

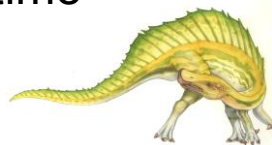
Chapter 3: Processes





Process Concept

- ❑ An operating system executes a variety of programs:
 - ❑ Batch system – **jobs**
 - ❑ Time-shared systems – **user programs** or **tasks**
- ❑ Textbook uses the terms **job** and **process** almost interchangeably
- ❑ **Process** – a program in execution; process execution must progress in sequential fashion
- ❑ Multiple parts
 - ❑ The program code, also called **text section**
 - ❑ Current activity including **program counter**, processor registers
 - ❑ **Stack** containing temporary data
 - ▶ Function parameters, return addresses, local variables
 - ❑ **Data section** containing global variables
 - ❑ **Heap** containing memory dynamically allocated during run time

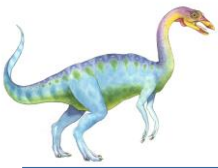




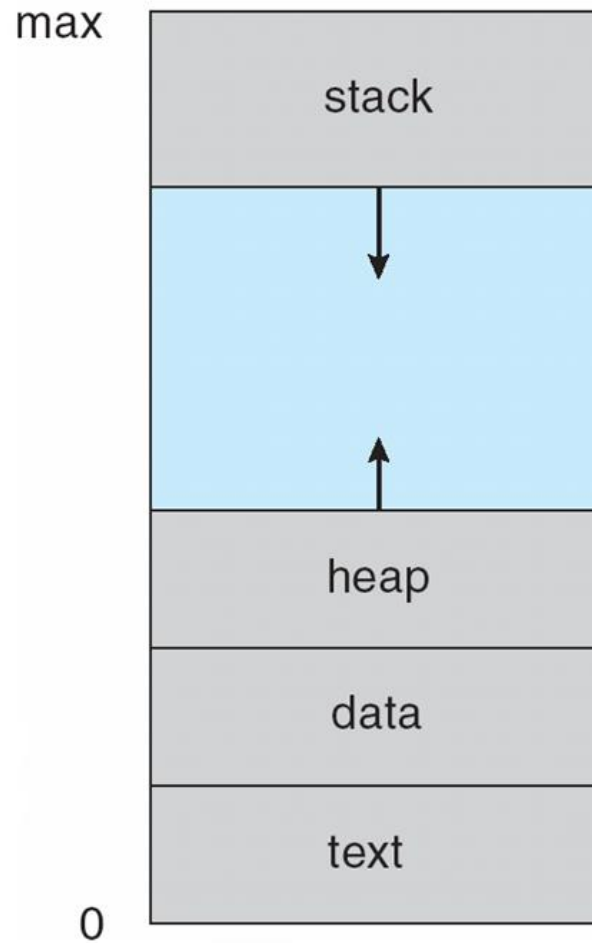
Process Concept (Cont.)

- Program is **passive** entity stored on disk (**executable file**), process is **active**
 - Program becomes process when executable file loaded into memory
- Execution of program started via GUI mouse clicks, command line entry of its name, etc
- One program can be several processes
 - Consider multiple users executing the same program



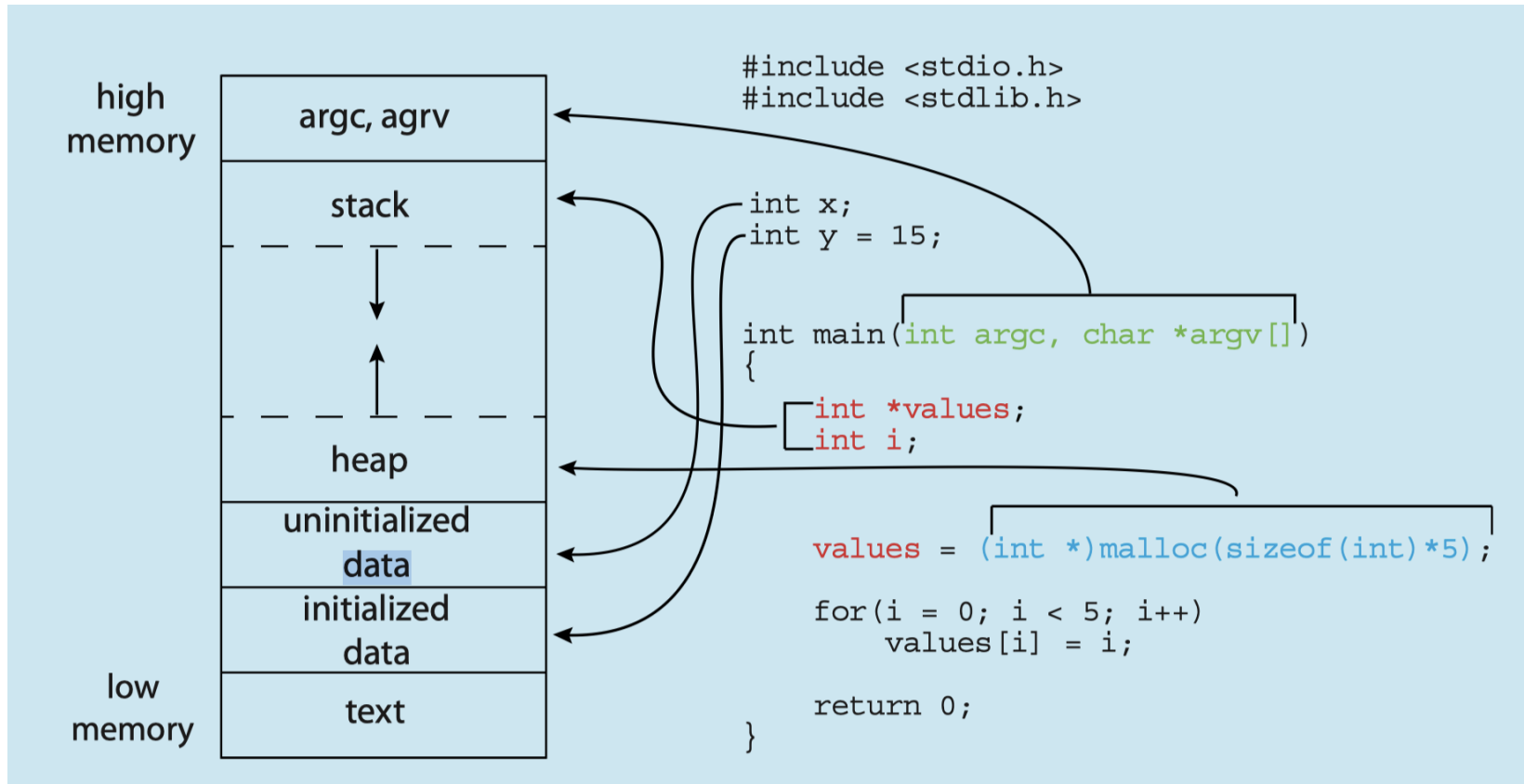


Process in Memory





Process in Memory





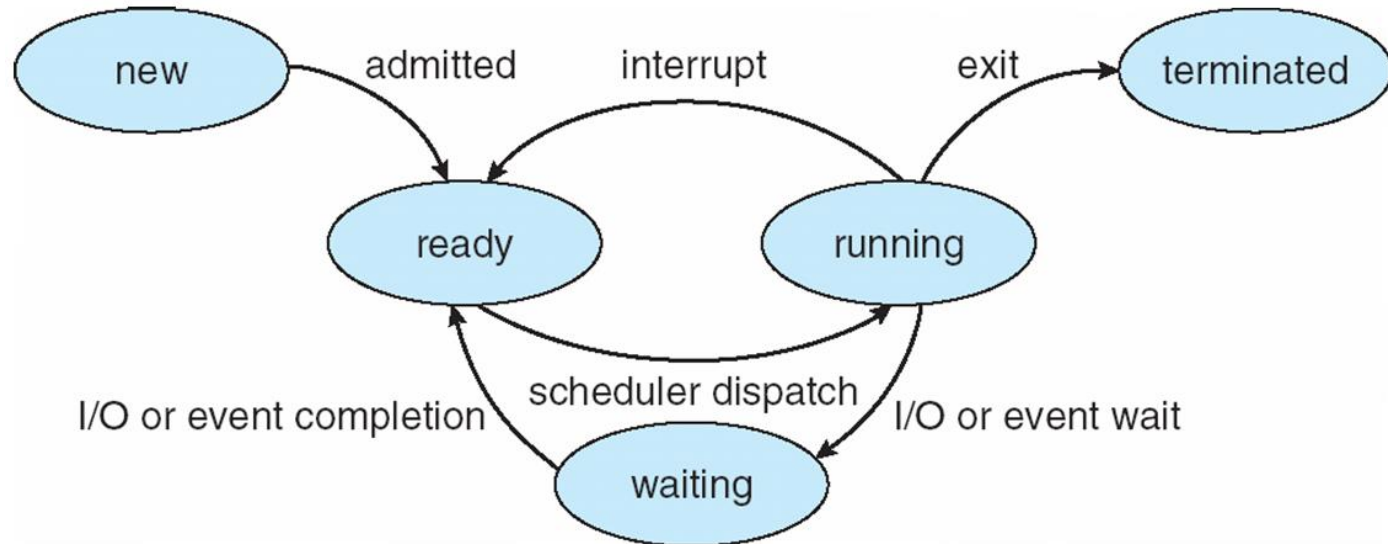
Process State

- As a process executes, it changes **state**
 - **new**: The process is being created
 - **running**: Instructions are being executed
 - **waiting**: The process is waiting for some event to occur
 - **ready**: The process is waiting to be assigned to a processor
 - **terminated**: The process has finished execution





Diagram of Process State

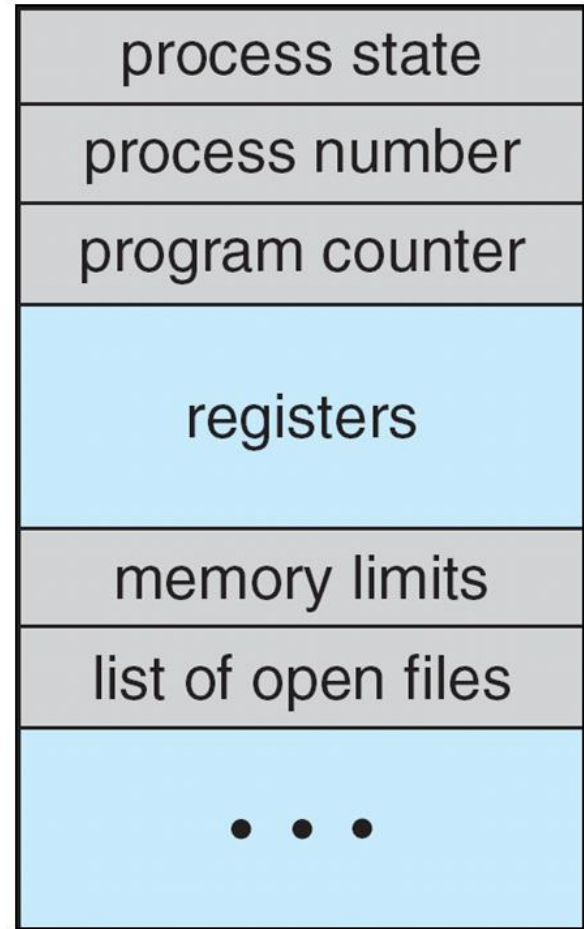




Process Control Block (PCB)

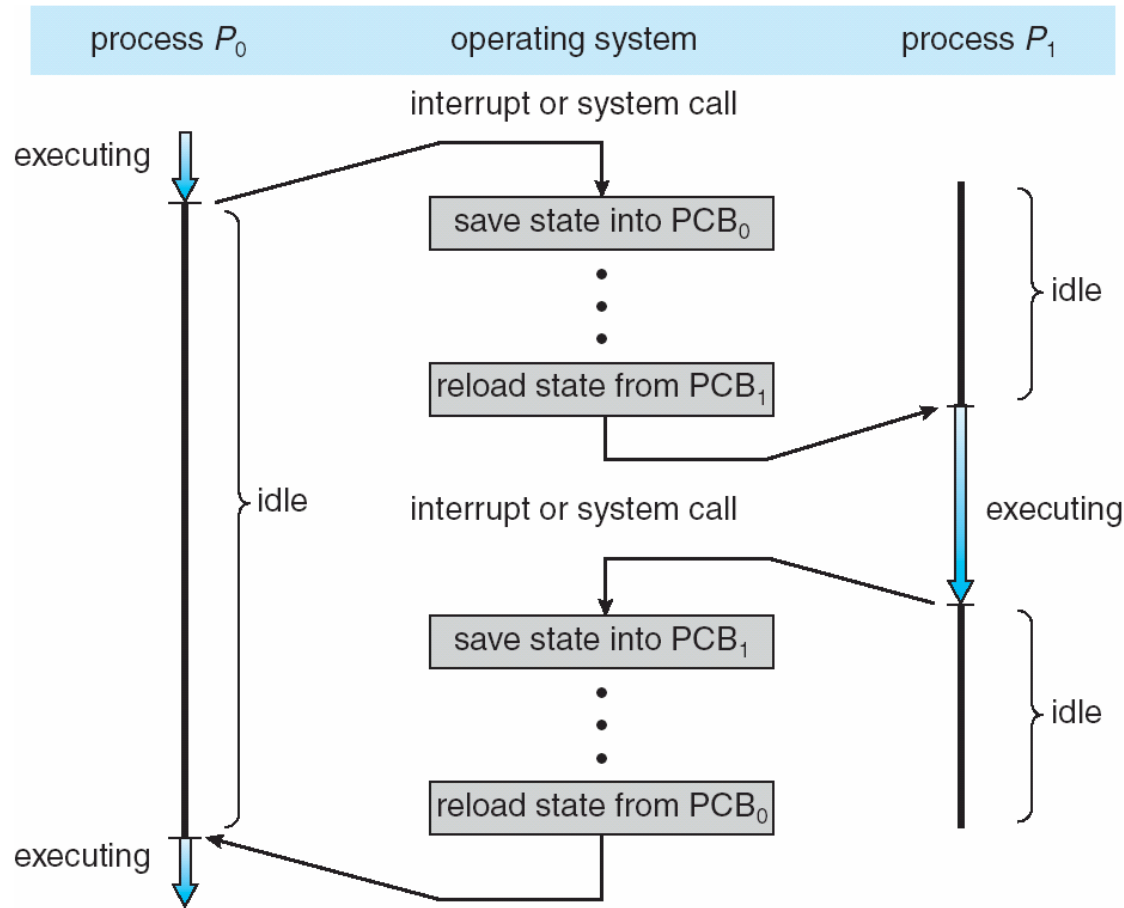
Information associated with each process
(also called **task control block**)

- ❑ Process state – running, waiting, etc
- ❑ Program counter – location of instruction to next execute
- ❑ CPU registers – contents of all process-centric registers
- ❑ CPU scheduling information- priorities, scheduling queue pointers
- ❑ Memory-management information – memory allocated to the process
- ❑ Accounting information – CPU used, clock time elapsed since start, time limits
- ❑ I/O status information – I/O devices allocated to process, list of open files





CPU Switch From Process to Process





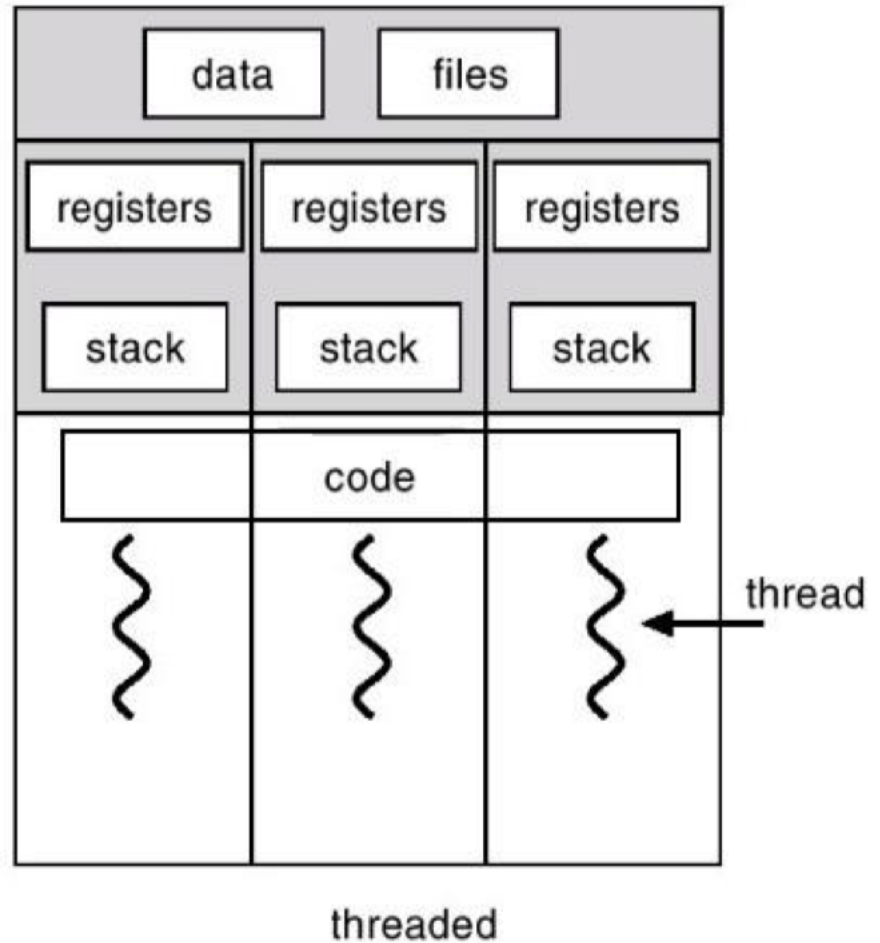
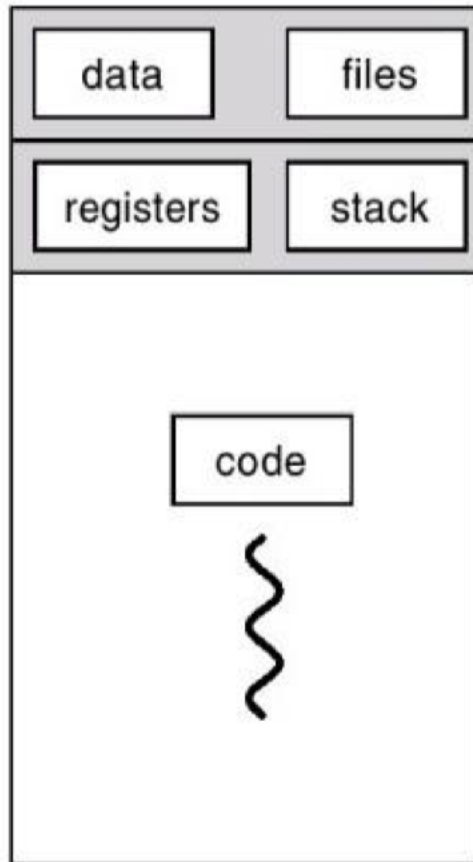
Threads

- ❑ So far, process has a single thread of execution
- ❑ Consider having multiple program counters per process
 - ❑ Multiple locations can execute at once
 - ▶ Multiple threads of control -> **threads**
- ❑ Must then have storage for thread details, multiple program counters in PCB





Threads

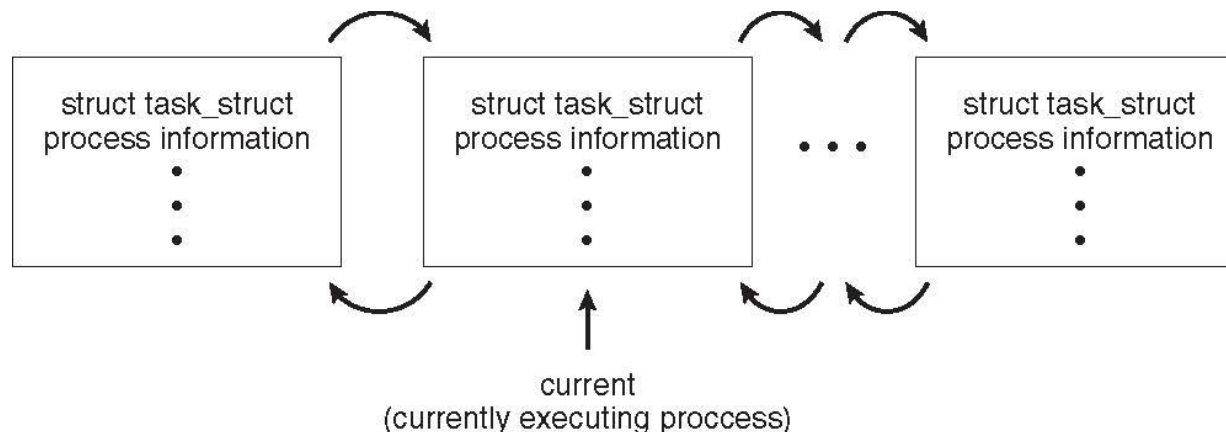




Process Representation in Linux

Represented by the C structure `task_struct`

```
pid t_pid; /* process identifier */
long state; /* state of the process */
unsigned int time_slice /* scheduling information */
struct task_struct *parent; /* this process's parent */
struct list_head children; /* this process's children */
struct files_struct *files; /* list of open files */
struct mm_struct *mm; /* address space of this process */
```





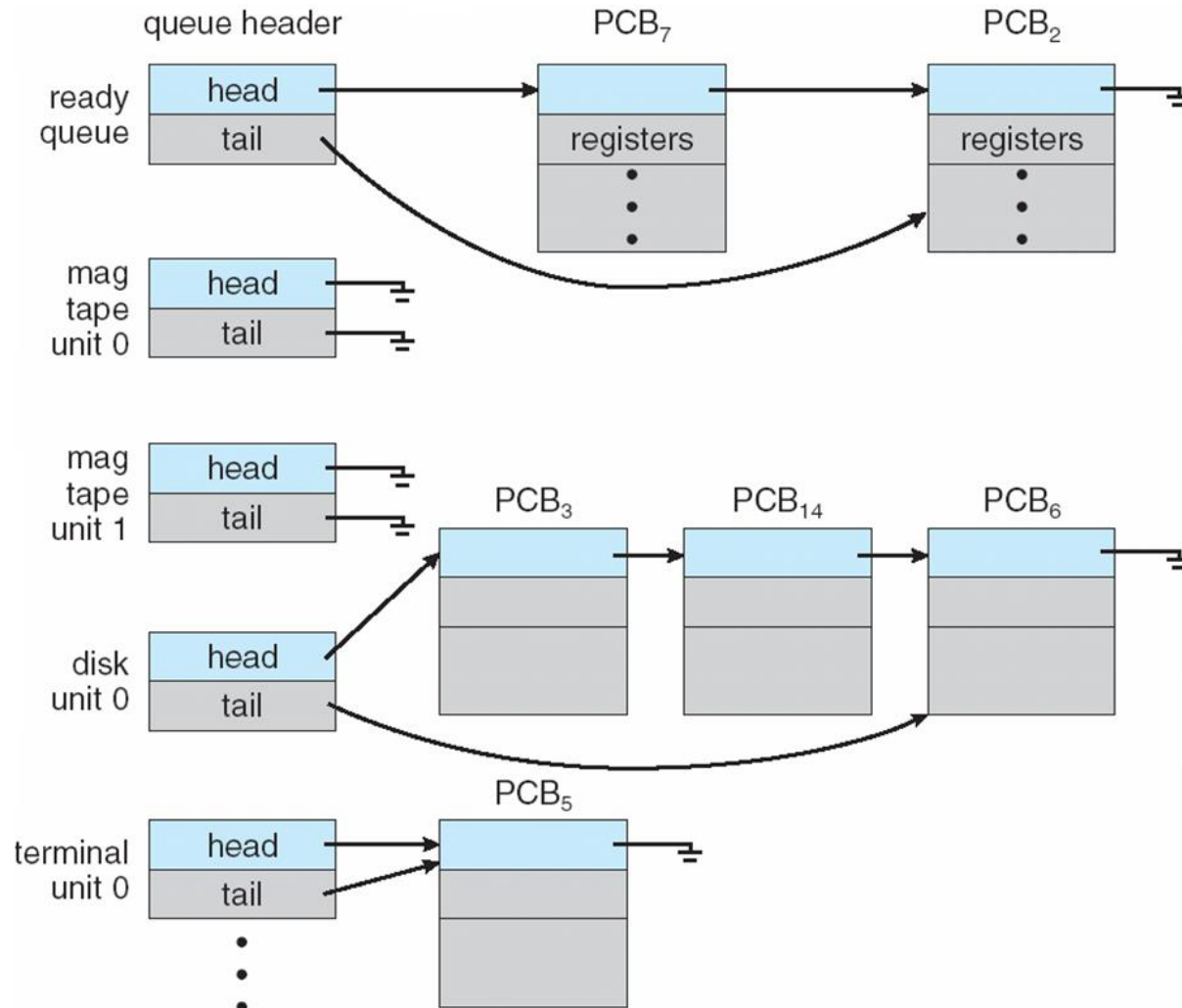
Process Scheduling

- Maximize CPU use, quickly switch processes onto CPU for time sharing
- **Process scheduler** selects among available processes for next execution on CPU
- Maintains **scheduling queues** of processes
 - **Job queue** – set of all processes in the system
 - **Ready queue** – set of all processes residing in main memory, ready and waiting to execute
 - **Device queues** – set of processes waiting for an I/O device
 - Processes migrate among the various queues





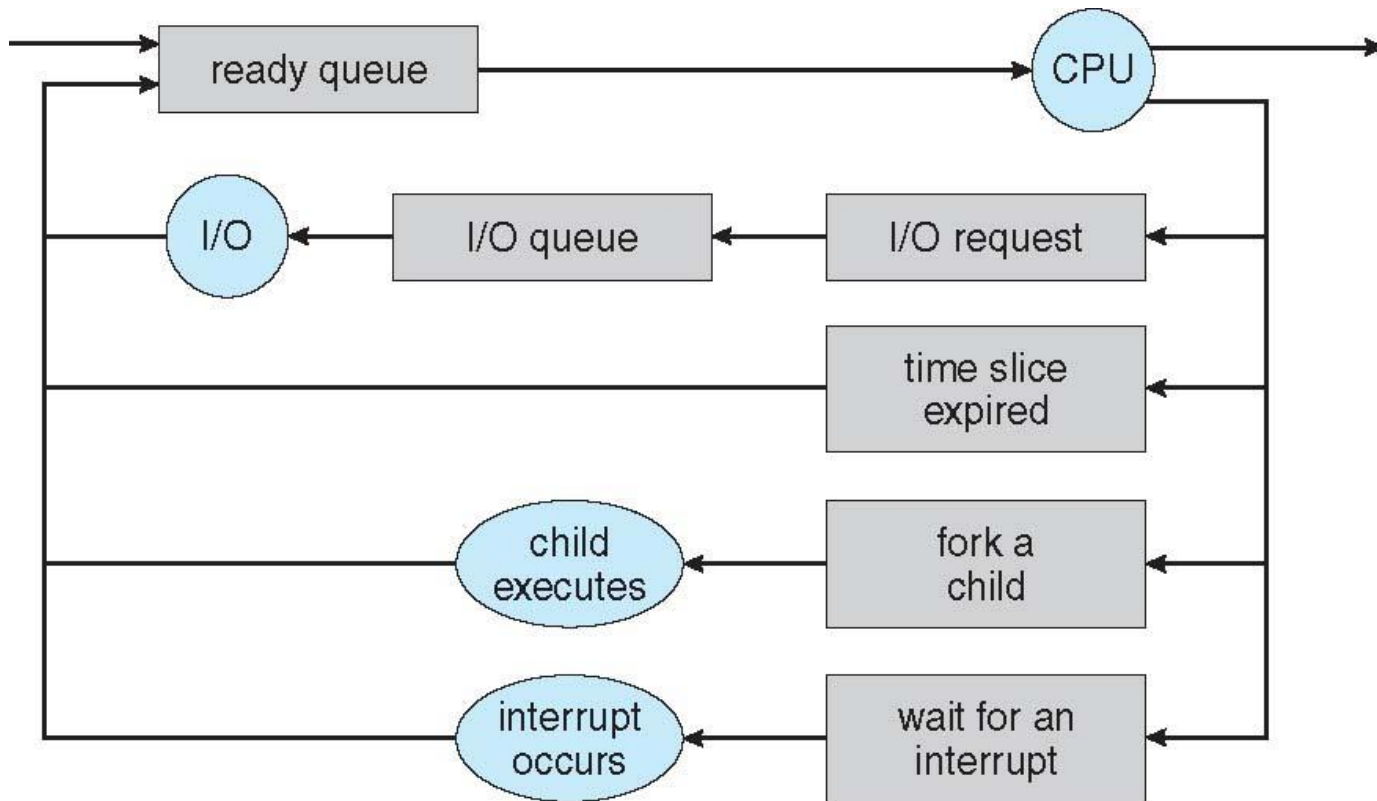
Ready Queue And Various I/O Device Queues





Representation of Process Scheduling

- **Queueing diagram** represents queues, resources, flows





Schedulers

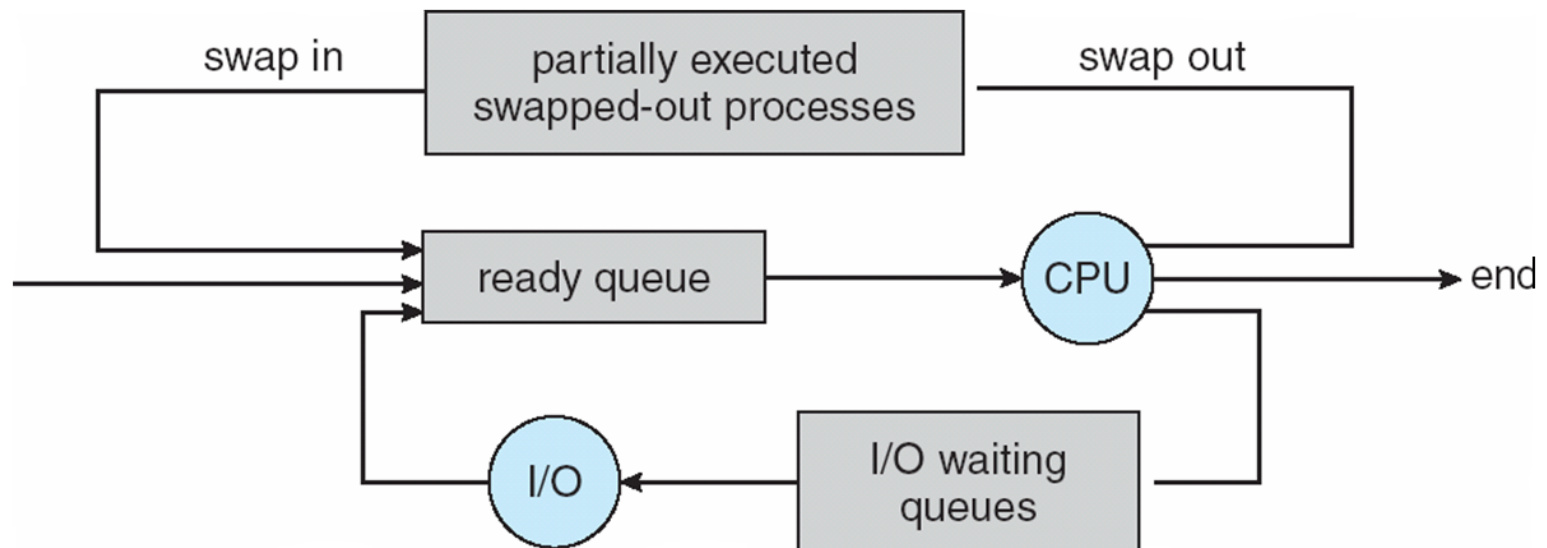
- **Short-term scheduler** (or **CPU scheduler**) – selects which process should be executed next and allocates CPU
 - Sometimes the only scheduler in a system
 - Short-term scheduler is invoked frequently (milliseconds) \Rightarrow (must be fast)
- **Long-term scheduler** (or **job scheduler**) – selects which processes should be brought into the ready queue
 - Long-term scheduler is invoked infrequently (seconds, minutes) \Rightarrow (may be slow)
 - The long-term scheduler controls the **degree of multiprogramming**
- Processes can be described as either:
 - **I/O-bound process** – spends more time doing I/O than computations, many short CPU bursts
 - **CPU-bound process** – spends more time doing computations; few very long CPU bursts
- Long-term scheduler strives for good ***process mix***





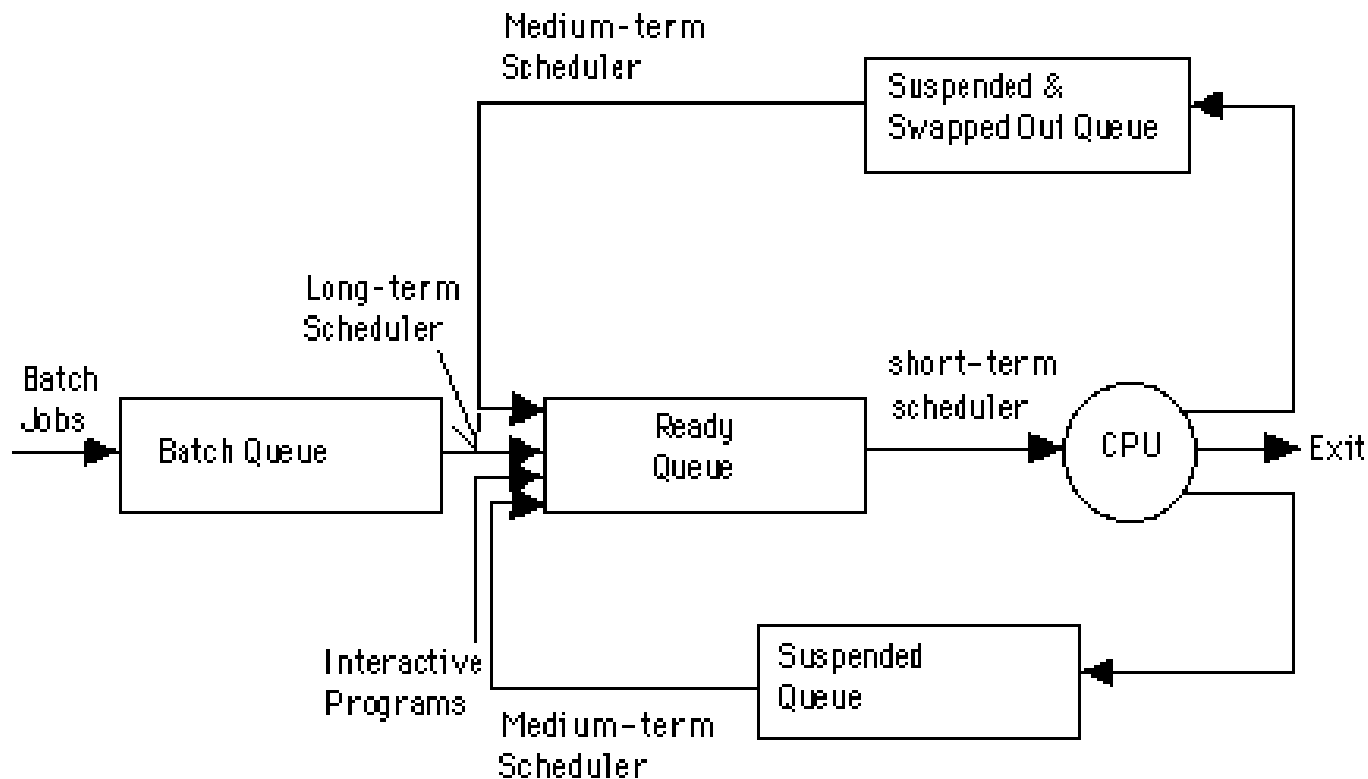
Addition of Medium Term Scheduling

- **Medium-term scheduler** can be added if degree of multiple programming needs to decrease
 - Remove process from memory, store on disk, bring back in from disk to continue execution: **swapping**





Addition of Medium Term Scheduling





Multitasking in Mobile Systems

- ❑ Some mobile systems (e.g., early version of iOS) allow only one process to run, others suspended
- ❑ Due to screen real estate, user interface limits, iOS provides for a
 - ❑ Single **foreground** process- controlled via user interface
 - ❑ Multiple **background** processes– in memory, running, but not on the display, and with limits
 - ❑ Limits include single, short task, receiving notification of events, specific long-running tasks like audio playback
- ❑ Android runs foreground and background, with fewer limits
 - ❑ Background process uses a **service** to perform tasks
 - ❑ Service can keep running even if background process is suspended
 - ❑ Service has no user interface, small memory use





Context Switch

- ❑ When CPU switches to another process, the system must **save the state** of the old process and load the **saved state** for the new process via a **context switch**
- ❑ **Context** of a process represented in the PCB
- ❑ Context-switch time is overhead; the system does no useful work while switching
 - ❑ The more complex the OS and the PCB → the longer the context switch
- ❑ Time dependent on hardware support
 - ❑ Some hardware provides multiple sets of registers per CPU → multiple contexts loaded at once





Operations on Processes

- System must provide mechanisms for:
 - process creation,
 - process termination,
 - and so on as detailed next





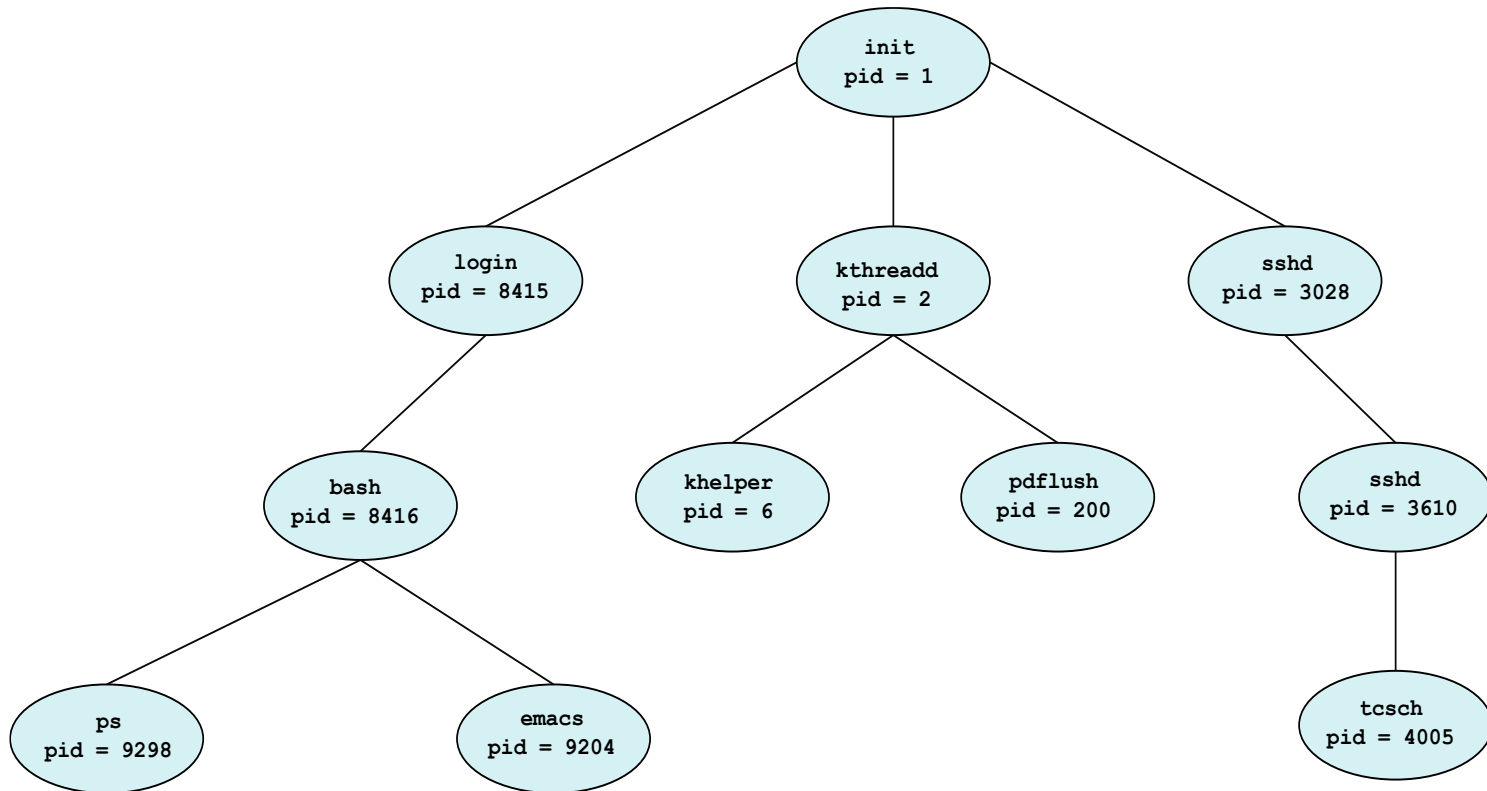
Process Creation

- ❑ **Parent** process create **children** processes, which, in turn create other processes, forming a **tree** of processes
- ❑ Generally, process identified and managed via a **process identifier (pid)**
- ❑ Resource sharing options
 - ❑ Parent and children share all resources
 - ❑ Children share subset of parent's resources
 - ❑ Parent and child share no resources
- ❑ Execution options
 - ❑ Parent and children execute concurrently
 - ❑ Parent waits until children terminate





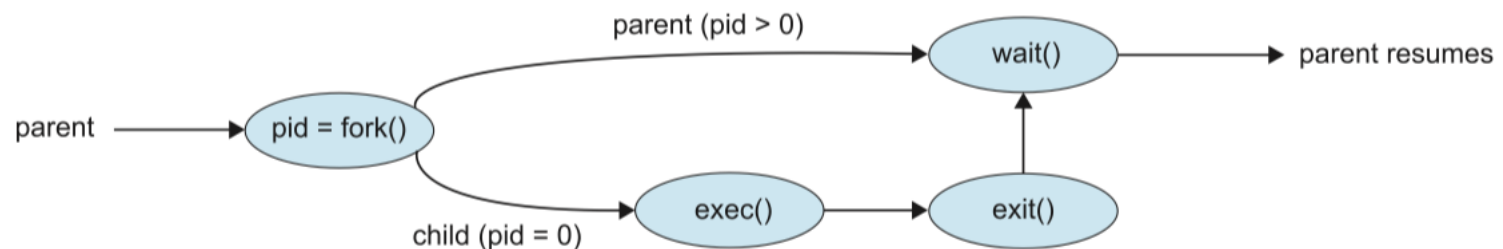
A Tree of Processes in Linux





Process Creation (Cont.)

- Address space options
 - Child duplicate of parent
 - Child has a program loaded into it
- UNIX examples
 - **fork()** system call creates new process
 - **exec()** system call used after a **fork()** to replace the process' memory space with a new program





C Program Forking Separate Process

```
#include <sys/types.h>
#include <stdio.h>
#include <unistd.h>

int main()
{
    pid_t pid;

    /* fork a child process */
    pid = fork();

    if (pid < 0) { /* error occurred */
        fprintf(stderr, "Fork Failed");
        return 1;
    }
    else if (pid == 0) { /* child process */
        execlp("/bin/ls", "ls", NULL);
    }
    else { /* parent process */
        /* parent will wait for the child to complete */
        wait(NULL);
        printf("Child Complete");
    }

    return 0;
}
```





Process Termination

- Process executes last statement and then asks the operating system to delete it using the **exit()** system call.
 - Returns status data from child to parent (via **wait()**)
 - Process' resources are deallocated by operating system
- Parent may terminate the execution of children processes using the **abort()** system call. Some reasons for doing so:
 - Child has exceeded allocated resources
 - Task assigned to child is no longer required
 - The parent is exiting and the operating systems does not allow a child to continue if its parent terminates





Process Termination

- Some operating systems do not allow child to exist if its parent has terminated. If a process terminates, then all its children must also be terminated.
 - **cascading termination.** All children, grandchildren, etc. are terminated.
 - The termination is initiated by the operating system.
 - The parent process may wait for termination of a child process by using the `wait()` system call. The call returns status information and the pid of the terminated process
- ```
pid = wait(&status);
```
- If no parent waiting (did not invoke `wait()`) process is a **zombie**
  - If parent terminated without invoking `wait`, process is an **orphan**





# Multiprocess Architecture – Chrome Browser

- Many web browsers ran as single process (some still do)
  - If one web site causes trouble, entire browser can hang or crash
- Google Chrome Browser is multiprocess with 3 different types of processes:
  - **Browser** process manages user interface, disk and network I/O
  - **Renderer** process renders web pages, deals with HTML, Javascript. A new renderer created for each website opened
    - ▶ Runs in **sandbox** restricting disk and network I/O, minimizing effect of security exploits
  - **Plug-in** process for each type of plug-in





# Interprocess Communication

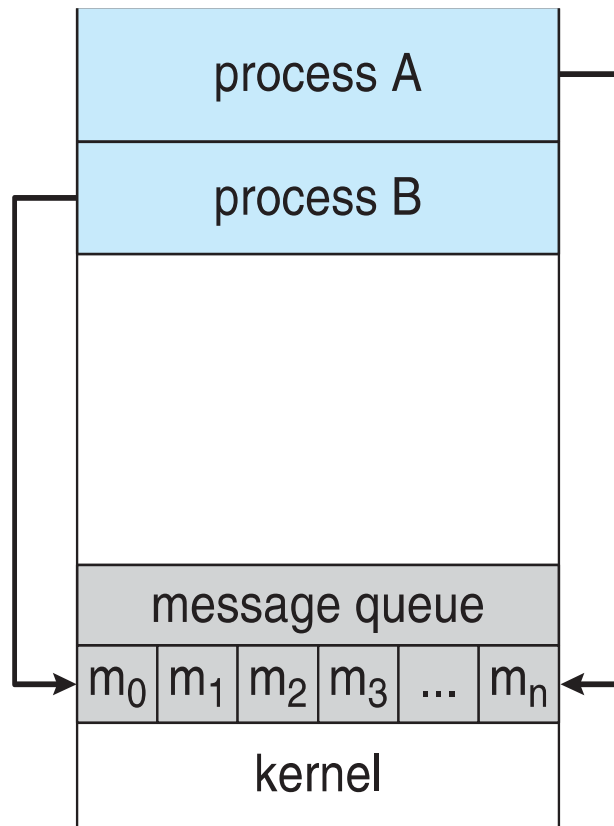
- Processes within a system may be ***independent*** or ***cooperating***
- ***Cooperating*** process can affect or be affected by other processes, including sharing data
- ***Independent*** process cannot affect or be affected by the execution of another process
- Reasons for cooperating processes:
  - Information sharing
  - Computation speedup
  - Modularity
- Cooperating processes need **interprocess communication (IPC)**
- Two models of IPC
  - **Shared memory**
  - **Message passing**



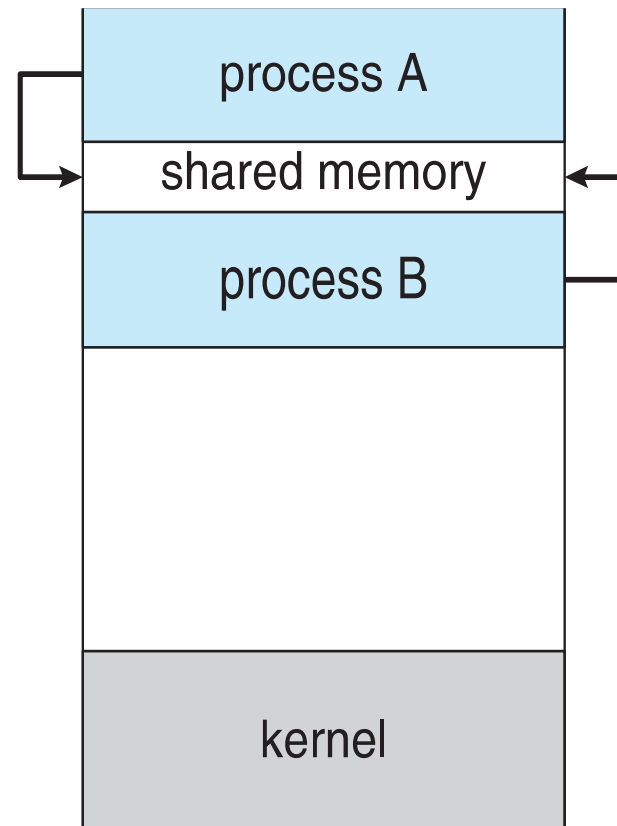


# Communications Models

(a) Message passing. (b) shared memory.



(a)



(b)





# Producer-Consumer Problem

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- Paradigm for cooperating processes, *producer* process produces information that is consumed by a *consumer* process
  - **unbounded-buffer** places no practical limit on the size of the buffer
  - **bounded-buffer** assumes that there is a fixed buffer size





# Bounded-Buffer – Shared-Memory Solution

## □ Shared data

```
#define BUFFER_SIZE 10
typedef struct {
 . . .
} item;

item buffer[BUFFER_SIZE];
int in = 0;
int out = 0;
```

## □ Solution is correct, but can only use BUFFER\_SIZE-1 elements







# Bounded-Buffer – Producer

---

```
item next_produced;
while (true) {
 /* produce an item in next produced */
 while (((in + 1) % BUFFER_SIZE) == out)
 ; /* do nothing */
 buffer[in] = next_produced;
 in = (in + 1) % BUFFER_SIZE;
}
```





# Bounded Buffer – Consumer

---

```
item next_consumed;
while (true) {
 while (in == out)
 ; /* do nothing */
 next_consumed = buffer[out];
 out = (out + 1) % BUFFER_SIZE;

 /* consume the item in next consumed */
}
```





# Interprocess Communication – Shared Memory

---

- ❑ An area of memory shared among the processes that wish to communicate
- ❑ The communication is under the control of the users processes not the operating system.
- ❑ Major issues is to provide mechanism that will allow the user processes to synchronize their actions when they access shared memory.
- ❑ Synchronization is discussed in great details in Chapter 5.





# Interprocess Communication – Message Passing

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- Mechanism for processes to communicate and to synchronize their actions
- Message system – processes communicate with each other without resorting to shared variables
- IPC facility provides two operations:
  - **send**(*message*)
  - **receive**(*message*)
- The *message* size is either fixed or variable





## Message Passing (Cont.)

---

- If processes  $P$  and  $Q$  wish to communicate, they need to:
  - Establish a **communication link** between them
  - Exchange messages via send/receive
- Implementation issues:
  - How are links established?
  - Can a link be associated with more than two processes?
  - How many links can there be between every pair of communicating processes?
  - What is the capacity of a link?
  - Is the size of a message that the link can accommodate fixed or variable?
  - Is a link unidirectional or bi-directional?





# Message Passing (Cont.)

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- Implementation of communication link
  - Physical:
    - ▶ Shared memory
    - ▶ Hardware bus
    - ▶ Network
  - Logical:
    - ▶ Direct or indirect
    - ▶ Synchronous or asynchronous
    - ▶ Automatic or explicit buffering





# Direct Communication

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- Processes must name each other explicitly:
  - **send** ( $P$ , *message*) – send a message to process  $P$
  - **receive**( $Q$ , *message*) – receive a message from process  $Q$
- Properties of communication link
  - Links are established automatically
  - A link is associated with exactly one pair of communicating processes
  - Between each pair there exists exactly one link
  - The link may be unidirectional, but is usually bi-directional





# Indirect Communication

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- ❑ Messages are directed and received from mailboxes (also referred to as ports)
  - ❑ Each mailbox has a unique id
  - ❑ Processes can communicate only if they share a mailbox
- ❑ Properties of communication link
  - ❑ Link established only if processes share a common mailbox
  - ❑ A link may be associated with many processes
  - ❑ Each pair of processes may share several communication links
  - ❑ Link may be unidirectional or bi-directional







# Indirect Communication

---

- Operations
  - create a new mailbox (port)
  - send and receive messages through mailbox
  - destroy a mailbox
- Primitives are defined as:
  - send**(*A, message*) – send a message to mailbox A
  - receive**(*A, message*) – receive a message from mailbox A





# Indirect Communication

- Mailbox sharing
  - $P_1$ ,  $P_2$ , and  $P_3$  share mailbox A
  - $P_1$ , sends;  $P_2$  and  $P_3$  receive
  - Who gets the message?
- Solutions
  - Allow a link to be associated with at most two processes
  - Allow only one process at a time to execute a receive operation
  - Allow the system to select arbitrarily the receiver. Sender is notified who the receiver was.





# Synchronization

- ❑ Message passing may be either blocking or non-blocking
- ❑ **Blocking** is considered **synchronous**
  - ❑ **Blocking send** -- the sender is blocked until the message is received
  - ❑ **Blocking receive** -- the receiver is blocked until a message is available
- ❑ **Non-blocking** is considered **asynchronous**
  - ❑ **Non-blocking send** -- the sender sends the message and continue
  - ❑ **Non-blocking receive** -- the receiver receives:
    - ❑ A valid message, or
    - ❑ Null message
- ❑ Different combinations possible
  - ❑ If both send and receive are blocking, we have a **rendezvous**





# Synchronization (Cont.)

---

## □ Producer-consumer becomes trivial

```
message next_produced;
while (true) {
 /* produce an item in next produced */
 send(next_produced);
}
```

```
message next_consumed;
while (true) {
 receive(next_consumed);

 /* consume the item in next consumed */
}
```





# Buffering

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- Queue of messages attached to the link.
- implemented in one of three ways
  1. Zero capacity – no messages are queued on a link.  
Sender must wait for receiver (rendezvous)
  2. Bounded capacity – finite length of  $n$  messages  
Sender must wait if link full
  3. Unbounded capacity – infinite length  
Sender never waits





# Communications in Client-Server Systems

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- Sockets
- Remote Procedure Calls
- Pipes
- Remote Method Invocation (Java)





# Sockets

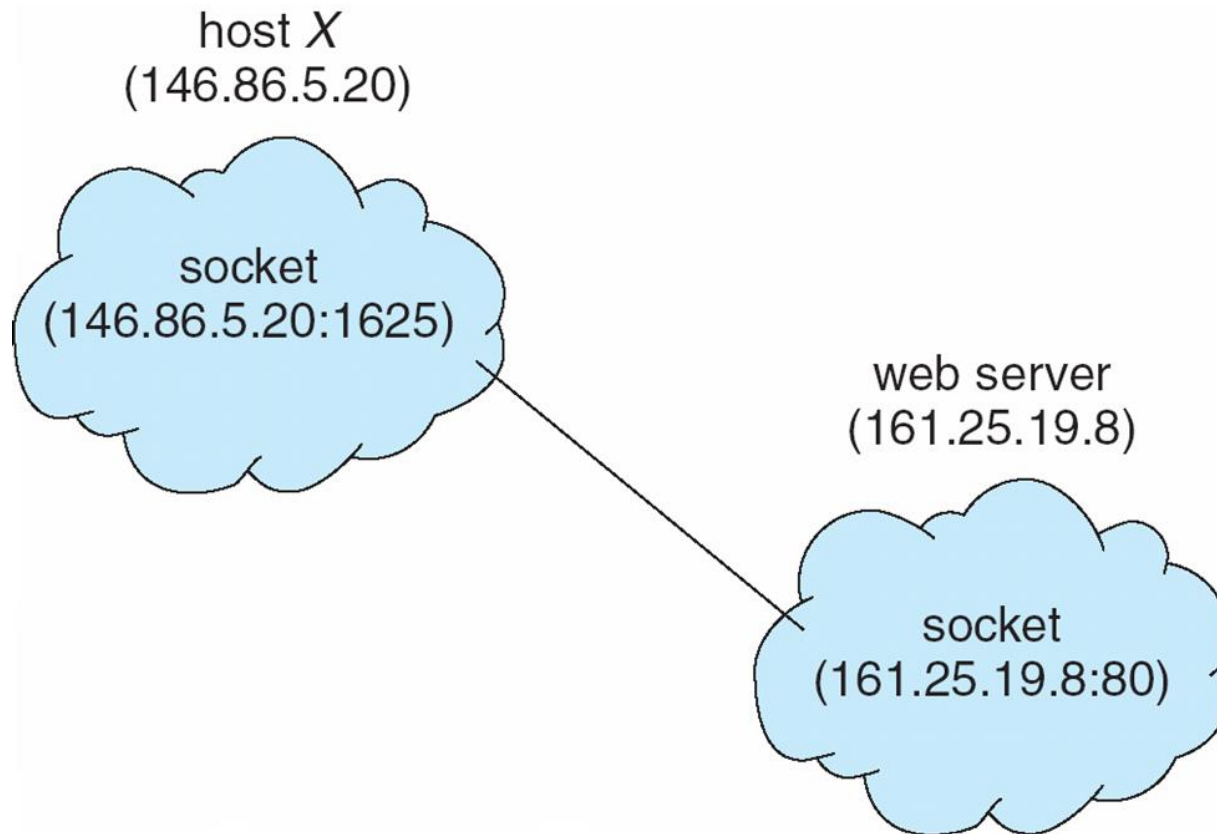
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- A **socket** is defined as an endpoint for communication
- Concatenation of IP address and **port** – a number included at start of message packet to differentiate network services on a host
- The socket **161.25.19.8:1625** refers to port **1625** on host **161.25.19.8**
- Communication consists between a pair of sockets
- All ports below 1024 are **well known**, used for standard services
- Special IP address 127.0.0.1 (**loopback**) to refer to system on which process is running





# Socket Communication







# Pipes

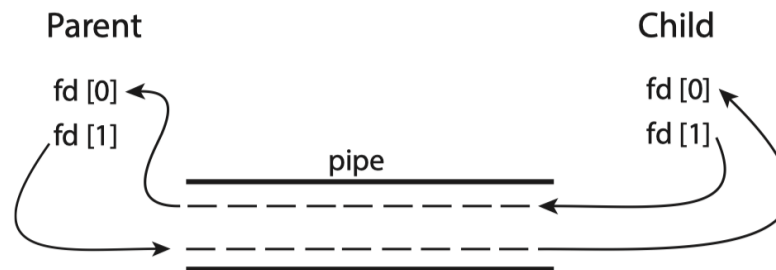
- Acts as a conduit allowing two processes to communicate
- Issues:
  - Is communication unidirectional or bidirectional?
  - In the case of two-way communication, is it half or full-duplex?
  - Must there exist a relationship (i.e., **parent-child**) between the communicating processes?
  - Can the pipes be used over a network?
- Ordinary pipes – cannot be accessed from outside the process that created it. Typically, a parent process creates a pipe and uses it to communicate with a child process that it created.
- Named pipes – can be accessed without a parent-child relationship.





# Ordinary Pipes

- ❑ Ordinary Pipes allow communication in standard producer-consumer style
- ❑ Producer writes to one end (the **write-end** of the pipe)
- ❑ Consumer reads from the other end (the **read-end** of the pipe)
- ❑ Ordinary pipes are therefore unidirectional
- ❑ Require parent-child relationship between communicating processes



- ❑ Windows calls these **anonymous pipes**
- ❑ See Unix and Windows code samples in textbook





# Named Pipes

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- ❑ Named Pipes are more powerful than ordinary pipes
- ❑ Communication is bidirectional
- ❑ No parent-child relationship is necessary between the communicating processes
- ❑ Several processes can use the named pipe for communication
- ❑ Provided on both UNIX and Windows systems



# End of Chapter 3

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