index]+=value;
           return;
17
18
       int mid = (left + right)/2;
19
       if (lazy [index]!=0)
20
21
           tree[2*index]+=(LLD)(mid-left+1)*lazy[index];
22
           tree[2*index+1]+=(LLD)(right-mid)*lazy[index];
           lazy [2*index]+=lazy [index];
24
           lazy [2*index+1]+=lazy [index];
25
           lazy[index]=0;
26
27
       if(x \le mid)
28
29
           update(left, mid,2*index,x,y,value);
30
31
       if (y>mid)
32
33
       {
           update(mid+1,right,2*index+1,x,y,value);
34
35
       tree[index] = tree[2*index] + tree[2*index+1];
36
37
38 LLD query(int left, int right, int index, int x, int y)
39
40
      LLD a1=0, a2=0;
       if (x<=left&&y>=right)
41
42
       {
           return tree[index];
43
       int mid = (left + right)/2;
45
       if (lazy [index]!=0)
46
47
           tree[2*index]+=(LLD)(mid-left+1)*lazy[index];
48
           tree[2*index+1]+=(LLD)(right-mid)*lazy[index];
49
           lazy [2*index] += lazy [index];
50
```

```
lazy [2*index+1]+=lazy [index];
51
            lazy[index]=0;
52
53
       if(x \le mid)
54
       {
            a1=query(left, mid, 2*index, x, y);
57
       if (y>mid)
58
59
       {
            a2=query(mid+1, right, 2*index+1, x, y);
60
61
       return (a1+a2);
62
63
64 int main()
65
66
       int test,t;
       scanf("%d",&test);
67
       for (t=1;t \le test;t++)
68
69
            memset(tree, 0, sizeof(tree));
70
            memset(lazy,0,sizeof*lazy);
71
            int s, q;
72
            scanf("%d %d",&s,&q);
73
            while (q--)
74
75
            {
                 int x,y,v,dec;
76
                 scanf("%d",&dec);
                 if (dec)
79
                 {
                      scanf("%d %d",&x,&y);
80
                     LLD ans=query (0, s-1, 1, x-1, y-1);
81
                      printf("%lld\n",ans);
82
                 }
83
                 else
84
                 {
85
                      scanf("%d %d %d",&x,&y,&v);
86
                      update (0, s-1, 1, x-1, y-1, v);
87
89
90
       return 0;
91
92
```

Lazy Propagration2

```
13 {
14
        if (left=right)
            Tree [indx][0]=1;
16
            Tree [indx][1] = Tree [indx][2] = lazy[indx] = 0;
17
            return;
18
19
       int mid = (left + right)/2;
20
       build (left, mid, 2*indx);
21
       build (mid+1, right, 2*indx+1);
22
       for (int i=0; i<3; i++)
23
       {
24
            Tree [indx][i] = Tree[2*indx][i] + Tree[2*indx+1][i];
25
26
27
28
   void update(int left, int right, int indx, int x, int y, int add)
29
        if (lazy [indx])
30
31
             int lazy_val=lazy[indx];
32
            lazy [2*indx] = (lazy [2*indx] + lazy_val) \%3;
33
            lazy [2*indx+1]=(lazy [2*indx+1]+lazy_val)%3;
34
            for (int i=0; i<3; i++)temp [(lazy_val+i)%3]=Tree [indx][i];
35
            for (int i=0; i<3; i++)Tree [indx][i]=temp[i];
36
            lazy[indx]=0;
37
38
        if (left >y | | right <x) return;</pre>
39
40
       if(x \le left \&\&right \le y)
41
42
             for (int i = 0; i < 3; i++)
43
            {
                 temp[(i+add)%3]=Tree[indx][i];
44
45
            for (int i = 0; i < 3; i++) Tree [indx] [i] = temp[i];
46
            lazy [2*indx] = (lazy [2*indx] + add) \%3;
47
            lazy [2*indx+1]=(lazy [2*indx+1]+add) \%3;
48
49
            return;
50
        int mid = (left + right)/2;
51
       update(left, mid, 2 * indx, x, y, add);
52
       update(mid+1, right, 2*indx+1, x, y, add);
       for (int i = 0; i < 3; i++)
54
            Tree [indx][i] = Tree[2*indx][i] + Tree[2*indx+1][i];
56
57
58
       query (int left, int right, int indx, int x, int y)
59
  int
60
   {
        if (lazy [indx])
61
       {
62
            int lazy_val=lazy[indx];
63
            lazy [2*indx] = (lazy [2*indx] + lazy_val) \%3;
64
            lazy [2*indx+1]=(lazy [2*indx+1]+lazy_val)%3;
65
            for (int i=0;i<3;i++)temp[(lazy_val+i)%3]=Tree[indx][i];
66
             for (int i=0; i<3; i++) Tree [indx] [i]=temp[i];
67
            lazy[indx]=0;
68
69
       if(left>y||right<x)return 0;
70
        if(x \le left \&\& right \le y) return Tree[indx][0];
71
       int mid = (left + right)/2;
72
       return query(left, mid, 2*indx,x,y)+query(mid+1,right,2*indx+1,x,y);
```

```
74 }
75 int main()
76
        int x,y;
77
        int test;
78
        scanf("%d",&test);
79
80
        for (int t=1; t \le t \in t; t++)
81
            memset(lazy,0, sizeof(lazy));
82
            int N,Q;
83
            scanf("%d %d",&N,&Q);
84
            build(0,N-1,1);
85
             printf("Case \%d: \n",t);
86
87
             for (int i=0; i<Q; i++)
88
             {
89
                 int d;
                 scanf("%d %d %d",&d,&x,&y);
91
                  if(d==0)
92
                      update(0, N-1, 1, x, y, 1);
93
94
                  else printf("%d \setminus n", query(0,N-1,1,x,y));
95
            }
96
97
98
        return 0;
99
```

Segment Tree Variant 1

```
1 /**
2 **Give a array Of N numbers. Finding Maximum cumulative number frequency in
      **the range.
3 **input:
4 **10 4
5 **1 1 1 3 3 3 3 2 2 2
6 **1 5
7 **1 6
8 **1 7
9 **Output:
10 **3
11 **3
12 **4
13 **2
14 */
15 #include < bits / stdc++.h>
16 using namespace std;
17 typedef long long LLD;
18 #define MAX 50005
19 struct info
20 {
       int Lcnt, Rcnt, Max, Lnum, Rnum;
21
       info(int Lcnt=0,int Rcnt=0,int Max=0,int Lnum=0,int Rnum=0):Lcnt(Lcnt),
22
      Rcnt(Rcnt), Max(Max), Lnum(Lnum), Rnum(Rnum) {};
23 };
24 info Tree[3*MAX];
25 int arr [MAX];
info marge (const info &L, const info &R)
27
  {
28
       info ret;
       if (L.Rnum=R.Lnum)
29
```

```
ret.Max=max(L.Rcnt+R.Lcnt, max(L.Max, R.Max));
31
32
        }
33
        else ret.Max=max(L.Max,R.Max);
        ret.Lnum=L.Lnum;
34
        ret.Rnum=R.Rnum;
35
        if(L.Lnum=R.Lnum) ret.Lcnt=L.Lcnt+R.Lcnt;
36
37
        else ret.Lcnt=L.Lcnt;
        if (L.Rnum=R.Rnum) ret.Rcnt=L.Rcnt+R.Rcnt;
38
        else ret.Rcnt=R.Rcnt;
39
        return ret;
40
41 }
  void build (int L, int R, int indx)
42
43
44
        if (L=R)
45
        {
46
              Tree [indx] = info (1,1,1,arr [L], arr [R]);
              return;
        int mid=(L+R)>>1;
49
        build (L, mid, 2*indx);
50
        \texttt{build} \, (\, \texttt{mid} \! + \! 1, \! R, 2 \! * \! \texttt{ind} x \! + \! 1) \, ;
        Tree [indx] = marge (Tree [2*indx], Tree [2*indx+1]);
53
   info query(int L, int R, int indx, int x, int y)
54
55
        if (L>=x&&R<=y) return Tree[indx];</pre>
56
57
        int mid=(L+R)>>1;
58
        info c1, c2;
59
        if(x \le mid) c1 = query(L, mid, 2 * indx, x, y);
60
        if(y>mid)c2=query(mid+1,R,2*indx+1,x,y);
61
        return marge(c1,c2);
62
   int main()
63
64
65
        int test;
        scanf("%d",&test);
66
        for (int t=1; t <= test; t++)
67
69
              int N,C,Q;
              scanf("%d %d %d",&N,&C,&Q);
70
              for (int i = 0; i < N; i++)
72
                   int x;
73
                   scanf("%d",&arr[i+1]);
74
75
              build (1,N,1);
76
              printf("Case \%d: \n",t);
77
              while (Q--)
79
              {
80
                   int x,y;
                   scanf("%d %d",&x,&y);
81
                   p \, \texttt{rintf} \, (\, \texttt{"%d} \backslash \texttt{n"} \, , \texttt{query} \, (\, 1 \, , N, 1 \, , x \, , y \,) \, . \, Max) \, ;
82
              }
83
84
        return 0;
85
86 }
```

Segment Tree Variant 2

```
_2 **You are given a sequence A of N (N <= 50000) integers between -10000 and
      10000.
_3 **On this sequence you have to apply M (M <= 50000) operations:
4 **modify the i-th element in the sequence or for given x y print max{Ai + Ai
      +1 + ... + Aj \mid x \le i \le j \le y }.
6 #include < bits / stdc++.h>
7 using namespace std;
8 typedef long long LLD;
9 template < class T> T MAX3(T a, T b, T c) {return max(a, max(b, c));}
10 LLD Inf = (111 << 60);
11 #define MN 50005
12 struct info
13 {
       LLD prefixSum;
14
       LLD suffixSum;
15
       LLD Total;
16
       LLD TotalMax;
17
       info(int pre=-Inf,int suff=-Inf,int total=-Inf,int totalmax=-Inf):
18
       prefixSum(pre), suffixSum(suff), Total(total), TotalMax(totalmax) { };
19
  };
20
  info marge (const info &a, const info &b)
21
       info ret;
22
       ret. Total=a. Total+b. Total;
23
       ret.prefixSum=max(a.prefixSum,a.Total+b.prefixSum);
2.4
       ret.suffixSum=max(a.suffixSum+b.Total,b.suffixSum);
25
       ret. TotalMax=MAX3(a. TotalMax, b. TotalMax, a. suffixSum+b. prefixSum);
26
       return ret;
27
28
29 LLD arr [MN];
30 info Tree [3*MN];
  void build (int L, int R, int indx)
32
       if (L=R)
33
34
       {
            Tree [indx] = info (arr [L], arr [L], arr [L], arr [L]);
35
           return;
36
37
       int mid=(L+R)>>1;
38
       build (L, mid, 2*indx);
39
       build (mid+1,R,2*indx+1);
       Tree [indx] = marge (Tree [2*indx], Tree [2*indx+1]);
41
42
  void update(int L, int R, int indx, int x, LLD val)
43
44
       if (L==R)
45
       {
46
           Tree [indx] = info (val, val, val, val);
47
           return;
48
49
       int mid=(L+R)>>1;
       if(x \le mid) update(L, mid, 2 * indx, x, val);
51
       else update (mid+1,R,2*indx+1,x,val);
       Tree [indx] = marge (Tree [2*indx], Tree [2*indx+1]);
53
54
55 info query(int L, int R, int indx, int x, int y)
56 {
```

```
if (L=x and y=R)return Tree[indx];
57
58
         int mid = (L+R) >> 1;
         if(y \le mid) return query(L, mid, 2 * indx, x, y);
59
         else if (x>mid) return query (mid+1,R,2*indx+1,x,y);
60
         return marge(query(L, mid, 2*indx, x, mid), query(mid+1,R,2*indx+1,mid+1,y));
61
62 }
63 int main()
64 {
        #ifdef _ANICK_
65
         //f_input;
66
        #endif // _ANICK_
67
         int N;
68
         scanf("%d",&N);
69
70
         for (int i=1; i <= N; i++) scanf ("%lld", & arr [i]);
71
         build (1,N,1);
72
         int Q;
         scanf("%d",&Q);
73
         while(Q--)
74
75
               int t,x,y;
76
               scanf("%d %d %d",&t,&x,&y);
77
               \begin{array}{l} \textbf{if} (t) \, \texttt{printf} \, (\text{"\%lld} \, \backslash \text{n"} \, , \texttt{query} \, (1 \, , \! N, 1 \, , \! x \, , y) \, . \, \\ \textbf{TotalMax}) \, ; \end{array}
78
               else update (1, N, 1, x, y);
79
80
81
         return 0;
82
```

Segment Tree Variant 3

```
1 /**
2 **Given a bracket sequence.
3 ** On a bracket word one can do the following operations:
4 **replacement -- changes the i-th bracket into the opposite one
5 **check — if the word is a correct bracket expression
7 #include < bits / stdc++.h>
8 using namespace std;
9 typedef long long LLD;
10 #define MAX 50005
11 struct info
12 {
       int sum, sub;
13
       info(int sum=0, int sub=0): sum(sum), sub(sub) {};
14
15 };
info Tree [4*MAX];
17 char inp [MAX];
info marge(const info &L, const info &R)
19 {
20
       info ret;
       ret.sum= L.sum+R.sum;
21
       ret.sub=L.sub;
22
       ret.sub=min(ret.sub,L.sum+R.sub);
23
      return ret;
24
25 }
  void build (int L, int R, int indx)
26
27
       if (L==R)
28
29
       {
30
           int x;
           if(inp[L]=='(')x=1;
31
           else x=-1;
32
```

```
Tree [indx] = info(x,x);
33
34
            return;
35
       int mid=(L+R)>>1;
36
       build (L, mid, 2*indx);
37
38
       build (mid+1,R,2*indx+1);
39
       Tree [indx] = marge (Tree [2*indx], Tree [2*indx+1]);
40
  void update(int L, int R, int indx, int x)
41
42
       if (L=R)
43
       {
44
            int x;
45
46
            if(inp[L]=='(')x=1;
47
            else x=-1;
48
            Tree [indx] = info(x,x);
49
            return;
50
       int mid=(L+R)>>1;
51
       if(x \le mid) update(L, mid, 2 * indx, x);
       else update (mid+1,R,2*indx+1,x);
53
       Tree [indx] = marge (Tree [2*indx], Tree [2*indx+1]);
54
55
  info query (int L, int R, int indx, int x, int y)
56
57
       if (L=x&R=y) return Tree [indx];
58
59
       int mid=(L+R)>>1;
60
       if(y \le mid) return query(L, mid, 2 * indx, x, y);
61
       else if (x>mid) return query (mid+1,R,2*indx+1,x,y);
62
       else return marge(query(L, mid, 2*indx, x, mid), query(mid+1, R, 2*indx+1, mid))
       +1,y));
63
  int main()
64
65
       int N, t=1;
66
       while (scanf("%d",&N)==1)
67
68
            scanf("%s", inp);
69
70
            build (0, N-1, 1);
            int Q;
            printf("Test %d:\n",t++);
72
            scanf("%d",&Q);
73
            while (Q--)
74
75
                 int x;
76
                 scanf("%d",&x);
77
78
                 if(x)
                 {
                      if(inp[x-1]=='(')inp[x-1]=')';
80
                      else inp[x-1]='(';
81
                      update(0, N-1, 1, x-1);
82
                 }
83
                 else
84
                 {
85
                      info y=query (0, N-1, 1, 0, N-1);
86
                      if(y.sum=0\&\&y.sub>=0)printf("YES\n");
87
                      else printf("NO\n");
88
                 }
89
90
91
       return 0;
92
```

93 }

1.2.5 Sliding Window RMQ

```
1 /**
2
       every K size window RMQ
       Calculate in O(N+K) time
3
4 **/
5 #include < bits / stdc++.h>
6 using namespace std;
7 vector < int > Sliding RMQ (int *A, int N, int k)
8 {
       /** Create a Double Ended Queue, Qi that will store indexes of array
9
      elements
           The queue will store indexes of useful elements in every window and
10
       it will
           maintain decreasing order of values from front to rear in Qi, i.e.,
11
           arr [Qi.front []] to arr [Qi.rear()] are sorted in increasing order
       **/
13
       vector < int > MinWindow;
14
       deque<int>Q;
       int i;
16
       /* Process first k (or first window) elements of array */
17
       for (i = 0; i < k; i++)
18
19
           /// For very element, the previous largest elements are useless so
20
           /// remove them from Qi
21
           while (!Q.empty() \text{ and } A[i] \le A[Q.back()])Q.pop_back();
22
           Q. push_back(i);
23
24
       /// Process rest of the elements, i.e., from arr[k] to arr[n-1]
25
       while (i<N)
26
27
           /// The element at the front of the queue is the smallest element of
28
           /// previous window, so insert it result
29
           MinWindow.push_back(A[Q.front()]);
30
31
           /// Remove the elements which are out of this window
32
           while (!Q.empty() \text{ and } Q.front() \le i-k)Q.pop_front();
33
34
           /// Remove all elements larger than the currently
35
           /// being added element (remove useless elements)
36
           while (!Q.empty() \text{ and } A[i] \le A[Q.back()])Q.pop_back();
37
38
           /// Add current element at the rear of Qi
39
           Q. push_back(i);
40
41
42
       /// insert the minimum element of last window
43
       MinWindow.push_back(A[Q.front()]);
44
       return MinWindow;
45
46 }
47 int main()
48 {
       int A[] = \{100, 10, -1, 2, -3, -4, 10, 1, 100, 20\};
49
       vector < int > a = SlidingRMQ(A, 10, 2);
50
       for (int i=0; i < a. size(); i++)cout << a[i] << "";
51
       return 0;
52
53 }
```

1.2.6 Sparse Table

```
1
      Compute sparse table in O(NlogN)
3
       query in O(1)
       Ref link: https://www.topcoder.com/community/data-science/data-science-
      tutorials/range-minimum-query-and-lowest-common-ancestor/
6 \frac{\text{\#include}}{\text{otts}} / \text{stdc} + + .h >
7 using namespace std;
8 #define Max 10000005
9 int rmq[24][Max];
10 int A[Max];
void Compute_ST(int N)
12 {
       for (int i = 0; i < N; ++i)rmq[0][i] = i;
13
       for (int k = 1; (1 << k) < N; ++k)
14
           for (int i = 0; i + (1 << k) <= N; i++)
16
17
               int x = rmq[k - 1][i];
18
               int y = rmq[k - 1][i + (1 << k - 1)];
19
20
               rmq[k][i] = A[x] \le A[y] ? x : y;
21
           }
22
       }
23
24
  int RMQ(int i, int j)
25
26
       int k = log2(j-i);
27
       int x = rmq[k][i];
28
       29
       return A[x] \ll A[y] ? x : y;
30
31 }
33 int main()
34 {
35
36
      return 0;
37 }
```

1.3 Ternary Bit Mask

```
2 int more_bit [10];
3 int get_bit(int mask , int pos)
4 {
      return (mask / more_bit[pos]) % 3;
5
6 }
7 int set_bit(int mask, int pos , int bit)
8 {
      int tmp = (mask / more_bit[pos]) % 3;
9
      mask -= tmp * more_bit[pos];
10
11
      mask += bit * more_bit[pos];
12
      return mask;
13 }
void init(void){
      more_bit[0] = 3;
15
      for(int i = 1; i < 10; i++) more_bit[i] = 3 * more_bit[i - 1];
16
17 }
```

1.4 Largest Rectangle in Histogram

1.4.1 Largest rectangle in histogram using Stack

```
int GetMaxArea(int *hist, int n)
2 {
       stack<int> s;
3
4
5
       int max_area = 0;
6
       int tp;
       int area_with_top;
       int i = 0;
       while (i < n)
9
10
            if (s.empty() || hist[s.top()] <= hist[i])</pre>
11
                s.push(i++);
12
            else
13
            {
14
                tp = s.top();
15
16
                s.pop();
                area_with_top = hist[tp] * (s.empty() ? i : i - s.top() - 1);
17
                if (max_area < area_with_top)</pre>
18
                    max_area = area_with_top;
19
           }
20
21
       while (s.empty() == false)
22
23
           tp = s.top();
           s.pop();
25
           area_with_top = hist[tp] * (s.empty() ? i : i - s.top() - 1);
26
27
28
            if (max_area < area_with_top)</pre>
                max_area = area_with_top;
31
       return max_area;
32
```

1.4.2 Largest rectangle in histogram using Segment Tree

```
1 #include < bits / stdc++.h>
2 using namespace std;
3 #define MN 1000005
4 typedef long long LLD;
5 LLD Histogram [MN], N;
6 template < class T> T MAX3(T a, T b, T c)
7 {
       return max(a, max(b,c)); ///maximum of 3 number
8
9 }
10 struct info
11
12
       int Num;
13
       int pos;
       info(int Num=0,int pos=0):Num(Num),pos(pos) {};
14
15 };
16 info Tree [3*MN];
17 info marge(const info &L, const info &R)
18 {
       if (L.Num<R.Num) return L;
19
20
       return R;
void build (int L, int R, int indx)
```

```
23 {
       if (L==R)
24
25
       {
            Tree [indx]=info (Histogram [L], L);
26
27
            return;
28
       int mid=(L+R)>>1;
29
       build (L, mid, 2*indx);
30
       build (mid+1,R,2*indx+1);
31
       Tree [indx] = marge(Tree[2*indx], Tree[2*indx+1]);
32
33
34 info query (int L, int R, int indx, int x, int y)
35
36
       if (L=x&&R=y) return Tree[indx];
37
       int mid = (L+R) >> 1;
38
       if(y \le mid) return query(L, mid, 2 * indx, x, y);
39
       else if (x>mid) return query (mid+1,R,2*indx+1,x,y);
40
       return marge (query (L, mid, 2*indx, x, mid), query (mid+1,R, 2*indx+1, mid+1,y));
41
  LLD GetMaximumArea(int L, int R)
42
43
       if (L>R) return -1;
44
       if (L=R) return Histogram [L];
45
       info p=query(1,N,1,L,R);
46
       LLD = GetMaximumArea(L, p. pos -1);
47
       LLD b=GetMaximumArea(p.pos+1,R);
48
       LLD c=(LLD) p.Num*(R-L+1);
49
50
       return MAX3(a,b,c);
51
52
  int main()
53
       while (scanf ("%d",&N) and N)
54
            for(int i=1; i<=N; i++)scanf("%lld",&Histogram[i]);
56
57
            build (1,N,1);
            printf("%lld\n",GetMaximumArea(1,N));
58
59
60
       return 0;
61
```

1.5 Least common ancestor (LCA)

```
vector < pp > Graph [2005];
2 int Level [2005];
3 int Lavel_cost [2005];
4 int sparse_table [2005][20];
5 int sparse_cost [2005][20];
6 void init()
7 {
      mem(parents, -1);
8
9
      mem(sparse\_table, -1);
10
      mem(sparse_cost,0);
11
  void dfs(int from, int u, int deep)
12
13
       Level [u]=deep;
14
       parents [u]=from;
15
       int sz=Graph[u].size();
16
17
       for(int i=0; i < sz; i++)
18
```

```
int v=Graph[u][i].second;
19
            int c=Graph[u][i].first;
20
            if (v==from) continue;
21
            Cost[v]=c;
22
            dfs(u,v,deep+1);
23
24
25
  }
  void Sparse (int N)
26
27
       for (int i=0; i< N; i++)
28
29
            sparse_cost [ i ] [0] = Cost [ i ];
30
            sparse_table[i][0] = parents[i];
31
32
33
       for (int j=1; (1 << j) < N; j++)
34
       {
            for (int i=0; i< N; i++)
36
            {
                if (sparse_table[i][j-1]!=-1)
37
                {
38
                     int a=sparse\_table[i][j-1];
39
                     sparse_table [i][j]=sparse_table [a][j-1];
40
                     sparse\_cost[i][j] = sparse\_cost[i][j-1] + sparse\_cost[a][j-1];
41
42
            }
43
44
45
  }
46
  int LCA(int p, int q)
47
       if(Level[p] < Level[q])swap(p,q);
48
       int Log=log 2 (Level[p]) + 1;
49
       int ans=0;
50
       for (int i=Log; i>=0; i--)
51
            if ((Level[p]-(1<<i))>=Level[q])
53
            {
54
                ans+=sparse_cost[p][i];
55
56
                p=sparse_table[p][i];
58
       if(p=q)return ans;
59
       for (int i=Log; i>=0; i--)
60
61
            if (sparse_table [p][i]!=-1&&sparse_table [p][i]!=sparse_table [q][i])
62
63
                ans+=sparse_cost[p][i]+sparse_cost[q][i];
64
                p=sparse_table[p][i],q=sparse_table[q][i];
65
67
       ans+=Cost[p]+Cost[q];
68
       ///Lca will be parent[p];
69
       return ans;
70
71 }
```

1.6 Union Find

```
int parent[10000];
2 void init (int N)
3 {
         for (int i=0; i \le N; i++)
4
            parent[i]=i;
5
6 }
7 int Find(int r)
8 {
         if (parent[r]==r)return r;
9
         return parent[r]=Find(parent[r]);
10
11 }
void make_parent(int p,int q)
13 {
         parent[p]=q;
14
15 }
void check(int x, int y)
17 {
         int p=Find(x);
18
         int q=Find(y);
19
         if(p!=q)
20
         {
21
               make_parent(p,q);
22
               printf("Make friend\n");
23
               return;
24
         }
25
         printf("they allready friend\n");
26
         return;
27
28 }
```

Chapter 2

Graph Theory

2.1 DFS

2.1.1 Bicoloring

```
///color will be initial with -1
int color[20005];
bool dfs(int u,int c)

{
    if(color[u]==c)return true;
    if(color[u]==(1-c))return false;
    color[u]=c;
    bool ret=true;
    for(auto v:graph[u])ret&=dfs(v,1-c);
    return ret;
}
```

2.1.2 Cycle Finding

```
int color [20005];
2 bool dfs(int u)
3 {
       color [u]=GREY;
4
5
       bool no_cycle=true;
6
       for (auto v:graph[u])
           if ( color [ v]==WHITE)
           {
9
                no_cycle=dfs(v);
10
11
           else if(color[v]==GREY)return false;
12
13
       color [u]=BLACK;
14
       return no_cycle;
15
16 }
```

2.2 Topological Sort

```
#include < bits / stdc ++.h>
using namespace std;
#define WHITE 0
#define GREY 1
#define BLACK 2
vector < int > graph [100005];
```

```
7 vector<int> ans;
8 int visit [100005];
9 bool dfs(int u)
10 {
11
       visit[u]=GREY;
12
       bool no_cycle=true;
13
       int sz=graph[u].size();
       for (int i = 0; i < sz; i++)
14
15
            int v=graph[u][i];
16
            if ( visit [v]==WHITE)
17
            {
18
                 no_cycle=dfs(v);
19
20
            }
21
            else if(visit[v]==GREY)return false;
22
       visit[u]=BLACK;
23
24
       ans.push_back(u);
       return no_cycle;
25
26
  bool topsort (int N)
27
28
       ans.clear();
29
       memset(visit, false, sizeof(visit));
30
       int no_cycle=true;
31
       for (int i = 0; i < N; i++)
32
33
       {
            if ( visit [ i]==WHITE) no_cycle&=dfs(i);
34
35
36
       return no_cycle;
37 }
  int main()
38
39
40
       return 0;
41
```

2.3 Strongly Connected Component

```
1 typedef vector<int> vint;
2 vint G[100000+5];
3 \text{ vint } G2[100000+5];
4 vint sorted;
5 vint cycle;
6 bool color [100000+5];
7
  void dfs(int u)
8
9
       color[u] = false;
       for (int i=0; i < G[u]. size(); i++)
10
11
            int v=G[u][i];
12
            if (color[v])
13
            {
14
                 dfs(v);
16
17
18
       sorted.pb(u);
19 }
void dfs2(int u)
21
  {
       cycle.pb(u);
```

```
color [u]=false;
23
          for (int i=0; i<G2[u].size(); i++)
24
25
                 int v=G2[u][i];
26
                if (color [v])
27
                {
                       dfs2(v);
                 }
30
          }
31
32 }
33 void SCC(int N)
34 {
35
          mem(color,true);
36
          for (int i=1; i \le N; i++)
37
                 if (color[i])
38
                       dfs(i);
40
41
42
         mem(color, true);
43
          \begin{array}{ll} \textbf{for} (\, \textbf{int} & \textbf{i} \! = \! \textbf{sorted} \, . \, \textbf{size} \, (\,) \, -1; \  \, \textbf{i} > = \! 0; \  \, \textbf{i} \, --) \end{array}
44
45
                 int u=sorted[i];
46
                 if (color [u])
47
                 {
48
                       cycle.clear();
50
                       dfs2(u);
                       print_vector(cycle);
51
                 }
52
53
54
```

2.4 Havel Hakimi

Given N degree d1,d2,d3.....dn. Is it possible to make a graph which have no cycle and different two node will be connected with one Edge?

$$S = d_1, d_2, d_3, \dots, d_n$$
$$d_i \ge d_{i+1}$$

- 1. If any $d_i \ge n$ then fail
- 2. If there is an odd number of odd degrees then fail
- 3. If there is a $d_i < 0$ then fail
- 4. If all $d_i = 0$ then report success
- 5. Reorder Sinto non -increasing order
- 6. Let $k = d_1$
- 7. Remove d_1 from S.
- 8. Subtract 1 from the first k terms remaining of the new sequence
- 9. Go to step 3 above

```
1 #include < stdio.h>
2 #include <queue>
3 #include < vector >
4 using namespace std;
5 int main()
6 {
7
       int N;
       while (scanf ("%d",&N) and N)
8
9
            priority_queue<int>Q;
            bool Ok=true;
11
            int Odd_Node=0;
            for (int i=0; i < N; i++)
13
14
15
                 int x;
                 scanf("%d",&x);
16
17
                 if(x>=N or x<0)Ok\&=false;
                 Odd_Node += (x\%2);
19
                Q. push(x);
20
            Ok\&=(Odd_Node\%2==0); ///Handshaking Theorem
21
            for (int i=0; i < N and Ok; i++)
22
            {
23
                 int k=Q. top();
24
                Q. pop();
25
                 vector <int> v;
26
                 for (int j=0; j < k and Ok; j++)
                 {
28
                     int x=Q. top();
29
30
                     Q. pop();
                     x--;
                     Ok\&=(x>=0);
32
                     v.push_back(x);
33
34
                 for (int j=0; j < k and Ok; j++)
35
36
                     Q. push (v[j]);
37
39
            if (Ok) printf("Possible\n");
40
            else printf("Not possible\n");
41
42
       return 0;
43
44
```

2.5 Articulation Point/Bridge

Articulation point: A vertex in an undirected connected graph is an articulation point (or cut vertex) iff removing it (and edges through it) disconnects the graph. Articulation points represent vulnerabilities in a connected network single points whose failure would split the network into 2 or more disconnected components. They are useful for designing reliable networks. For a disconnected undirected graph, an articulation point is a vertex removing which increases number of connected components. Following are some example graphs with articulation points encircled with red color.

A O(V+E) algorithm to find all Articulation Points:

The idea is to use DFS (Depth First Search). In DFS, we follow vertices in tree form called DFS tree. In DFS tree, a vertex u is parent of another vertex v, if v is discovered

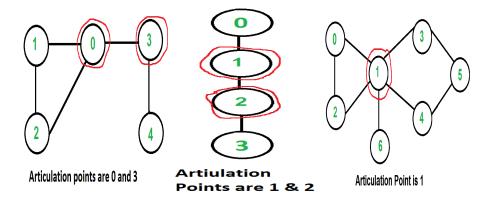


Figure 2.1: Articulation Point

by u (obviously v is an adjacent of u in graph). In DFS tree, a vertex u is articulation point if one of the following two conditions is true.

- 1. u is root of DFS tree and it has at least two children.
- 2. u is not root of DFS tree and it has a child v such that no vertex in subtree rooted with v has a back edge to one of the ancestors (in DFS tree) of u.

u is not root of DFS tree and it has a child v such that no vertex in subtree rooted with v has a back edge to one of the ancestors (in DFS tree) of u.

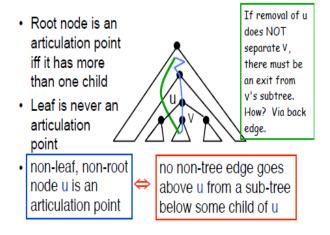


Figure 2.2: Articulation Point

We do DFS traversal of given graph with additional code to find out Articulation Points (APs). In DFS traversal, we maintain a parent[] array where parent[u] stores parent of vertex u. Among the above mentioned two cases, the first case is simple to detect. For every vertex, count children. If currently visited vertex u is root (parent[u] is NIL) and has more than two children, print it. How to handle second case? The second case is trickier. We maintain an array disc[] to store discovery time of vertices. For every node u, we need to find out the earliest visited vertex (the vertex with minimum discovery time) that can be reached from subtree rooted with u. So we maintain an additional array low[] which is defined as follows.

```
low[u] = min(disc[u], disc[w]) where w is an ancestor of u and there is a back edge from some descendant of u to w.
```

2.5.1 Find Articulation Point:

```
vector < int > Graph [10000];
2 bool visit [10000];
3 int arti[100000];
4 int discover [100000], Back [100000];
5 int predfn;
6 int source;
7 int child_of_root;
8 int cnt = 0;
9 void reset()
10 {
       memset(visit, false, sizeof(visit));
11
12
       memset(arti, false, sizeof(arti));
13
       predfn=child_of_root=0;
14
  void articulation (int v)
15
16
       visit[v] = true;
17
       predfn++;
18
       discover[v]=Back[v]=predfn;
19
       for (int i=0;i<Graph[v].size();i++)</pre>
20
21
            int w=Graph[v][i];
22
            if (! visit [w])
23
            {
24
25
                 articulation (w);
26
                 Back[v] = min(Back[v], Back[w]);
27
                 if (Back [w]>=discover [v]&&v!=source)
28
                 {
                      arti[v] = true;
29
                 }
30
                 else if (v==source)
31
32
                      child_of_root++;
33
                      if (child_of_root == 2)
34
                      {
                          arti[v] = true;
36
                 }
38
            }
39
            else
40
            {
41
                 Back[v]=min(Back[v], discover[w]);
42
            }
43
       }
44
45
```

2.5.2 Find Bridge version 1:

```
vector < int > Graph [200];
int Back [205], Discover [205];
bool visit [205];
bool bridge [205][205];
int brount;
void reset (int n)
{
```

```
for (int i=0; i <=n; i++)Graph[i].clear();
8
       memset(visit, false, sizeof(visit));
9
       memset(bridge, false, sizeof(false));
10
       brcount=0;
11
12
  void find_bridge(int u, int parent, int depth)
13
14
       visit[u] = true;
15
       Discover [u] = Back [u] = depth;
16
17
       for (int i=0 ; i<Graph[u].size() ; i++)</pre>
18
19
            int v = Graph[u][i];
20
21
22
            if (visit[v] && v!=parent)
23
            {
                Back[u] = min(Back[u], Discover[v]);
24
25
               (! visit [v])
            i f
26
27
                find_bridge(v, u, depth+1);
28
                Back[u] = min(Back[u], Back[v]);
29
                if (Back[v]>Discover[u])
30
                {
31
                     brcount++;
32
                     bridge[u][v] = bridge[v][u] = true;
33
                }
34
35
            }
36
       }
37
38
```

2.5.3 Find Bridge version 2:

```
void find_bridge(int node, int parent)
2
3
       discovery_time [node] = bedge [node] = ++T;
       int to, i, connected = adj[node].size();
4
       for (i = 0; i < connected; i++)
5
6
           to = adj [node][i];
           if (to = parent) continue;
8
           if (! discovery_time[to])
9
           {
10
                printf("%d %d\n", node, to);
11
               find_bridge(to, node);
12
               bedge [node] = min(bedge [node], bedge [to]);
13
                if (bedge [to] > discovery_time [node]) printf ("%d %d\n", to, node)
14
15
           else if (discovery_time[node] > discovery_time[to])
16
17
                printf("%d %d\n", node, to);
18
                bedge [node] = min(bedge [node], discovery_time [to]);
19
20
21
       }
22 }
```

2.6 Stable Marriage Problem

```
1 const int Max=400;
2 int position [Max][Max];
3 int Husband [Max];
4 queue<int>womenlist [Max];
5 void stable_marage(int N)
6 {
       mem(Husband, -1);
7
       queue<int>freeman;
8
       for (int i=1; i \le N; i++) freeman.push(i);
9
       while (!freeman.empty())
11
            int man=freeman.front();
12
           int woman=womenlist[man].front();
13
           womenlist [man].pop();
14
            if (Husband [woman]==-1)
16
                Husband [woman]=man;
17
18
                freeman.pop();
           }
19
           else
20
            {
                int herhusband=Husband[woman];
22
                if (position [woman] [man] < position [woman] [herhusband])
23
24
                     Husband [woman]=man;
25
                     freeman.pop();
                     freeman.push(herhusband);
27
28
           }
29
       }
30
31
```

2.7 2 SAT

```
1 /*
2 A or B
з A =>В
4 Algo Steps:
_{5} 1. Construct Graph with A \rightarrow\!\! B & B \rightarrow\!\! A Edges
6 2. Find the Strongly Connected component of Graph & ranked them.
7 3. Check A & A is not in the same connected component. If in same group
      then, there is no solution
  4. If there is solution exist, then solution will be higher rank between A &
       A node.
10 #include < bits / stdc++.h>
11 using namespace std;
12 #define pb push_back
13 \#define mem(x,y) memset(x,y,sizeof(x))
14 typedef vector<int> vint;
15 vint Graph1 [5 * 8000];
16 vint Graph2[5*8000];
17 bool visit [5*8000];
18 int color [5*8000];
19 vint ans;
20 struct info
21 {
int u, v;
```

```
23
       info(int u, int v): u(u), v(v) \{\};
24 };
25 vector<info>Edges;
vint topsort;
void init (int N)
28
       Edges.clear();
29
       for (int i=0; i <= 3*N; i++)
30
31
            Graph1[i].clear();
32
            Graph2[i].clear();
33
34
35
36
  void construct_graph(int N)
37
38
       int sz=Edges.size();
       for(int i=0; i < sz; i++)
39
40
            int u=Edges[i].u;
41
            int v=Edges[i].v;
42
            {\color{red} \textbf{int}} \quad a \ , b \ , a comp \ , b comp \ ;
43
            if(u>0)
44
            {
45
                 a=2*u;
46
                 acomp=2*u+1;
47
            }
48
49
            else
50
            {
51
                 a=2*(-u)+1;
52
                 acomp=2*(-u);
53
            if(v>0)
54
            {
55
                 b=2*v;
56
                 bcomp=2*v+1;
57
            }
58
            else
59
60
            {
                 b=2*(-v)+1;
61
                 bcomp=2*(-v);
62
63
            Graph1[acomp].pb(b);
64
            Graph1 [bcomp].pb(a);
65
            Graph2[b].pb(acomp);
66
            Graph2 [a].pb(bcomp);
67
68
69
  void dfs(int u)
70
71
  {
       if ( visit [u]) return;
72
        visit[u] = true;
73
       int sz=Graph1[u].size();
74
       for (int i=0; i < sz; i++)
75
76
77
            int v=Graph1[u][i];
            dfs(v);
78
79
       topsort.pb(u);
80
81 }
82 void dfs2(int u, int Rank)
```

```
if ( visit [u]) return;
84
85
        visit[u] = true;
        color [u]=Rank;
86
        int sz=Graph2[u].size();
87
        for (int i=0; i < sz; i++)
88
89
             int v=Graph2[u][i];
90
             dfs2(v,Rank);
91
92
93
   void SCC(int N)
94
95
        topsort.clear();
96
97
        mem(visit, false);
98
        for (int i=2; i <= 2*N+1; i++)
99
             if(visit[i] == false)dfs(i);
100
101
        int r=1;
102
       mem(visit, false);
103
        for (int i=topsort.size()-1; i>=0; i--)
             if(visit[topsort[i]] == false)dfs2(topsort[i], r++);
106
107
108
   bool is 2 sat possible (int N)
109
110
111
        for(int i=1; i \le N; i++)
112
113
             int a=2*i;
             int acomp=2*i+1;
114
             if (color[a]==color[acomp]) return false;
115
116
        return true;
117
118
   void solution (int N)
119
120
121
        ans.clear();
        for (int i=1; i \le N; i++)
122
123
             int a=2*i;
124
             int acomp=2*i+1;
             if (color[a]>color[acomp]) ans.pb(i);
126
127
128
   int main()
129
130
131
        int test;
        scanf("%d",&test);
132
133
        for (int t=1; t <= test; t++)
134
             int M,N;
135
             scanf("%d %d",&M,&N);
136
             init(N);
137
             for(int i=0; i \triangleleft M; i++)
138
139
             {
140
                  int u, v;
                  scanf("%d %d",&u,&v);
141
                  Edges.pb(info(u, v));
143
             construct_graph(N);
144
```

```
SCC(N);
145
             printf("Case %d: ",t);
146
             if (is2satpossible (N))
147
148
                  cout << "Yes" << endl;
149
                  solution (N);
150
                  int sz=ans.size();
151
                  printf("%d",sz);
                  for (int i=0; i < sz; i++) printf(" %d", ans[i]);
153
                  printf("\n");
154
155
             else cout << "No" << endl;
157
        return 0;
158
159
```

2.8 Heavy Light Decomposition

```
1 #include < bits / stdc++.h>
2 using namespace std;
3 #define pp pair<int,int>
4 #define pb push_back
5 const int Max=10000;
6 struct info
7
  {
       int v, cost;
9
       info(int v=0, int cost=0): v(v), cost(cost)  {};
10 };
vector<pp>edges;
vector < info>Graph [Max + 5];
  int Tree [5*Max+5], BaseArray [Max+5], SubTreeSize [Max+5];
14 int ChainHead [Max+5], ChainNum [Max+5], PosInBaseArray [Max+5], ChainNo;
int Level [Max+5], Parent [Max+5], SparseTable [Max+5][16];
16 int ptr;
17 void init (int N)
18 {
19
       for (int i=0; i \le N; i++)
20
           Graph[i].clear(), ChainHead[i]=-1;
21
           for (int j=0; j \le 15; j++)SparseTable [i][j]=-1;
22
23
       edges.clear();
24
       ptr=ChainNo=0;
25
26
  void buildSegmentTree(int 1, int r, int indx)
27
28
29
       if ( l==r )
30
       {
           Tree [indx]=BaseArray[l];
31
32
           return;
33
       int mid=(1+r)>>1;
34
       int lindx=indx << 1;
35
36
       int rindx=lindx | 1;
37
       buildSegmentTree(l, mid, lindx);
38
       buildSegmentTree(mid+1,r,rindx);
39
       Tree [indx]=max(Tree [lindx], Tree [rindx]);
40 }
  void updateSegmentTree(int l, int r, int indx, int update_indx, int value)
```

```
if ( l==r )
43
        {
44
45
            Tree [indx]=value;
46
            return;
47
        int mid=(1+r)>>1;
48
        int \lim dx = \inf dx <<1;
49
50
        int rindx=lindx | 1;
        if (update_indx <= mid) updateSegmentTree(l, mid, lindx, update_indx, value);</pre>
51
        else updateSegmentTree(mid+1,r,rindx,update_indx,value);
        Tree [indx]=max(Tree [lindx], Tree [rindx]);
53
54
       querySegmentTree(int 1, int r, int indx, int x, int y)
55
   int
56
57
        if (1>y | | r<x) return 0;
58
        if (x<=l&&y>=r) return Tree[indx];
59
        int mid=(1+r)>>1;
60
        int lindx = indx << 1;
        int rindx=lindx | 1;
61
        int c1=0, c2=0;
62
        if (x<=mid) c1=querySegmentTree(l, mid, lindx, x, y);</pre>
63
        if (y>mid) c2=querySegmentTree(mid+1,r,rindx,x,y);
64
        return max(c1,c2);
65
66
   void dfs (int from, int u, int depth)
67
68
   {
        Level [u]=depth;
69
70
        Parent [u]=from;
71
        SubTreeSize[u]=1;
72
        int sz=Graph[u].size();
73
        for (int i=0; i < sz; i++)
74
        {
             int v=Graph[u][i].v;
75
             if (v==from) continue;
76
77
             dfs(u,v,depth+1);
            SubTreeSize[u]+=SubTreeSize[v];
78
79
80
   void sparseTable(int N)
81
82
        for (int i=0; i \le N; i++)SparseTable[i][0] = Parent[i];
83
        for (int j=1; (1 << j) <= N; j++)
84
85
             for(int i=0; i <= N; i++)
86
             {
87
                 if(SparseTable[i][j-1]!=-1)
88
89
                 {
                      int a=SparseTable[i][j-1];
                      SparseTable [i] [j] = SparseTable [a] [j-1];
91
92
            }
93
        }
94
95
   int LCA(int p, int q)
96
97
        if (Level[p] < Level[q]) swap(p,q);
98
        int Log=log 2 (Level[p]) + 1;
99
        for (int i=Log; i>=0; i--)
100
        {
             if ((Level[p]-(1<<i))>=Level[q])p=SparseTable[p][i];
103
```

```
if(p==q)return p;
104
       for (int i=Log; i>=0; i--)
106
            if (SparseTable [p][i]!=-1&&SparseTable [p][i]!=SparseTable [q][i])
107
108
            {
                p=SparseTable[p][i],q=SparseTable[q][i];
109
111
       return Parent[p];
113
114
    * Actual HL-Decomposition part
115
    * Initially all entries of chainHead[] are set to -1.
116
    * So when ever a new chain is started, chain head is correctly assigned.
117
    * As we add a new node to chain, we will note its position in the baseArray
118
    * In the first for loop we find the child node which has maximum sub-tree
       size.
    * The following if condition is failed for leaf nodes.
120
    * When the if condition passes, we expand the chain to special child.
    * In the second for loop we recursively call the function on all normal
      nodes.
    * chainNo++ ensures that we are creating a new chain for each normal child.
   **/
124
   void heavyLightDecompositon(int from, int curNode, int cost)
125
126
   {
       if (ChainHead [ChainNo]==-1)ChainHead [ChainNo]=curNode; /// Assign chain
       head
       ChainNum [curNode]=ChainNo;
128
       PosInBaseArray[curNode]=ptr; /// Position of this node in baseArray
129
       which we will use in Segtree
       BaseArray [ ptr++]=cost;
130
       int sc=-1,nextCost;
131
       int sz=Graph[curNode].size();
       for (int i=0; i < sz; i++) /// Loop to find special child
133
134
            int v=Graph[curNode][i].v;
135
            if (v==from) continue;
136
            if (sc==-1||SubTreeSize[sc]<SubTreeSize[v])
137
138
139
                nextCost=Graph[curNode][i].cost;
140
141
       if (sc!=-1)heavyLightDecompositon(curNode, sc, nextCost); /// Expand the
143
       for (int i=0; i < sz; i++)
144
145
       {
            int v=Graph[curNode][i].v;
146
            int cost=Graph[curNode][i].cost;
147
148
            if(v = from | | sc = v) continue;
           ChainNo++;
149
           heavyLightDecompositon(curNode, v, cost);
151
152
   void updateTree(int ith, int val)
153
154
       pp a=edges[ith];
       int u=a.first ,v=a.second;
       int indx=PosInBaseArray[u];
       if (Level[u] < Level[v]) indx=PosInBaseArray[v];</pre>
```

```
updateSegmentTree(0, ptr-1, 1, indx, val);
159
160 }
161
162
   * query_up:
    * It takes two nodes u and v, condition is that v is an ancestor of u
163
    * We query the chain in which u is present till chain head, then move to
       next chain up
    * We do that way till u and v are in the same chain, we query for that part
165
        of chain and break
    **/
166
   int queryUp(int u,int v)
167
168
       if(u=v)return 0;
169
170
       int uchain, vchain=ChainNum[v], ans=-1;
171
       while (true)
172
       {
            uchain=ChainNum[u];
174
            if (uchain=vchain)
175
            {
                                  /// Both u and v are in the same chain, so we
                if (u==v)
176
       need to query from u to v, update answer and break.
                    break;
                                  /// We break because we came from u up till v,
177
       we are done
                ans=max(ans, querySegmentTree(0, ptr-1, 1, PosInBaseArray[v]+1,
178
       PosInBaseArray[u]));
                break;
179
            }
180
181
            int uchainhead=ChainHead[uchain];
            ans=max(ans, querySegmentTree(0, ptr-1,1, PosInBaseArray[uchainhead],
182
       PosInBaseArray[u]));
            /// Above is call to segment tree query function. We do from
183
       chainHead of u till u. That is the whole chain from
           u=Parent [uchainhead];
184
185
186
       return ans;
187
   int queryTree(int u,int v)
188
189
       int lca=LCA(u,v);
190
       return max(queryUp(u,lca),queryUp(v,lca));
191
192
   int main()
193
194
       int test;
195
       cin>>test;
196
       while (test --)
197
198
            int N;
199
200
            cin >> N;
201
            init(N);
            for (int i=0; i < N-1; i++)
202
            {
203
                int u, v, c;
204
                cin>>u>>v>>c;
205
                u--,v--:
206
                Graph[u].pb(info(v,c));
207
                Graph[v].pb(info(u,c));
208
                edges.pb(pp(u,v));
209
            dfs(-1,0,0);
211
            sparseTable(N);
212
```

```
{\tt heavyLightDecompositon} \, (\, -1\,, 0\,, -1) \, ;
213
             buildSegmentTree(0,ptr-1,1);
214
             string ch;
215
216
             int x, y;
             while (true)
217
             {
218
                  cin>>ch;
                   if(ch[0] == 'D')break;
220
                  \verb|cin>>x>>y|;
221
                   if(ch[0]== Q') printf("%d\n", queryTree(x-1,y-1));
222
                  else if (ch[0] == 'C') updateTree (x-1,y);
223
             }
224
225
         return 0;
226
227 }
```

Flow networks, matching

3.1 Max Flow

```
1 #include < bits / stdc++.h>
2 using namespace std;
3 #define pb push_back
4 #define MN 1000
5 typedef vector < vector <int> > vint2D;
6 const int inf=(1 < < 29);
7 vint2D graph;
8 int Cost [MN] [MN];
9 int parent [MN+5];
10 int flow;
void init (int N)
12 {
       graph=vint2D(N);
13
       memset (Cost, 0, size of (Cost));
14
15 }
  void AddEdge(int u, int v, int cost)
17 {
18
       graph[u].pb(v);
19
       graph[v].pb(u);
       Cost[u][v]+=cost;
20
       Cost[v][u]+=cost;
21
22 }
  bool augmenting_path(int source, int sink)
23
24
       memset (parent, -1, size of (parent));
26
       queue<int>Q;
       Q. push (source);
       while (!Q. empty())
28
20
           int u=Q. front();
30
           Q. pop();
31
           int sz=graph[u].size();
32
            for (int i=0; i < sz; i++)
33
34
                int v=graph[u][i];
35
                if (parent[v]==-1 \text{ and } Cost[u][v]>0)
37
                {
                     parent [v]=u;
38
                     Q. push(v);
39
                     if (v==sink)return true;
40
                }
41
```

```
43
          return false;
44
45 }
void path (int v, int source)
47 {
48
          int u=parent[v];
          flow=min(flow, Cost[u][v]);
49
          if (source!=u) path(u, source);
50
          Cost[u][v]-=flow;
51
         Cost\left[\,v\,\right]\left[\,u\right] += flow\;;
          return;
53
54 }
int max_flow(int source, int sink)
56 {
57
          int ret = 0;
58
          while (augmenting_path (source, sink))
59
         {
                flow=inf;
60
                path(sink, source);
61
                ret+=flow;
62
63
          return ret;
64
   }
65
   int main()
66
67
   {
68
          int test;
69
          scanf("%d",&test);
70
          while (test --)
71
                int P,S,C,M;
72
                scanf ("%d %d %d %d",&P,&S,&C,&M);
73
                init(P+S+5);
74
                int superSource=0,SuperSikn=P+S+1;
75
                \begin{array}{ll} \textbf{for} \; (\; \textbf{int} \quad i = 1; \; \; i < = P; \; \; i + +) \\ \text{AddEdge} (\; \text{superSource} \; , i \; , 1) \; ; \end{array}
76
                \begin{array}{ll} \text{for} \; (\; \text{int} \quad i = 1; \; \; i < = S \; ; \quad i + +) \\ \text{AddEdge} \left(P + 1, \text{SuperSikn} \; , C \right) \; ; \end{array}
77
                for (int i=0; i \triangleleft M; i++)
78
79
                      int x, y;
80
                      scanf("%d %d",&x,&y);
                      AddEdge (x, P+y, (1 < < 30));
82
83
                printf("%d\n", max_flow(superSource, SuperSikn));
84
85
          return 0;
86
87 }
```

3.2 Maximum Bipartite Matching

```
bool khun(int u)
2 {
3
       int sz=graph[u].size();
       for (int i=0; i < sz; i++)
4
            int v=graph[u][i];
6
            if (seen[v]) continue;
            seen[v] = true;
8
            if(Right[v] == -1||khun(Right[v]))
9
10
                 Right[v]=u;
11
                Left[u]=v;
12
                return true;
13
14
15
       return false;
16
17 }
  int bipartite_matching(int m)
18
19
       mem(Left, -1);
20
       mem(Right, -1);
21
       int cnt = 0;
22
       for (int i=0; i \triangleleft m; i++)
23
24
           mem(seen, false);
25
            if (khun(i))cnt++;
26
27
       return cnt;
28
29 }
```

Dynamic programming

4.1 Longest Increasing sub sequence

4.1.1 LIS1

```
1 #include < bits / stdc++.h>
2 using namespace std;
3 const int inf = 2000000000; // a large value as infinity
5 int n; // the number of items in the sequence
6 int Sequence[32]; // the sequence of integers
7 int \ L[32]; \ // \ L[] as described in the algorithm
8 int I[32]; // I[] as described in the algorithm
9 vector <int>LIS;
10 void takeInput()
11 {
      scanf("%d", &n); // how many numbers in the sequence ?
12
      13
          scanf("%d", &Sequence[i]);
14
15 }
16
17 int LisNlogK()
                  // which runs the NlogK LIS algorithm
18
      int i; // auxilary variable for iteration
19
      I[0] = -inf; // I[0] = -infinite
21
      for (i = 1; i \le n; i++) // observe that i \le n are given
22
          I[i] = inf; // I[1 \text{ to } n] = infinite
23
      int LisLength = 0; // keeps the maximum position where a data is
25
      inserted
26
      for (i = 0; i < n; i++) // iterate left to right
27
28
          int low, high, mid; // variables to perform binary search
          low = 0; // minimum position where we to put data in I[]
          high = LisLength; // the maximum position
31
32
          while (low <= high) // binary search to find the correct position
33
34
              mid = (low + high) / 2;
35
              if ( I [mid] < Sequence[i] )</pre>
36
37
                  low = mid + 1;
              else
38
                  high = mid - 1;
```

```
40
            // observe the binary search carefully, when the binary search ends
41
            // low > high and we put our item in I[low]
42
            I [low] = Sequence [i];
43
            L[i]=low;
44
            if (LisLength < low ) // LisLength contains the maximum position
45
                LisLength = low;
46
47
       }
48
49
       i = 0;
50
       for (int j=0; j< n; j++)
51
53
            if (L[j]>L[i])
            {
                 i=j;
56
57
       LIS.push_back(Sequence[i]);
58
       for (int j=i-1; j>=0; j--)
59
60
            if (Sequence [j] < Sequence [i] & L[j] == L[i] - 1)
61
            {
62
63
                 i=i:
                 LIS. push_back (Sequence [j]);
64
65
66
       reverse(LIS.begin(),LIS.end());
67
68
       return LisLength; // return the result
69 }
  int main()
70
71
   {
       takeInput();
72
       int result = LisNlogK();
73
       printf("\,The\ LIS\ length\ is\ \%d\backslash n"\,,\ result);
74
       for(int i=0; i<LIS.size(); i++)printf("%d",LIS[i]);</pre>
       return 0;
76
77 }
```

4.1.2 LIS(Stanford)

```
1 // Given a list of numbers of length n, this routine extracts a
2 // longest increasing subsequence.
3 //
4 // Running time: O(n log n)
5 //
6 //
       INPUT: a vector of integers
       OUTPUT: a vector containing the longest increasing subsequence
9 #include <iostream>
10 #include <vector>
11 #include <algorithm>
12 using namespace std;
13 typedef vector <int> VI;
14 typedef pair <int , int > PII;
15 typedef vector <PII > VPII;
16
17 #define STRICTLY_INCREASNG
19 VI LongestIncreasingSubsequence(VI v)
```

```
VPII best;
21
22
       VI dad(v.size(), -1);
23
       for (int i = 0; i < v.size(); i++)
24
25
       {
  #ifdef STRICTLY_INCREASING
27
           PII item = make_pair(v[i], 0);
           VPII:: iterator it = lower_bound(best.begin(), best.end(), item);
28
           item.second = i;
29
30 #else
           PII item = make_pair(v[i], i);
31
           VPII::iterator it = upper_bound(best.begin(), best.end(), item);
32
  #endif
33
           if (it == best.end())
34
35
           {
36
               dad[i] = (best.size() == 0 ? -1 : best.back().second);
37
                best.push_back(item);
           }
38
           else
39
           {
40
               dad[i] = dad[it->second];
41
               *it = item;
42
           }
43
       }
44
45
       VI ret;
46
       for (int i = best.back().second; i >= 0; i = dad[i])
47
48
           ret.push_back(v[i]);
49
       reverse (ret.begin(), ret.end());
50
       return ret;
51
```

4.2 Edit Distance

```
1 #include < bits / stdc++.h>
2 using namespace std;
3 int dp[88][88];
4 int N,M, step;
5 char S1[88], S2[88];
6 int solve(int i, int j)
7
8
       if (i = N \text{ and } j = M) \text{ return } 0;
       if (i=N) return M-j;
9
       if (j=M) return N-i;
10
       int &ret=dp[i][j];
11
       if (ret!=-1)return ret;
12
       ret = (1 < <28);
13
       if(S1[i]==S2[j])ret=solve(i+1,j+1);
14
       else
       {
16
            ret=min(ret, solve(i, j+1)+1);
17
            ret=min(ret, solve(i+1,j)+1);
18
            ret = min(ret, solve(i+1, j+1)+1);
19
20
       }
21
       return ret;
22 }
  void pathPrint(int i, int j, int del, int ins, int st)
24
  {
       if(i=N&kj=M) return ;
25
       if (i==N)
```

```
27
             for (int k=j; k < M; k++, i++)
28
29
                  printf("%d Insert %d,%c\n",st++,i-del+1+ins,S2[k]);
30
31
32
             return ;
33
        if ( j==M)
34
35
             for (; i \lt N; i++)
36
             {
37
                  printf("%d Delete %d\n", st++,i-del+1+ins);
38
                  del++;
39
40
41
             return ;
42
43
        int ret = solve(i,j);
        int tmp;
44
        if (S1[i]==S2[j])
45
46
            tmp = solve(i+1,j+1);
47
             if(ret = tmp)
48
             {
49
                  pathPrint(i+1,j+1,del,ins,st);
50
                  return ;
51
52
             }
53
54
       tmp = solve(i, j+1)+1;
55
        if (tmp=ret)
56
             printf("\%d Insert \%d,\%c\n",st,i-del+1+ins,S2[j]);
57
             pathPrint(i, j+1, del, ins+1, st+1);
58
             return ;
59
60
       tmp = solve(i+1,j)+1;
61
        if (tmp=ret)
62
63
             printf("%d Delete %d\n",st,i-del+1+ins);
64
             pathPrint(i+1,j,del+1,ins,st+1);
65
             return ;
66
67
       tmp = solve(i+1, j+1)+1;
68
        if (tmp==ret)
69
        {
70
             printf("%d Replace %d,%c\n", st, i-del+1+ins, S2[j]);
71
             pathPrint(i+1,j+1,del,ins,st+1);
72
73
             return ;
74
75
        return ;
76
77 int main()
78
        bool New=false;
79
        while (gets (S1))
80
81
             gets(S2);
82
             if (New) printf("\n");
83
            New=true;
84
            N=strlen(S1);
85
            M≡strlen (S2);
86
            \operatorname{memset}(\operatorname{dp}, -1, \operatorname{sizeof}(\operatorname{dp}));
87
```

```
step=solve(0,0);
printf("%d\n", step);
pathPrint(0,0,0,0,1);
}
return 0;
}
```

4.3 Distinct Sub Sequence

```
1 LLD solve()
2
  {
3
       dp[0] = 1;
4
      mem(Last_occurance,0);
       for (int i=1;s[i];i++)
5
6
           dp[i]=2*dp[i-1];
           int x=s[i]-A';
8
           if (Last\_occurance[x]) dp[i] = dp[Last\_occurance[x] - 1];
9
10
           Last_occurance [x]=i;
           dp[i]=dp[i];
11
12
13
       return dp[strlen(s+1)];
14
```

4.4 Knuth Optimization

```
1 #include < bits / stdc++.h>
2 using namespace std;
3 int dp[505][505];
4 int positionK [505][505];
5 int Csum[505];
  int main()
6
7
       int N;
8
       while (scanf("%d",&N)==1)
9
10
            for(int i=1; i \le N; i++)
11
            {
12
                 scanf("%d",&Csum[i]);
13
14
                 Csum [i] += Csum [i-1];
15
16
            for (int L=1; L \leftarrow N; L++)
17
                 for(int i=1; i \le N; i++)
18
19
                      int j=L+i-1;
20
                      if (j>N) continue;
21
                      if (L==1)
22
                      {
23
                          dp[i][j]=0;
24
                           positionK[i][j]=i;
25
                      }
26
                      else
2.7
                      {
2.8
                          dp[i][j]=(1<<28);
29
                           for (int k=positionK[i][j-1]; k \le positionK[i+1][j]; k++)
30
                           {
31
                               int cost=dp[i][k-1];
32
33
                               cost += dp[k+1][j];
                               cost += (Csum[j] - Csum[k]);
```

```
cost+=(Csum[k-1]-Csum[i-1]);
35
                                 if ( cost < dp[i][j])</pre>
36
37
                                      dp[i][j] = cost;
38
                                      positionK[i][j]=k;
39
                                 }
40
                            }
41
                       }
42
                  }
43
44
             printf("%d\n",dp[1][N]);
45
46
47
        return 0;
48
```

4.5 Max Sum

4.5.1 2D Max Sum 1

Complexity $O(n^4)$

```
precalculation:
for(int i=1;i<=N;i++)

{
    for(int j=1;j<=N;j++)
    {
        Mat[i][j]=Mat[i-1][j]+Mat[i][j-1]-Mat[i-1][j-1]+A[i][j];
    }
}
Query:
Sum=Mat[x2][y2]-Mat[x1-1][y2]-Mat[x2][y1-1]+Mat[x1-1][y1-1]</pre>
```

4.5.2 2D Max sum 2

Complexity n^3

```
int arr[SZ][SZ];
1 int input(void);
3 int max_sum(int n);
4 int main()
5 {
6
       int test, n;
       scanf("%d", \&test);
7
       while (test --)
8
9
           n = input();
            printf("%d\n", max_sum(n));
11
            if (test)
12
                printf("\n");
13
14
       return 0;
15
16
17
  int input (void)
18
19
       int i, len, r = 1, n;
20
       char str [26];
21
       scanf(" %s", &str);
22
       len = strlen(str);
23
       for (i = 1; i \le len; i++)
24
```

```
n = str[i - 1] - '0';
26
27
            if(n)
                arr[r][i] = n;
28
            else
29
                arr[r][i] = -700;
30
31
32
       r++;
       while (r <= len)
33
34
            scanf("%s", &str);
35
            for (i = 1; i \le len; i++)
36
37
           n = str[i - 1] - '0';
38
            if(n)
39
40
                arr[r][i] = n;
41
            else
                arr[r][i] = -700;
43
           r++;
44
45
       return len;
46
  }
47
48
  int max_sum(int n)
49
50
  {
       int i, j, c1, c2, r, sum = 0,m = 0;
51
52
       for (i = 0; i \le n; i++)
53
            arr[i][0] = arr[0][i] = 0;
54
       for(i = 1; i \le n; i++)
55
            for (j = 1; j \le n; j++)
                arr[i][j] += arr[i][j - 1];
56
       for(c1 = 1; c1 \le n; c1++)
57
            for(c2 = c1; c2 \le n; c2++)
58
            {
59
                sum = 0;
60
                for (r = 1; r \le n; r++)
61
62
63
                     sum += arr[r][c2] - arr[r][c1 - 1];
64
                     if (sum < 0)
                         sum = 0;
65
                     else if (sum > m)
66
                         m = sum;
67
                }
68
            }
69
70
       return m;
71 }
```

4.5.3 3D Max sum

```
Complexity O(n^4)
```

Cumulative Sum formula: S(x,y,z) = value(x,y,z) + S(x-1,y,z) + S(x,y-1,z) + S(x,y,z-1) - S(x-1,y-1,z) - S(x,y-1,z-1) - S(x-1,y,z-1) + S(x-1,y-1,z-1)

Query Formula: S([x1,y1,z1] to [x2,y2,z2]) = S(x2,y2,z2) + S(x1,y1,z2) + S(x1,y2,z1) + S(x1,y1,z1) - S(x2,y2,z1) - S(x2,y1,z2) - S(x1,y2,z2) - S(x1,y1,z1)

```
1 for (int i = n; i <= N; i++)
```

```
2 {
3
       for (int j = n; j \ll N; j++)
4
5
6
            for (int k = n; k \le N; k++)
            {
                int x2=i, y2=j, z2=k;
                int x1=i-n, y1=j-n, z1=k-n;
9
                int temp=[Query formul]
                          ans=max(ans,temp);
11
           }
12
       }
13
14 }
```

4.6 Minimum Vertex Cover

```
1 #define MAXN 100002
_{2} int dp [MAXN] [5];
3 int par [MAXN];
vector<int>edges[MAXN];
5 int f(int u, int isGuard)
6
       if (edges[u].size()==0)return 0;
       if(dp[u][isGuard]!=-1) return dp[u][isGuard];
8
       int sum=0;
9
       for(int i=0; i<(int)edges[u].size(); i++)
10
11
           int v=edges[u][i];
12
           if(v!=par[u])
13
           {
14
                par[v]=u;
15
                if(isGuard==0) sum+=f(v,1);
16
17
                else sum+=\min(f(v,1),f(v,0));
18
19
       return dp[u][isGuard]=sum+isGuard;
20
21 }
  int main()
22
23
       memset(dp, -1, sizeof(dp));
24
25
       int n;
       scanf("%d",&n);
26
       for (int i=1; i < n; i++)
27
28
29
           int u, v;
           scanf("%d%d",&u,&v);
30
           edges [u].push_back(v);
31
           edges[v].push_back(u);
32
33
       int ans=0;
34
       ans=\min(f(1,1), f(1,0));
35
       printf("%d\n", ans);
36
       return 0;
37
38 }
```

4.7 Nth Permutation Of a String

```
1 #include < bits / stdc++.h>
2 using namespace std;
\# define mem(x,y) memset(x,y,sizeof(x))
4 typedef long long LLD;
5 LLD FAC[22];
6 int Freq [27];
7 string S;
8 void FACTORIAL()
9 {
       FAC[0]=1;
       for (int i=1; i \le 20; i++)FAC[i]=FAC[i-1]*i;
11
12 }
13 LLD Permutation(int n)
14
       LLD per=FAC[n];
15
       for (int i=0; i<26; i++)per/=FAC[Freq[i]];
16
17
       return per;
18
  string NthPermutation(int nth)
19
20
       string ans="";
21
       int len=S.size();
22
       while (len)
23
24
           LLD upto=0;
25
26
            for (int i=0; i<26; i++)
27
            {
                if (Freq[i])
28
                {
29
                     Freq[i]-=1;
30
                     LLD now=Permutation(len -1);
31
                     if(upto+now>=nth)
32
33
                          nth-=upto;
34
                          ans+=('a'+i);
35
36
                          len --;
37
                          break;
                     }
38
                     else
39
                     {
40
                          upto+=now;
41
                          Freq[i]+=1;
42
43
                }
44
45
       }
46
47
       return ans;
48 }
49
  int main()
50
       FACTORIAL();
51
       int test;
       scanf("%d",&test);
53
       for (int t=1; t <= test; t++)
54
55
56
           mem(Freq, 0);
            char inp [100];
           LLD Nth;
```

```
scanf("%s %11d", inp,&Nth);
59
60
                  S=inp;
                   \label{eq:formula} \begin{array}{ll} \mbox{for} \, (\, \mbox{int} \  \, i = \! 0; \, \, \mbox{inp} \, [\, i \, ]\,; \, \, \, i + \! +) \mbox{Freq} \, [\, \mbox{inp} \, [\, i \, ] - \, \mbox{`a\,'}] + +; \end{array}
61
                  string Ans="Impossible";
62
                  if (Permutation(S.size())>=Nth)
63
                  {
64
                          Ans=NthPermutation(Nth);
65
                  }
66
                  printf("Case %d: %s\n",t,Ans.c_str());
67
68
69
           return 0;
70 }
```

Strings

5.1 KMP

Tutorial

```
1 #include < bits / stdc++.h>
using namespace std;
3 char TXT[10000000], ptr [10000000];
4 vector<int> compute_prefix(const char *p)
5
6
       int m=strlen(p+1);
       vector < int > prefix (m+1);
        prefix[1]=0;
       int k=0;
       for (int i=2; i \leqslant m; i++)
11
            while (k>0 \text{ and } p[k+1]!=p[i]) k=prefix[k];
            if(p[k+1]==p[i])k=k+1;
13
            prefix[i]=k;
14
15
       return prefix;
16
17 }
  vector <int > KMP_match(const char *txt, const char *ptrn)
18
19
20
       int n=strlen(txt+1);
21
       int m=strlen(ptrn+1);
22
       vector<int> Prefix=compute_prefix(ptrn);
23
       vector<int>Match_position;
24
       int q=0;
       for(int i=1; i \le n; i++)
25
26
            while (q>0 \text{ and } ptrn[q+1]!=txt[i]) q=Prefix[q];
27
            if(ptrn[q+1]==txt[i])q=q+1;
28
            i f ( q==m)
            {
30
                 Match_position.push_back(i-m);
32
                 q=Prefix[q];
33
34
       return Match_position;
35
36
  int main()
37
38
  {
       \operatorname{scanf}(\text{"%s \%s"},\operatorname{TXT+1},\operatorname{ptr+1});
39
       vector <int > Match_position=KMP_match(TXT, ptr);
40
       for (int i=0; i<Match_position.size(); i++)
```

5.2 Aho Corasick

Algorithm: We can divide the algorithm into 2 steps:

- 1. constructing a finite state automaton
 - constructing a Trie with input words
 - finding the right fall function for each node of trie
- 2. moving through automaton reading each character from input, storing the results. If an algorithm finds right words in some step, (more words can end in the same place), the nodes light up and the numbers by their indices are incremented.

5.2.1 Aho Corasick with Dynamic Trie

```
1 #include < bits / stdc++.h>
2 using namespace std;
з #define Max 26
4 int getID (char c)
5 {
       return c>='a'?c-'a':c-'A';
6
7 }
8 char inp[1000005];
9 char text[1000005];
10 int ans [5000];
11 map<string , int>Map;
  vector < int > v;
  struct Trie
13
14
       Trie *next[26], * fail;
       int stringMap;
16
       Trie()
17
18
       {
           stringMap=0;
19
            for (int i=0; i<Max; i++)next[i]=NULL;
20
            fail=NULL;
21
22
  };
23
  Trie *root;
  void Insert (const char *str, int M)
25
26
       Trie *p=root;
27
       for (int i=0; str[i]; i++)
28
29
            int id=getID(str[i]);
30
            if (p->next[id]==NULL)p->next[id]=new Trie();
31
32
           p=p->next[id];
33
34
       p->stringMap=M;
35
```

```
36 void computeFailure()
37
       Trie *u, * prefix;
38
       queue<Trie*>Q;
39
       Q. push (root);
40
41
       while (!Q. empty())
42
            u=Q.front(); ///Take a new node
43
44
            Q. pop();
            \quad \quad \text{for} \; (\; \text{int} \quad i = 0; \quad i < \text{Max} \; ; \quad i + +)
45
46
                 if(u->next[i]!=NULL) ///select fail position of ith node of
47
       parent u
48
                      prefix=u->fail; /// Going to u node fail position/ prefix
49
       position
                      while ( prefix!=NULL)
50
                           if ( prefix -> next[i]!=NULL) ///if match found
52
53
                           {
                                u->next[i]->fail=prefix->next[i];
54
                                break;
55
56
                           prefix=prefix->fail; /// match not found, going to upper
57
        child prefix position
58
59
                      if (prefix=NULL)u->next[i]->fail=root;
60
                      Q. push(u\rightarrow next[i]);
61
            }
62
63
64 }
  void AhoCorasick(const char *str)
65
66
       Trie *p=root;
67
       for (int i = 0; str[i]; i++)
68
69
            int id=getID(str[i]);
70
            while (p->next [id]==NULL&&p!=root)p=p->fail;
71
            if (p->next[id]!=NULL)p=p->next[id];
72
            Trie *tp=p;
73
            while (tp!=root)
74
75
                 if (tp->stringMap >0) ans [tp->stringMap]++;
76
                 tp=tp->fail;
77
            }
78
       }
79
80
  }
  void Delete (Trie *u)
81
82 {
       if (u=NULL) return;
83
       for (int i = 0; i < Max; i++) Delete (u->next[i]);
84
       delete u;
85
86 }
87
  int main()
88
89
90
       int test;
       scanf("%d",&test);
91
       for (int t=1; t <= test; t++)
92
93
```

```
Map. clear();
94
95
                 v.clear();
                 memset(ans,0,sizeof(ans));
96
                 root=new Trie();
97
                  int N;
98
                 \operatorname{scanf}("%d",&N);
99
                 scanf("%s", text);
100
101
                 int cnt=1;
                 for (int i=0; i< N; i++)
102
                        scanf("%s",inp);
104
                        \begin{array}{l} \textbf{if} \ (\mathrm{Map.\,find} \ (\mathrm{\,inp}\,) \\ \end{array} \\ = \hspace{-0.5cm} \mathrm{Map.\,end} \ () \ ) \\ \mathrm{Map} \ [\, \mathrm{inp}\,] \\ = \hspace{-0.5cm} \mathrm{cnt} \\ + +; \end{array}
                        Insert(inp,Map[inp]);
106
                        v.push_back(Map[inp]);
107
108
109
                 computeFailure();
                 AhoCorasick(text);
                  printf("Case %d:\n",t);
111
                  for (int i=0; i < N; i++)
113
                        printf("%d\n", ans[v[i]]);
114
115
                 Delete (root);
117
           return 0;
118
119
```

5.2.2 Aho Corasick with Static Trie

```
1 #include < bits / stdc++.h>
2 using namespace std;
3 #define root 0
4 #define NuLL −1
_{5} #define Max 248878
6 #define MC 26
7 int ans[10000];
8 char text[1000005];
9 char inp[100000];
10 map<string , int >Map;
  vector < int > v;
11
  int getID(const char c)
12
13
       return c>='a'?c-'a':c-'A';
14
15 }
16 struct Trie
17 {
       struct node
18
19
       {
20
            int Next[26], fail;
21
            int stringMap;
            void clear()
22
23
                memset(Next, -1, sizeof(Next));
24
                fail=-1;
25
                stringMap = 0;
26
27
       } T[Max];
28
       int ptr;
29
       void clear()
31
       {
32
            ptr=1;
```

```
T[0].clear();
33
34
35
       void Insert(char *str, int M)
36
            int p=0;
37
            for (int i=0; str[i]; i++)
            {
                int id=getID(str[i]);
40
                if(T[p].Next[id]==-1)
41
42
                     T[p].Next[id]=ptr;
43
                     T[ptr++].clear();
44
45
46
                int q=p;
47
                p=T[p]. Next[id];
48
49
            T[p].stringMap=M;
50
       void ComputeFailure()
51
52
            queue<int>Q;
53
            Q. push (root);
54
            int u, prefix;
            while (!Q.empty())
56
57
            {
                u=Q.front();
58
                Q. pop();
60
                for (int i=0; i < MC; i++)
61
                     if(T[u].Next[i]!=NuLL)
62
                     {
63
                          int now=T[u].Next[i];
64
                          prefix=T[u].fail;
65
                          while (prefix!=NuLL)
66
67
                               if (T[prefix]. Next[i]!=NuLL)
68
69
                                   T[now]. fail=T[prefix]. Next[i];
                                   break;
72
                               prefix=T[prefix].fail;
73
                          if ( prefix=NuLL)T[now ]. fail=root;
75
                          Q. push (now);
76
                     }
77
                }
78
            }
79
       }
80
81 };
  void AhoCorasick(const Trie &A, const char *str)
83
       int p=root;
84
       for (int i=0; str[i]; i++)
85
86
            int id=getID(str[i]);
87
            while (A.T[p].Next[id] == NuLL \&p! = root) p = A.T[p]. fail;
88
            if (p!=NuLL&&A.T[p]. Next[id]!=NuLL)p=A.T[p]. Next[id];
89
            int tp=p;
90
91
            while (tp!=root)
92
            {
                 if (A.T[tp].stringMap>0)ans [A.T[tp].stringMap]++;
93
```

```
tp=A.T[tp].fail;
94
             }
95
96
97
   Trie A;
98
99
  int main()
100
   #ifdef _ANICK_
101
        freopen("input.txt","r",stdin);
   #endif // _ANICK_
        int test;
104
        scanf("%d",&test);
        for(int t=1; t \le test; t++)
106
107
108
            Map. clear();
109
             v.clear();
             memset (ans, 0, size of (ans));
111
            A. clear ();
             int N;
             scanf("%d",&N);
113
             scanf("%s", text);
114
             int cnt=1;
115
             for (int i=0; i< N; i++)
117
             {
                  scanf("%s", inp);
118
                  if (Map. find (inp) = Map. end ()) Map [inp] = cnt ++;
119
                 A. Insert (inp, Map[inp]);
121
                 v.push_back(Map[inp]);
122
            A. ComputeFailure();
123
             A hoCorasick\left(A, text\right);
             printf("Case \%d: \n",t);
             for(int i=0; i<N; i++)
126
127
                  printf("%d\n", ans[v[i]]);
128
129
130
131
        return 0;
132
```

5.3 Manacher's Algorithm

```
1 #include < bits / stdc++.h>
2 using namespace std;
3 string s, t;
4 char str [1000005];
5
  void prepare_string()
6
       int i;
7
       t="\hat{\ }\#";
8
       for (i=0; i < s. size(); i++)
9
            t+=s[i], t+="#";
10
       t+="$";
11
12 }
13
14
  int manacher()
15 {
16
       prepare_string();
       int P[t.size()], c=0, r=0, i, i\_mirror, n=t.size()-1;
17
       for (i=1; i < n; i++)
```

```
19
            i_m irror = (2*c)-i;
20
           P[i]=r>i?min(r-i,P[i_mirror]):0;
21
            while (t[i+1+P[i]]==t[i-1-P[i]]) P[i]++;
22
            if(i+P[i]>r)
23
24
            {
                c=i;
25
                r=i+P[i];
26
            }
27
28
       return *max_element(P+1,P+n);
29
30 }
31 int main()
32 {
33
       int kase=1;
       while (scanf (" %s", str)&&str[0]!= 'E')
34
35
36
            s=str;
            printf("Case %d: %d\n", kase++, manacher());
37
38
       return 0;
39
40
```

5.4 Suffix Array

5.4.1 Suffix Array $Nlog^2N$ complexity:

```
1 #include < bits / stdc++.h>
2 using namespace std;
3 const int MAX = 400004;
4 struct info
5 {
       int tup[2], indx; //tup[0] = prev rank, tup[1] = new rank
6
  } arr [MAX];
  int sorted[18][MAX], LCP[MAX], step;
10 string text;
11
  bool comp(const info &a, const info &b)
       return a.tup[0] != b.tup[0]? a.tup[0] < b.tup[0] : a.tup[1] < b.tup[1];
14
15 }
16
  void build_suffix_array(void)
17
18
       int i, n = text.size(), jump;
19
       for (i = 0; i < n; i++)
20
21
           sorted [0][i] = text[i] - 'A'; ///rank suffixes according to 1st char
22
           memset(arr[i].tup, 0, sizeof(arr[i].tup));
23
       for (step = 1, jump = 1; jump \leq n; step++, jump \leq 1)
25
26
           for(i = 0; i \le n; i++)
2.7
2.8
               arr[i].indx = i;
29
               arr[i].tup[0] = sorted[step - 1][i]; /// what i was in prev step
30
               arr[i].tup[1] = i + jump < n? sorted[step - 1][i + jump] : -1;
31
32
           sort(arr, arr + n, comp);
33
           sorted[step][arr[0].indx] = 0;
```

```
for (i = 1; i < n; i++)
35
                 sorted[step][arr[i].indx] = arr[i].tup[0] = arr[i-1].tup[0]
36
      &&
                 arr[i].tup[1] = arr[i-1].tup[1]? sorted[step][arr[i-1].
37
       indx] : i;
38
       }
       cout << "Suffix Array : \n\n";</pre>
39
       for (i = 0; i < n; i++)
40
            cout<<arr[i].indx<<' '<<text.substr(arr[i].indx)<<endl;</pre>
41
42
43
44
  void build_LCP_array(void)
45
46
47
       LCP[0] = 0;
48
       int n = text.size(), i, j, id1, id2;
       for (i = 1; i < n; i++)
49
50
            id1 = arr[i - 1].indx;
51
            id2 = arr[i].indx;
            LCP[i] = 0;
53
            \label{eq:cout} $$ // \cout << "id" << id1 << ' ' << id2 << endl; \\ // \cout << sorted[0][id1] << ' ' << sorted[0][id2] << endl; \\ $$
54
55
            for (j = step - 1; j >= 0; j--)
56
                 if(sorted[j][id1] = sorted[j][id2] && sorted[j][id2])
57
                 {
58
59
                     LCP[i] += (1 << j);
60
                     id1 += (1 << j);
61
                     id2 += (1 << j);
62
            cout << arr[i - 1].indx << '' ' << arr[i].indx << '' ' << LCP[i] <<
63
       endl;
64
       for(i = 0; i < n; i++)
cout << i << ' ' << LCP[i] << endl;
65
66
67
68
69
  int main()
70
       text = "helloworld";
71
       build_suffix_array();
72
       build_LCP_array();
73
       return 0;
74
75 }
```

5.4.2 Suffix Array *NlogN* complexity:

```
string text;
2 char str [10001];
_3 int revSA [MAX], SA [MAX];
4 int cnt [MAX] , nxt [MAX];
5 bool bh [MAX], b2h [MAX];
6 int lcp [MAX];
7 bool cmp(int i,int j)
8 {
        return text[i]<text[j];</pre>
9
10 }
11
12
13
void sortFirstChar(int n)
15
        /// sort for the first char
16
        for(int i = 0 ; i < n ; i++)
17
            SA[i] = i;
18
        sort (SA, SA+n ,cmp);
19
20
        ///indentify the bucket ......
21
22
        for(int i=0 ; i< n ; i++)
23
             bh[i] = (i==0 | | text[SA[i]]! = text[SA[i-1]]);
24
            b2h[i] = false;
25
26
        return;
27
  }
28
29
  int CountBucket(int n)
30
31
   {
32
        int bucket = 0;
33
        for (int i = 0, j; i < n; i = j)
34
        {
35
            j = i+1;
             while (j < n \&\& bh[j] == false) j++;
36
            \mathrm{nxt}\,[\;i\;]\;=\;j\;;
37
            bucket++;
38
39
40
        return bucket;
41 }
42
   void SetRank(int n)
43
44
        for(int i = 0 ; i < n ; i = nxt[i])
45
46
             cnt[i] = 0;
47
             for (int j =i ; j<nxt[i] ; j++)</pre>
48
49
                 \mathrm{revSA}\left[\mathrm{SA}\left[\:j\:\right]\:\right] \;=\; i\;;
50
             }
51
       }
52
53
        return;
54 }
55
void findNewRank(int l, int r, int step)
57 {
       for (int j = l ; j < r ; j++)
58
59
```

```
int pre = SA[j] - step;
60
61
            if(pre>=0)
62
                 int head = revSA[pre];
63
                 revSA[pre] = head+cnt[head]++;
64
65
                 b2h[revSA[pre]] = true;
66
67
       return;
68
69
70
   void findNewBucket(int 1, int r, int step)
71
72
       for(int j = l ; j < r ; j++)
73
74
75
            int pre = SA[j] - step;
            if (pre>=0 && b2h [revSA [pre]])
76
77
                 for (int k = \text{revSA}[\text{pre}]+1; b2h[k] && !bh[k]; k++) b2h[k] =
78
       false;
79
       }
80
       return;
81
82
83
   void buildSA(int n)
84
85
   {
86
       ///start sorting in logn step ...
87
       sortFirstChar(n);
       for(int h = 1 ; h < n ; h < <=1)
88
89
            if(CountBucket(n)=n) break;
90
            SetRank(n);
                          /// cause n-h suffix must be sorted
91
            b2h[revSA[n-h]] = true;
92
            cnt[revSA[n-h]]++;
93
            for (int i = 0 ; i < n ; i = nxt[i])
94
95
96
                 findNewRank(i,nxt[i], h);
                 findNewBucket(i , nxt[i] , h);
97
98
            ///set the new sorted suffix array ...
99
            for(int i = 0 ; i < n ; i++)
100
101
                 SA[revSA[i]] = i;
                 bh[i] = b2h[i]; //new bucket ....
103
104
105
       return;
106
107
108
109
   void buildLCP(int n)
110
111
       int len = 0;
112
       for (int i = 0 ; i < n ; i++)
113
            revSA[SA[i]] = i;
114
       for (int i = 0 ; i < n ; i++)
116
            int k = revSA[i];
117
            if(k==0)
118
119
```

```
lcp[k] = 0;
120
                  continue;
             int j = SA[k-1];
123
             while (\text{text} [i+\text{len}] = \text{text} [j+\text{len}]) len++;
124
             lcp[k] = len;
125
             if (len) len --;
126
127
        return;
128
129
130
131 void printSA()
132
133
        for (int i=0; i < text. size (); i++) printf ("%d %d %s %d\n", i, SA[i], text.
        substr(SA[i]).c_str(), lcp[i]);
        puts("");
134
        for (int i=1;i<text.size();i++) printf("%d ",lcp[i]);</pre>
135
        puts("");
136
137
        return ;
138
   int main()
139
140
        int test;
141
        scanf("%d", \&test);
142
        while (test --)
143
144
        {
             scanf(" %s", str);
146
             text = str;
147
             text += text;
148
             buildSA (text.size());
             buildLCP(text.size());
149
             printSA();
151
        return 0;
152
153 }
```

5.4.3 Longest common sub sequence:

```
1 #include < bits / stdc++.h>
2 using namespace std;
3 #define pp pair<int, int>
4 const int MAX = 3*5005;
5 struct info
6 {
       int tup[2], indx; //tup[0] = prev rank, tup[1] = new rank
8
  } arr [MAX];
int sorted [18] [MAX], LCP [MAX], step;
11 string text;
12
13 bool comp(const info &a, const info &b)
14 {
       return a.tup[0] \stackrel{!}{=} b.tup[0]? a.tup[0] < b.tup[0] : a.tup[1] < b.tup[1];
15
16 }
17
18 void build_suffix_array(void)
19 {
       int i, n = text.size(), jump;
20
21
       for (i = 0; i < n; i++)
22
       {
           sorted [0][i] = text[i] - 'A'; ///rank suffixes according to 1st char
```

```
memset(arr[i].tup, 0, sizeof(arr[i].tup));
24
25
26
       for (step = 1, jump = 1; jump \leq n; step++, jump \leq 1)
27
           for (i = 0; i \le n; i++)
28
           {
29
                arr[i].indx = i;
30
                arr[i].tup[0] = sorted[step - 1][i]; /// what i was in prev step
31
                arr[i].tup[1] = i + jump < n? sorted[step - 1][i + jump] : -1;
32
33
           sort(arr, arr + n, comp);
34
           sorted[step][arr[0].indx] = 0;
35
           for (i = 1; i < n; i++)
36
37
                sorted[step][arr[i].indx] = arr[i].tup[0] = arr[i - 1].tup[0]
      &&
38
                arr[i].tup[1] = arr[i-1].tup[1]? sorted[step][arr[i-1].
      indx] : i;
39
         cout << "Suffix Array : \n\n";</pre>
40
         for (i = 0; i < n; i++)
41
             cout<<arr[i].indx<<' '<<text.substr(arr[i].indx)<<endl;</pre>
42
43
44
45
  void build_LCP_array(void)
46
47
      LCP[0] = 0;
48
49
       int n = text.size(), i, j, id1, id2;
50
       for (i = 1; i < n; i++)
51
           id1 = arr[i - 1].indx;
52
           id2 = arr[i].indx;
53
           LCP[i] = 0;
54
           //cout << "id" << id1 << ' ' << id2 << endl;
           //cout << sorted[0][id1] << ' ' << sorted[0][id2] << endl;
56
           for (j = step - 1; j >= 0; j--)
57
                if(sorted[j][id1] = sorted[j][id2] && sorted[j][id2])
58
59
                    LCP[i] += (1 << j);
60
                    id1 += (1 << j);
61
                    id2 += (1 << j);
62
63
             cout << arr[i - 1].indx << ' ' << arr[i].indx << ' ' << LCP[i] <<
64
      endl;
       }
65
         for (i = 0; i < n; i++)
66
             cout << i << ' ' ' << LCP[i] << endl;
67
68 }
69 pp Index [3];
70 int Position (int indx)
71 {
       for (int i=0; i<3; i++)
72
73
           if (indx>=Index[i]. first and indx<Index[i]. second)return i;</pre>
74
75
76
       return -1;
77
  bool Ok(int val)
78
79 {
       for (int i=1;i<text.size();i++)</pre>
80
81
```

```
if(LCP[i]>=val)
82
83
                 bool a,b,c;
84
                 a=b=c=false;
85
                 int indx=Position(arr[i].indx);
86
                 if (indx = -1)continue;
                 indx==0?a=true:indx==1?b=true:indx==2?c=true:c=false;
                 indx=Position(arr[i-1].indx);
89
                 indx==0?a=true:indx==1?b=true:indx==2?c=true:c=false;
90
                 int r=i+1;
91
                 while (r<text.size() and LCP[r]>=val and !(a and b and c))
92
93
                     indx=Position(arr[r].indx);
94
                     indx==0?a=true:indx==1?b=true:indx==2?c=true:c=false;
95
96
                     r++;
97
                 if (a and b and c) return true;
99
100
       return false;
101
102
   int main()
103
104
       int test;
       scanf("%d",&test);
106
       for (int t=1; t \le test; t++)
107
108
       {
109
            char s[5005];
110
            text.clear();
111
            for (int i=0; i<3; i++)
                 scanf("%s",s);
113
                 Index[i].first=text.size();
114
                 text+=s;
                 Index[i].second=text.size();
116
                 if(i!=2) text = "#";
117
118
            build_suffix_array();
119
            build_LCP_array();
120
            int ans=0;
121
            int low=0,high=text.size();
            while (low<=high)
124
                 int mid = (low + high) >> 1;
                 if (Ok(mid))
126
                 {
127
                     ans=mid;
128
129
                     low=mid+1;
130
131
                 else high=mid-1;
            }
132
            printf("Case %d: %d\n",t,ans);
134
       return 0;
135
136
```

5.4.4 Number of Distinct sub sequence:

```
1 The solution consists of constructing the suffix array and then finding the
      number of distinct substrings based on the Longest Common Prefixes.
2 One key observation here is that:
3 If you look through the prefixes of each suffix of a string, you have
      covered all substrings of that string.
5 Let us take an example: BANANA
6 Suffixes are:
7 0) BANANA
8 1) ANANA
9 2) NANA
10 3) ANA
11 4) NA
12 5) A
13 It would be a lot easier to go through the prefixes if we sort the above set
       of suffixes, as we can skip the repeated prefixes easily.
14 Sorted set of suffixes:
15 5) A
16 3) ANA
17 1) ANANA
18 0) BANANA
19 4) NA
20 2) NANA
22 From now on,
_{\rm 23} LCP = Longest Common Prefix of 2 strings.
24 Initialize
ans = length (first suffix) = length ("A") = 1.
27 Now consider the consecutive pairs of suffixes, i.e, [A, ANA], [ANA, ANANA],
       [ANANA, BANANA], etc. from the above set of sorted suffixes.
29 We can see that,
^{30} LCP("A", "ANA") = "A".
31 All characters that are not part of the common prefix contribute to a
      distinct substring. In the above case, they are 'N' and 'A'. So they
      should be added to ans.
32 So we have,
33 1.ans += length("ANA") - LCP("A", "ANA")
34 \ 2. \, ans = ans + 3 - 1 = ans + 2 = 3
36 do same thing for other pairs.
```

Math

6.1 Pick's Theorem:

```
Area = \frac{B}{2} + I - 1
```

6.2 Reduce Ratio

```
\left(\frac{A}{B}\right) ratio reduce to \left(\frac{x}{y}\right)
```

```
int main()

int main()

int A,B,x,y;

int A,B,x,y;

cin>>A>>B>>x>>y;

int g=_-gcd(x,y);

x/=g,y/=g;

int t=min(A/x,B/y);

cout<<x*t<<" "<<y*t<<endl;

return 0;

}</pre>
```

6.3 Floyd's Cycle Finding algorithm

```
1 #include < bits / stdc++.h>
2 using namespace std;
3 #define pp pair<int,int>
_{4} int Z,L,M,I;
5 int f(int L)
6 {
       return (Z*L+I)%M;
7
8 }
9 pp CycleFinding()
10
       ///L here initial seed
11
       int hare, tortoise, lambda, meu;
12
       bool cyclefind=false;
13
       hare=tortoise=L;
14
       while (!cyclefind)
15
16
           tortoise=f(tortoise);
17
           hare=f(hare);
18
19
           hare=f(hare);
           if (hare=tortoise) cyclefind=true;
```

```
21
       hare=L;
22
23
       meu=0;
       while (hare!=tortoise)
24
25
           meu++;
           hare=f(hare);
27
           tortoise=f(tortoise);
28
29
       int i=0;
30
       hare=L;
31
       while (i<=meu)
32
33
34
           i++;
35
           hare=f(hare);
36
       tortoise=f(hare);
37
38
       lambda=1;
       while (hare!=tortoise)
39
40
            tortoise=f(tortoise);
41
           lambda++;
42
43
       return {meu,lambda}; ///meu is starting index and lambda is cycle length
44
45 }
  int main()
46
47
  {
48
       int t=1;
       while (scanf("%d %d %d %d",&Z,&I,&M,&L) and (Z or I or M or L))
49
50
           pp a=CycleFinding();
51
           cout << "Cycle starts from index "<< a.first << "\nCycle length is "<< a.
      second << endl;
       }
53
54
       return 0;
55 }
```

6.4 Angle between clock hand

```
double angle(int h, int m)

double hAngle = 0.5D * (h * 60 + m);

double mAngle = 6 * m;

double angle = abs(hAngle - mAngle);

angle = min(angle, 360 - angle);

return angle;

}
```

6.5 Big Integer

```
struct Bigint

// representations and structures
string a; // to store the digits
int sign; // sign = -1 for negative numbers, sign = 1 otherwise

// constructors
Bigint() {} // default constructor
Bigint( string b )
{
```

```
(*this) = b; // constructor for string
11
12
13
       Bigint (long long num)
14
           if (\text{num} < 0) sign = -1;
           else sign=1;
           if (num==0)a.push_back('0');
17
           while (num)
18
19
               a.push_back(num\%10+'0');
20
               num/=10;
21
           }
22
23
       // some helpful methods
24
25
       int size() // returns number of digits
26
27
           return a.size();
28
                              // changes the sign
       Bigint inverseSign()
29
30
           sign *= -1;
31
           return (*this);
32
33
       Bigint normalize ( int newSign ) // removes leading 0, fixes sign
34
35
           for ( int i = a. size () - 1; i > 0 && a[i] == '0'; i- )
36
               a.erase(a.begin() + i);
37
           sign = (a.size() = 1 && a[0] = '0') ? 1 : newSign;
38
39
           return (*this);
40
       }
41
       // assignment operator
42
       void operator = ( string b ) // assigns a string to Bigint
43
44
           a = b[0] = '-'; b.substr(1) : b;
45
           reverse(a.begin(), a.end());
this->normalize(b[0] = '-' ? -1 : 1);
46
47
48
49
       // conditional operators
50
       bool operator < ( const Bigint &b ) const // less than operator
51
52
           if( sign != b.sign ) return sign < b.sign;</pre>
53
           if ( a.size() != b.a.size() )
               return sign == 1 ? a.size() < b.a.size() : a.size() > b.a.size()
55
           for ( int i = a. size () - 1; i >= 0; i --- ) if ( a[i] != b.a[i] )
56
                    return sign = 1 ? a[i] < b.a[i] : a[i] > b.a[i];
57
           return false;
58
       bool operator == ( const Bigint &b ) const // operator for equality
60
61
           return a == b.a && sign == b.sign;
62
63
     mathematical operators
64
       Bigint operator + (Bigint b) // addition operator overloading
65
66
           if( sign != b.sign ) return (*this) - b.inverseSign();
67
68
           for (int i = 0, carry = 0; i < a.size() || i < b.size() || carry; i++)
69
```

```
carry += (i < a. size() ? a[i] - 48 : 0) + (i < b.a. size() ? b.a[i] - 48 : 0)
71
72
               c.a += (carry \% 10 + 48);
               carry /= 10;
73
74
           return c.normalize(sign);
75
76
       Bigint operator - ( Bigint b )
                                       // subtraction operator overloading
77
78
       {
           if(sign != b.sign) return(*this) + b.inverseSign();
79
           int s = sign;
80
           sign = b.sign = 1;
81
           if((*this) < b) return ((b - (*this)).inverseSign()).normalize(-s)
82
83
           for ( int i = 0, borrow = 0; i < a.size(); i++)
84
               borrow = a[i] - borrow - (i < b.size() ? b.a[i] : 48);
86
               c.a += borrow >= 0? borrow + 48: borrow + 58;
87
               borrow = borrow >= 0 ? 0 : 1;
88
89
           return c.normalize(s);
90
91
       Bigint operator * (Bigint b ) // multiplication operator overloading
92
93
           Bigint c("0");
94
           for (int i = 0, k = a[i] - 48; i < a.size(); i++, k = a[i] - 48)
95
96
               while (k--) c = c + b; // ith digit is k, so, we add k times
97
               b.a.insert(b.a.begin(), '0'); // multiplied by 10
98
99
           return c.normalize(sign * b.sign);
100
101
       Bigint operator / (Bigint b) // division operator overloading
103
           if (b.size() = 1 \&\& b.a[0] = '0') b.a[0] /= (b.a[0] - 48);
           Bigint c("0"), d;
           106
           int dSign = sign * b.sign;
107
           b.sign = 1;
108
           for ( int i = a.size() - 1; i >= 0; i - )
109
               c.a.insert(c.a.begin(), '0');
111
               c = c + a.substr(i, 1);
112
               while (!(c < b)) c = c - b, d.a[i]++;
113
114
           return d. normalize (dSign);
115
       Bigint operator % (Bigint b) // modulo operator overloading
117
118
           if(b.size() = 1 \&\& b.a[0] = '0') b.a[0] /= (b.a[0] - 48);
119
           Bigint c("0");
120
           b.sign = 1;
           for (int i = a.size() - 1; i >= 0; i-)
               c.a.insert(c.a.begin(), '0');
124
               c = c + a.substr(i, 1);
               while (!(c < b)) c = c - b;
126
           return c.normalize(sign);
128
129
```

Game Theory

7.1 Min Max Algorithm

```
1 /**
      There is N stone of pile, every player in each move can remove Move[i]
      stone. If Any move there is no Stone remain
      these player will be lost.
3
4 **/
5 #include < bits / stdc++.h>
6 using namespace std;
7 int Move[100], k, n;
9 bool is_valid(int x)
10
       return x>=0;
11
12 }
  bool solve (int sum)
13
15
       if (sum==0)return false;
16
       if(dp[sum]!=-1)return dp[sum];
17
       for (int i=1; i \le k; i++)
18
           if(is_valid(sum-Move[i]) and !solve(sum-Move[i]))return dp[sum]=true
19
20
       return dp[sum]=false;
21
22 }
23
  int main()
24
       while (scanf("%d",&n)==1)
25
26
           scanf("%d",&k);
27
           for (int i=1; i \le k; i++)
2.8
2.9
               scanf("%d",&Move[i]);
30
31
           memset(dp, -1, sizeof(dp));
32
           if(solve(n))printf("A\n");
33
           else printf("B\n");
35
       return 0;
36
37 }
```

7.2 Grundy Number

```
1 \text{ const int } MX = 1000005 ;
_{2} int grunday [MX] , rev [6] = { 1 , 4 , 27, 256, 3125, 46656} ; // in a single
      move n*n can be removed
3 int Cal(int x)
4 {
       if ( x = 0 ) return 0 ; // base case , grunday is always zero
5
       int &ret = grunday[x];
6
       if (ret != -1) return ret;
7
       set < int > s;
8
       int i ;
9
       for (i = 0 ; i < 6 ; i++)
10
11
            if(x - rev[i] >= 0)
12
13
                s. insert (Cal(x-(rev[i])));
14
15
       int ans = 0;
       while(s.find(ans) != s.end()) ans++;
16
       return ret = ans ;
17
18 }
  int main()
19
20
       // ios_base::sync_with_stdio(0); cin.tie(0);
21
       memo(grunday, -1);
22
       int cs , t ;
23
       scanf("%d",&t);
24
25
       for (cs = 1; cs \ll t; cs \leftrightarrow )
26
           int n , i , ans = 0 , x ;
27
           \operatorname{scanf}("%d",\&n);
28
           29
           {
30
                scanf("%d",&x);
31
                ans \hat{} = Cal(x);
32
33
           if( ans ) puts("Little Chef"); // first mover win
else puts("Head Chef");
34
37
       return 0;
38
39 }
```

Chapter 8

Number Theory

8.1 Greatest common divisor (GCD)

```
int gcd(int a, int b)
2 {
3    return b == 0 ? a : gcd(b, a % b);
4 }
```

8.2 Least common multiple (LCM)

```
int lcm(int a, int b)

{
    return (a / gcd(a, b)) * b;
}
```

8.3 Big Mod

8.3.1 Iterative version

```
LLD bigmod (LLD B, LLD P, LLD M)

LLD R=1;

while (P>0)

f

if (P%2==1)

R=(R*B)%M;

P/=2;
B=(B*B)%M;

B=(B*B)%M;

return R;
```

8.3.2 Recursive version

```
int big_mod(int n, int p, int m)

{
    if(!n) return 1;
    int x=(n, p>>1 , m);
    x= (x*x)%m;
    if(p&1) x=(x*n)%m;
    return x;
}
```

8.4 Sieve of Eratosthenes

8.4.1 Sieve

```
1 const int Max=1000000;
2 bool prime_check [Max+5];
з vint prime;
4 void sieve()
5
6
       int i,j;
       int sq=sqrt(Max);
       for (i = 2; i * i < = Max; i++)
8
            if (!prime_check[i])
10
            {
11
                 prime.pb(i);
12
                 for (j=i*i;j \le Max;j+=i)
13
14
                      prime_check[j]=true;
            }
17
18
       for (i=sq+1; i \le Max; i++)
19
       {
20
            if (! prime_check[i]) prime.pb(i);
21
22
23
```

8.4.2 Bitwise Sieve

```
1 #include < stdio.h>
2 #include <math.h>
\frac{\text{#define check}(N, pos)}{\text{(N&(1<<(pos)))}}
4 #define Set (N, pos) N|(1 << (pos))
5 const int M=100000;
6 int status [M/32+2]=\{0\};
  void sieve()
8
        int sqN=(int)sqrt(M);
9
        \begin{array}{ll} \textbf{int} & i\ , j\ , cnt \!=\! 1; \end{array}
        puts("2");
        for (i=3; i \le sqN; i+=2)
12
13
              if(check(status[i>>5], i\&31)==0)
14
              {
15
                   for (j=i*i;j \le M;j+=2*i)
16
17
                       status[j>>5]=Set(status[j>>5],j&31);
18
19
              }
20
        for (i = 3; i < M; i + = 2)
22
23
              if(check(status[i>>5], i\&31)==0)
24
              {
25
                   printf("%d\n",i);
27
28
29
```

8.5 Euler's totient function

Euler's totient function (or Euler's phi function), denoted as $\varphi(n)or\phi(n)$, is an arithmetic function that counts the positive integers less than or equal to n that are relatively prime to n. (These integers are sometimes referred to as totatives of n.) Thus, if n is a positive integer, then $\varphi(n)$ is the number of integers k in the range $1 \le k \le n$ for which the greatest common divisor $\gcd(n, k) = 1$.

8.5.1 Euler phi sieve version

```
int phi[10000];
2 const int M=1000;
3 void Generate_phi()
4 {
5
         \quad \quad i\, n\, t \quad i \ ,\, j \ ; \quad
         phi[1]=1;
6
         for (i=2; i < M; i++)
7
8
               if (! phi [ i ])
9
10
                    phi[i]=i-1;
11
                    for(j=i+i; j<M; j+=i)
13
                          if (! phi [ j ] ) phi [ j ] = j;
14
                          phi[j] = phi[j] / i*(i-1);
15
              }
17
18
        }
19
```

8.5.2 Euler phi function

```
1 int phi (int n)
2 {
3
      for (int i = 2; i * i <= n; i++)
4
5
           if (n \% i == 0)
6
               while (n \% i == 0)
8
9
                   n /= i;
11
               ret = ret / i;
12
           }
13
14
         this case will happen if n is a prime number
      // in that case we won't find any prime that divides n
16
      // that's less or equal to sqrt(n)
17
      if (n > 1) ret -= ret / n;
18
      return ret;
19
20 }
```

8.6 Number of Divisors

If a number

$$N = P1^{k1} * P2^{k2} * P3^{k3}.....Pn^{kn}$$

Number of divisor:

$$f(n) = (k1+1) * (k2+1) * (k3+1) * \dots (kn+1)$$

8.6.1 NOD sieve version

```
1 \text{ int } nod[100000+5];
2 void Generate()
3 {
       nod[1] = 1;
4
       for (int i=2; i \le 100000; i++)
5
6
                           //here checking i is prime or not ???
            if (!nod[i])
8
            {
                 nod[i]=2;
9
                 for (int j=i+i; j <=100000; j+=i)
10
                      if (!nod[j])nod[j]=1;
12
                      int n=j, cnt=0;
13
                      while (!(n%i))
14
15
                          cnt++;
16
                          n/=i;
17
18
                     nod[j]*=(cnt+1);
20
            }
21
       }
22
23 }
```

8.6.2 NOD function

```
int find_divisor(int n){
       int i, ans=1, count=1;
       while (n\%2==0)
4
            n/=2;
5
            count++;
6
       ans *= count;
8
       for (i=3; i*i <= n; i+=2)
9
10
            count=1;
11
            while (n\%i == 0)
12
13
                 n/=i;
14
15
                 count++;
16
            ans*=count;
17
18
       if (n>1)
19
20
            ans*=2;
21
22
23
       return ans;
24 }
```

8.7 Summation of Divisors

If the prime factorization of an integer is:

$$P1^{x1} * P2^{x2} * \dots * Pn^{xn}$$

Where P1, P2, Pn are primes, then the summation of divisors is

$$\frac{(P1^{x1}-1)}{(P1-1)}*\frac{(P2^{x2}-1)}{(P2-1)}*.....*\frac{(Pn^{xn}-1)}{(Pn-1)}$$

8.8 Extended Euclid GCD

```
typedef pair < int , int > pii;
#define x first
#define y second

pii extendedEuclid(int a, int b)  // returns x, y | ax + by =gcd(a,b)

if (b == 0) return pii(1, 0);
else

pii d = extendedEuclid(b, a % b);
return pii(d.y, d.x - d.y * (a / b));
}
```

8.9 Modular multiplicative inverse

8.9.1 Modular inverse using EGCD

```
int modularInverse(int a, int n)
{
    pii ret = extendedEuclid(a, n);
    return ((ret.x % n) + n) % n;
}
```

8.9.2 Modular multiplicative inverse using Euler Theorem

According to Euler's theorem, if a is coprime to m, that is, gcd(a, m) = 1, then

$$a^{\varphi(m)} \equiv 1 \pmod{m}$$

where (m) is Euler's totient function. This follows from the fact that a belongs to the multiplicative group $(\frac{Z}{mZ})$ iff a is coprime to m. Therefore the modular multiplicative inverse can be found directly:

$$a^{\varphi(m)-1} \equiv a^{-1} \pmod{m}$$

. In the special case when m is a prime, the modular inverse is given by the below equation as:

$$a^{-1} \equiv a^{m-2} \pmod{m}$$

.

8.10 NCR

8.10.1 Lucas Theorem

```
1 /**
       Fine NCR \% M when N C M are large number.
2
3
       using Lucas theorem.
4 **/
5 \#include < bits / stdc ++.h >
6 using namespace std;
7 typedef long long LLD;
8 LLD mod=1000003;
9 LLD big_mod(LLD n,LLD p,LLD m)
10 {
       if (p==0)return (LLD) 1;
11
       LLD x=big_mod(n, p/2, m);
12
       x=(x*x)\%m;
13
       if(p&1)x=(x*n)\%m;
14
       return x;
15
16
17 LLD inverse_modulo(LLD t,LLD m)
18 {
       return big_mod(t,m-2,m);
19
20
  }
21 LLD combi (LLD n, LLD k, LLD m)
22 {
       if (n<k)
23
           return 0;
24
       if (n-k<k)
25
           return combi(n, n-k, m);
26
       LLD i, p=1, t=1;
27
       for (i=n-k+1; i \le n; i++)
28
29
           p=(p*i)\%m;
30
       for (i=1; i \le k; i++)
           t = (t * i) \%m;
31
       return (p*inverse_modulo(t,m))%m;
32
33
34 LLD lucas (LLD n, LLD k, LLD m)
35
       if (n<k)
36
           return 0;
37
       if (k==0 || n==k)
38
           return 1;
39
       return (lucas (n/m, k/m,m) *combi (n/m, k/m,m))/m;
40
41 }
42 int main()
43 {
44
       return 0;
45 }
```

Chapter 9

Computational geometry

9.1 Geo Library (Stanford)

```
1 // C++ routines for computational geometry.
3 #include <iostream>
4 #include <vector>
5 #include <cmath>
6 #include <cassert>
7 #define M_PI (2*acos(0))
8 using namespace std;
double INF = 1e100;
double EPS = 1e-12;
13 struct PT {
    double x, y;
14
    PT() {}
15
    PT(double x, double y) : x(x), y(y) \{\}
16
    PT(const PT \& p) : x(p.x), y(p.y)
17
                                        {}
    PT operator + (const PT &p) const { return PT(x+p.x, y+p.y); }
18
    PT operator - (const PT &p) const { return PT(x-p.x, y-p.y); }
    PT operator * (double c)
                                 const \{ return PT(x*c, y*c); \}
    PT operator / (double c)
                                  const { return PT(x/c,
22 };
23
double dot (PT p, PT q)
                            { return p.x*q.x+p.y*q.y; }
double dist2(PT p, PT q)
                            \{ return dot(p-q,p-q); \}
double cross (PT p, PT q) { return p.x*q.y-p.y*q.x; }
ostream & operator << (ostream & os, const PT & p) {
    os << "(" << p.x << "," << p.y << ")";
28
29 }
31 // rotate a point CCW or CW around the origin
32 PT RotateCCW90(PT p)
                        \{ return PT(-p.y, p.x) ;
33 PT RotateCW90(PT p)
                         \{ return PT(p.y,-p.x); \}
34 PT RotateCCW(PT p, double t) {
    return PT(p.x*cos(t)-p.y*sin(t), p.x*sin(t)+p.y*cos(t));
36 }
37
38 // project point c onto line through a and b
39 // assuming a != b
40 PT ProjectPointLine(PT a, PT b, PT c) {
return a + (b-a)*dot(c-a, b-a)/dot(b-a, b-a);
```

```
// project point c onto line segment through a and b
44
45 PT ProjectPointSegment (PT a, PT b, PT c) {
     double r = dot(b-a, b-a);
46
     if (fabs(r) < EPS) return a;</pre>
47
     r = dot(c-a, b-a)/r;
     if (r < 0) return a;
     if (r > 1) return b;
50
     return a + (b-a)*r;
51
52 }
53
_{\rm 54} // compute distance from c to segment between a and b
55 double DistancePointSegment (PT a, PT b, PT c) {
     return sqrt(dist2(c, ProjectPointSegment(a, b, c)));
56
57 }
58
59 // compute distance between point (x,y,z) and plane ax+by+cz=d
   double DistancePointPlane(double x, double y, double z,
                                 double a, double b, double c, double d)
61
62
     return fabs(a*x+b*y+c*z-d)/sqrt(a*a+b*b+c*c);
63
64 }
65
66 // determine if lines from a to b and c to d are parallel or collinear
67 bool LinesParallel(PT a, PT b, PT c, PT d) {
     return fabs (cross(b-a, c-d)) < EPS;
69 }
70
71 bool LinesCollinear (PT a, PT b, PT c, PT d) {
     return LinesParallel(a, b, c, d)
         && fabs(cross(a-b, a-c)) < EPS
73
         && fabs (cross(c-d, c-a)) < EPS;
74
75 }
76
   // determine if line segment from a to b intersects with
77
   // line segment from c to d
   bool SegmentsIntersect (PT a, PT b, PT c, PT d) {
79
     if (LinesCollinear(a, b, c, d)) {
        \begin{array}{ll} \mbox{if} & (\mbox{ dist2}\,(a,\ c) < \mbox{EPS} \ || \ \mbox{dist2}\,(a,\ d) < \mbox{EPS} \ || \\ \mbox{dist2}\,(b,\ c) < \mbox{EPS} \ || \ \mbox{dist2}\,(b,\ d) < \mbox{EPS}) \ \mbox{return} \ \mbox{true}\,; \end{array}
81
82
        if (dot(c-a, c-b) > 0 \&\& dot(d-a, d-b) > 0 \&\& dot(c-b, d-b) > 0)
83
          return false;
84
        return true;
85
86
     if (cross(d-a, b-a) * cross(c-a, b-a) > 0) return false;
87
     if (cross(a-c, d-c) * cross(b-c, d-c) > 0) return false;
88
89
     return true;
90 }
91
92 // compute intersection of line passing through a and b
93 // with line passing through c and d, assuming that unique
94 // intersection exists; for segment intersection, check if
95 // segments intersect first
96 PT ComputeLineIntersection (PT a, PT b, PT c, PT d) {
     b=b-a; d=c-d; c=c-a;
97
     assert(dot(b, b) > EPS \&\& dot(d, d) > EPS);
98
99
     return a + b*cross(c, d)/cross(b, d);
100
  // compute center of circle given three points
103 PT ComputeCircleCenter(PT a, PT b, PT c) {
```

```
b = (a+b)/2;
104
     c = (a+c)/2;
     return ComputeLineIntersection(b, b+RotateCW90(a-b), c, c+RotateCW90(a-c))
106
107 }
109 // determine if point is in a possibly non-convex polygon (by William
110 // Randolph Franklin); returns 1 for strictly interior points, 0 for
111 // strictly exterior points, and 0 or 1 for the remaining points.
112 // Note that it is possible to convert this into an *exact* test using
113 // integer arithmetic by taking care of the division appropriately
114 // (making sure to deal with signs properly) and then by writing exact
115 // tests for checking point on polygon boundary
   bool PointInPolygon(const vector <PT> &p, PT q) {
116
117
     bool c = 0;
118
     for (int i = 0; i < p.size(); i++){
119
       int j = (i+1)\%p. size();
120
       if ((p[i].y \le q.y \&\& q.y < p[j].y ||
121
         p[j].y \le q.y \&\& q.y < p[i].y) \&\&
         q.x < p[i].x + (p[j].x - p[i].x) * (q.y - p[i].y) / (p[j].y - p[i].y)
         c = !c;
     }
124
     return c;
126
127
   // determine if point is on the boundary of a polygon
128
   bool PointOnPolygon(const vector <PT> &p, PT q) {
     for (int i = 0; i < p.size(); i++)
       if (dist2(ProjectPointSegment(p[i], p[(i+1)\%p.size()], q), q) < EPS)
131
132
         return true;
133
       return false;
134
135
   // compute intersection of line through points a and b with
136
   // circle centered at c with radius r > 0
137
   vector <PT> CircleLineIntersection (PT a, PT b, PT c, double r) {
138
     vector<PT> ret;
139
     b = b-a;
140
     a = a-c;
141
     double A = dot(b, b);
142
     double B = dot(a, b);
143

double C = dot(a, a) - r*r;

144
     double D = B*B - A*C;
145
     if (D < -EPS) return ret;
146
     ret.push_back(c+a+b*(-B+sqrt(D+EPS))/A);
147
     if (D > EPS)
148
       ret.push_back(c+a+b*(-B-sqrt(D))/A);
149
     return ret;
150
151 }
   // compute intersection of circle centered at a with radius r
153
   // with circle centered at b with radius R
   vector <PT> CircleCircleIntersection (PT a, PT b, double r, double R) {
155
     vector <PT> ret;
     double d = sqrt(dist2(a, b));
     if (d > r+R \mid | d+min(r, R) < max(r, R)) return ret;
158
159
     double x = (d*d-R*R+r*r)/(2*d);
160
     double y = sqrt(r*r-x*x);
     PT v = (b-a)/d;
     ret.push_back(a+v*x + RotateCCW90(v)*y);
     if (y > 0)
163
```

```
ret.push_back(a+v*x - RotateCCW90(v)*y);
164
      return ret;
165
166 }
167
   // This code computes the area or centroid of a (possibly nonconvex)
168
   // polygon, assuming that the coordinates are listed in a clockwise or
   // counterclockwise fashion. Note that the centroid is often known as
   // the "center of gravity" or "center of mass".
   double ComputeSignedArea(const vector <PT> &p) {
      double area = 0;
173
      for(int i = 0; i < p.size(); i++) {
174
        int j = (i+1) \% p. size();
175
        a\,rea \,\,+\!\!=\,p\,[\,i\,]\,.\,x\!*\!p\,[\,j\,]\,.\,y\,\,-\,\,p\,[\,j\,]\,.\,x\!*\!p\,[\,i\,]\,.\,y\,;
      }
177
178
      return area / 2.0;
179
    double ComputeArea(const vector <PT> &p) {
      return fabs(ComputeSignedArea(p));
182
183
184
   PT ComputeCentroid(const vector<PT> &p) {
185
      PT c(0,0);
186
      double scale = 6.0 * ComputeSignedArea(p);
187
      for (int i = 0; i < p.size(); i++){
188
        int j = (i+1) \% p. size();
189
        c = c + (p[i]+p[j])*(p[i].x*p[j].y - p[j].x*p[i].y);
191
      return c / scale;
192
193
194
    // tests whether or not a given polygon (in CW or CCW order) is simple
195
    bool IsSimple(const vector <PT> &p) {
196
      for (int i = 0; i < p.size(); i++)
197
         for (int k = i+1; k < p.size(); k++) {
198
           int j = (i+1) \% p.size();
199
           int l = (k+1) \% p. size();
200
           if (i = l \mid | j = k) continue;
201
           if (SegmentsIntersect(p[i], p[j], p[k], p[l]))
202
              return false;
203
204
205
      return true;
206
207
208
   int main() {
209
210
      // expected: (-5,2)
211
      cerr \ll RotateCCW90(PT(2,5)) \ll endl;
212
213
      // expected: (5,-2)
214
      cerr << RotateCW90(PT(2,5)) << endl;
215
216
      // expected: (-5,2)
217
      cerr \ll RotateCCW(PT(2,5), M_PI/2) \ll endl;
218
219
      // expected: (5,2)
220
      \operatorname{cerr} \ll \operatorname{ProjectPointLine}(\operatorname{PT}(-5,-2), \operatorname{PT}(10,4), \operatorname{PT}(3,7)) \ll \operatorname{endl};
221
222
      // expected: (5,2) (7.5,3) (2.5,1)
223
      \operatorname{cerr} << \operatorname{ProjectPointSegment}(\operatorname{PT}(-5,-2), \operatorname{PT}(10,4), \operatorname{PT}(3,7)) << ""
```

```
<< ProjectPointSegment(PT(7.5,3), PT(10,4), PT(3,7)) << ""
225
            << ProjectPointSegment (PT(-5,-2), PT(2.5,1), PT(3,7)) <math><< endl;
226
227
      // expected: 6.78903
228
      \operatorname{cerr} << \operatorname{DistancePointPlane}(4, -4, 3, 2, -2, 5, -8) << \operatorname{endl};
229
230
      // expected: 1 0 1
231
      cerr << LinesParallel(PT(1,1) \;,\; PT(3,5) \;,\; PT(2,1) \;,\; PT(4,5)) \;<< \text{"} \;\text{"}
            << \  \, LinesParallel\left(PT(1\,,1)\;,\; PT(3\,,5)\;,\; PT(2\,,0)\;,\; PT(4\,,5)\,\right)\;<<\;"\;\;"
233
            << LinesParallel(PT(1,1), PT(3,5), PT(5,9), PT(7,13)) << endl;
234
235
      // expected: 0 0 1
236
      {\tt cerr} \, <\! < \, {\tt LinesCollinear} \, ({\tt PT}(1\,,\!1) \,, \, {\tt PT}(3\,,\!5) \,, \, {\tt PT}(2\,,\!1) \,, \, {\tt PT}(4\,,\!5) \,) \, <\! < \, " \, "
237
            << LinesCollinear (PT(1,1), PT(3,5), PT(2,0), PT(4,5)) <math><< ""
238
            << LinesCollinear(PT(1,1), PT(3,5), PT(5,9), PT(7,13)) << endl;</pre>
239
240
      // expected: 1 1 1 0
      cerr << \ SegmentsIntersect (PT(0\,,0)\;,\; PT(2\,,4)\;,\; PT(3\,,1)\;,\; PT(-1,3)\,) <<\ "\;\;"
242
            << SegmentsIntersect(PT(0,0), PT(2,4), PT(4,3), PT(0,5)) << ""
243
            << \  \, {\rm SegmentsIntersect}\left(PT(0\,,\!0)\;,\; PT(2\,,\!4)\;,\; PT(2\,,\!-1)\;,\; PT(-2\,,\!1)\,\right)\;<<\;"\;\;"
244
            << SegmentsIntersect (PT(0,0), PT(2,4), PT(5,5), PT(1,7)) <math><< endl;
245
246
      // expected: (1,2)
247
      cerr << ComputeLineIntersection\left(PT(0\,,\!0)\,,\;PT(2\,,\!4)\,,\;PT(3\,,\!1)\,,\;PT(-1\,,\!3)\right) <<
248
        endl;
249
      // expected: (1,1)
250
      cerr \ll ComputeCircleCenter(PT(-3,4), PT(6,1), PT(4,5)) \ll endl;
251
252
253
      vector < PT > v;
      v.push_back(PT(0,0));
254
      v.push_back(PT(5,0));
255
      v.push_back(PT(5,5));
256
      v.push_back(PT(0,5));
257
258
      // expected: 1 1 1 0 0
259
      cerr << PointInPolygon(v, PT(2,2)) << ""
260
            << PointInPolygon(v, PT(2,0)) << " "</pre>
261
            << PointInPolygon(v, PT(0,2)) <<
262
            << PointInPolygon(v, PT(5,2)) << ""
263
            << PointInPolygon(v, PT(2,5)) << endl;</pre>
264
265
      // expected: 0 1 1 1 1
266
      cerr << PointOnPolygon(v, PT(2,2)) << ""
267
            << PointOnPolygon(v, PT(2,0)) << ""
268
            << PointOnPolygon(v, PT(0,2)) << ""
269
            << PointOnPolygon(v, PT(5,2)) << ""
270
            << PointOnPolygon(v, PT(2,5)) << endl;
271
272
          expected: (1,6)
273
                      (5,4) (4,5)
274
                      blank line
275
                      (4,5) (5,4)
276
                      blank line
277
                      (4,5) (5,4)
278
      vector < PT > u = CircleLineIntersection(PT(0,6), PT(2,6), PT(1,1), 5);
279
      for (int i = 0; i < u.size(); i++) cerr << u[i] << ""; cerr <math><< endl;
280
      u = CircleLineIntersection(PT(0,9), PT(9,0), PT(1,1), 5);
281
      for (int i = 0; i < u.size(); i++) cerr << u[i] << "; cerr << endl;
      u = CircleCircleIntersection(PT(1,1), PT(10,10), 5, 5);
      for (int i = 0; i < u.size(); i++) cerr << u[i] << " "; cerr << endl;
```

```
u = CircleCircleIntersection(PT(1,1), PT(8,8), 5, 5);
285
     for (int i = 0; i < u.size(); i++) cerr << u[i] << ""; <math>cerr << endl;
286
     u = CircleCircleIntersection(PT(1,1), PT(4.5,4.5), 10, sqrt(2.0)/2.0);
287
     288
     u = CircleCircleIntersection(PT(1,1), PT(4.5,4.5), 5, sqrt(2.0)/2.0);
289
     for (int i = 0; i < u.size(); i++) cerr << u[i] << " "; cerr << endl;
291
     // area should be 5.0
292
     // centroid should be (1.1666666, 1.166666)
293
     PT pa[] = \{ PT(0,0), PT(5,0), PT(1,1), PT(0,5) \};
294
     {\tt vector} <\!\! {\tt PT}\!\! > \; p\left(\, {\tt pa} \,,\;\; {\tt pa} \!+\! 4\right);
295
     PT c = ComputeCentroid(p);
296
     cerr << "Area: " << ComputeArea(p) << endl; \\
297
     cerr << "Centroid: " << c << endl;</pre>
298
300
     return 0;
301 }
```

9.2 Convex Hull

9.2.1 Convex Hull 1

```
1 typedef long long LLD;
2 struct POINT
3 {
       double x, y;
4
5
      POINT(LLD x=0,LLD y=0):x(x),y(y) {};
6 };
  POINT P0;
  int area2 (POINT p0, POINT p1, POINT p2)
10
11
       return (p2.x-p0.x)*(p1.y-p0.y)-(p1.x-p0.x)*(p2.y-p0.y);
12
13 }
int dist2 (POINT a, POINT b)
15 {
       int dx=a.x-b.x, dy=a.y-b.y;
16
       return dx*dx+dy*dy;
17
18
  }
19
  ///Compare version 1
  bool angle_comp (POINT a, POINT b)
21
22
       if(area2(P0,a,b)<0)return true;
23
       else if (area2(P0,a,b)==0)return dist2(P0,a)<dist2(P0,b);
24
       return false;
25
26
27
  ///Compare version 2
28
  bool angle_comp2 (POINT a, POINT b)
30
31
       if(area2(P0,a,b)==0)
32
33
           return dist2 (P0, a) < dist2 (P0, b);
34
35
       int d1x=a.x-P0.x, d1y=a.y-P0.y;
36
       int d2x=b.x-P0.x, d2y=b.y-P0.y;
37
       return (atan2((double)d1y,(double)d1x)-atan2((double)d2y,(double)d2x))
```

```
39 }
40
vector <POINT> Find_convex_hull(vector <POINT>P)
42
       int i,N=(int)P.size();
43
44
       int po=0;
45
       for (i=0; i< N; i++)
46
       {
            if(P[i].y < P[po].y \mid | (P[i].y = P[po].y && P[i].x < P[po].x))po=i;
47
48
       P0 = P[po];
49
       sort(P.begin(),P.end(),angle\_comp);
50
       POINT prev(0,0), now(0,0);
52
       stack < POINT > S;
53
       S. push (P0);
54
       i = 1;
55
       while (i < N)
56
       {
            if (S. size()<2)
57
58
            {
                 S.push(P[i++]);
59
            }
60
            else
61
            {
62
                 now = S.top();
63
                 S.pop();
64
65
                 prev=S.top();
66
                 S.push(now);
67
                 if ( area2 ( prev , now , P[ i ] ) < 0)</pre>
68
                      S.push(P[i++]);
                 else S.pop();
69
            }
70
       }
71
72
       vector < POINT > ConvexHull;
73
       ConvexHull.push_back(P0);
74
       while (!S.empty())
75
76
77
            ConvexHull.push_back(S.top());
78
            S.pop();
79
80
       return ConvexHull;
81
82
```

9.2.2 Convex Hull (Stanford)

```
1 // Compute the 2D convex hull of a set of points using the monotone chain
2 // algorithm. Eliminate redundant points from the hull if REMOVEREDUNDANT
     is
3 // #defined.
4 //
5 // Running time: O(n log n)
6 //
                a vector of input points, unordered.
       INPUT:
7 //
       OUTPUT:
                a vector of points in the convex hull, counterclockwise,
8 //
      starting
                with bottommost/leftmost point
9
11 #define REMOVE.REDUNDANT
```

```
13 typedef double T;
14 const T EPS = 1e-7;
15 struct PT
16 {
       T x, y;
17
       PT() {}
18
       PT(T x, T y) : x(x), y(y) \{ \}
19
20
       bool operator < (const PT &rhs) const
21
           return make_pair(y,x) < make_pair(rhs.y,rhs.x);</pre>
22
23
       bool operator == (const PT &rhs) const
24
       {
25
26
           return make_pair(y,x) == make_pair(rhs.y,rhs.x);
27
28
  };
29
  T cross (PT p, PT q)
30
31
       return p.x*q.y-p.y*q.x;
32
33
34 T area2 (PT a, PT b, PT c)
35
       return cross(a,b) + cross(b,c) + cross(c,a);
36
37
38
  #ifdef REMOVE.REDUNDANT
40 bool between (const PT &a, const PT &b, const PT &c)
41
       return (fabs(area2(a,b,c)) < EPS && (a.x-b.x)*(c.x-b.x) <= 0 && (a.y-b.y)
42
      )*(c.y-b.y) <= 0);
43
44 #endif
45
  void ConvexHull(vector<PT> &pts)
46
47
       sort(pts.begin(), pts.end());
48
       pts.erase(unique(pts.begin(), pts.end()), pts.end());
49
50
       vector < PT > up, dn;
       for (int i = 0; i < pts.size(); i++)
51
           while (up.size() > 1 \&\& area2(up[up.size()-2], up.back(), pts[i]) >=
53
        0) up.pop_back();
           while (dn.size() > 1 \&\& area2(dn[dn.size()-2], dn.back(), pts[i]) \le
54
        0) dn.pop_back();
           up.push_back(pts[i]);
55
           dn.push_back(pts[i]);
56
57
58
       pts = dn;
       for (int i = (int) up. size() - 2; i \ge 1; i--) pts.push_back(up[i]);
59
60
  #ifdef REMOVE.REDUNDANT
61
       if (pts.size() <= 2) return;
62
       dn.clear();
63
       dn.push_back(pts[0]);
64
65
       dn.push_back(pts[1]);
       for (int i = 2; i < pts.size(); i++)
66
67
       {
            if (\text{between}(\text{dn}[\text{dn.size}()-2], \text{dn}[\text{dn.size}()-1], \text{pts}[i])) \text{dn.pop-back}()
           dn.push_back(pts[i]);
69
```

9.3 Determine a point inside polygon or not

```
1 #include < bits / stdc++.h>
2 using namespace std;
3 struct point
4 {
5
       int x,y;
6
       point() {}
       point(int a, int b)
8
9
            x=a, y=b;
10
11 };
12
   int trun (point p0, point p1, point p2)
13
14
       int result = (p2.x-p0.x)*(p1.y-p0.y)-(p1.x-p0.x)*(p2.y-p0.y);
15
       if(result > 0)return 1;
       if (result < 0) return -1;
16
       return result;
17
18
19
  bool on_segment (point pi, point pj, point pk)
20
       int minx=min(pi.x,pj.x);
21
       int maxx=max(pi.x,pj.x);
22
23
       int miny=min(pi.y,pj.y);
       int maxy=max(pi.y,pj.y);
24
       \label{eq:continuous_pk.y} \begin{subarray}{ll} if (pk.x)=minx\&kpk.y<=maxy) return & true; \\ \end{subarray}
25
26
27 }
  bool Segment_intersect (point p1, point p2, point p3, point p4)
28
29
       int d1=trun(p3,p4,p1);
30
31
       int d2=trun(p3,p4,p2);
       int d3=trun(p1,p2,p3);
32
       int d4=trun(p1,p2,p4);
33
       if ((d1*d2<0)&&(d3*d4<0)) return true;
34
       if(d1=0\&\&on\_segment(p3,p4,p1))return true;
       if(d2=0\&\&on\_segment(p3,p4,p2))return true;
36
       if(d3=0\&\&on\_segment(p1,p2,p3))return true;
37
       if(d4=0\&\&on\_segment(p1,p2,p4))return true;
38
       return false;
39
  }
40
  bool in_poligon(vector<point>p, point given, point inf2)
41
42
  {
43
       int intersect = 0;
44
       int N=p.size();
45
       for (int i=0; i < N; i++)
46
47
            point p1=p[i];
            point p2=p[(i+1)\%N];
```

```
if (Segment_intersect (p1, p2, given, inf2))
49
50
51
                 intersect++;
            }
52
53
       return intersect%2==1;
54
55 }
56 int main()
57 {
       int N;
58
       while (scanf ("%d",&N)&&N)
59
60
            int Max = -10005;
61
            vector<point>poligon;
62
63
            for (int i=0; i< N; i++)
64
                 int x,y;
                 scanf("%d %d",&x,&y);
                 poligon.push_back(point(x,y));
67
                Max=max(x, Max);
68
69
            int x, y;
70
            scanf("%d %d",&x,&y);
71
            if(in\_poligon(poligon, point(x,y), point(Max,y)))
72
73
            {
                 printf("T\n");
74
            }
75
76
            else
77
            {
                 printf("F\n");
78
79
80
       return 0;
81
82
```

9.4 Integer point in a Segment or lattice point

```
int IntegerPointsOnSegment(const point &P1, const point &P2)

point P;

point P;

P.x=abs(P1.x-P2.x);

P.y=abs(P1.y-P2.y);

if(P.x==0) return P.y-1;

if(P.y==0) return P.x-1;

return (--gcd(P.x,P.y))-1;

}
```

9.5 Segment Intersection

```
q.y \le max(p.y, r.y) \&\& q.y >= min(p.y, r.y)
10
11
            return true;
12
       return false;
13 }
14 bool Segment_Intersect(point_int p1, point_int p2, point_int q1, point_int
15 {
       int d1 = Direction(p1, p2, q1);
16
       int d2 = Direction(p1, p2, q2);
17
       int d3 = Direction(q1, q2, p1);
18
       int d4 = Direction(q1, q2, p2);
19
       if (d1 != d2 && d3 != d4) return true;
20
       if (!d1 && On_Segment(p1, q1, p2)) return true;
21
        if (!d2 \&\& On\_Segment(p1, q2, p2)) \ \underline{return} \ \underline{true}; \\
22
       if (!d3 \&\& On\_Segment(q1, p1, q2)) return true;
24
       if (!d4 && On_Segment(q1, p2, q2)) return true;
25
       return false;
26
```

9.6 Verify Convex polygon

```
int trun (point p0, point p1, point p2)
2 {
       int result = (p2.x-p0.x)*(p1.y-p0.y)-(p1.x-p0.x)*(p2.y-p0.y);
3
4
       return result;
  }
5
  bool isConvex(int n, vector<point>v)
6
7
8
       int pos=0, neg = 0;
       for (int i = 0; i < n; i++)
9
10
           int prev = (i + n - 1) \% n, next = (i + 1) \% n;
11
12
           point A=v[i];
13
           point B=v[prev];
14
           point C=v[next];
15
           int pv=trun(A,B,C);
           if(pv > 0) pos++;
16
           else
17
18
                if(pv < 0) neg++;
19
20
21
22
       return (pos == 0) || (neg == 0);
```

9.7 Closest Pair Algorithm

```
const int MAX=100005;
struct point
{
    int x, y, i;
};
point arr [MAX], sortedY [MAX];
bool flag [MAX];
template <class T> int getdist(T a, T b)
{
    return max(abs(a.x - b.x), abs(a.y - b.y));
}
bool compareX(const point &a, const point &b)
{
```

```
14
      return a.x < b.x;
15 }
bool compareY(const point &a, const point &b)
17 {
      return a.y < b.y;
18
19 }
20
      closest_pair (point X[], point Y[], int n)
_{21} int
22 {
      int left_call , right_call , mindist;
23
       if(n == 1) return inf;
24
       if(n = 2) return getdist(X[0], X[1]);
25
      26
27
      point xL[m + 1], xR[m + 1], yL[m + 1], yR[m + 1], Xm = X[m - 1], yS[n];
28
      for (i = 0; i < m; i++)
29
      {
          xL[i] = X[i];
30
           flag[X[i].i] = 0;
31
32
      for (; i < n; i++)
33
34
          xR[i - m] = X[i];
35
          flag[X[i].i] = 1;
36
37
      for (i = n2 = n1 = 0; i < n; i++)
38
39
      {
40
           if(!flag[Y[i].i]) yL[n1++] = Y[i];
41
           else yR[n2++] = Y[i];
42
      left_call = closest_pair(xL, yL, n1);
43
      right_call = closest_pair(xR, yR, n2);
44
      mindist = min(left_call, right_call);
45
      for (i = ns = 0; i < n; i++)
46
           if((Y[i].x - Xm.x) < mindist)
47
               yS[ns++] = Y[i];
48
       for(i = 0; i < ns; i++)
49
           for (j = i + 1; j < ns && (yS[j].y - yS[i].y) < mindist; j++)
50
               mindist = min(mindist, getdist(yS[i], yS[j]));
52
      return mindist;
53
54
  int Find_closestpair(int n)
55
56
57
      sort(arr, arr + n, compareX);
      sort(sortedY, sortedY + n, compareY);
58
      int ans = closest_pair(arr, sortedY, n);
59
60
      return ans;
61 }
```