

**Computer Science and Engineering Discipline
Khulna University, Khulna - 9208**

Roll#:190208

Name: Tanvir Hassan

email: tanvir1908@cseku.ac.bd

Mobile: 01989567104

PROBLEM :) $T(n) = 2 T(n/2) + n \log_2 n$

Here $a = 2$, $b = 2$ and $f(n) = n \log_2 n$, $p = 1$.

$x = \log_b a = \log_2 2 = 1$. Here $x = p$, but $n^x \neq f(n)$,

Find out the solution of the recurrence in term of asymptotic notation.

Solution :

$$n^{\log_b a} = n^{\log_2 2} = n^1 = n$$

Comparing $n^{\log_b a} = n$ and $f(n) = n \log n$

Does not satisfy either Case 1 or 2 or 3 of the Master's theorem .

Case 3 states that $f(n)$ should be polynomially larger but here it is asymptotically larger than $n \log b a$ only by a factor of $\log n$

Let us take $n = 2^m$. Then we have the recurrence

$$m) = 2T(2^{m-1}) + 2^m \log_2(2^m) = 2T(2^{m-1}) + m2^m$$

Calling $T(2^m)$ as $f(m)$, we get that

$$\begin{aligned} f(m) &= 2f(m-1) + m2^m \\ &= 2(2f(m-2) + (m-1)2^{m-1}) + m2^m \\ &= 4f(m-2) + (m-1)2^m + m2^m \\ &= 4(2f(m-3) + (m-2)2^{m-2}) + (m-1)2^m + m2^m \\ &= 8f(m-3) + (m-2)2^m + (m-1)2^m + m2^m \end{aligned}$$

Proceeding on these lines, we get that

$$\begin{aligned}
 f(m) &= 2^m f(0) + 2^m (1+2+3+\dots+m) \\
 &= 2^m f(0) + 2^m (1+2+3+\dots+m) = 2^m f(0) + \frac{m(m+1)}{2} 2^m + \\
 &= 2^m f(0) + m(m+1) 2^{m-1}
 \end{aligned}$$

$$\text{Hence, } T(n) = n(T_1) + n \left(\frac{\log_2(n)(1+\log_2(n))}{2} \right) = \Theta(n \log^2 n).$$

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