Persistent Homology: Filtration building techniques

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Outline for studying persistent homology

- 1. Intro to persistent homology
 - Build intuitions of persistent homology: what it does, what it produces
- 2. Formalizing persistent homology
 - Introduce its input (filtration) and study an algorithm for computation
- 3. Different ways for building filtrations
 - Vietoris-Rips filtration, sub-levelset filtration
 - Cubical complexes (for images)
- 4. Interpretation and stability of persistence diagram

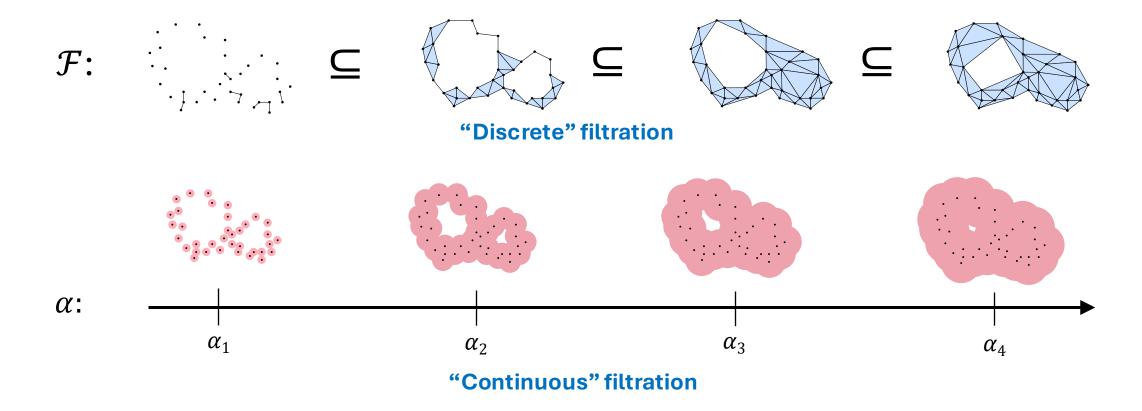
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- But we haven't formally defined PD for a continuous filtration, where we have a space varying over $\alpha \in [0, \infty)$ (technically, there're infinitely many of them)



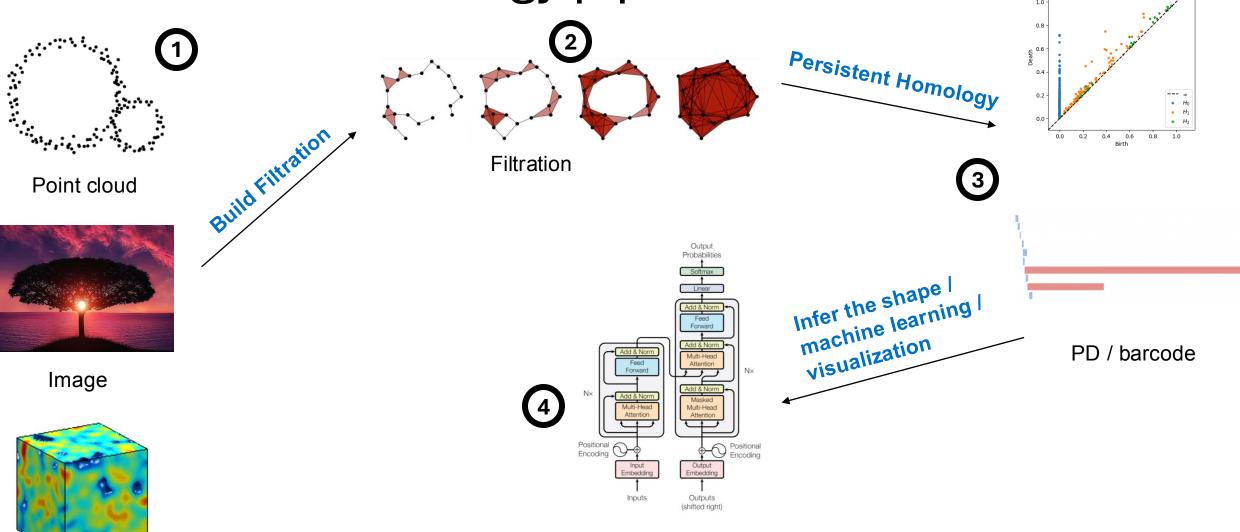
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- So to harness the power of persistence, you have to do this
- So we shall not only formally define PD on input data (which are typically continuous at least theoretically), but also learn ways to preprocess the data into filtrations to feed into the persistent homology pipeline

Persistent homology pipeline

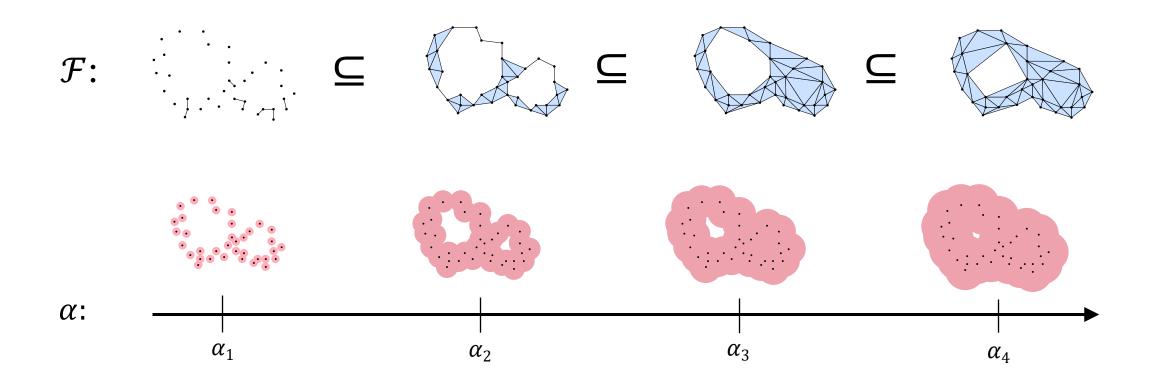


3D volume data

Some img from: AATRN; https://quantdare.com/understanding-the-shape-of-data-ii/; https://pixabay.com/photos/new-year-background-tree-sunset-736885/; Adler et al. Persistent homology for random fields and complexes.; https://builtin.com/artificial-intelligence/transformer-neural-network

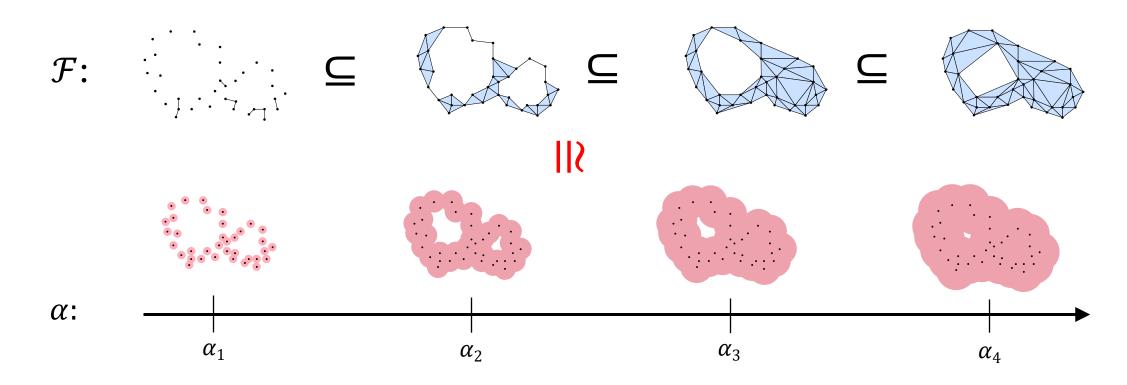
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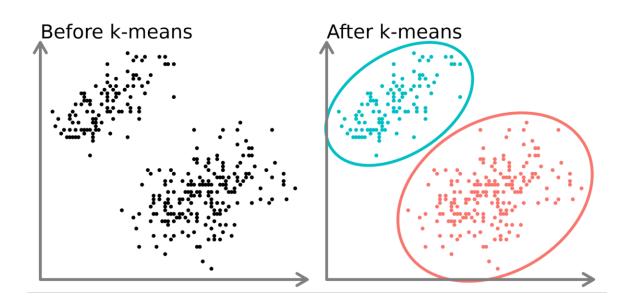
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- We shall eventually show that continuous filtration is in some sense "equivalent" to the discrete one



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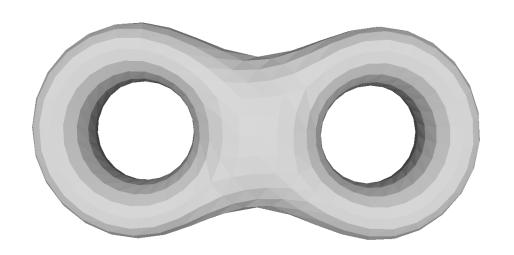
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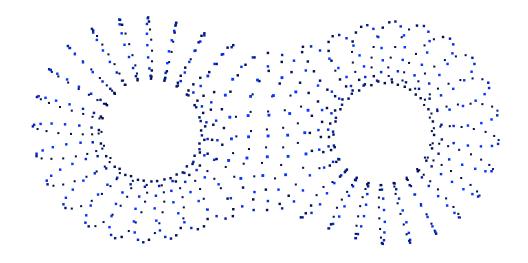


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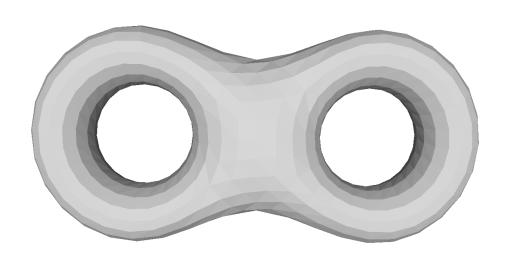
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 - Even for supervised learning (another type of more popular? machine learning), if you ignore the "labels" for the data, then the data become point clouds
 - After all, each element in your data is in some sense a "point"

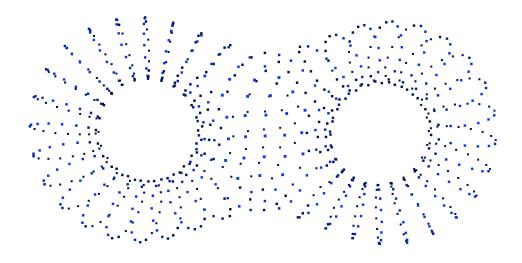
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- Trying to infer the shape of point cloud is indeed a major motivation for topological data analysis





- Recall that each space in the "growing of balls" filtration is to take a ball of the same radius α centering in each point, and then take the union of the α -radius balls of all points
- To get the (continuous) filtration, we then let the radius α increase from 0 to
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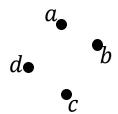
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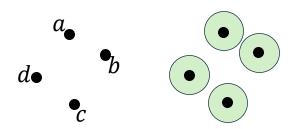
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- The remaining task is to build high-dimensional simplices out of the points (vertices)

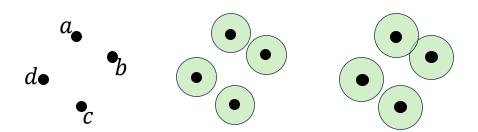
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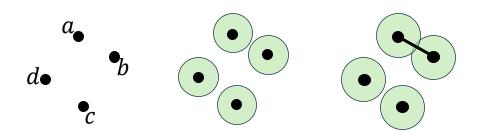
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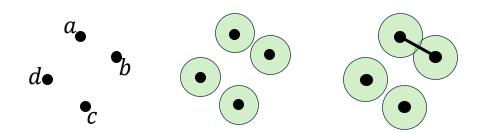
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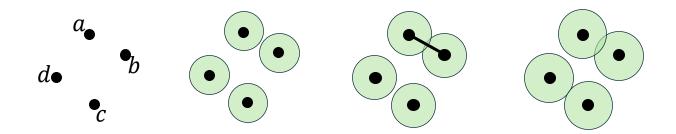
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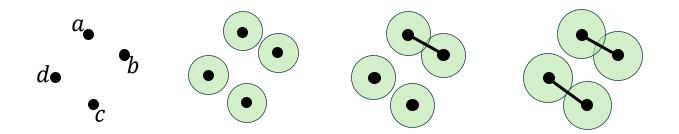
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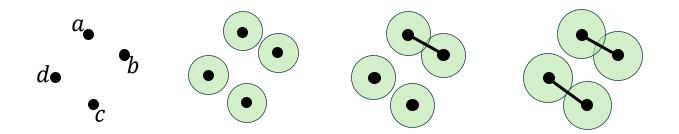
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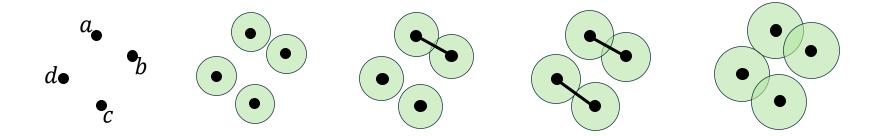
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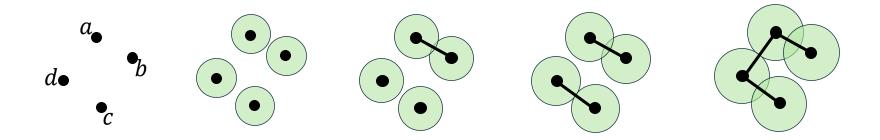
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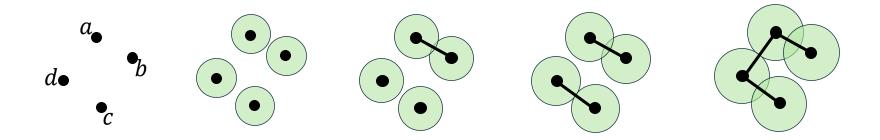
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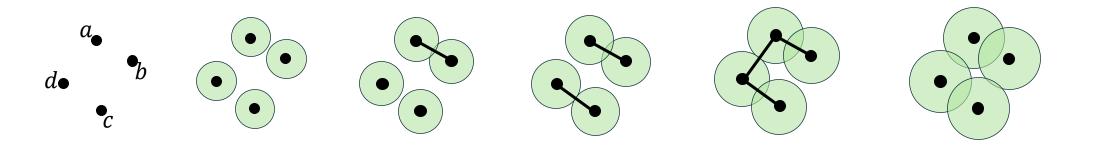
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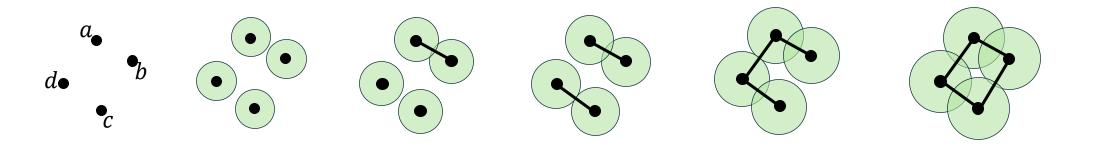
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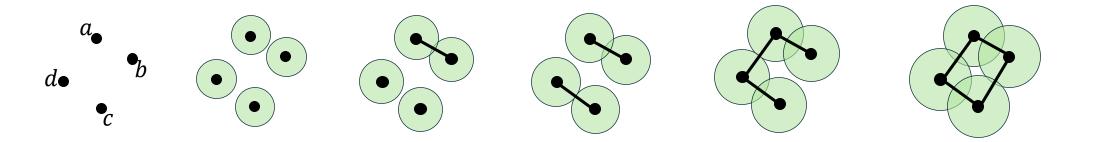
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- a ab
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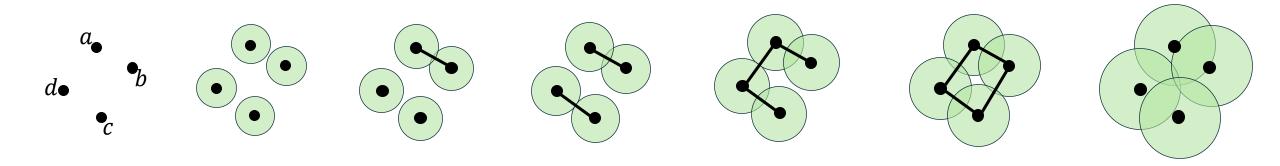


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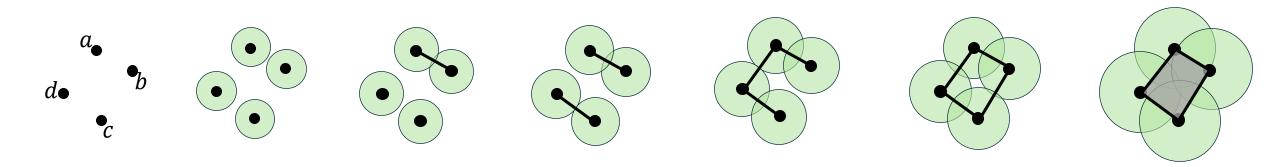
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- a ab
- *b cd*
- c ad
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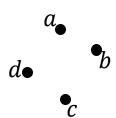
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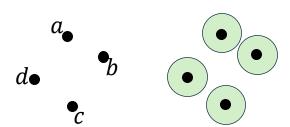
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- Of course, in practice, the algorithms that people use to compute Čech Complexes are much more efficient ones

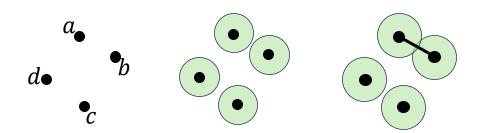
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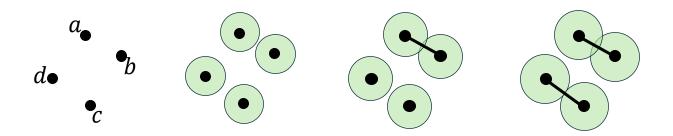
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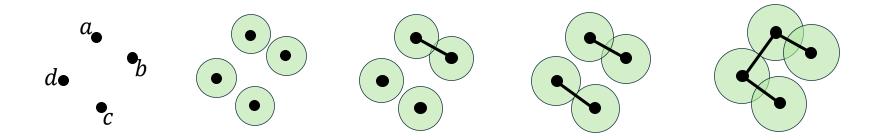
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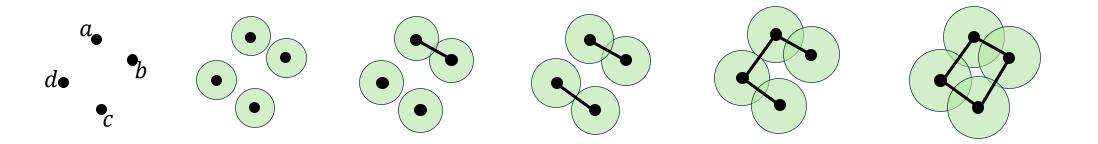
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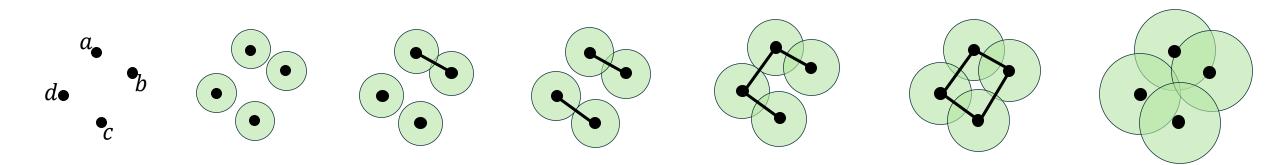
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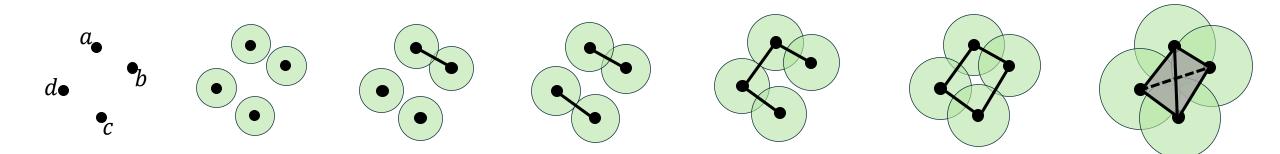
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- a ab
- b cd
- c ad
- *d bc*



- a ab abcd with all its faces
- b cd (abc, abd, acd, and
- *c ad bcd*).
- *d bc*



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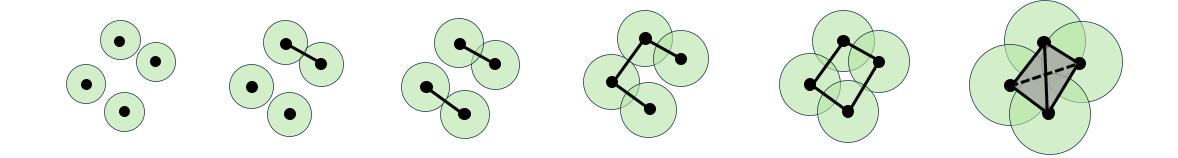
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- E.g., if three balls intersect, then any two balls also intersect. So the edges of the corresponding triangle are also in the Cech complex.

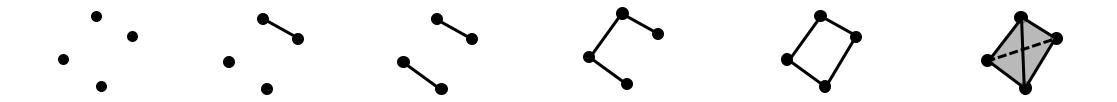
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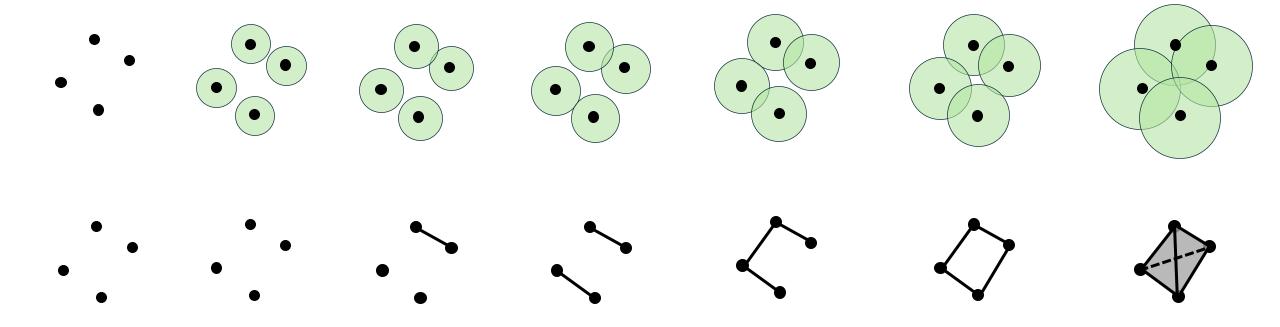


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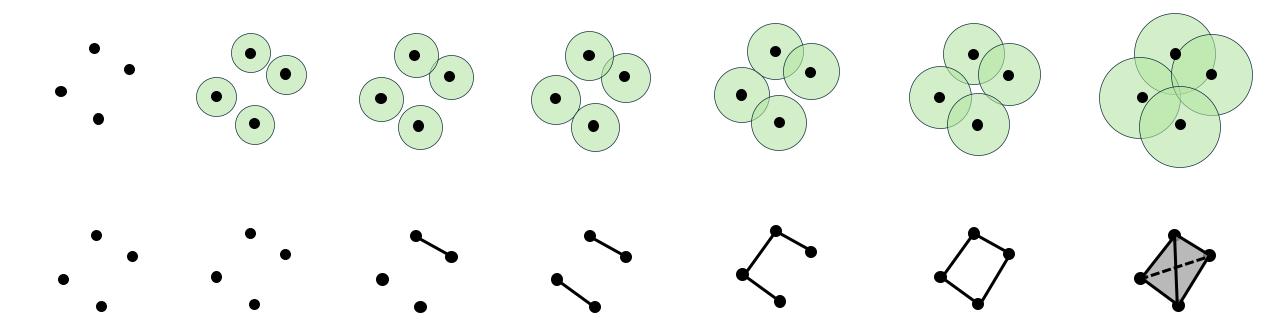


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- The above "equivalence" is called the "homotopy equivalence" in algebraic topology, whose definition is beyond the scope



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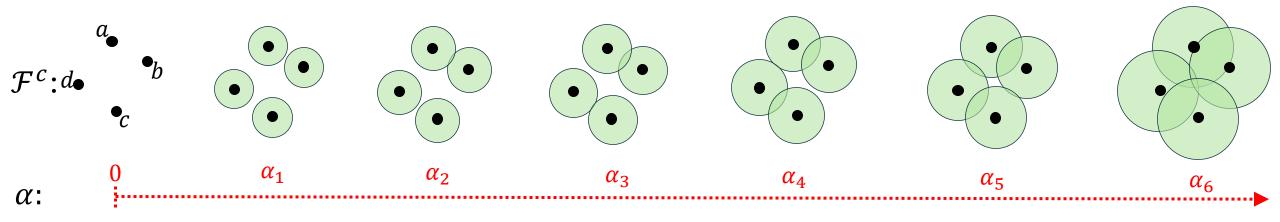
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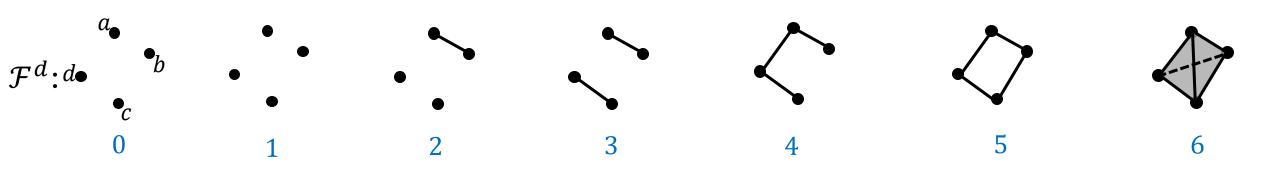
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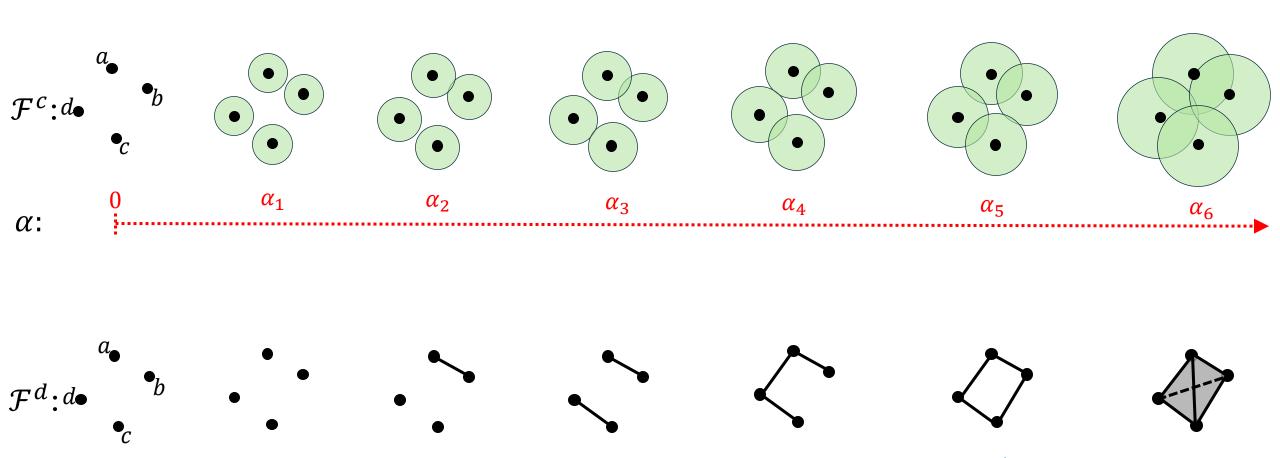
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 - First, we have that each complex in the discrete filtration corresponds to a range of α -values (and a range of spaces) in the continuous filtration
 - For each interval $[b,d) \in PD(\mathcal{C}(P))$, consider the corresponding complexes in $\mathcal{C}(P)$, which are $K_b, K_{b+1}, \dots, K_{d-1}$.

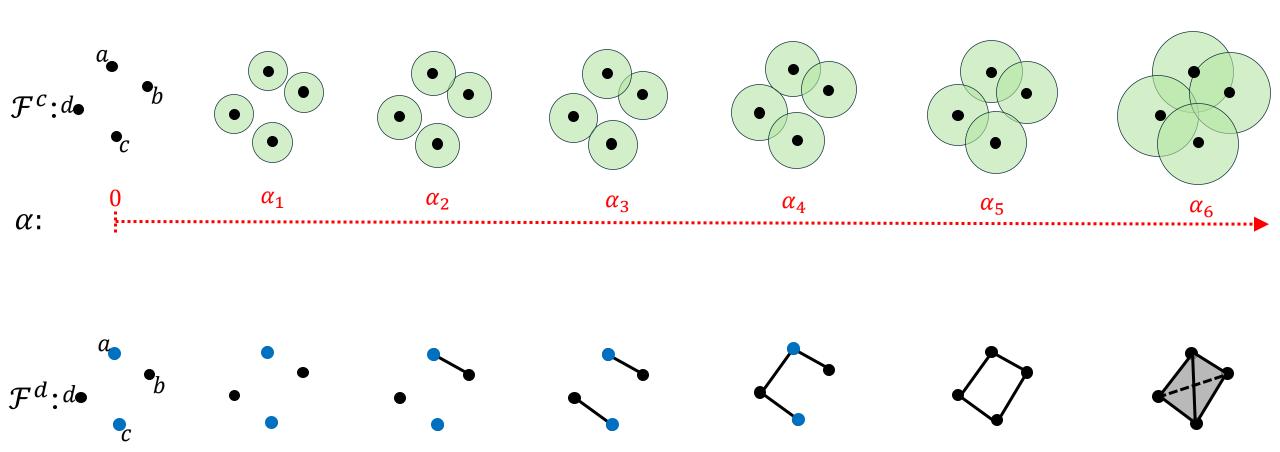
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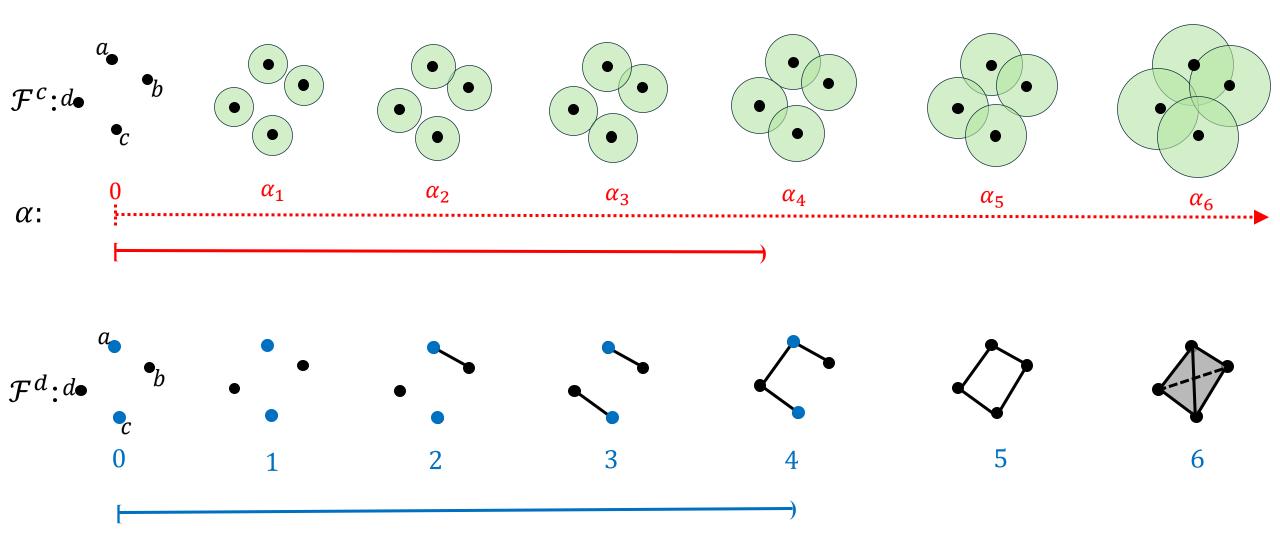
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 - Notice that an interval of $PD(\mathcal{F}^c)$ is actually a continuous interval over the real line \mathbb{R} .



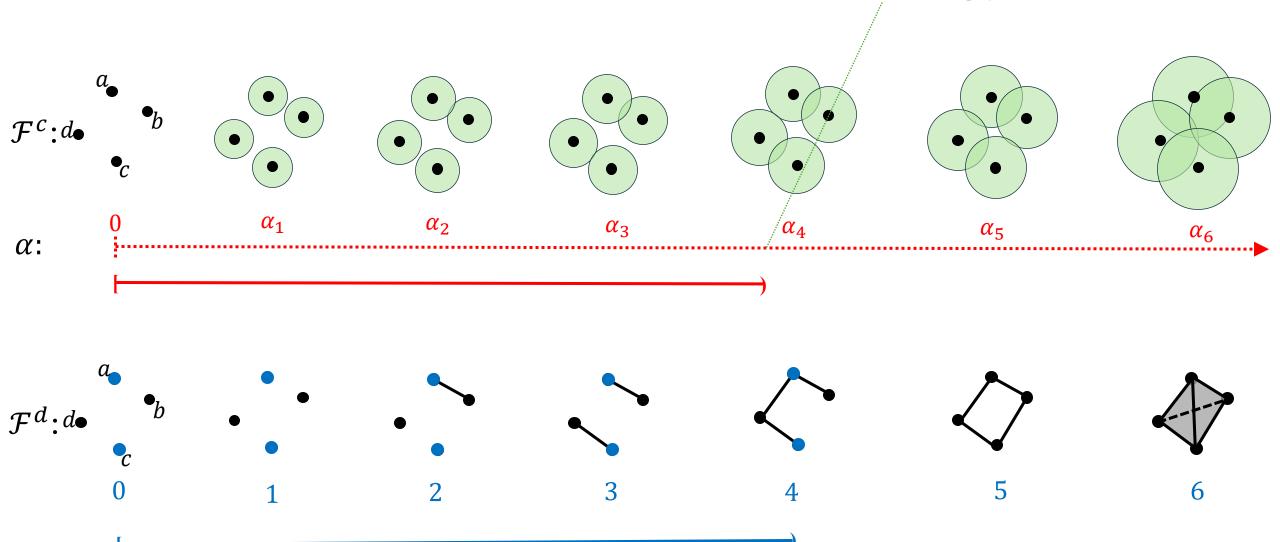


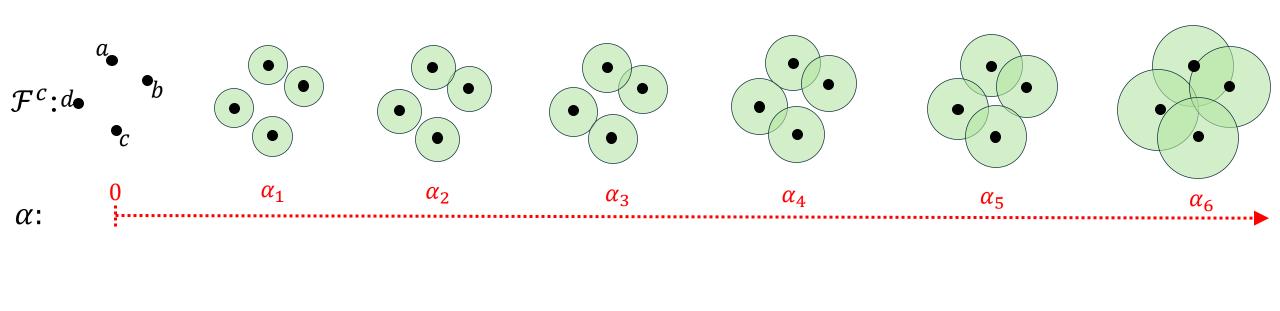


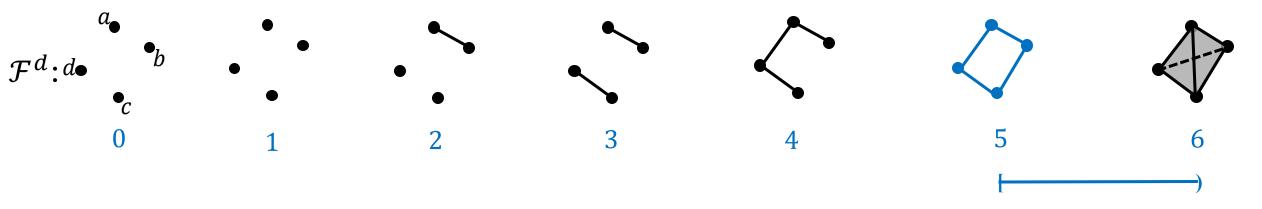


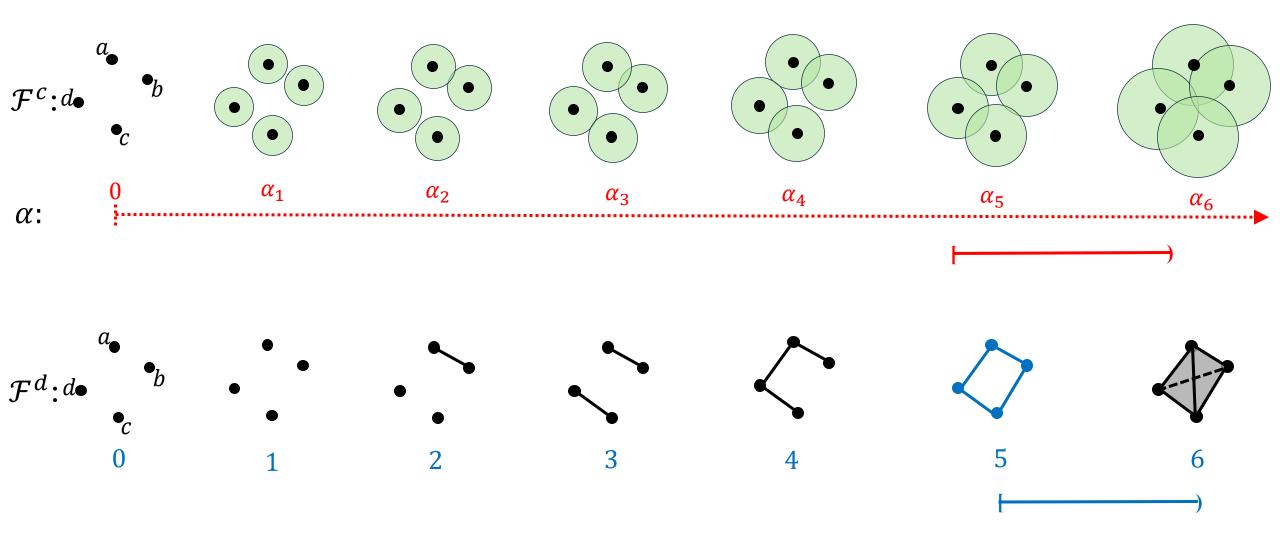


Between α_3 and α_4 , balls of a and d touch so the gap between a and d filled

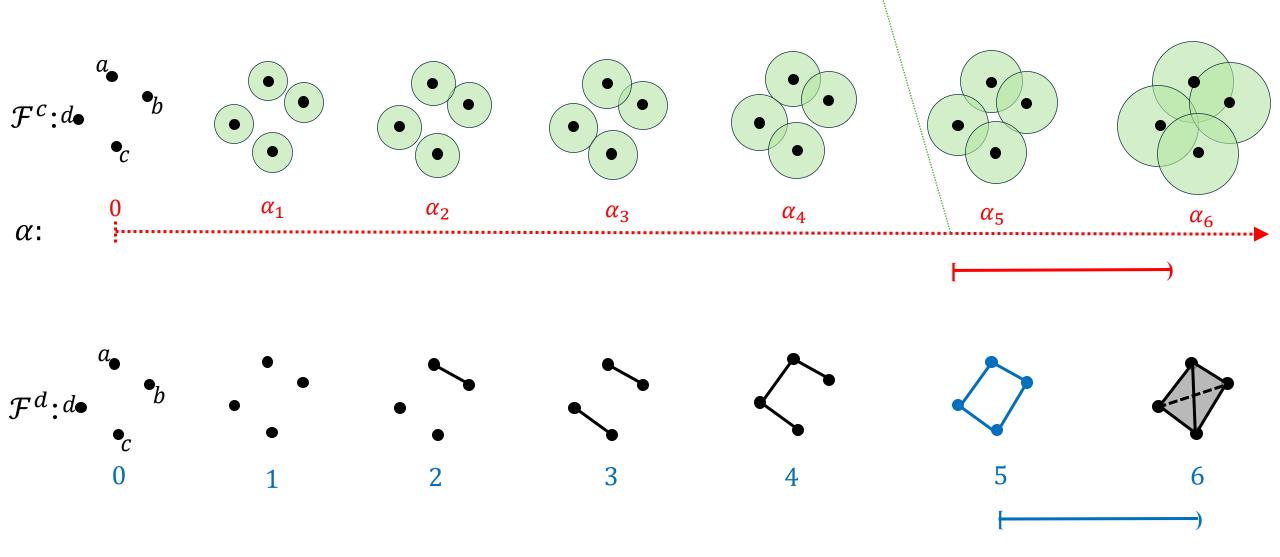




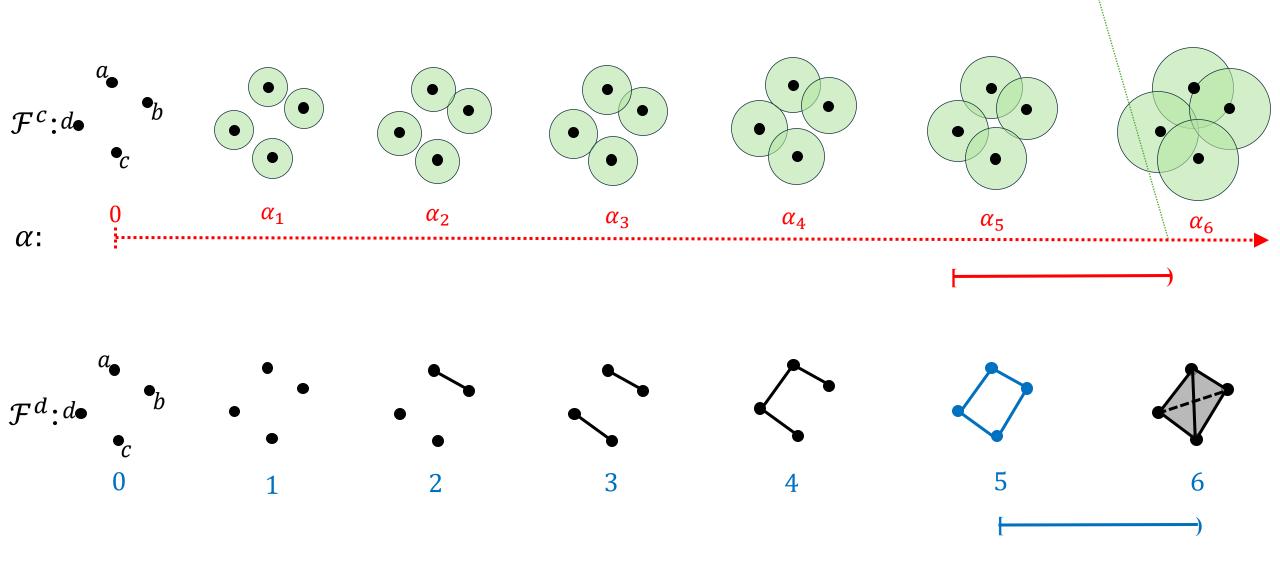




Between α_4 and α_5 , balls of b and c touch so the blue 1-cycle forms



Between α_5 and α_6 , the central hole gets filled so 1-cycle become trivial



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- After all, trade-offs were made everywhere in computer science between efficiency and quality

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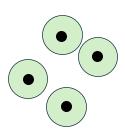
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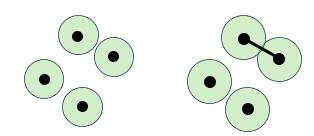
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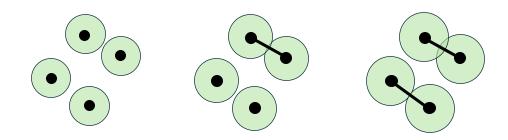
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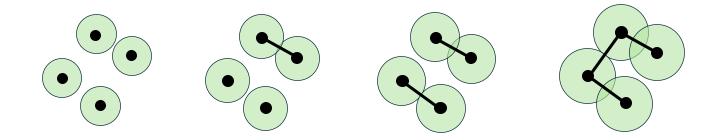
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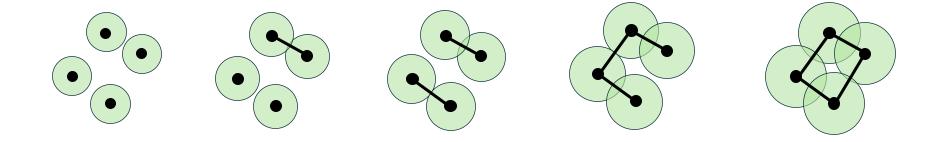
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- There are again finitely many α where the Rips complexes change, so the Rips filtration is a discrete finite filtration

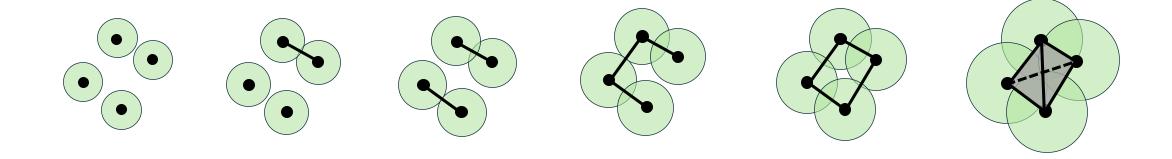


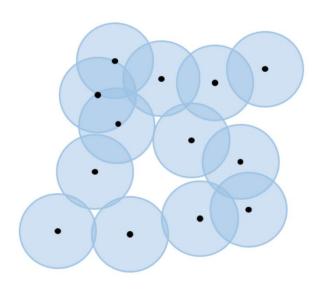


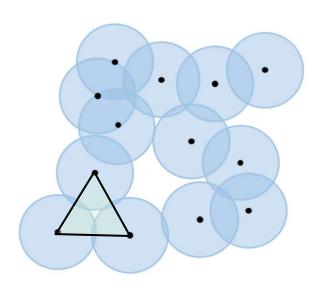


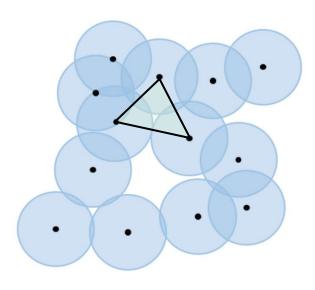




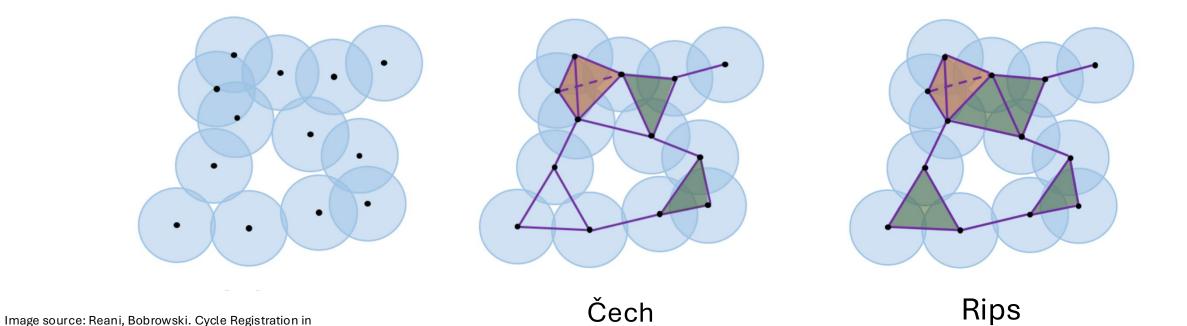




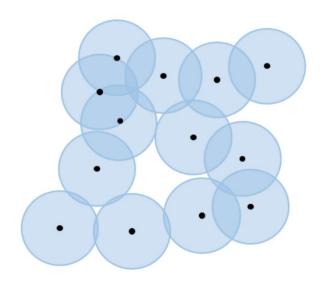


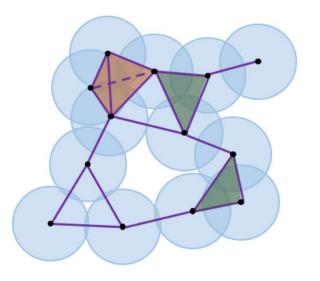


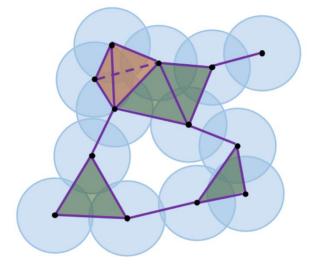
Persistent Homology with Applications in Topological Bootstrap



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- But notice that other than the two triangles the remaining simplices in the two complexes are the same

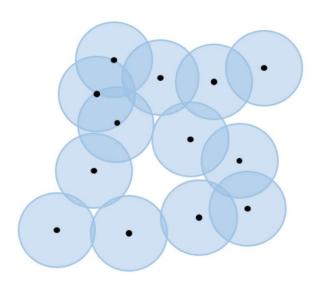


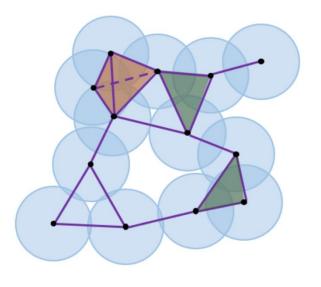


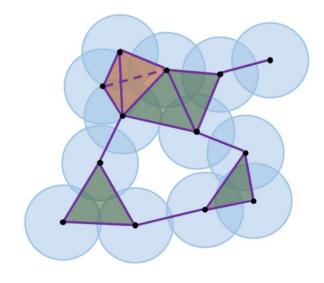


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- Furthermore, if we increase the radius for Čech complex, the two missing two triangles will come into picture ---- in some sense, the sequences of Čech and Rips complexes are "interleaved with each other"







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"Similarity" of Čech and Vietoris-Rips

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- Reason: Edges are formed by two points. If you check the definition of Čech and Vietoris-Rips ("all balls for a set of points intersect" and "each pair of balls for a set of points intersect"), when we only have two point, the two criteria become the same.

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- The benefit of the above alternative definition is that we can completely
 eliminate balls and define Rips complexes / filtration by only considering the
 pair-wise distances between points

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- They are abstract objects but we have some pair-wise "distances" between these objects.

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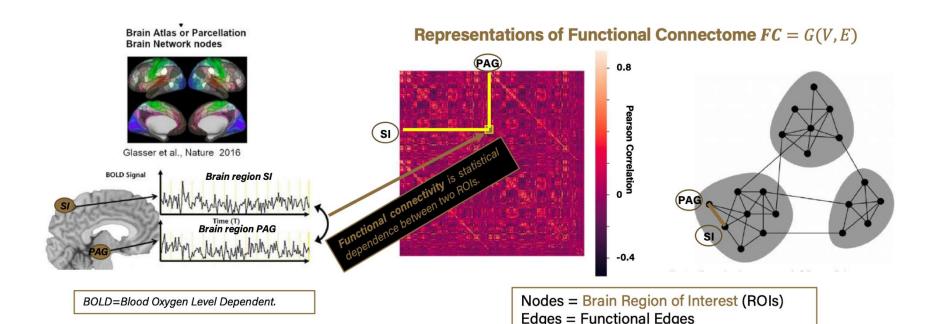


Figure courtesy of: Duy Duong-Tran

- In this example, the data points are "regions of the brain", and we can calculate their "similarity" by measuring the correlation of their blood oxygen fluctuation over time (a time series data).
- We then measure distances of two brain regions by taking inverse of similarity

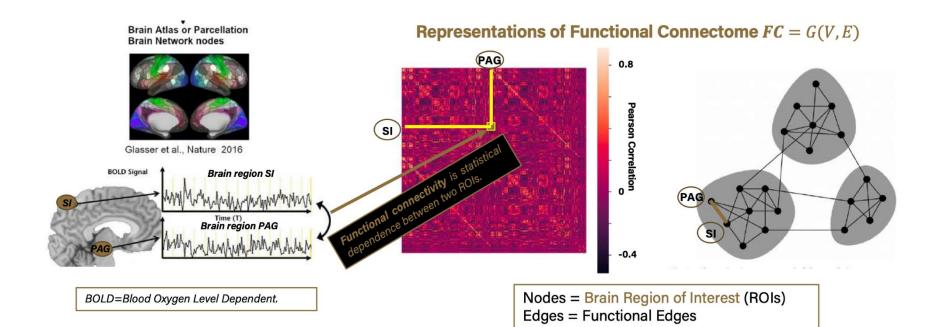
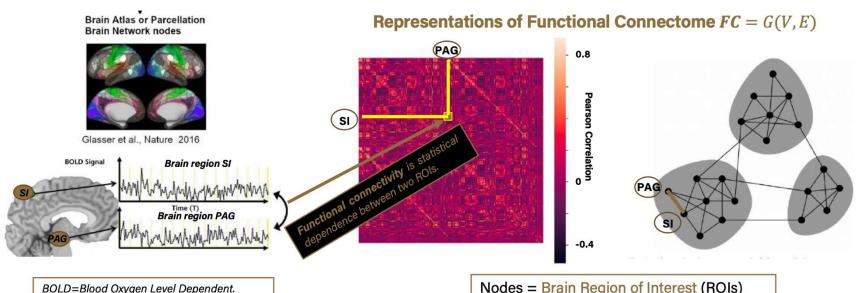


Figure courtesy of: Duy Duong-Tran

- In this example, the data points are "regions of the brain", and we can calculate their "similarity" by measuring the correlation of their blood oxygen fluctuation over time (a time series data).
- We then measure distances of two brain regions by taking inverse of similarity
- These regions are not really technically having a position (each region is represented by a blood oxygen level function over time), but we have distances between the regions



Nodes = Brain Region of Interest (ROIs)
Edges = Functional Edges

• For this data, we still can build Rips filtration on these brain regions

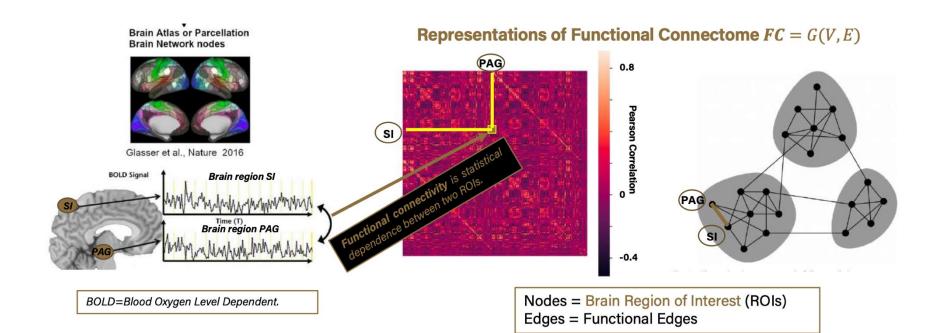


Figure courtesy of: Duy Duong-Tran

- We shall briefly look at some facts concerning computing Rips Filtration.
- Recall:
- **Definition**: Given a point set P and a distance r, the **Vietoris-Rips** complex of P corresponding to the distance r, denoted $\mathbb{VR}^r(P)$, is a simplicial complex whose vertices are points in P such that
 - A subset of points $p_0, p_1, ..., p_d$ of P form a d-simplex if and only for each pair of points p_i, p_i in the subset, their distance is no more than r.

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- Since a pair of points form an edge, this also means that a $\mathbb{VR}^r(P)$ can be completely determined once we have figured out the edges (1-simplices) for $\mathbb{VR}^r(P)$

- So, for a certain r, to compute $\mathbb{VR}^r(P)$, our first thing to do: Enumerate each pair of points in P and check whether their distance is no more than r.
 - If this is true, we let the pair form an edge in $\mathbb{VR}^r(P)$
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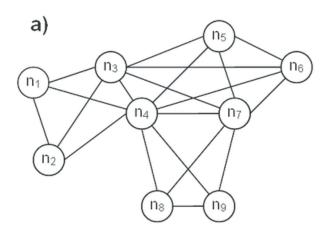
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- $\mathbb{VR}^r(P)$ is then the Clique complex of the graph $\mathbb{G}^r(P)$.

Clique

- **Definition**: A **clique** of a graph G = (V(G), E(G)) is a subset S of V(G) such that each pair of vertices of S form an edge in G.
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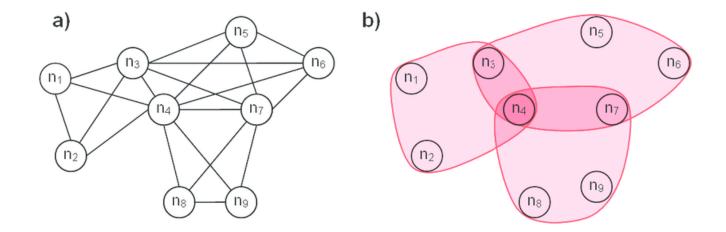


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- E.g, the following graph has three maximal cliques (a clique not contained in another clique)

$$G_1=\{n_1, n_2, n_3, n_4\}$$

 $G_2=\{n_3, n_4, n_5, n_6, n_7\}$
 $G_3=\{n_4, n_7, n_8, n_9\}$

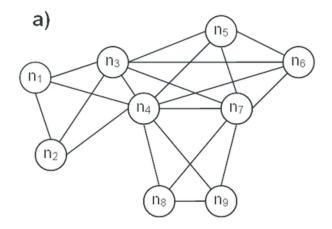


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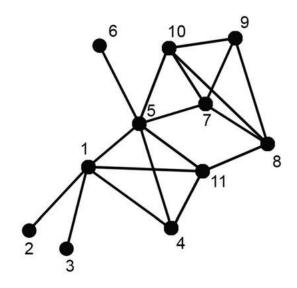
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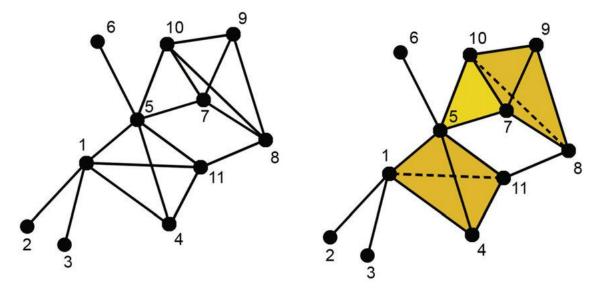


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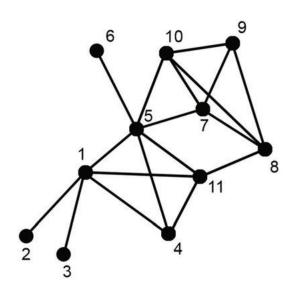
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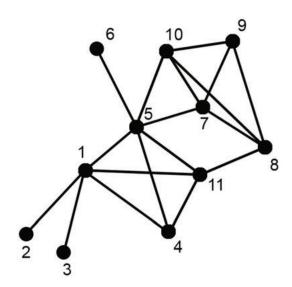
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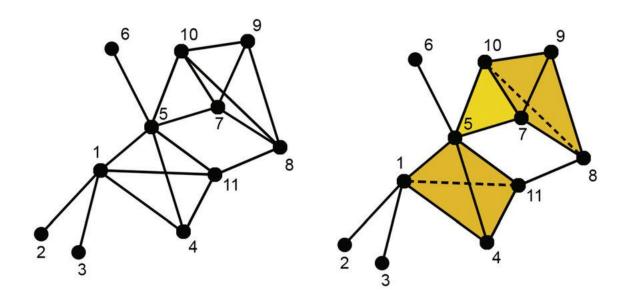
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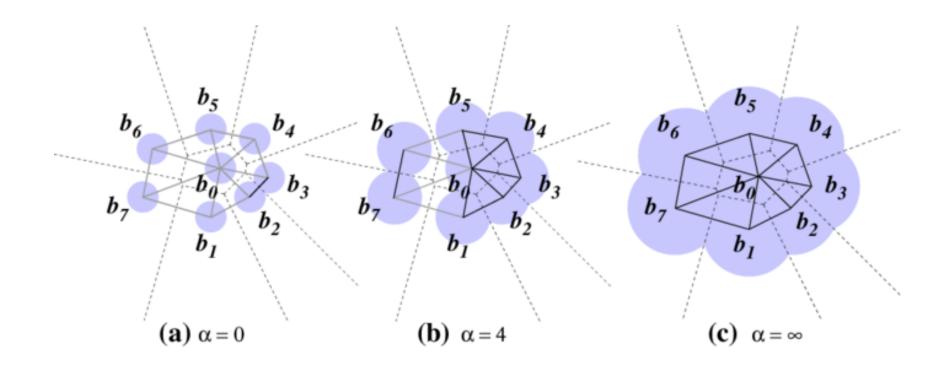
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- Finding these values takes $O(n^2 \log n)$ time dominated by the sorting, where n is the number of points in P

Another type of filtration

- **Delauney** complexes / filtrations: growing the balls around points, construct a simplex whenever their set of balls intersect (the same as Cech complexes)
- Difference: the part of a ball stops growing when touching another ball.



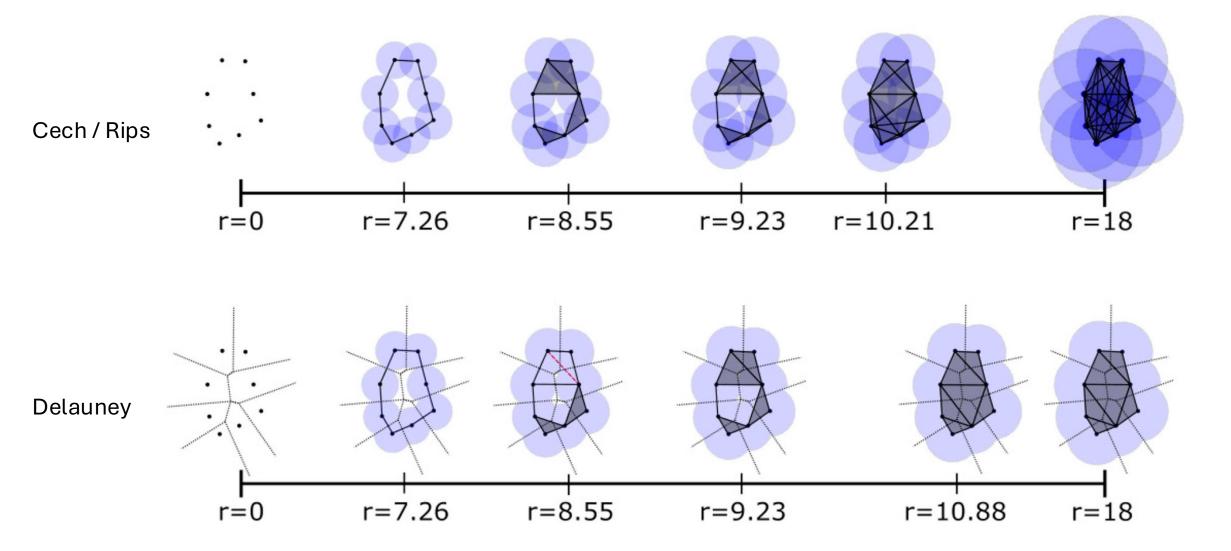


Image source: Mishra, Motta. Stability and machine learning applications of persistent homology using the Delaunay-Rips complex

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- Disadvantage of Delauney complexes: costly to compute, especially when the dimension of the points in the point cloud is high
- The go-to filtration for point cloud is **Rips filtration** because of (1) its computational efficiency and (2) the fact that it still faithfully recover the shape of the data (despite data loss)

Other types of complexes

- There are other types of complexes:
 - Witness complex
 - Graph-induced complex
 - Tangential complex
 - ...
- Will not cover them at least for the time being

Data as a function

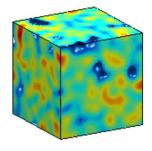
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Image



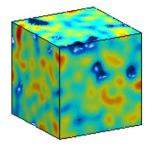
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Image



3D volume data

• E.g., all pixels in an image form the domain of a function and the color value on each pixel is basically the function value on a point of the domain

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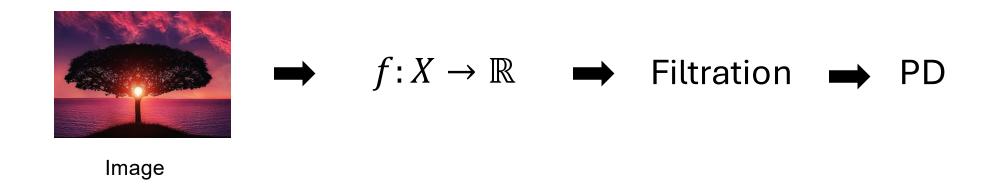
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 - Even if the range of the function is more than a single real value, say again, a colored image, we can take each channel (RGB), this will give you three individual real-valued functions. We can analyze each individually using persistence

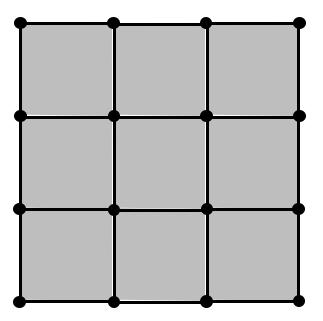
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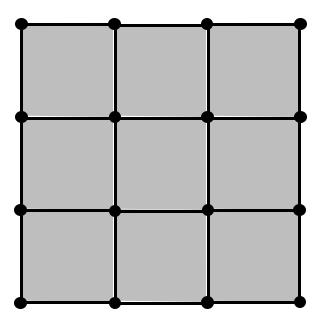


• We visualize the domain X of a 2d image as a regular grid, where pixels are grid points (below an example of 4x4 image)

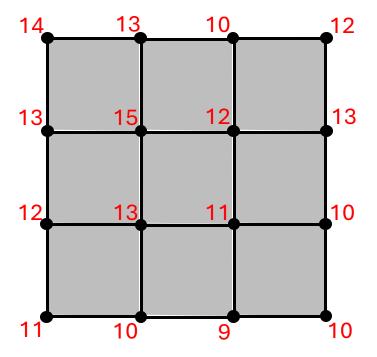
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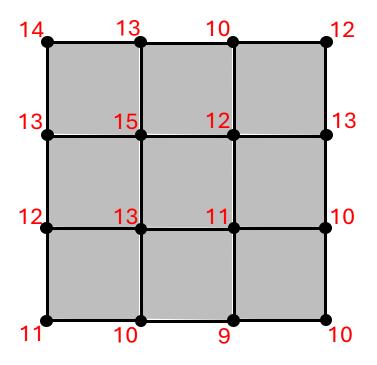
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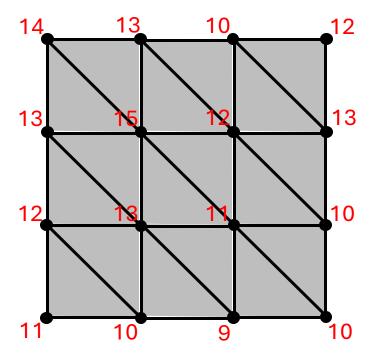
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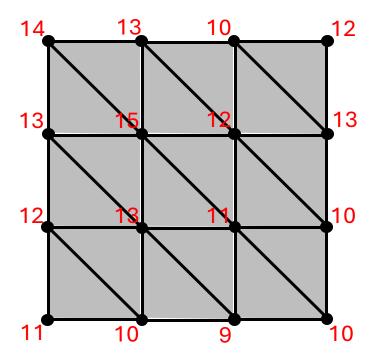
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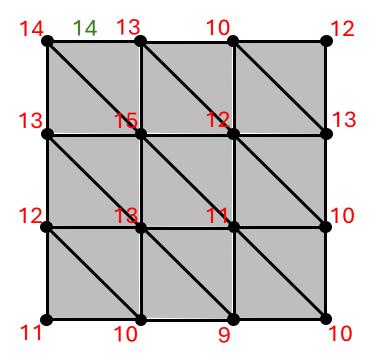
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- So we subdivide the grid to be consisting of triangles, so \boldsymbol{X} becomes a simplicial complex



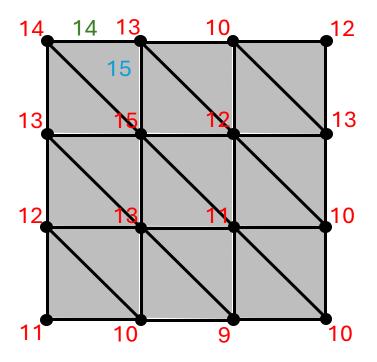
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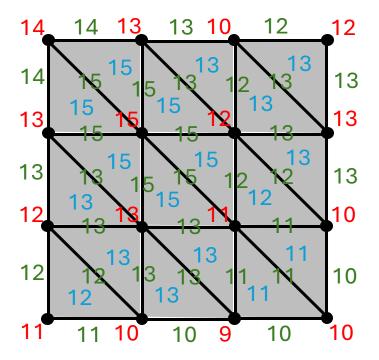
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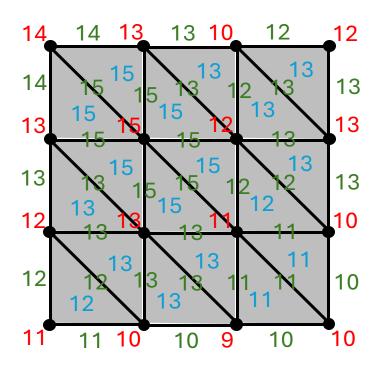
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• It should be esay to verify that $f^{-1}(-\infty, \alpha_i] \subseteq f^{-1}(-\infty, \alpha_{i+1}]$ for any i

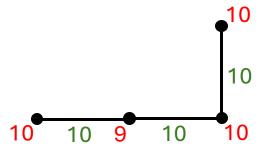
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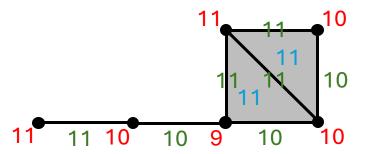
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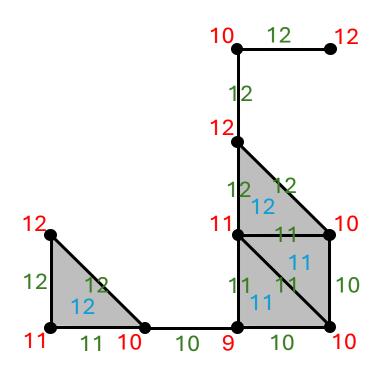


• $f^{-1}(9] \subseteq f^{-1}(10] \subseteq f^{-1}(11]$

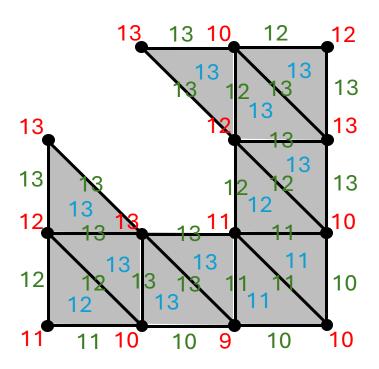




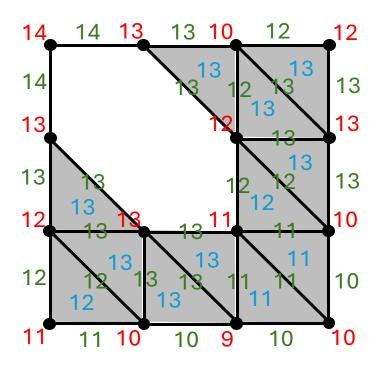
• $f^{-1}(9] \subseteq f^{-1}(10] \subseteq f^{-1}(11] \subseteq f^{-1}(12]$



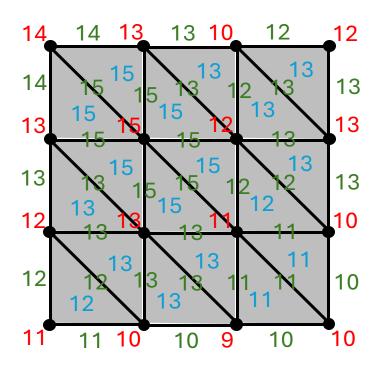
• $f^{-1}(9] \subseteq f^{-1}(10] \subseteq f^{-1}(11] \subseteq f^{-1}(12] \subseteq f^{-1}(13]$



• $f^{-1}(9] \subseteq f^{-1}(10] \subseteq f^{-1}(11] \subseteq f^{-1}(12] \subseteq f^{-1}(13] \subseteq f^{-1}(14]$



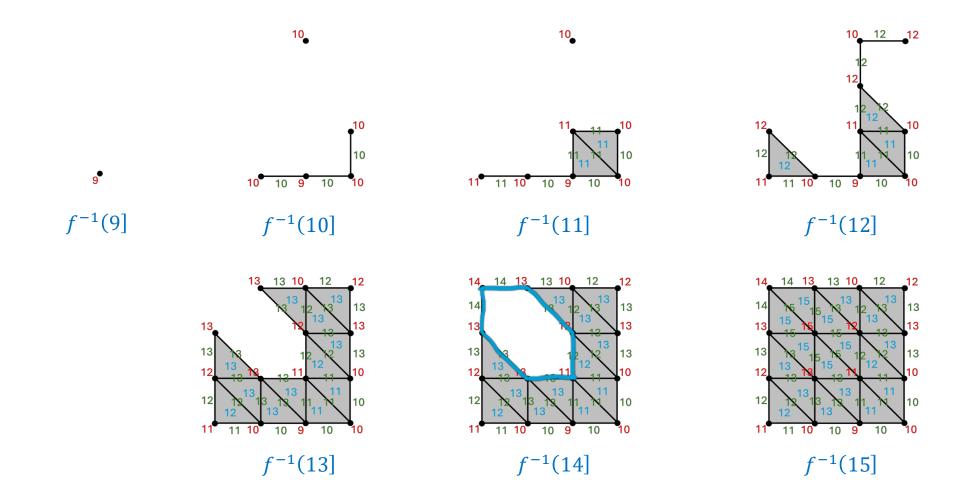
• $f^{-1}(9] \subseteq f^{-1}(10] \subseteq f^{-1}(11] \subseteq f^{-1}(12] \subseteq f^{-1}(13] \subseteq f^{-1}(14] \subseteq f^{-1}(15]$



PD for the sublevelset filtration

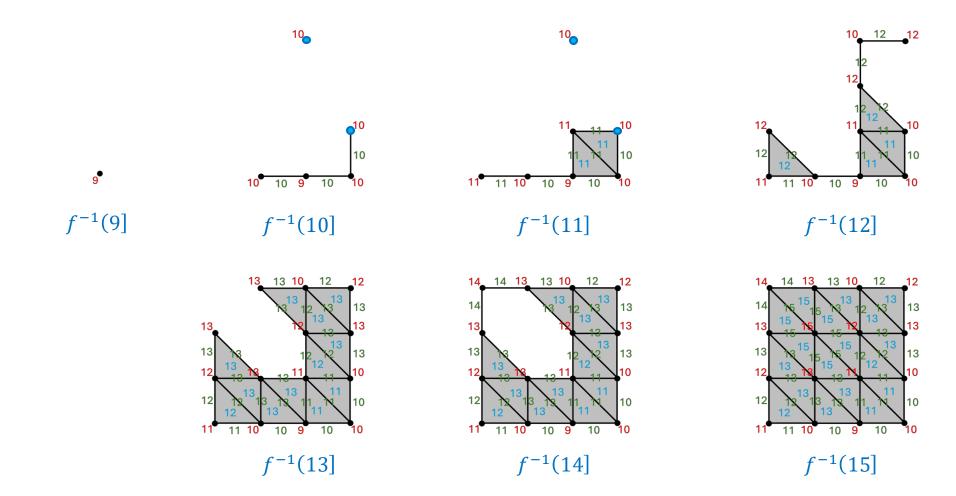
PD for the sublevelset filtration

• There is a 1-dimensional bar [14,15) in the PD



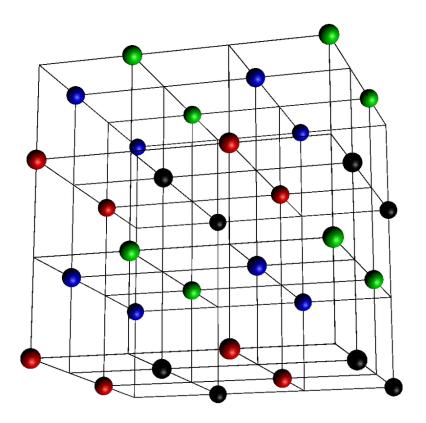
PD for the sublevelset filtration

• There is a 0-dimensional bar [10,12) in the PD



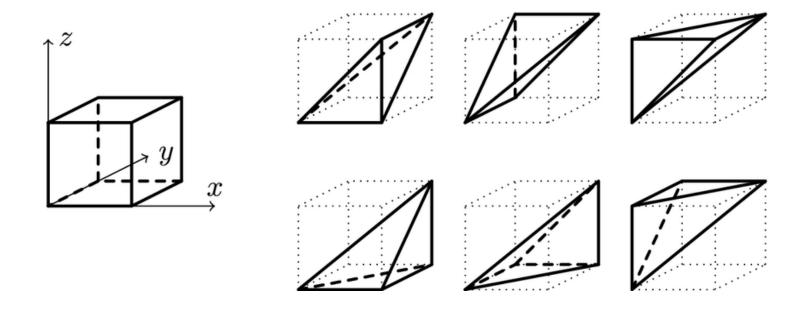
3D image

 We view the domain X for a 3D image as a 3D grid, and we have a function value on each grid point



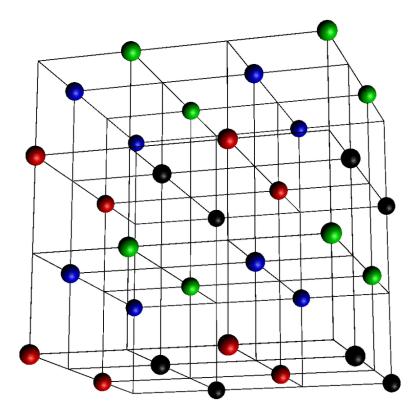
3D image

 We also need to subdivide the cube into (six) tetrahedra to make the domain a simplicial complex



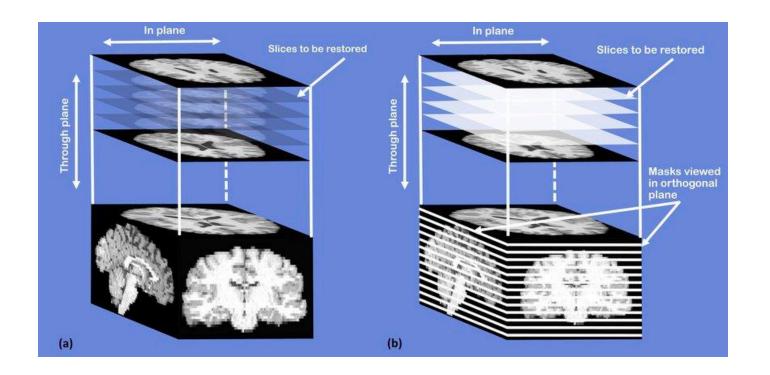
3D image

- And then we only need to assign value to each edge, triangle, tetrahedron based on the maximum values of their vertices
- The sublevelset filtration can then be defined similarly



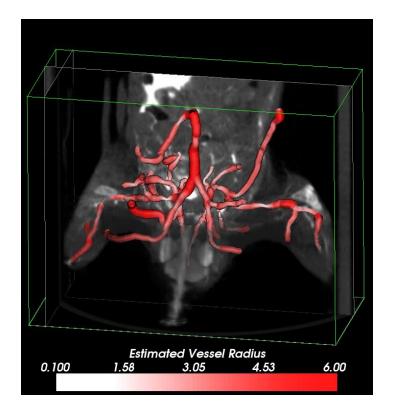
More about 3D images

• 3D images can be considered as a stacking of several 2D images, and are commonly used in medical imaging (e.g., CT-scans, MRI)

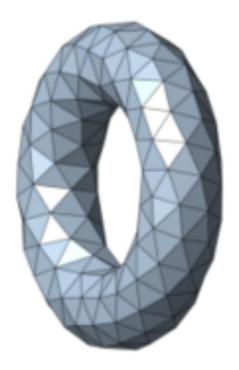


More about 3D images

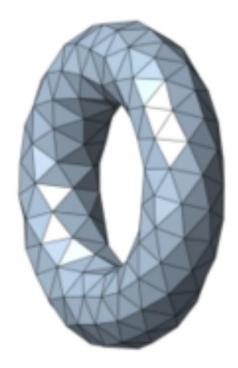
- 3D images can be considered as a stacking of several 2D images, and are commonly used in medical imaging (e.g., CT-scans, MRI)
- Analyzing medical images is a hot and important in image processing

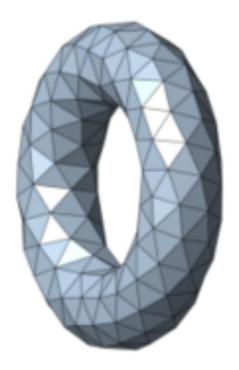


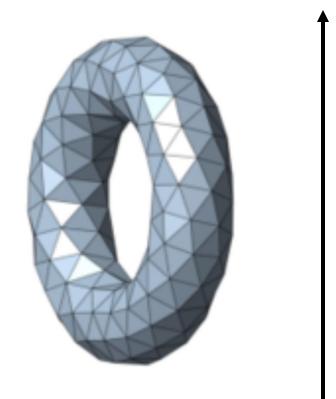
• Naturally, we could also define sublevelset filtrations on triangular meshes by assigning function values to the vertices (edges / triangles are then induced)

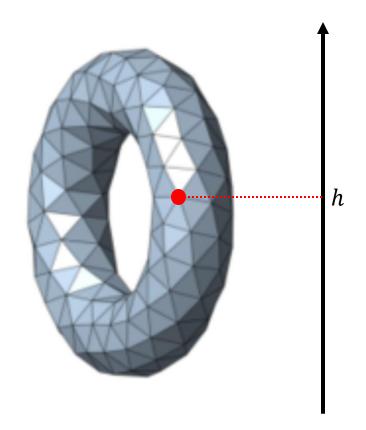


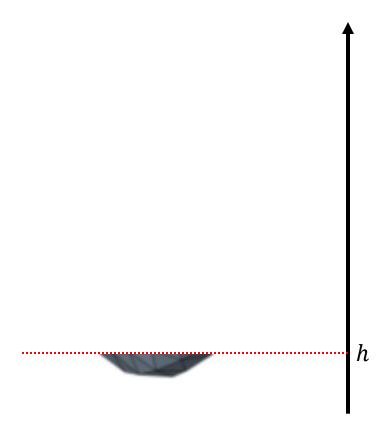
- Naturally, we could also define sublevelset filtrations on triangular meshes by assigning function values to the vertices (edges / triangles are then induced)
- There is a natural way to assign values to the vertices which is to use the "height function"

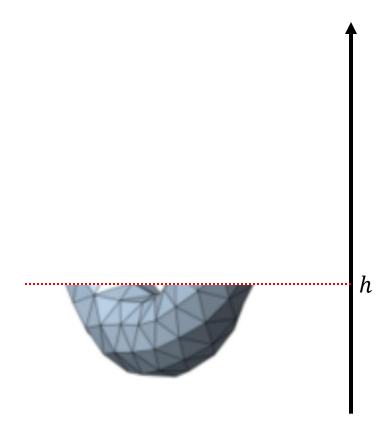


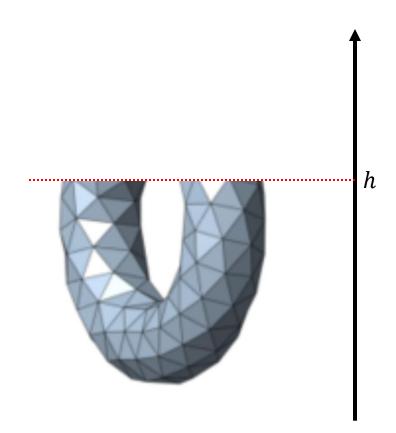


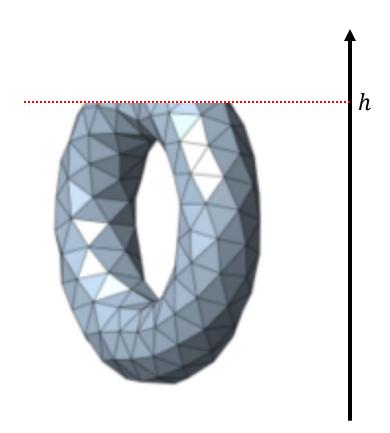












More sublevelset filtrations

- Indeed we have also seen sublevelset filtrations in previous slides
- An interactive example: https://iuricichf.github.io/ICT/filtration.html

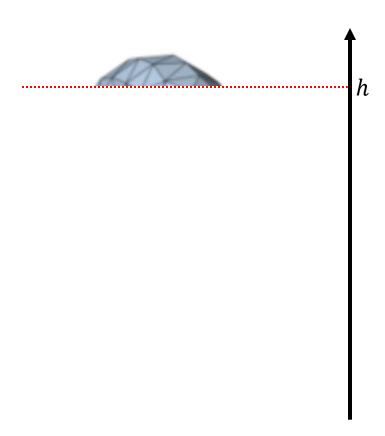
- There is a counterpart of sublevelset filtration called superlevelset filtration
- A superlevelset of is the subset of X whose function values are greater than or equal to a value α , and we denote it as $f^{-1}[\alpha, \infty)$
- We then take all possible functions values and descreasingly sort them (i.e, start with the height value):

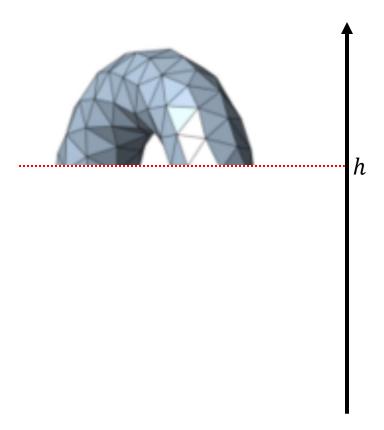
$$\alpha_0 > \alpha_1 > \dots > \alpha_m$$

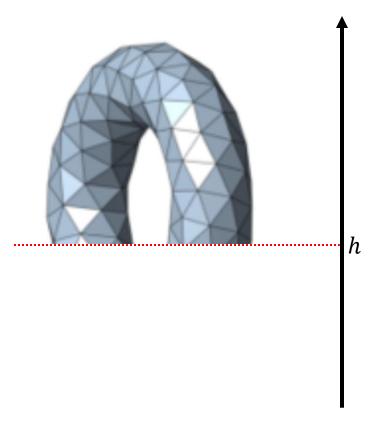
• The superlevelset filtration is then the superlevelsets over the above values:

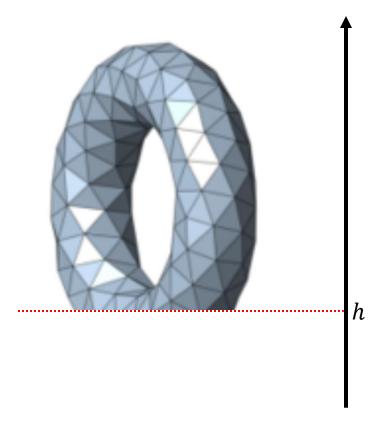
$$f^{-1}[\alpha_0,\infty)\subseteq f^{-1}[\alpha_1,\infty)\subseteq\cdots\subseteq f^{-1}[\alpha_m,\infty)$$











• Previously when we define PD by computing it from a discrete filtration:

$$\mathcal{F}\colon K_0\subseteq K_1\subseteq\cdots\subseteq K_m=K$$

intervals in the PD are "integer intervals" (e.g., $[3, 6) = \{3, 4, 5\}$).

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- In practice, filtrations are built from different types of data. Each complex in the discrete filtration is associate with a real value (or a bunch of them)

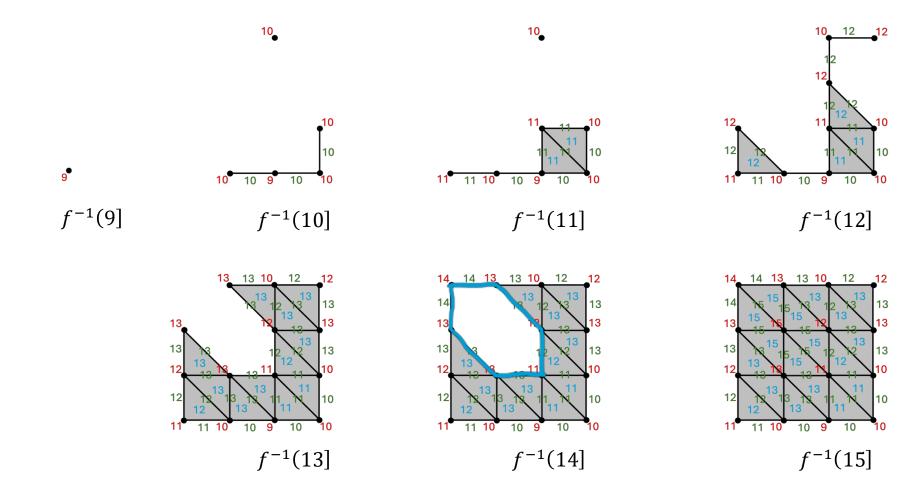
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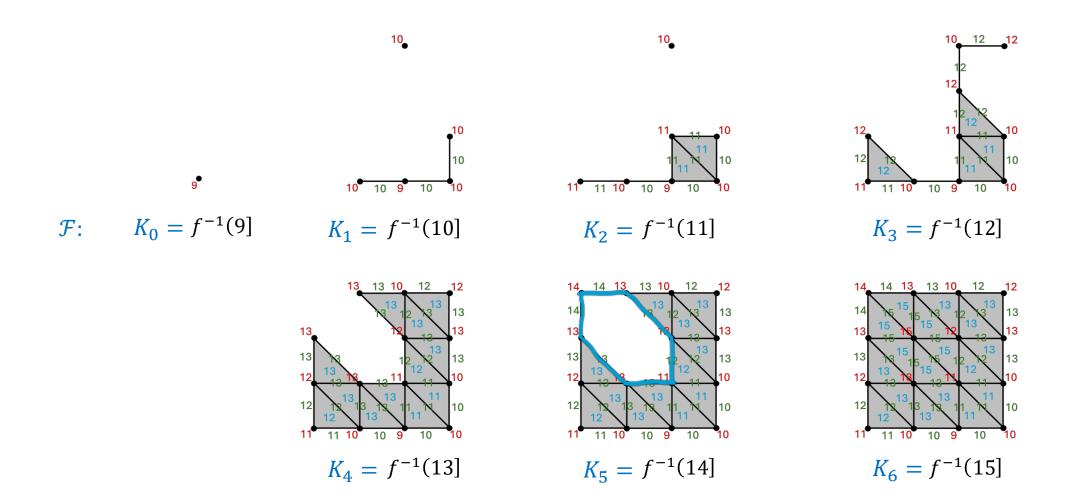
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- This applies when it is not clear where the discrete filtration is built from
- In practice, filtrations are built from different types of data. Each complex in the discrete filtration is associate with a real value (or a bunch of them)
- Intervals in the PD for such a filtration (when we know source data) is then continuous intervals of real values (e.g., [3.52, 6.37))

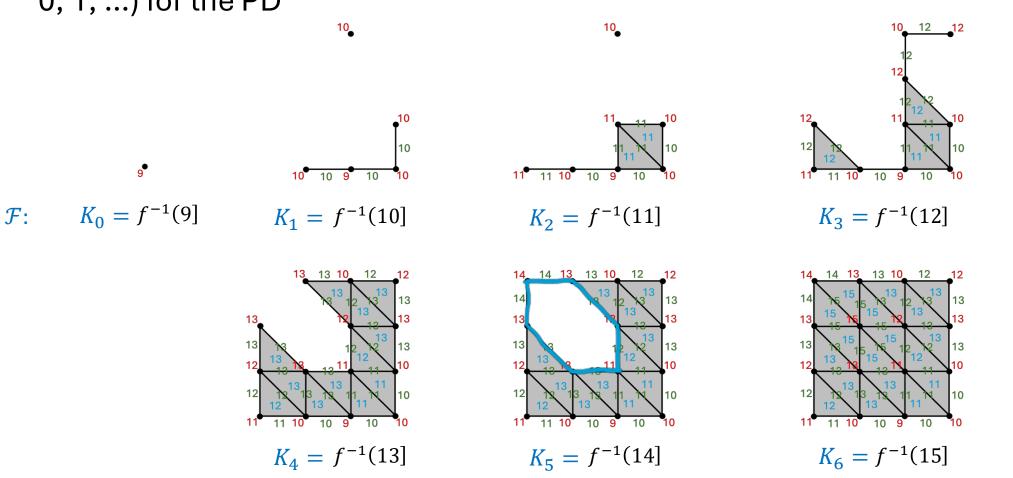
• For the previous sublevelset filtration for image, we can also number each complex in the filtration from 0 to 6



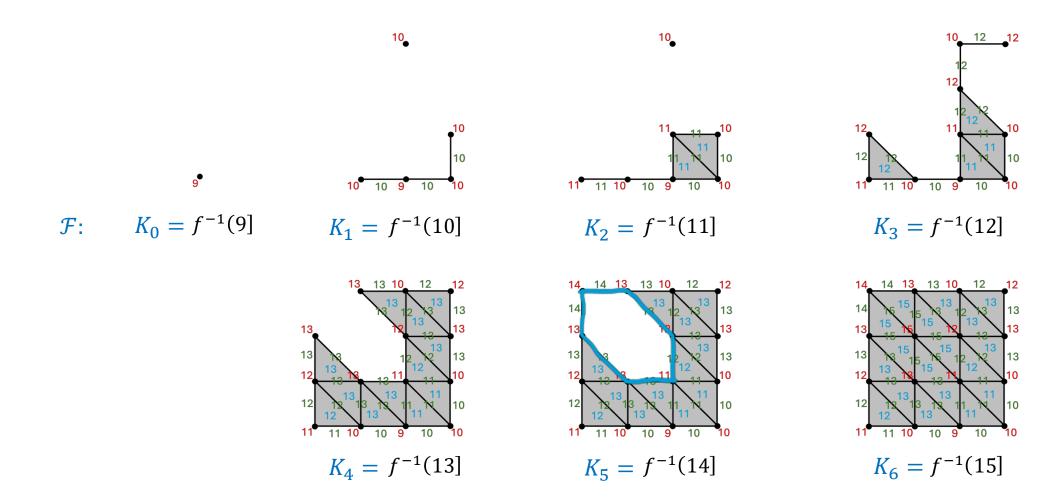
• For the previous sublevelset filtration for image, we can also number each complex in the filtration from 0 to 6



- For the previous sublevelset filtration for image, we can also number each complex in the filtration from 0 to 6
- But we use the pixel values (e.g., 9, 10, ...) instead of the integer indices (e.g., 0, 1, ...) for the PD



• E.g., the below 1d interval is [14,15) rather than [5,6)



• The below 0d interval is [10,12) rather than [1,3)

