Robustness of Empirical Revenue Maximization in Auction Learning

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Joint with

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"A Game-Theoretic Analysis of Empirical Revenue Maximization with Endogenous Sampling"

NeurIPS 2020

Key Points

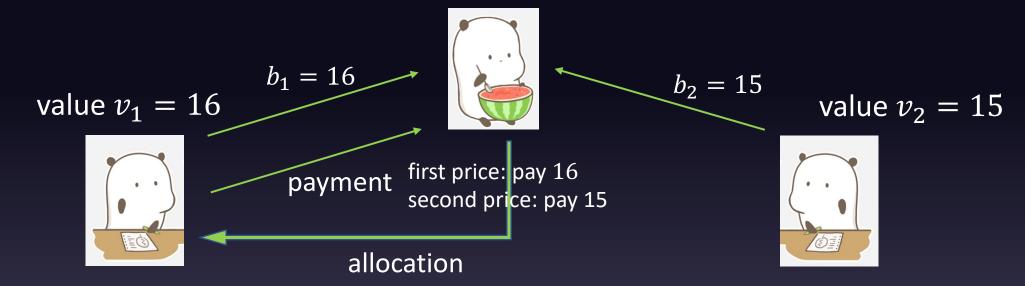
Auction + incentive-aware learning

Main Result

Empirical Revenue Maximization (ERM) is a *robust* learning algorithm against sample manipulation, as long as

- 1) the number of manipulated samples is **small**;
- 2) robustness is measured *in expectation* (rather than worst-case)
- Corollaries: a) in repeated auctions, ERM is an approximately truthful learning algorithm;
 b) uniform price auction is "group-strategyproof in the large".

Auctions



Examples

- 1) Second price auction: highest bidder wins, pays the 2nd highest bid.
- 2) Second price auction with reservation price p: highest bidder wins (if bid > p), pays

$$\max\{2^{\text{nd}} \text{ bid}, p\}$$

3) Posted-price auction: single-bidder version of 2); the bidder pays p if $v \ge p$

The optimal auction (Myerson's)

- One item for sell
- K bidders with i.i.d. values $v_i \sim F$ for the item

Theorem (Myerson, 1981)

The revenue-optimal auction is the second price auction with reservation price

$$p^* = \underset{p \in \mathbf{R}}{\operatorname{argmax}} \{ Rev_{\mathbf{F}}(p) \}, \qquad Rev_{\mathbf{F}}(p) \coloneqq p \left(1 - \mathbf{F}(p) \right) = p \operatorname{Pr}_{v \sim \mathbf{F}}[v \geq p]$$
expected revenue of the auction with posted price p

Assuming distributional knowledge is problematic!

Game theory has a great advantage in explicitly analyzing the consequences of trading rules that presumably really are common knowledge; it is deficient to the extent it assumes other features to be common knowledge, such as one agent's probability assessment about another's preferences or information.

Learning optimal auction from samples

- Suppose we have N samples $\mathbf{s} = (s_1, s_2, ..., s_N) \sim \mathbf{F}$
- We want to learn a good auction (reservation price) from s
 - Cole & Roughgarden (2014), Dhangwatnotai et al (2015), Huang et al (2015), Morgenstern & Roughgarden (2015), Gonczarowski & Nisan (2017), Guo et al (2019), ...

- *N* samples $s = (s_1, s_2, ..., s_N) \sim F$
- Output the optimal reservation price for the empirical distribution $E = \text{Uniform}\{s_1, s_2, ..., s_N\}$

$$ERM(s) = \underset{p \in \mathbb{R}}{\operatorname{argmax}} \{ \operatorname{Rev}_{\underline{E}}(p) \} = \underset{p \in \mathbb{R}}{\operatorname{argmax}} \{ p (1 - \underline{E}(p)) \}$$

ERM is important!

• ERM is the Empirical Risk Minimization in the auction learning problem:

Learning a classifier: Empirical Risk Minimization: Empirical Revenue Maximization: $c^* = \operatorname*{argmin} \{ \ loss_F(c) \ \} \qquad c^*_s = \operatorname*{argmin} \{ \ loss_E(c) \ \} \qquad p^*_s = \operatorname*{argmax} \{ \ Rev_E(p) \ \} \qquad p \in \mathbb{R}$

• ERM has good performance: $\left(1 - O\left(\sqrt{\frac{\log N}{N}}\right)\right)$ -optimal expected revenue [*Huang et al* '15]

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But, assuming sample access is still problematic!

Where do we get samples?

• Exogenous – by market research, by magic?

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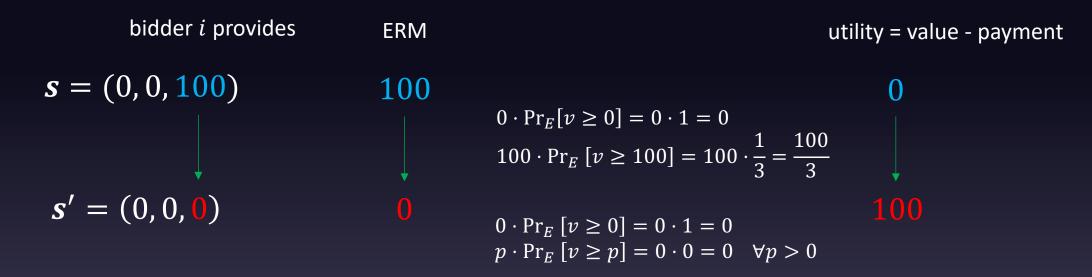
- Exogenous by market research, by magic?
- *Endogenous* from bidders e.g., repeated auctions

Incentive issue: bidders may provide false samples to lower ERM(s)!

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Example 1



- *N* samples $s = (s_1, s_2, ..., s_N) \sim F$
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$$\operatorname{ERM}(\boldsymbol{s}) = \operatorname{argmax} \{ \operatorname{Rev}_{\boldsymbol{E}}(p) \} = \operatorname{argmax} \{ \operatorname{pPr}_{\boldsymbol{E}}[v \ge p] \}$$

$$p \in \mathbf{R}$$

Example 2

bidder *i* provides ERM
$$\mathbf{s} = (100, 100, 100)$$

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$$\mathbf{s}' = \left(100, 100, \frac{200}{3}\right)$$

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Main result: **ERM** is robust in expectation!

Informal Theorem 1 (this work)

For any bounded F, if at most $m = o(\sqrt{N})$ samples are manipulated, then the expected change of ERM(s) is at most $\Delta = O\left(m^{2/3} \frac{\log^2 N}{N^{1/3}}\right) \to 0$ (as $N \to \infty$)

Main result: **ERM** is robust in expectation!

Formally:

- Suppose we have N values v_1, \ldots, v_N
- where N-m values $v_{-I}=\{v_{m+1},\ldots,v_N\}$ are i.i.d. samples from F.
- other m values $x_I = \{v_1, \dots, v_m\}$ can be changed to any m values $b_I \in \mathbf{R}_+^m$

• $ERM(x_I, v_{-I}) \longrightarrow ERM(b_I, v_{-I}).$

$$\Delta_{N,m}^F = \mathrm{E}_{v_{-I} \sim F^{N-m}} \left[\sup_{x_I, b_I \in \mathbf{R}_+^m} \frac{\mathrm{ERM}(x_I, v_{-I}) - \mathrm{ERM}(b_I, v_{-I})}{\mathrm{ERM}(x_I, v_{-I})} \right]$$

Related works & our work

- Lavi et al. (WWW'19) proposed $\Delta_{N,1}^F$ and showed $\Delta_{N,1}^F \to 0$ for discrete F
- Yao (arXiv'18) showed $\Delta_{N,1}^F \to 0$ for any [1,D]-bounded F No convergence rate is given!

Theorem 1 (our): Upper bound on $\Delta_{N,m}^F$

Assuming $m = o(\sqrt{N})$

- Universally for any [1,D]-bounded F, $\Delta_{N,m}^F = O\left(m^{2/3} \frac{\log^2 N}{N^{1/3}}\right) \to 0$
- Universally for any MHR F, $\Delta_{N,m}^F = O\left(m\frac{\log^3 N}{\sqrt{N}}\right) \to 0$ MHR: $\frac{f(x)}{1-F(x)}$ increases

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- For MHR F, $\Delta_{N,m}^F = \Omega(?)$

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Corollary: Repeated auctions

(cf. differential privacy techniques in [LHW'18] [ACKTM'19])

- *T* rounds of second price auctions
 - Each bidder participates in $\leq m$ rounds,
 - $m \ll T$
 - Each round has one item, K bidders (e.g., $K = 1 \ or \ 2$), with iid values

Corollary 1. If $m = o(\sqrt{T_1K})$, the two-phase ERM algorithm is ϵ_1 -BIC and $(1 - \epsilon_2)$ -revenue optimal, with

•
$$\epsilon_1 = O\left(m^{5/3} \frac{\log^2 T_1 K}{(T_1 K)^{1/3}}\right) \to 0$$
, and $\epsilon_2 = O\left(\frac{T_1}{T} + \sqrt{\frac{\log T_1 K}{T_1 K}}\right)$

Two-phase ERM: Divide the T rounds into two phases:

- T_1 rounds of exploration: set reservation price to 0.
- T_2 rounds of exploitation: set reservation price to ERM(s), where
 - $s = (b_1, ..., b_{T_1 K})$ are the first round bids.

Corollary 2: Uniform-price auctions

one-shot uniform-price auction:

("strategyproofness in the large" [Azevedo & Budish '18]])

- N copies of a good
- N unit-demand bidders with i.i.d. values v from F submit bids b.
- Auctioneer sets a price p = P(b)
- Each bidder i with bid $b_i \ge p$ receives a copy of the good and pays p, obtaining utility $v_i p$; otherwise the utility is 0.
- Assumption (joint deviation): we allow any m ≥ 1 bidders to coordinate and submit non-truthful bids jointly!

Corollary: (m, ϵ) -group BIC: No bidder gains ϵ more utility from joint deviation, where, for P = ERM and any bounded F, $\epsilon = D\Delta_{N,m}^{F,P} \to 0$

Summary

See details at

- (Single-dimensional) Auction + (incentive-aware) Learning
- Bidders have incentives to misreport samples to fool the **Empirical Revenue Maximization/Empirical Risk Minimization algorithm.**

Takeaway

Empirical Risk Minimization is a *robust* learning algorithm in the auction learning problem, as long as

- 1) robustness is measured *in expectation* (rather than worst-case)
- 2) misreported samples are *few*
- Future works: a) more complex auction environment?
 - b) tighten the bounds on $\Delta_{N,m}^F$