

## Recursive Formulation

Define  $f(k) = F(k, 1) + (1-\delta)k$   
 $\Rightarrow \sum_{t=0}^{\infty} \beta^t u(f(k_t) - k_{t+1}) \rightarrow \max_{\{k_{t+1}\}_{t=0}^{\infty}}$

s.t.  $\begin{cases} f(k_t) \geq k_{t+1} \geq 0 & \forall t \\ k_0 \text{ given} \end{cases}$

$$\max_{\substack{\{k_{t+1}\}_{t=0}^{\infty} \\ 0 \leq k_{t+1} \leq f(k_t)}} \sum_{t=0}^{\infty} \beta^t u(f(k_t) - k_{t+1}) =$$

$$= \max_{\{k_{t+1}\}_{t=0}^{\infty}} \left[ u(f(k_0) - k_1) + \beta \sum_{t=1}^{\infty} \beta^{t-1} (u(f(k_t) - k_{t+1})) \right]$$

$$= \max_{k_1, 0 \leq k_1 \leq f(k_0)} \left[ u(f(k_0) - k_1) + \beta \cdot \max_{\{k_{t+1}\}_{t=1}^{\infty}} \left( \sum_{t=1}^{\infty} \beta^{t-1} (u(f(k_t) - k_{t+1})) \right) \right]$$

$$0 \leq k_{t+1} \leq f(k_t)$$

$$= \max_{0 \leq k_1 \leq f(k_0)} \left[ u(f(k_0) - k_1) + \beta \cdot \max_{\{k_{t+2}\}_{t=0}^{\infty}} \left( \sum_{t=0}^{\infty} \beta^t (f(k_{t+1}) - k_{t+2}) \right) \right]$$

$$U(k_0) = \max_{\{k_{t+1}\}_{t=0}^{\infty}} \sum_{t=0}^{\infty} \beta^t u(f(k_t) - k_{t+1})$$

$$= \max_{\{k_1\}} [u(f$$

21 September

Sequential Problem of SP

$$\sum_{t=0}^{\infty} \beta^t u(f(k_t) - k_{t+1}) \rightarrow \max_{\{k_{t+1}\}_{t=0}^{\infty}}$$

s.t.  $f(k_t) \geq k_{t+1} \geq 0$   
 $k_0$  is given.

Recursive Problem of the SP

$$v(k_0) = \max_{0 \leq k_1 \leq f(k_0)} [u(f(k_0) - k_1) + \beta v(k_1)]$$



$k_0$  is given.

$k_1$  that maximizes  
utility for pd I / all the future

The problem is we don't know  $v(\cdot)$ . yet.

Jargon:

$v(\cdot)$  : the value function

$k$  : state variable (if you know this, you

$k'$  : choice variable - know everything)

(or control)  $\rightarrow$  in the example above

$$\begin{pmatrix} k \rightarrow k_0 \\ k' \rightarrow k_1 \end{pmatrix}$$

$k' = g(k) \rightarrow$  Policy function (decision rule)

Algorithm to solve the Bellman Equation

(Value function iteration) :

1) "Guess"  $v_0(k)$

2) Use BE to update your guess

$$v_{j+1}(k) = \max_{0 \leq k' \leq f(k)} [u(f(k) - k') + \beta v_j(k')]$$

3) If  $\|v_{j+1}(k) - v_j(k)\| < \epsilon$ ,

then  $v_{j+1}(k)$  is the solution.

Assume:

$$u = \log(c)$$

$$\delta = 1$$

(by definition)

$$y_t = \theta k_t^\alpha$$

$$v(k) = \max_{0 \leq k' \leq \theta k^\alpha} [\log(c) + \beta \cdot v(k')]$$

s.t.  $c + k' = \theta k^\alpha$

$k_0$  is given

Guess (and later we will verify), that

$$v(k) = a_0 + a_1 \log k$$

$$L = \log(c) + \beta(a_0 + a_1 \log k') + \lambda [\theta k^\alpha - c - k']$$

$$\frac{\partial L}{\partial c} : \frac{1}{c} = \lambda \quad \frac{\partial L}{\partial k} : \frac{\cancel{\theta a_1}}{k'} =$$

$$\frac{1}{c} = \frac{\cancel{\beta a_1}}{k'}$$

$$\Rightarrow \frac{1}{\theta k^\alpha - k'} = \frac{\beta a_1}{k'}$$

$$\Rightarrow \beta a_1 \theta k^\alpha - \beta a_1 k' = k'$$

$$\Rightarrow k' (1 + \beta a_1) = \beta a_1 \theta k^\alpha$$

$$\Rightarrow \boxed{k' = \frac{\beta a_1}{1 + \beta a_1} \theta k^\alpha}$$

$$c = \theta k^\alpha - k' = \frac{\theta k^\alpha - \beta a_1}{1 + \beta a_1} \theta k^\alpha$$

$$= \theta k^\alpha \left[ 1 - \frac{\beta a_1}{1 + \beta a_1} \right]$$

$$\boxed{c = \theta k^\alpha \cdot \left( \frac{1}{1 + \beta a_1} \right)}$$

Method of indeterminate coeffs:

$$v(k) = \max [ \log(c) + \beta v(k') ]$$

$$a_0 + a_1 \log(k) = \log \left( \frac{\theta k^\alpha}{1 + \beta a_1} \right) + \beta$$

$$+ \beta \cdot \left( a_0 + a_1 \log \left[ \frac{\theta k^\alpha - \beta a_1}{1 + \beta a_1} \right] \right)$$

$$a_0 + a_1 \log k = -\log(1 + \beta a_1) + \log \theta + \alpha \log(k) + \beta a_0 - \beta$$

$$- \beta a_1 \log(1 + \beta a_1) + \beta a_1 \log(\theta k^\alpha) + \beta a_1 \alpha \log(k)$$

$$a_1 = \alpha + \beta a_1 \alpha \iff a_1 (1 - \beta \alpha) = \alpha$$

$$\Rightarrow \boxed{a_1 = \frac{\alpha}{1-\beta\alpha}}$$

## Decentralizing Equilibrium

Representative agent,  $v(c)$

Firm with technology  $F(k, l)$

Agent owns the firm

Agent is endowed with  $k_0$ .

Def:

An ADE is allocations for the consumer

$\{(c_t, k_{t+1}, l_t)\}_{t=0}^{\infty}$ , allocation for firm  $\{y_t, k_t^f, l_t^f\}_{t=0}^{\infty}$   
and a price system  $\{\hat{p}_t, \hat{w}_t, \hat{i}_t\}$  s.t.

1) given the prices,  $\{\hat{c}_t, \hat{k}_{t+1}, \hat{p}_t\}$  solves

$$\max \sum_{t=0}^{\infty} \beta^t u(c_t)$$

$$\text{s.t. } \sum_{t=0}^{\infty} (\hat{p}_t c_t + \hat{p}_t \cdot i_t) \leq \sum_{t=0}^{\infty} \hat{w}_t l_t + \hat{\delta}_t k_t$$

$$i_t = k_{t+1} - (1-\delta)k_t$$

$$c_t \geq 0, k_t \geq 0 \quad \forall t$$

2) Given prices, the firm chooses optimally  $\{y_t, k_t^f, l_t^f\}$   
 s.t.  $p_t^f(k_t^f, l_t^f) - \hat{w}_t l_t^f - \hat{\gamma}_t k_t \rightarrow \max$

3) Markets clear.

$$\hat{c}_t + \hat{i}_t = \hat{y}_t$$

$$\hat{k}_t^f = \hat{k}_t$$

$$\hat{l}_t^f = \hat{l}_t$$

"all the"

Suppose hh sell  $\hat{k}$  capital to the firm in the first time period.  $t=1$ .

Hh:

$$\sum_{t=0}^{\infty} \hat{p}_t c_t \leq \sum_{t=0}^{\infty} \hat{w}_t l_t + \gamma_0 k_0 + p_0 (1-\delta) k_0$$

The firm:

$$\max \sum_{t=0}^{\infty} \hat{p}_t (F(k_t, l_t) - i_t) - \gamma_0 k_0 - p_0 (1-\delta) k - \sum_{t=0}^{\infty} \hat{w}_t l_t$$

$$i_t = k_{t+1} - (1-\delta)k_t$$

Def: An SME is allocations  $\{\tilde{c}_t, \tilde{b}_{t+1}, \tilde{k}_{t+1}, \tilde{l}_t\}$ ,  
 allocation for firms  $\{\tilde{y}_t^f, \tilde{k}_t^f, \tilde{l}_t^f\}$ , prices  
 $\{\tilde{w}_t, \tilde{x}_t^k, \tilde{x}_t^b\}$ . s.t.

1) Given prices, HHP solves:

$$\max_{\{c_t, l_t, b_{t+1}, k_{t+1}\}} V(\{c_t\})$$

s.t.

$$c_t + b_{t+1} + i_t = \tilde{w}_t l_t + \tilde{x}_t^k k_t + (1 + \tilde{x}_t^b) b_t$$

$$i_t = k_{t+1} - (1 - \delta) b_t$$

→  
 convention  
 write  $\underline{(1+\delta)}$   
 for interest  
 rates from  
 bond.

Tildas represent numbers we already have. Variables w/o tildas are the ones we are choosing.

2) Firms maximize profits

$$y_t - \tilde{w}_t l_t^f - \tilde{x}_t^k k_t^f \rightarrow \max$$

$$\text{s.t. } y_t \leq F(k_t^f, l_t^f)$$

3) Markets clear

$$\tilde{c}_t + \tilde{i}_t = y_t$$

$$\tilde{k}_t^f = \tilde{k}_t$$

$$\tilde{l}_t^f = \tilde{l}_t$$

$\tilde{b}_t = 0$

$\left\{ \begin{array}{l} \text{we still include} \\ \text{bonds even though} \\ \text{it is 0, to find the} \\ \text{equilibrium price} \\ \text{of assets to convince} \\ \text{HHS to not buy bonds} \end{array} \right.$

[Sep 26]

### Social Planner's Problem

$$\sum_{t=0}^{\infty} \beta^t u(c_t) \rightarrow \max$$

$$\text{s.t. } c_t + k_{t+1} - (1-\delta)k_t = y_t$$

$$c_t \geq 0, k_{t+1} \geq 0, k_0 \text{ is given}$$

Euler Equation:

$$u'_c(t) = \beta u'_c(t+1) [1 - \delta + f'_k(k_{t+1})]$$

\* Consumer (in SMO) seeks  $\{c_t, k_{t+1}, b_{t+1}\}_{t=0}^{\infty}$  to solve  
the following problem

$$\sum_{t=0}^{\infty} \beta^t u(c_t, l_t) \rightarrow \max$$

$$\text{s.t. } c_t + k_{t+1} + b_{t+1} \leq w_t l_t + (1+r_t^b) b_t + (1-\delta)k_t + r_t^k k_t$$

$$\mathcal{L} = \sum_{t=0}^{\infty} \beta^t u(c_t, l_t) + \sum_{t=0}^{\infty} \lambda_t [w_t l_t + (1+r_t^b) b_t + (1-\delta)k_t + r_t^k k_t - c_t - k_{t+1} - b_{t+1}]$$

FOC:

$$(c_t)$$

$$\beta^t u'_c = \lambda_t$$

$$(k_{t+1})$$

$$-\lambda_t + \lambda_{t+1} [(1-\delta) + r_{t+1}^k] = 0$$

$$(b_{t+1})$$

$$-\lambda_t + \lambda_{t+1} [1 + r_{t+1}^b]$$

Combine FOCs wrt  $k$  and  $b$ :

$$1 - \delta + \gamma_{t+1}^k = 1 + \gamma_{t+1}^b$$
$$\Leftrightarrow \boxed{\gamma_{t+1}^b = \gamma_{t+1}^k - \delta} \quad [\text{"No arbitrage cond"}]$$

Transversality Condition:

$$\lim_{t \rightarrow \infty} p^t u'_c(t) \cdot k_{t+1} = 0$$

$$\lim_{t \rightarrow \infty} p^t u'_c(t) \cdot b_{t+1} = 0$$

Combine FOCs wrt  $c$  and  $k_{t+1}$ :

$$u'_c(t) = \beta u'_c(t+1) [1 - \delta + \gamma_{t+1}^k]$$

Firms  $\stackrel{\circ}{=}$  (Maximizing profits)

$$\Pi = y_t - w_t l_t - r_t^k k_t$$

$$y_t \leq F(k_t, l_t)$$

$$\Pi = F(k_t, l_t) - w_t l_t - r_t^k k_t$$

$$\frac{\partial \Pi}{\partial k} \Rightarrow F'_k = r_t^k$$

$$\frac{\partial \Pi}{\partial l} \Rightarrow F'_l = w_t$$

$$F(k, l) \text{ is CRS} \Leftrightarrow \lambda F(k, l) = f(\lambda k, \lambda l)$$

Take derivative w.r.t  $\lambda$ .

$$F(k, l) = \left. \frac{\partial F}{\partial k} \right|_{\lambda k, \lambda l} \cdot k + \left. \frac{\partial F}{\partial l} \right|_{\lambda k, \lambda l} \cdot l$$

Consider  $\lambda = 1$

$$F(k, l) = F'_k \cdot k + F'_l \cdot l$$

From FOC:

$$F(k, l) = r_t^t k_t + w_t l_t$$

$\pi = 0 \rightarrow$  if  $F(k, l)$  is CRS.

EE :-

$$u'_c(t) = \beta u'_c(t+1) \left[ 1 - \delta + \underbrace{r_{t+1}^k}_{\rightarrow F'_k(t+1)} \right] \quad (\text{from FOC})$$

EE,  $k_0$  & transversality are the same as in SPP.

$\Rightarrow$  CE is Pareto Optimal // First Welfare Theorem holds in this environment.

## Neoclassical Growth Model with Pop. Growth.

$N_0 = 1 \rightarrow$  starting mass of consumers.

$n \rightarrow$  growth rate of population

$$N_t = (1+n)^t N_0 = (1+n)^t$$

Assume  $F(K_t, N_t (1+g)^t)$

$\underbrace{(1+g)^t}$   
labor augmented tech  
progress

- small  $\rightarrow$  per capita letters

- Capital  $\rightarrow$  Aggregate

Notation: —

$c_t, k_t \rightarrow$  aggregate cons<sup>m</sup> & Capital

$c_t, k_t \rightarrow$  per capita  $\downarrow \downarrow \downarrow \downarrow$

Suppose  $c_t$  maximizes utility per capita: —

$$\sum_{t=0}^{\infty} \beta^t u(c_t) \rightarrow \max$$

Resource constraint:

$$c_t + k_{t+1} - (1-s)k_t = F(k_t, N_t (1+g)^t)$$

Define  $\tilde{c}_t = \frac{c_t}{(1+g)^t}$   $\tilde{k}_t = \frac{k_t}{(1+g)^t}$

Divide  $R_C$  by  $(1+g)^t (1+n)^t$

$$\frac{c_t}{(1+g)^t (1+n)^t} + \frac{k_{t+1}}{(1+g)^t (1+n)^t} - (1-\delta) \frac{k_t}{(1+g)^t (1+n)^t}$$

$$= F\left(\frac{k_t, n_t (1+g)^t}{(1+g)^t (1+n)^t}\right)$$

Can take this  
inside the  
fn as  
~~CBS~~.

$$\tilde{c}_t + (1+g)(1+n)\tilde{k}_{t+1} - (1-\delta)\tilde{k}_t = F(\tilde{k}_t, 1)$$

$$\sum_{t=0}^{\infty} \beta^t u(c_t) = \sum_{t=0}^{\infty} \beta^t \cdot \frac{c_t^{1-\alpha}}{1-\alpha} = \sum_{t=0}^{\infty} \beta^t \frac{(\tilde{c}_t (1+g)^t)^{1-\alpha}}{1-\alpha}$$

$u(c_t) = \frac{c_t^{1-\alpha}}{1-\alpha}$   
(CRRA utility)

$$= \sum_{t=0}^{\infty} \left( \beta (1+g)^t \right)^t \cdot \frac{\tilde{c}_t^{1-\alpha}}{1-\alpha}$$

$$= \sum_{t=0}^{\infty} \tilde{\beta}^t \frac{\tilde{c}_t^{1-\alpha}}{1-\alpha}$$

5<sup>th</sup> Oct.

Bellman Equation

$$v(k) = \max_{0 \leq k' \leq f(k)} [u(f(k) - k') + \beta v(k')]$$

More generally :

$$v(x) = \max_{y \in F(x)} [F(x, y) + \beta v(y)]$$

$$(Tv)(x) = \max_{y \in F(x)} [F(x, y) + \beta v(y)]$$

Solution to Bellman is a fixed point of operator T

$$v^* = Tv^*$$

Question:-

1) When  $v^*$  exists?

2) Is  $v^*$  unique?

3) When is  $v^* = \lim_{j \rightarrow \infty} T^j(v_0)$ ?

} contraction  
Mapping  
Theorem

Def : A metric space is a set  $S$  along with a metric  $d(x, y)$ , satisfying certain properties :-

for any  $x, y \in S$  :-

- 1)  $d(x, y) \geq 0$
  - 2)  $d(x, y) = 0$  iff  $x = y$
  - 3)  $d(x, z) \leq d(x, y) + d(y, z)$   $\forall x, y, z \in S$
  - 4)  $d(x, y) = d(y, x)$
- Δ inequality

Example:  $S = \mathbb{R}$

$$d(x, y) = \begin{cases} 1 & \text{if } x \neq y \\ 0, & \text{otherwise} \end{cases}$$

This is a metric space. [Add now Δ inequality holds]

Ex 2 Let  $X \subseteq \mathbb{R}^e$ ,  $\mathcal{C} = C(X)$  set of all  
continuous & bounded functions  $f: X \rightarrow \mathbb{R}$   
equipped w/ supnorm  $d$

$$d(f, g) = \sup_{x \in X} |f(x) - g(x)|$$

This is a metric space.

Δ Inequality :-

Fix some  $x$ ,

$$|f(x) - g(x)| \leq |f(x) - h(x)| + |h(x) - g(x)|$$

Take sup norm on both sides.  $\blacksquare$

Def : Let  $(\mathcal{S}, d)$  be a metric space. A sequence  $\{x_n\}$  with  $x_n \in \mathcal{S}$   $\forall n$  converges to  $x$  if  $\forall \varepsilon > 0 \exists N \in \mathbb{N} . d(x, x_n) < \varepsilon \forall n > N$

Exercise : Prove that if limit exists  $\Rightarrow$  it's unique  
(Proof by contradiction)

Cauchy Sequences : Let  $(S, d)$  be a metric space.

$\{x_n\}$  with  $x_n \in S \forall n$  is Cauchy if  $\forall \epsilon > 0$

$\exists N : \text{s.t. } \forall n, m \geq N, d(x_n, x_m) < \epsilon$ .

Thm: If  $\{x_n\}$  converges  $\Rightarrow$  it's a Cauchy sequence

Proof: Pick  $\epsilon > 0, \Rightarrow \exists M_{\frac{\epsilon}{2}} \text{ s.t.}$

$$d(x_n, x) < \frac{\epsilon}{2} \quad \forall n > M_{\frac{\epsilon}{2}}$$

$$d(x_n, x_m) \leq d(x_n, x) + d(x_m, x) < \frac{\epsilon}{2} + \frac{\epsilon}{2} < \epsilon$$

[ $\Delta$  inequality  
applies as  
it metric  
space.]

Ex.  $S = \mathbb{R}, d(x, y) = \begin{cases} 1, & x \neq y \\ 0, & x = y \end{cases}$

$$\{x_n\}, \text{ where } x_n = \frac{1}{n}$$

The distance b/w the elements would always be 1  $\Rightarrow$  not convergent  $\Rightarrow$  not Cauchy.

Def: A metric space  $(S, d)$  is complete if +  
Cauchy sequence  $\{x_n\}$  converges to any element  
 $x \in S$ . ( $x_n \in S \forall n$ )

(limits might be different but all converges to a limit in set  $S$ ).

Ex:  $X = [1, 2]$

$S = C(X)$  set of all continuous, strictly decreasing  
on  $[1, 2]$

$d(\cdot)$  is supnorm. This is not a complete metric space. (why?)

Counter example:

$\{f_n\} = \frac{1}{nx}$  Is it cauchy?

Pick  $\epsilon > 0$ , & let  $N_\epsilon = \frac{2}{\epsilon}$ ,  $m > n > N_\epsilon$

$$d(f_n, f_m) = \sup_{x \in [1, 2]} \left| \frac{1}{nx} - \frac{1}{mx} \right|$$

$$= \sup_{x \in [1, 2]} \frac{\frac{1}{nx} - \frac{1}{mx}}{\frac{1}{20} \frac{1}{4}}$$

$$= \sup_{x \in [1, 2]} \frac{m-n}{mn} \quad \left( \text{Max achieved at } x=1 \right)$$

$$= \frac{m-n}{mn}$$

$$= \frac{1 - \frac{n}{m}}{mn} \leq \frac{1}{n} \leq \frac{1}{N_\epsilon} = \frac{\epsilon}{2} < \epsilon$$

(Stolz & Cesaro)

## Contraction :

Def: Let  $(S, d)$  be a metric space &  $T: S \rightarrow S$ .

An operator  $T$  is a contraction if

$$d(Tx, Ty) \leq \beta d(x, y) \quad \forall x, y \in S, \beta \in (0, 1)$$

$\beta$  is called modulus of contraction.