

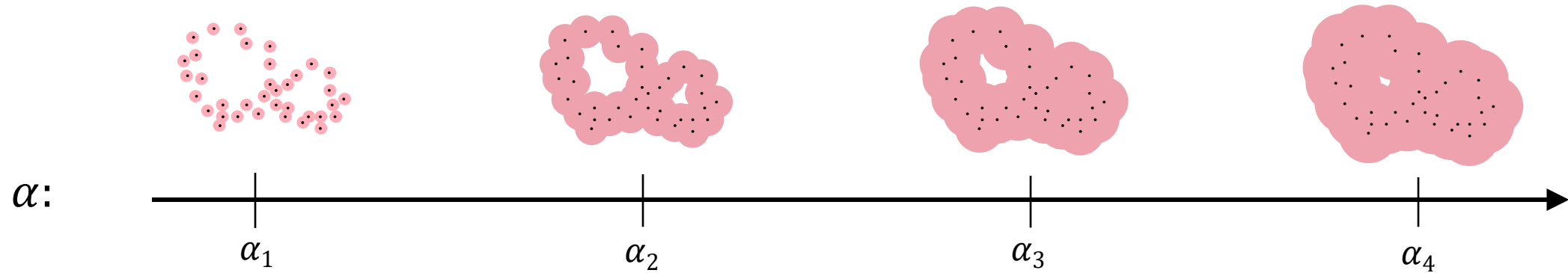
Persistent Homology: Formalization

Tao Hou, University of Oregon

Outline for studying persistent homology

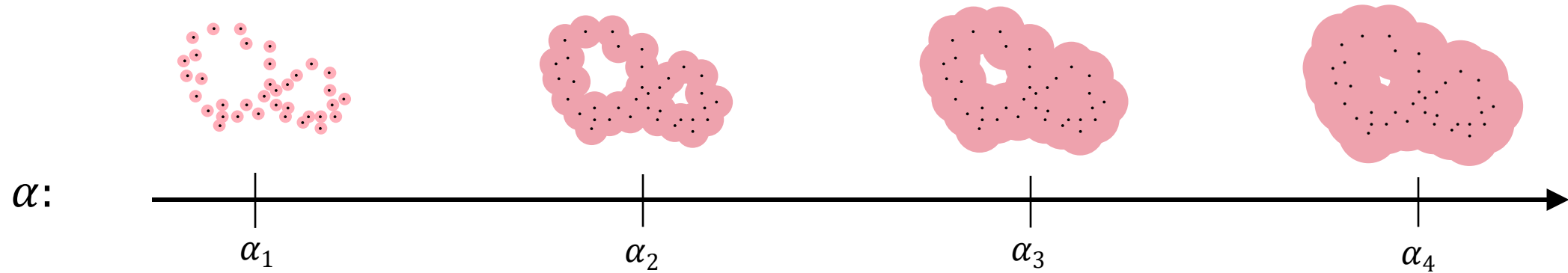
1. Intro to persistent homology
 - Build intuitions of persistent homology: what it does, what it produces
2. Formalizing persistent homology
 - Introduce its input (filtration) and study an algorithm for computation
3. Different ways for building filtrations
 - Vietoris-Rips filtration, sub-levelset filtration
 - Cubical complexes (for images)
4. Interpretation and stability of persistence diagram

“The growing space”



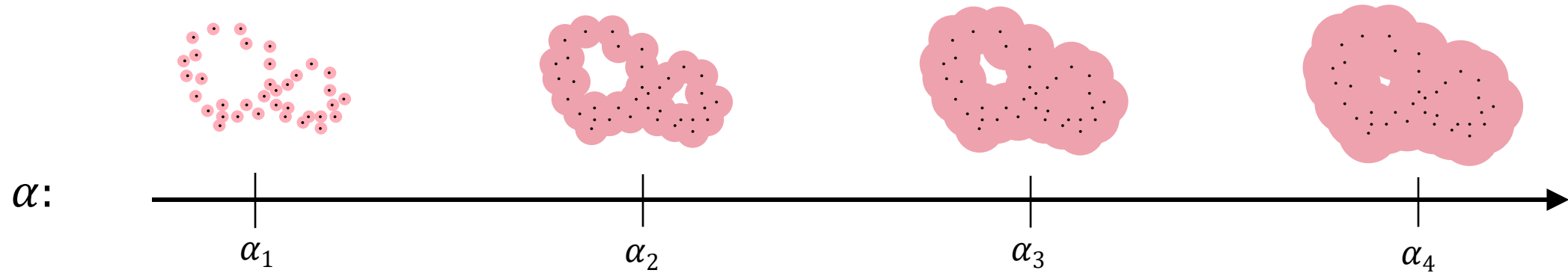
- Recall the growing space:
 - We have a value α ranging within an interval, say, from 0 to ∞
 - Let each value α corresponds to a topological space so that
 - The topological space grows as α increases from 0 to ∞

“The growing space”



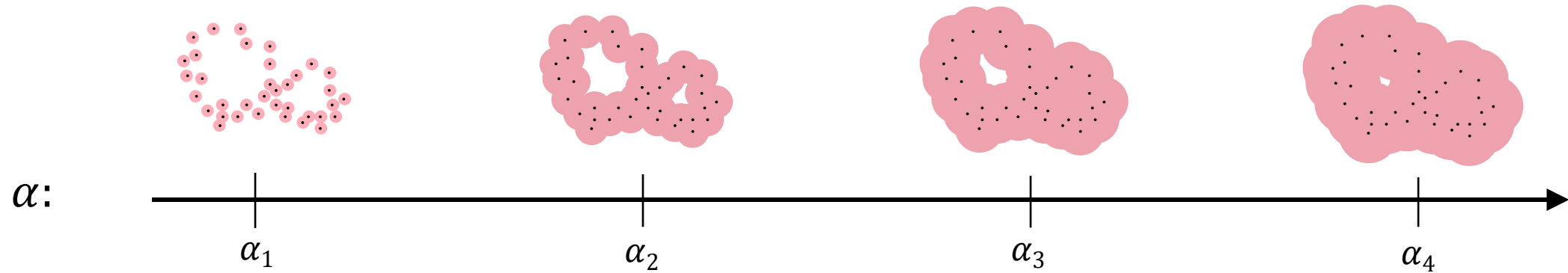
- Recall the growing space:
 - We have a value α ranging within an interval, say, from 0 to ∞
 - Let each value α corresponds to a topological space so that
 - The topological space grows as α increases from 0 to ∞
- Suppose I ask you to represent such a growing space in the computer, can you think of any problems?

“The growing space”



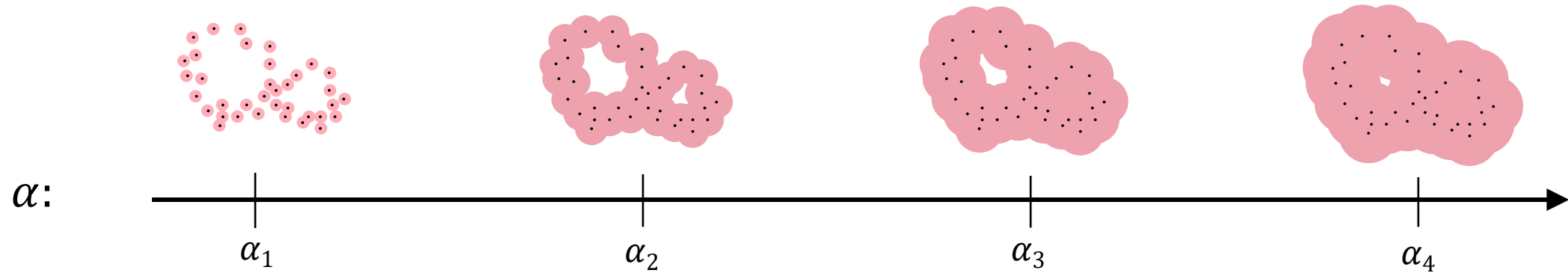
- Problem 1:
 - When α ranges within an interval $[s, f]$, no matter how small the interval is, there are always **infinitely many values** within the interval

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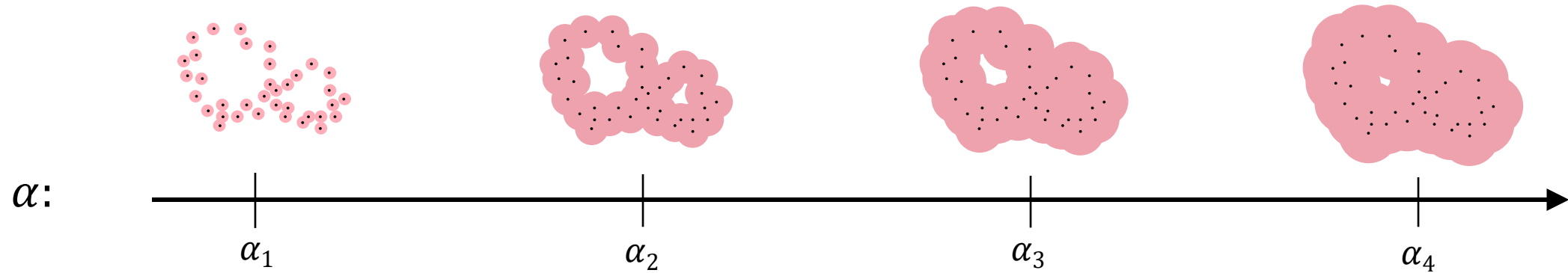
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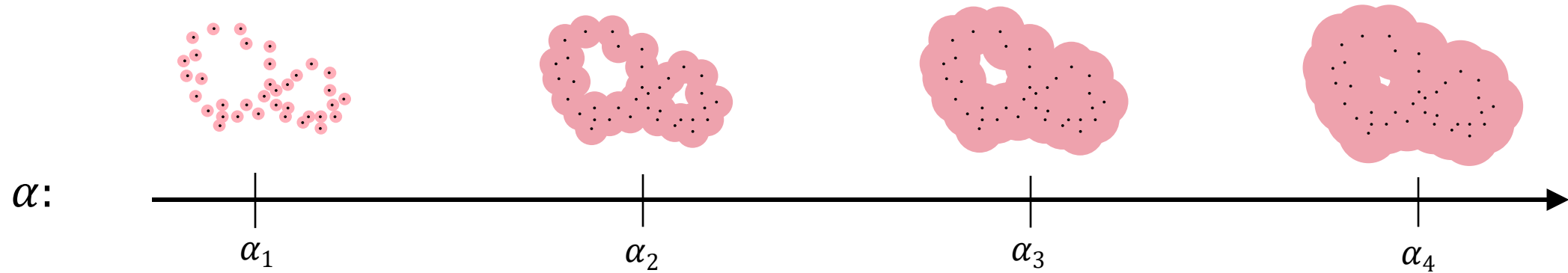
- Problem 1:
 - When α ranges within an interval $[s, f]$, no matter how small the interval is, there are always **infinitely many values** within the interval
 - Each α value may correspond to a possibly different space
 - This means there could be **infinitely many spaces** that we need to store in the computer, which is impossible

“The growing space”



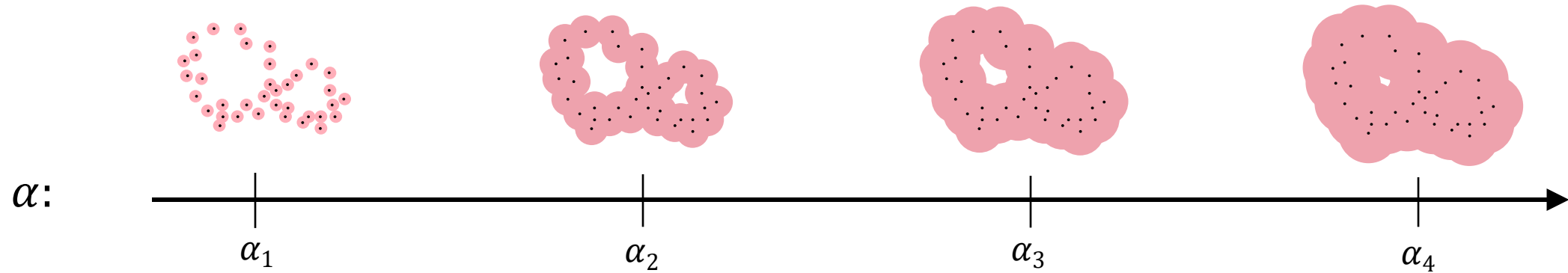
- Solution:
 - While there are infinitely many values for α , our data is still “finite” (e.g., the above point cloud contains finitely many points)

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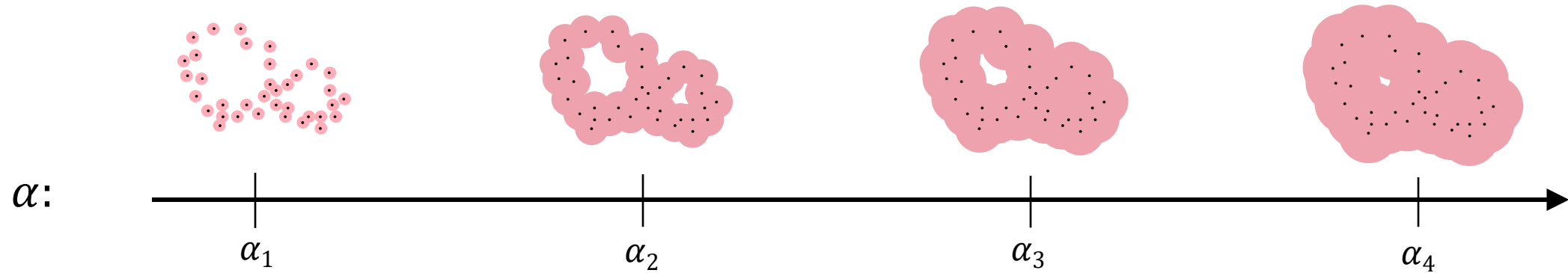
- Solution:
 - While there are infinitely many values for α , our data is still “finite” (e.g., the above point cloud contains finitely many points)
 - This means that there are only finitely many values of α where the topological space “essentially changes”

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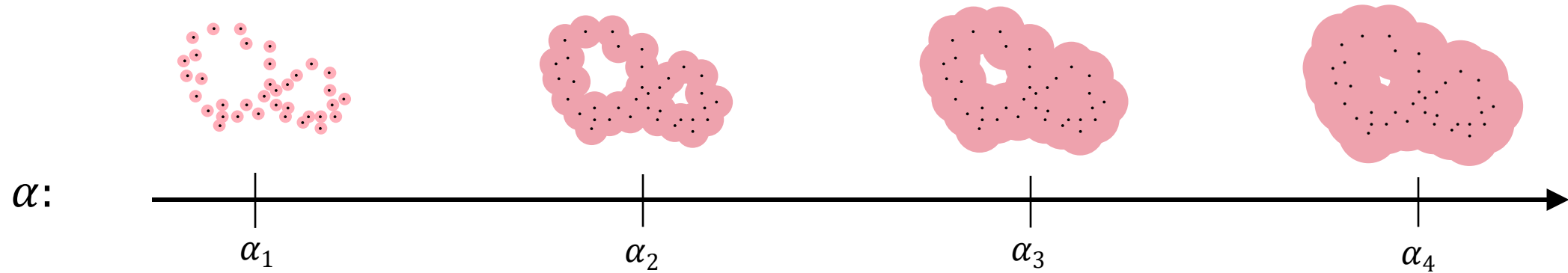
- Solution:
 - While there are infinitely many values for α , our data is still “finite” (e.g., the above point cloud contains finitely many points)
 - This means that there are only finitely many values of α where the topological space “essentially changes”
 - So we only need to record finitely many spaces in computer

“The growing space”



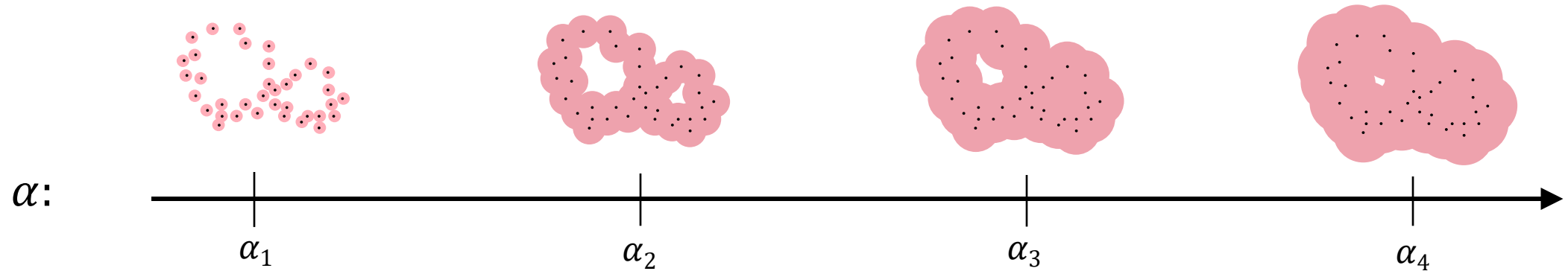
- Remark
 - We will not be very accurate on what the “essential changes” mean here (should be clearer later)

“The growing space”



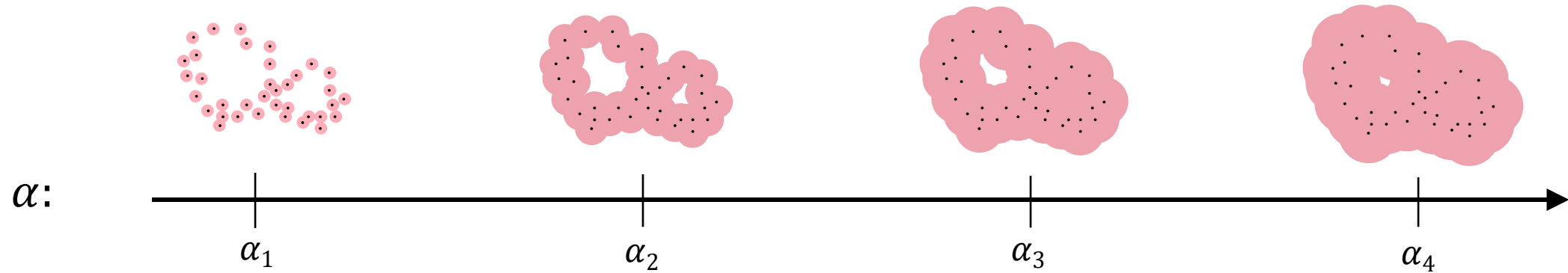
- Remark
 - We will not be very accurate on what the “essential changes” mean here (should be clearer later)
 - BTW, these values where topological space “essentially changes” are called **critical values**
 - Critical values are important concepts in “Morse theory”, but we will not go very deep on it in this course

“The growing space”



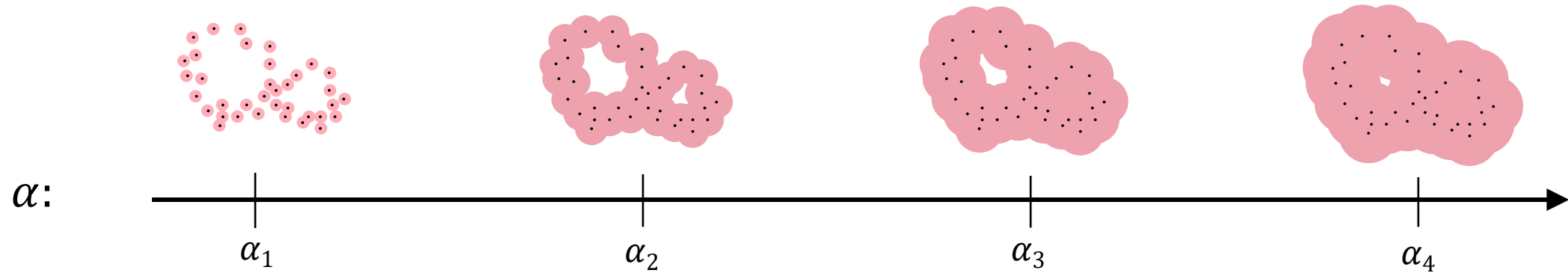
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“The growing space”



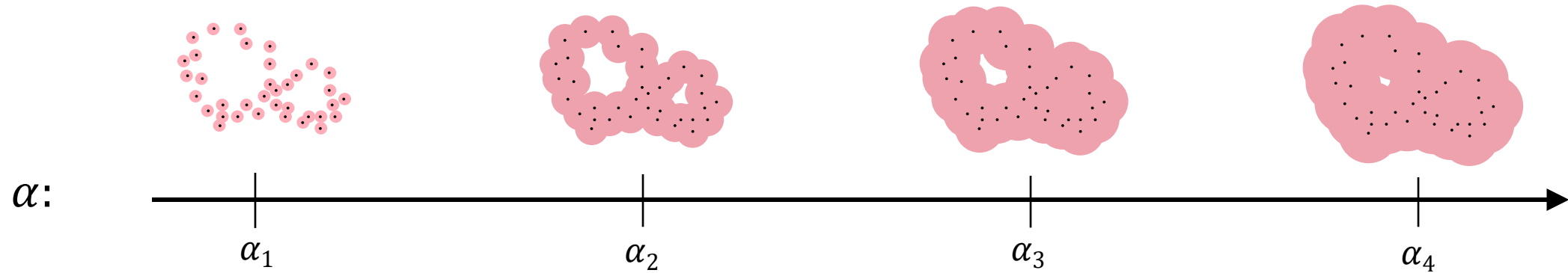
- Problem 2:
 - Even there are finitely many spaces to record, we still need a way to represent each topological space in computer

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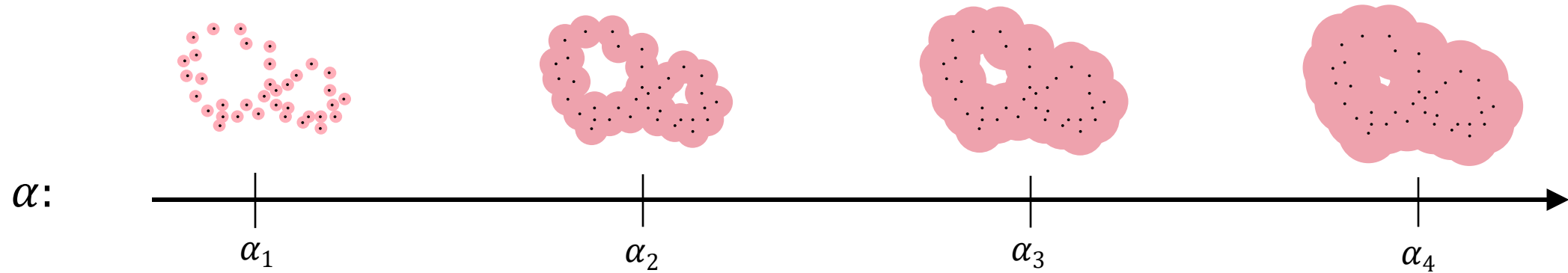
- Problem 2:
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- Solution:

“The growing space”



- Problem 2:
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- Solution:
 - Using simplicial complexes!

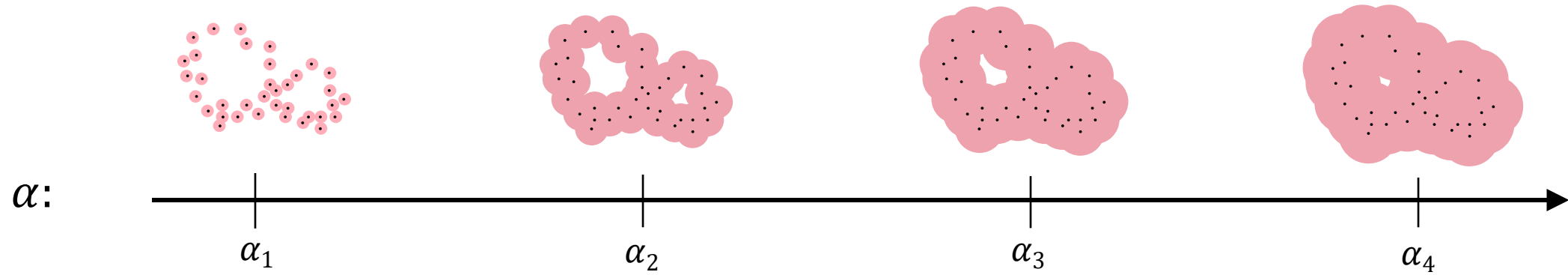
Filtration



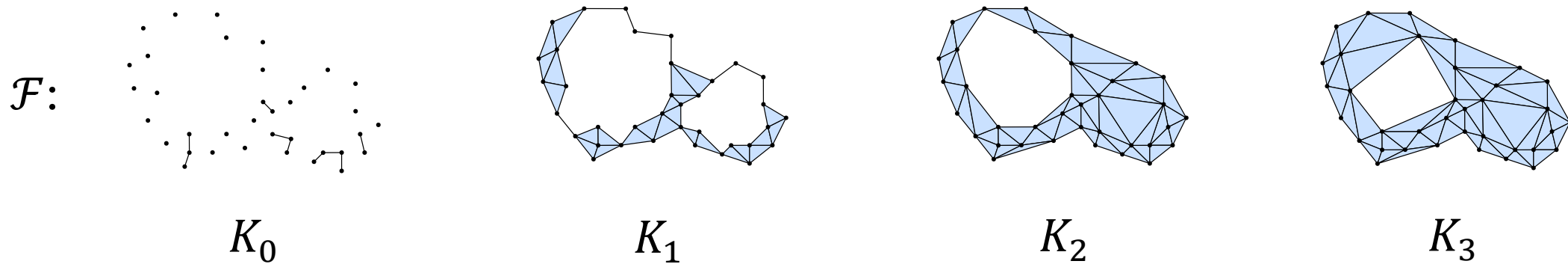
- Hence, the “growing space” in computer is represented by a **finite sequence of simplicial complexes**, called a **filtration**, which is typically denoted by a calligraphic letter \mathcal{F} ,

$$\mathcal{F}: K_0, K_1, \dots, K_m$$

Filtration

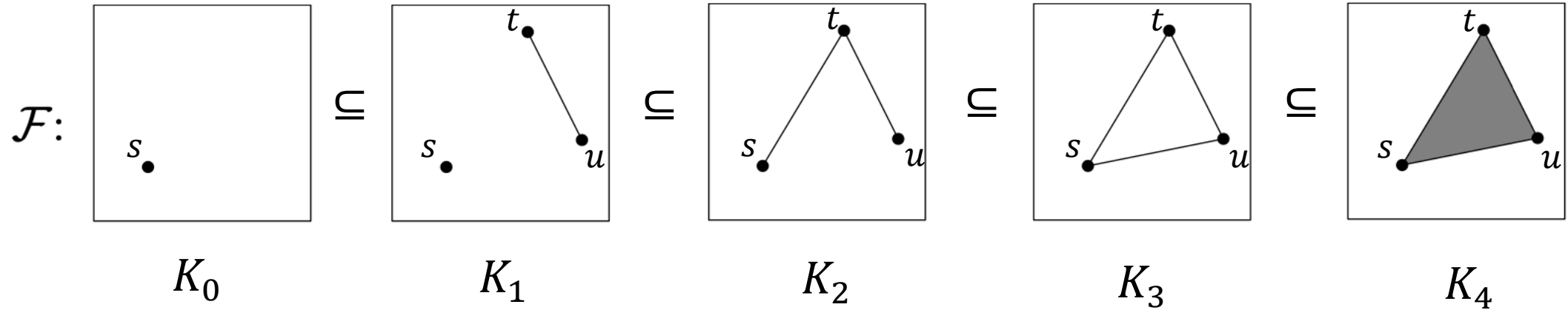


- **Below** is an example of a filtration:



Filtration

- Another example:



Filtration

- Question: In previous definition, a filtration is only a sequence of complexes.
- How do we account for the fact that the spaces (complexes) **grow**?

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- Answer: We make sure the complexes “grow” by making sure the previous complex is a “**subset**” (subcomplex) of the next complex.

Filtration

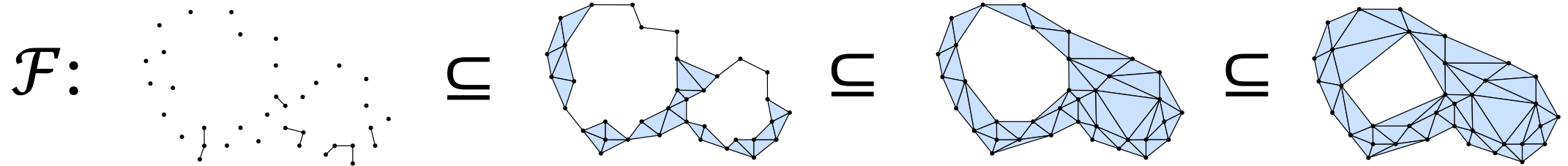
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- How do we account for the fact that the spaces (complexes) **grow**?
- Answer: We make sure the complexes “grow” by making sure the previous complex is a “**subset**” (subcomplex) of the next complex.
- **Definition:** A **filtration** is a nested sequence of simplicial complexes

$$\mathcal{F}: K_0 \subseteq K_1 \subseteq \cdots \subseteq K_m$$

such that each K_i is a subcomplex of K_{i+1} .

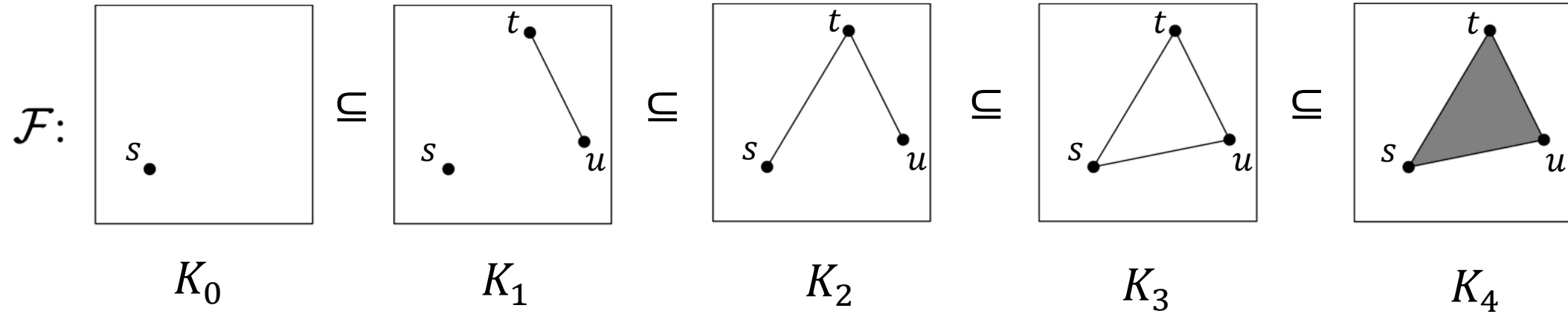
Filtration

- Example:

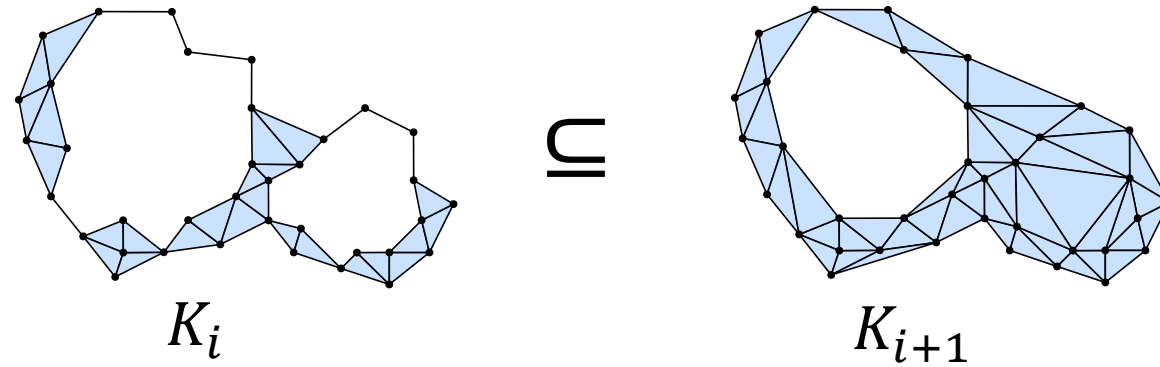


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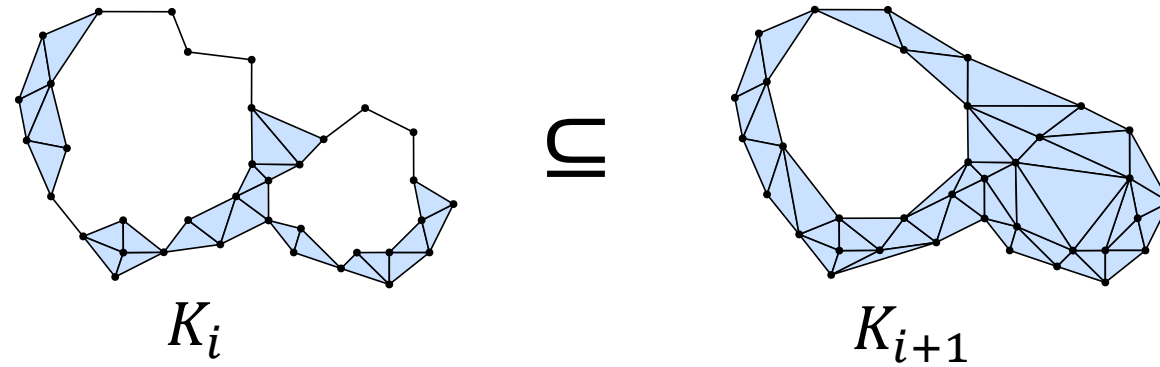


Filtration



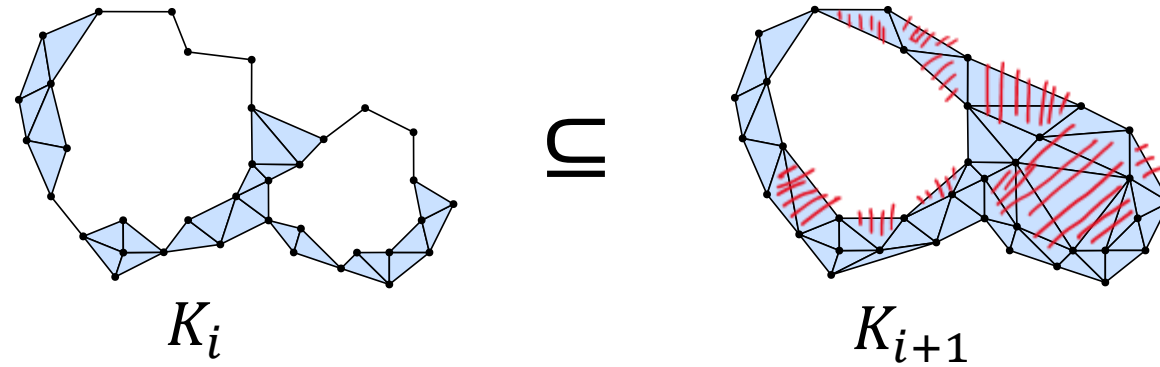
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Filtration



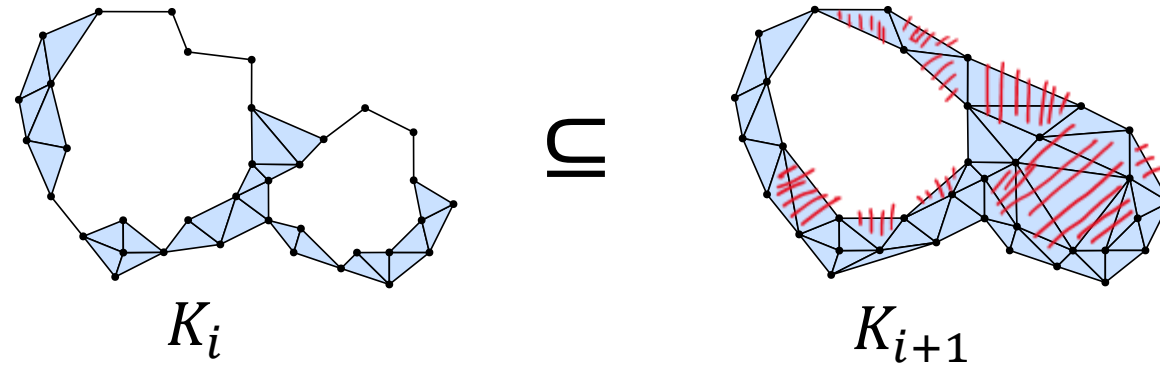
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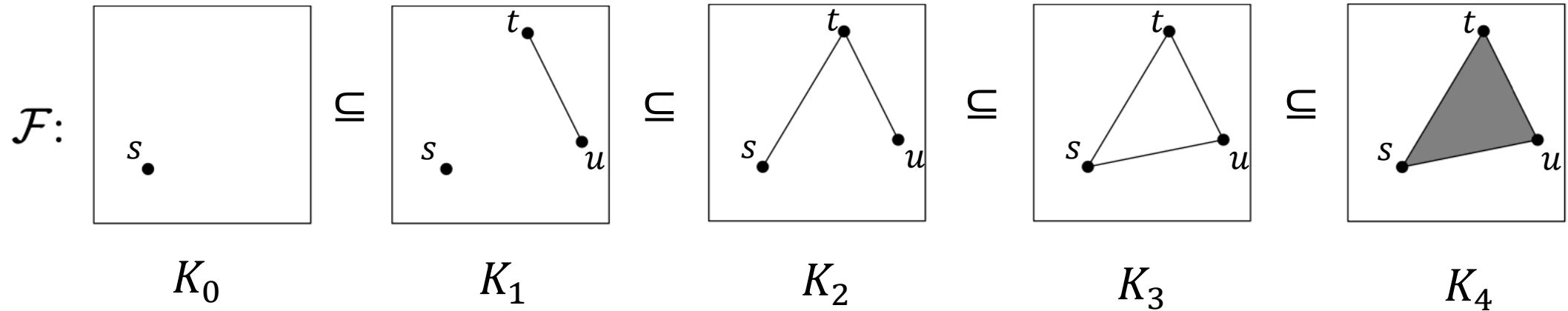
- Now we want to further interpret a filtration
- For this, we focus on a single inclusion in a filtration
- Since it's an inclusion, the difference of the two complexes is that K_{i+1} has some additional simplices than K_i
- So we can consider each inclusion $K_i \subseteq K_{i+1}$ in a filtration

$$\mathcal{F}: K_0 \subseteq K_1 \subseteq \cdots \subseteq K_m$$

as an **insertion of a bunch of simplices**

Filtration

For the example:



- K_0 to K_1 : insert vertices t and u and edge tu
- K_1 to K_2 : insert edge st
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Filtration

- More **regulations**: For a filtration

$$\mathcal{F}: K_0 \subseteq K_1 \subseteq \cdots \subseteq K_m$$

we typically let the first complex K_0 be empty, and call the last complex K_m the “**total complex**” (because it contains all simplices) and denote it as K .

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- **Observation**: (1). Any simplex of K is **added exactly once** in \mathcal{F}
(2). For any two simplices σ and τ in K such that σ is a face of τ , we have σ **cannot be added later** than τ .
- (1) is easy to see. To see (2), suppose that σ is added later than τ . Then at a certain time, τ is already added to a complex K_i but σ is not in K_i yet. This contradicts the fact that any face of a simplex in the complex is also in the complex.

PD for Filtration

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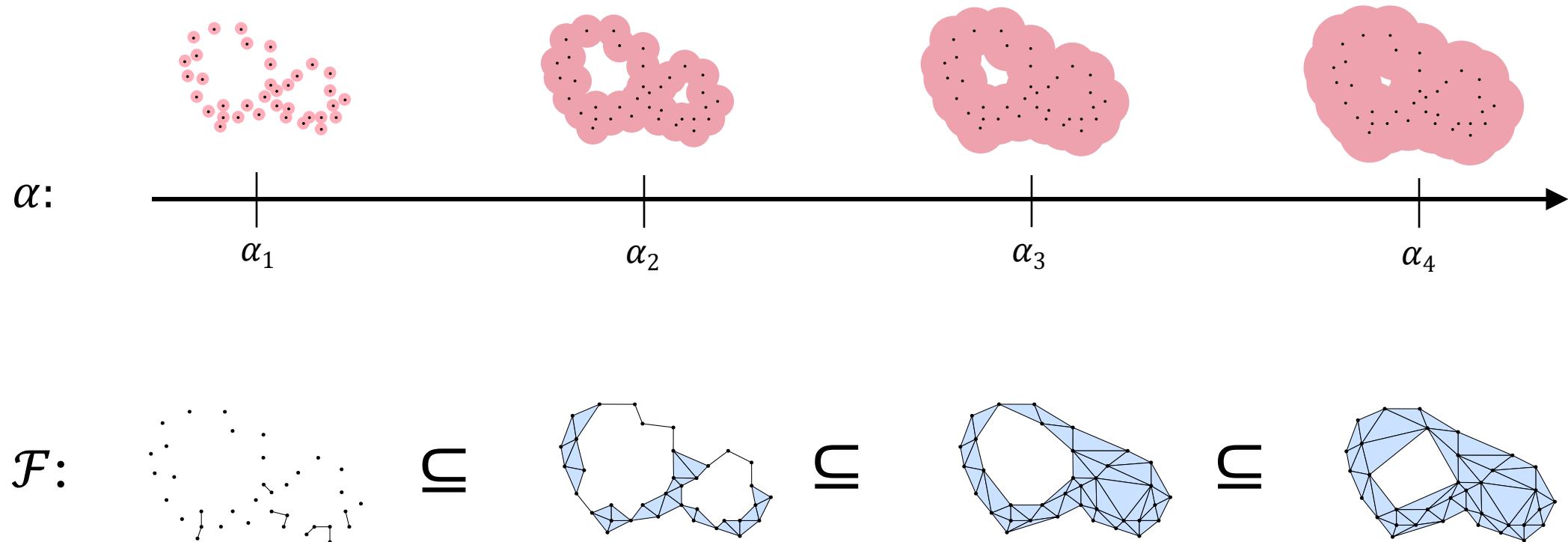
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- Previously, we only saw some examples of PD on a sequence of “growing spaces”, which are not exactly a filtration of complexes.

PD for Filtration

- Filtrations are inputs to the persistent homology pipeline that we want to formalize
- But still we need to **formally define** a PD on a filtration of simplicial complexes
- Previously, we only saw some examples of PD on a sequence of “growing spaces”, which are not exactly a filtration of complexes.
- Moreover, we haven’t really formally defined a PD on a growing space other than showing some examples

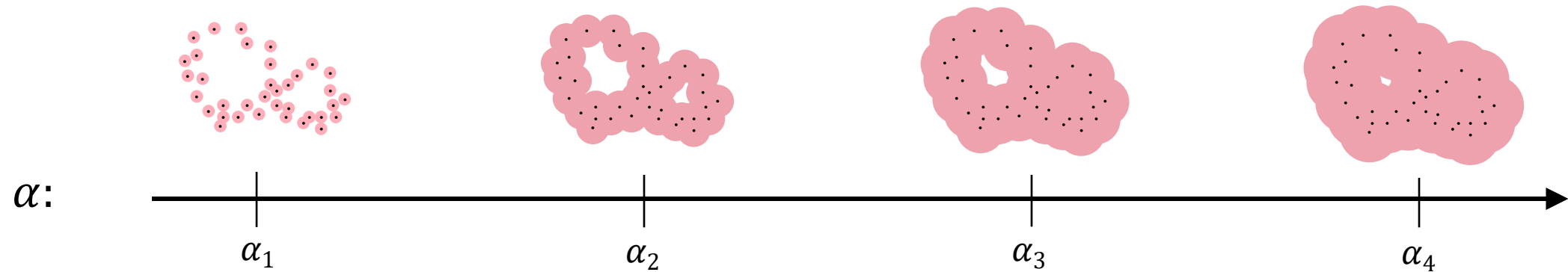
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- Eventually, we will show that, PDs can be formally defined on both a “growing space” (which is **continuous**) and a “filtration of complexes” (which is **discrete**).



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- Eventually, we will show that, PDs can be formally defined on both a “growing space” (which is **continuous**) and a “filtration of complexes” (which is **discrete**).
- We sometimes call the former one a “**continuous**” filtration and latter a “**discrete**” filtration (by default, a “filtration” without modifiers is **always a discrete one**).



“Continuous” filtration



“Discrete” filtration

PD for Filtration

- However, formally defining PD on a continuous or a discrete filtration needs a lot of mathematics (a lot of algebra, category theory, or quiver theory), which is beyond the scope of the course.
- So to understand the definition of a PD, we shall see how to [compute a PD on a discrete filtration](#).
- Things can get a bit technical from now on.

Simplex-wise Filtration

- For computing persistence diagram, we focus on a special type of filtration.
- **Definition:** A **simplex-wise filtration** is a filtration such that each consecutive complexes differ by only a single simplex, i.e., in

$$\mathcal{F}: \emptyset = K_0 \subseteq K_1 \subseteq \cdots \subseteq K_m = K$$

for each inclusion $K_{i-1} \subseteq K_i$, we have that K_i is derived from K_{i-1} by inserting a single simplex typically denoted σ_i .

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- Because of the constructions, we can also consider a simplex-wise filtration

$$\mathcal{F} : \emptyset = K_0 \xrightarrow{\sigma_1} K_1 \xrightarrow{\sigma_2} \cdots \xrightarrow{\sigma_{m-1}} K_{m-1} \xrightarrow{\sigma_m} K_m = K$$

as a sequence of simplices $\sigma_0, \sigma_1, \dots, \sigma_{m-1}$ inserted one by one following the order.

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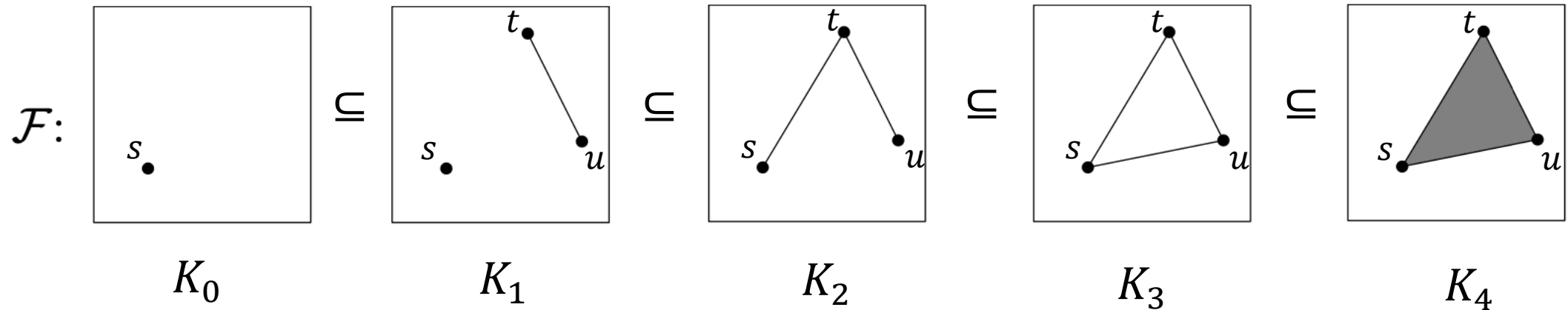
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as a sequence of simplices $\sigma_0, \sigma_1, \dots, \sigma_{m-1}$ inserted one by one following the order.

- **Fact:** Each general filtration (not necessarily simplex-wise) can be made into a simplex-wise one by **padding additional complexes** (or **expanding the inclusions**)

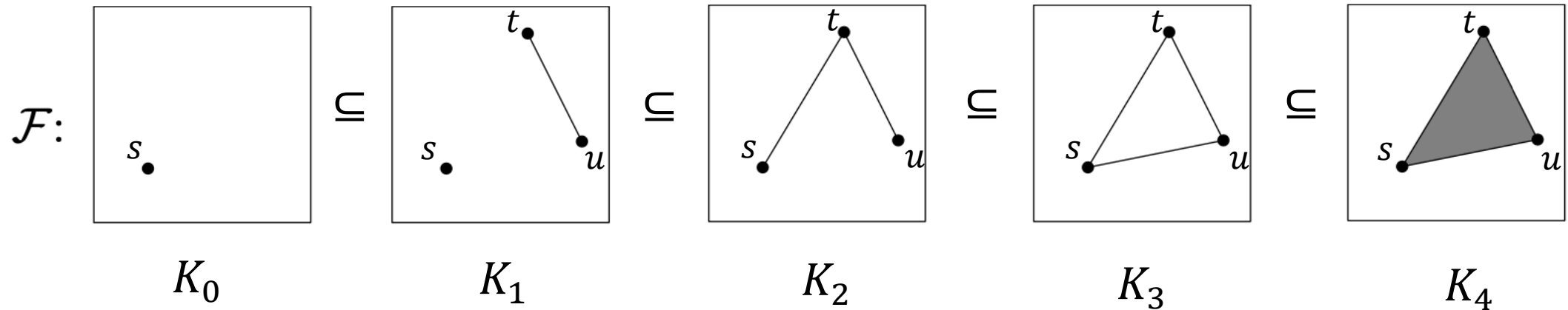
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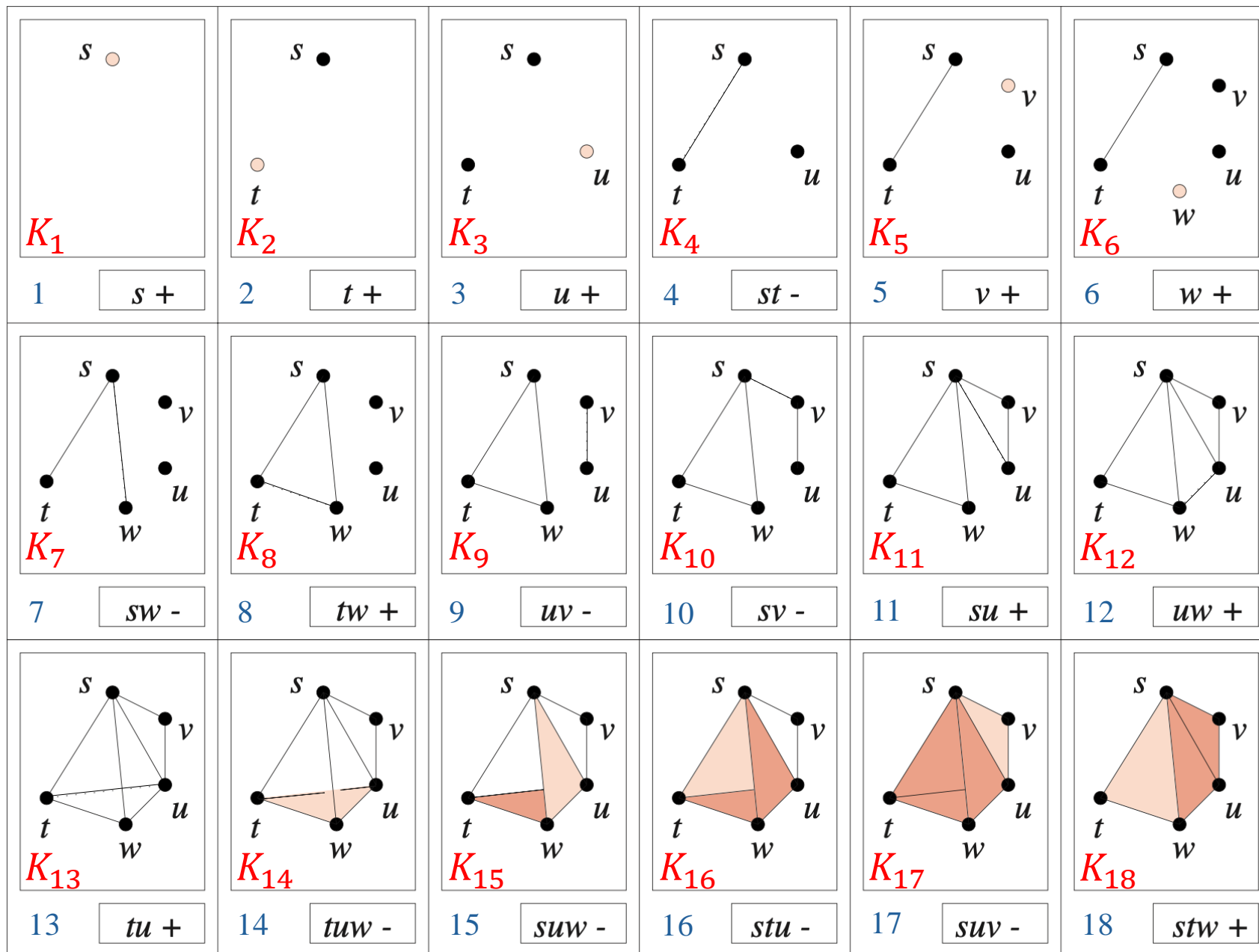


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- To convert to simplex-wise, only need to add an empty complex at the beginning and insert two additional complexes between K_0 to K_1 .



Algorithm

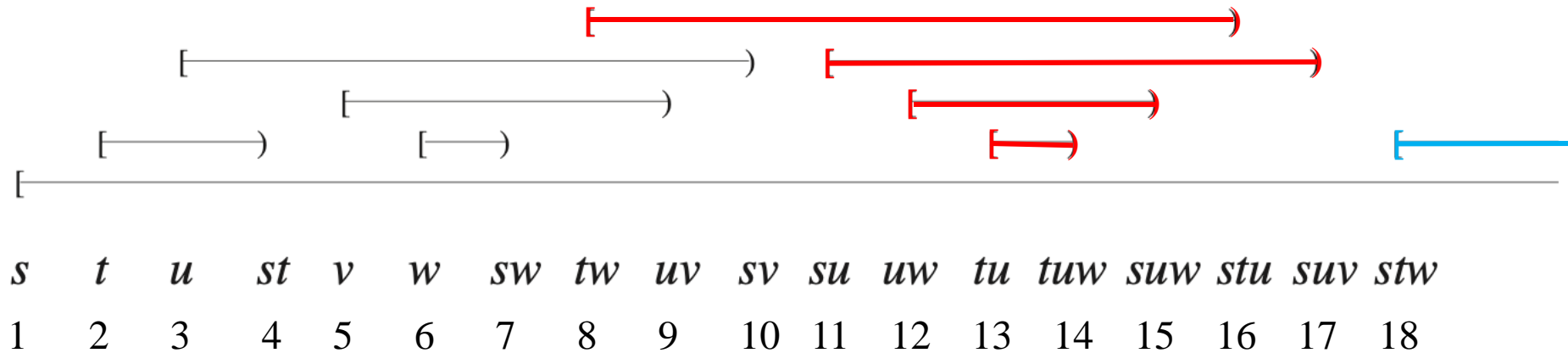
- Notice that the input filtration \mathcal{F} **must be simplex-wise**

Input: a filtration \mathcal{F} as a sequence of simplices $\sigma_1, \sigma_2, \dots, \sigma_m$

Output: p -th PD of \mathcal{F} , $\text{PD}_p(\mathcal{F})$, for each dimension p

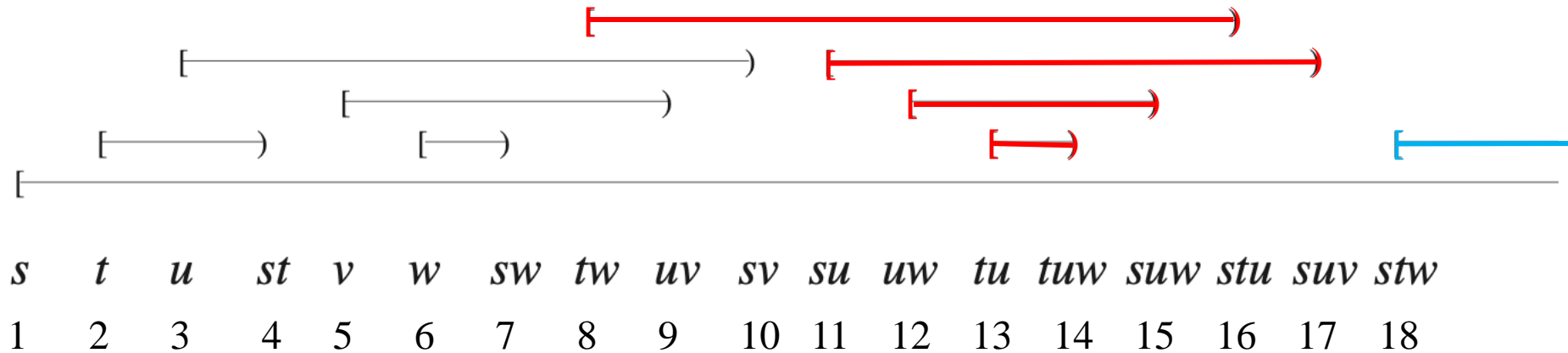
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1: set each  $\sigma_i$  in  $\mathcal{F}$  as “unpaired”
2:  $\zeta$  = a table mapping each  $\sigma_i$  to a cycle  $\zeta(\sigma_i)$  initially undefined
3: for  $\sigma_i = \sigma_1, \sigma_2, \dots, \sigma_m$  do
4:      $z = \partial(\sigma_i)$ 
5:     while  $z \neq 0$  do
6:         let  $\sigma_j$  be the simplex with maximum index in  $z$ 
7:         if  $\sigma_j$  is unpaired then break
8:          $z = z + \zeta(\sigma_j)$ 
9:     if  $z \neq 0$  then
10:        pair  $\sigma_j$  with  $\sigma_i$  and set  $\sigma_j, \sigma_i$  as “paired”
11:         $\zeta(\sigma_j) = z$ 
12:         $p = \text{dimension of } \sigma_j$ 
13:        add  $(j, i)$  to  $\text{PD}_p(\mathcal{F})$ 
14: for each each unpaired  $\sigma_i$  do
15:      $p = \text{dimension of } \sigma_i$ 
16:     add  $(i, \infty)$  to  $\text{PD}_p(\mathcal{F})$ 
```

Resulting PD



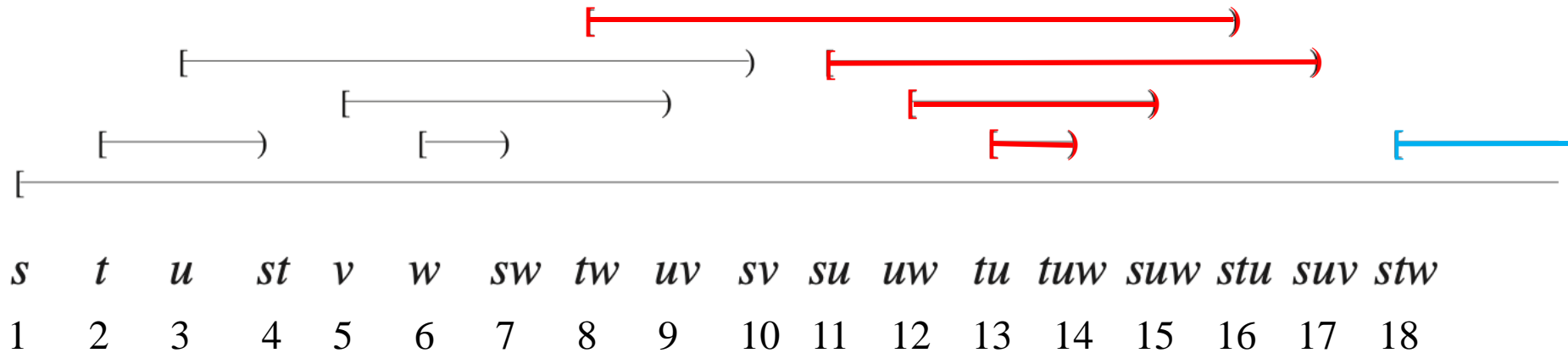
- Black: PD_0
- Red: PD_1
- Blue: PD_2

Resulting PD



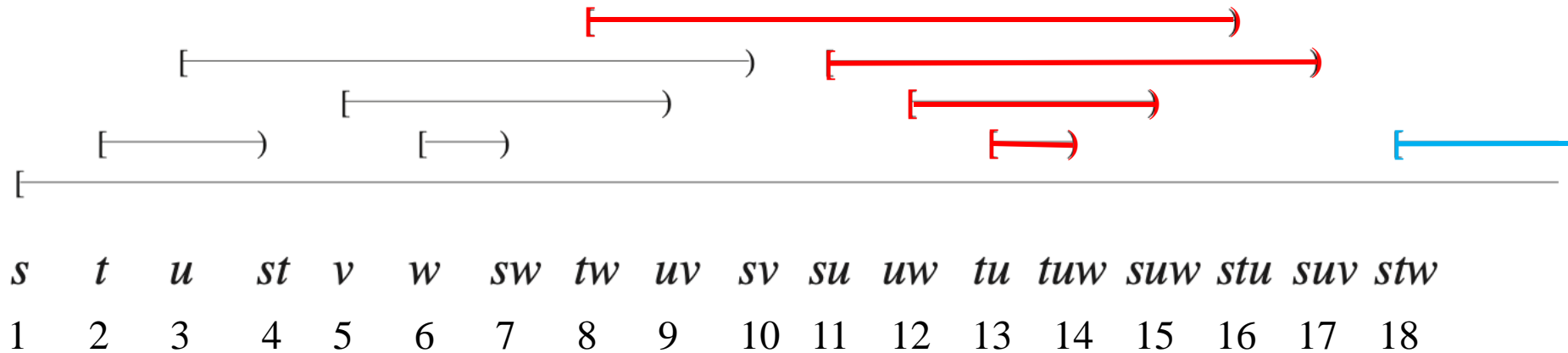
- Notice: instead of drawing each pair of birth / death as a point on 2D plane, we just let each pair of birth and death **form an interval**, indicating the “time” in which a certain homology hole persists (will see examples later)

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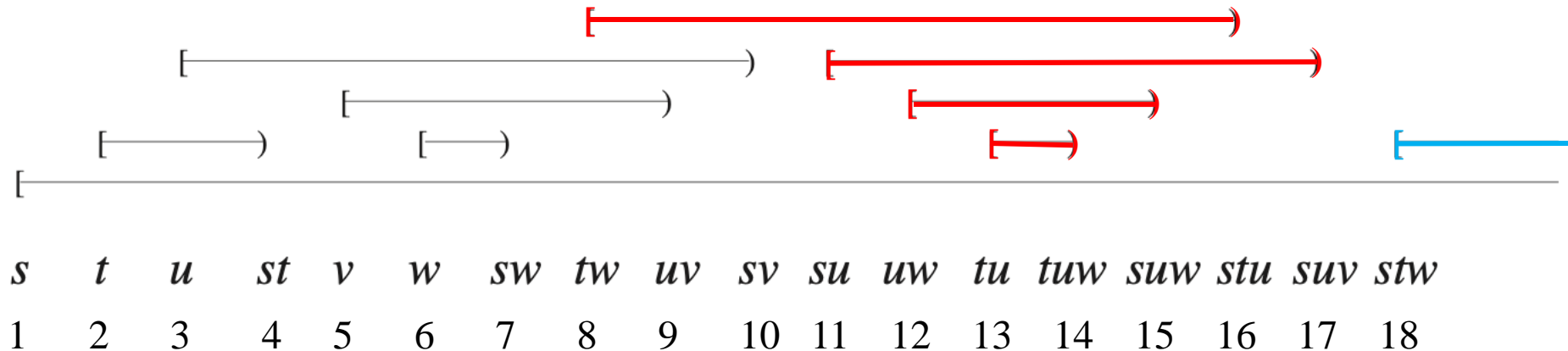
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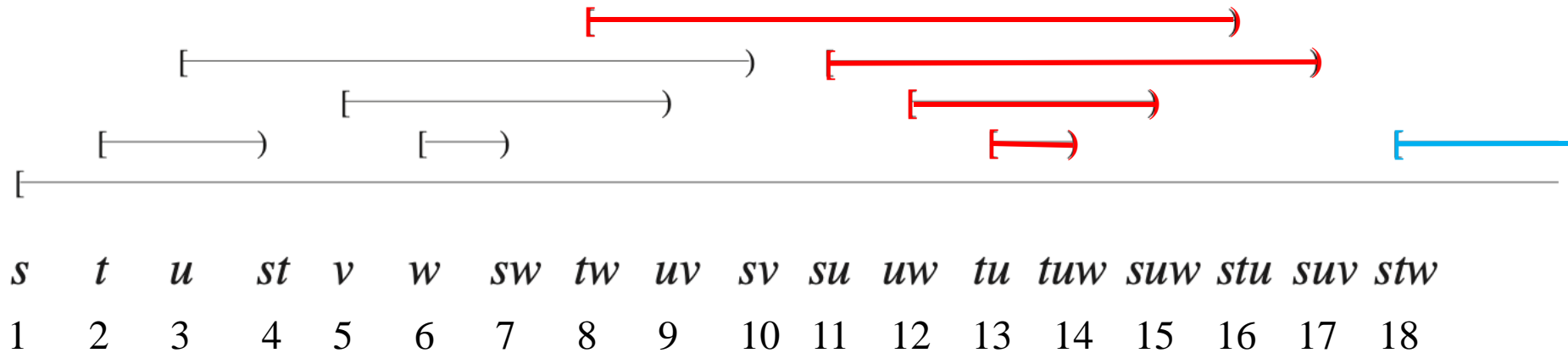
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- So persistence barcodes and persistence diagrams are just the same things displayed in different ways (we sometimes also use the two terms **interchangeably**)

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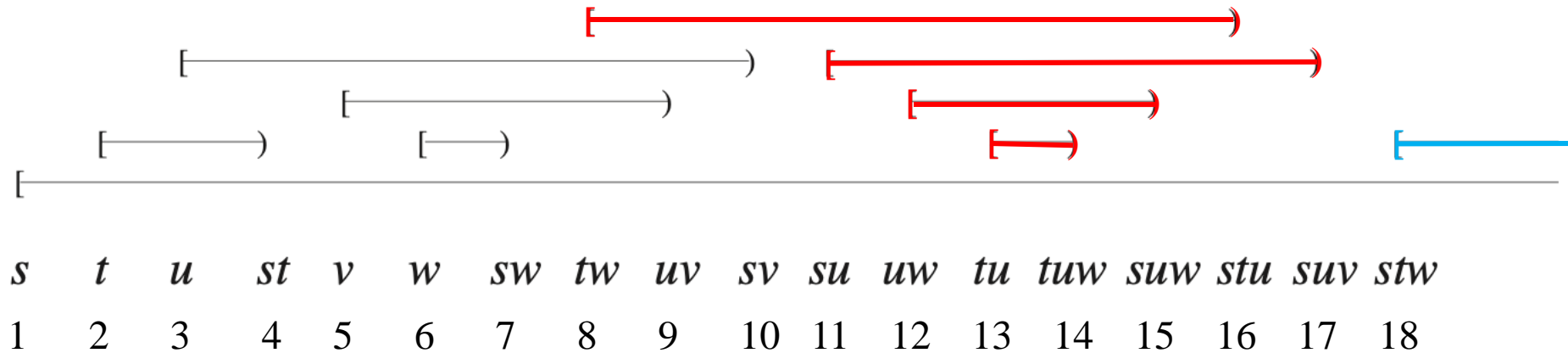
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- The above is also called the **persistence barcode**
- So persistence barcodes and persistence diagrams are just the same things displayed in different ways (we sometimes also use the two terms **interchangeably**)
- Also notice: In persistence barcode, we always draw each interval as **left-closed, right open** (there is a technical reason for this but explaining this a little beyond scope)

Resulting PD



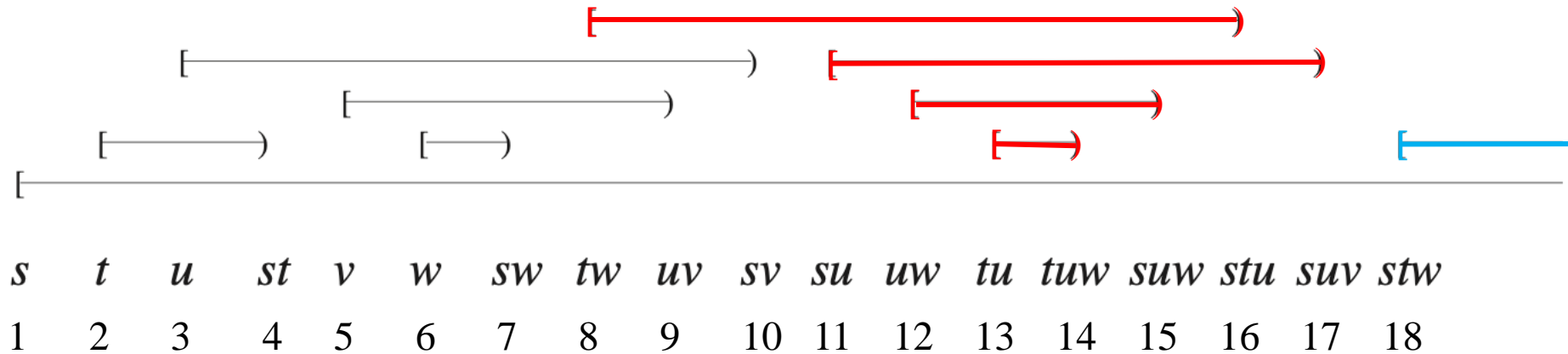
- We also notice that the cycle recorded in the “ ζ table” indeed captures the homology hole born and died with a birth-death interval in the barcode (point in the PD)

Resulting PD

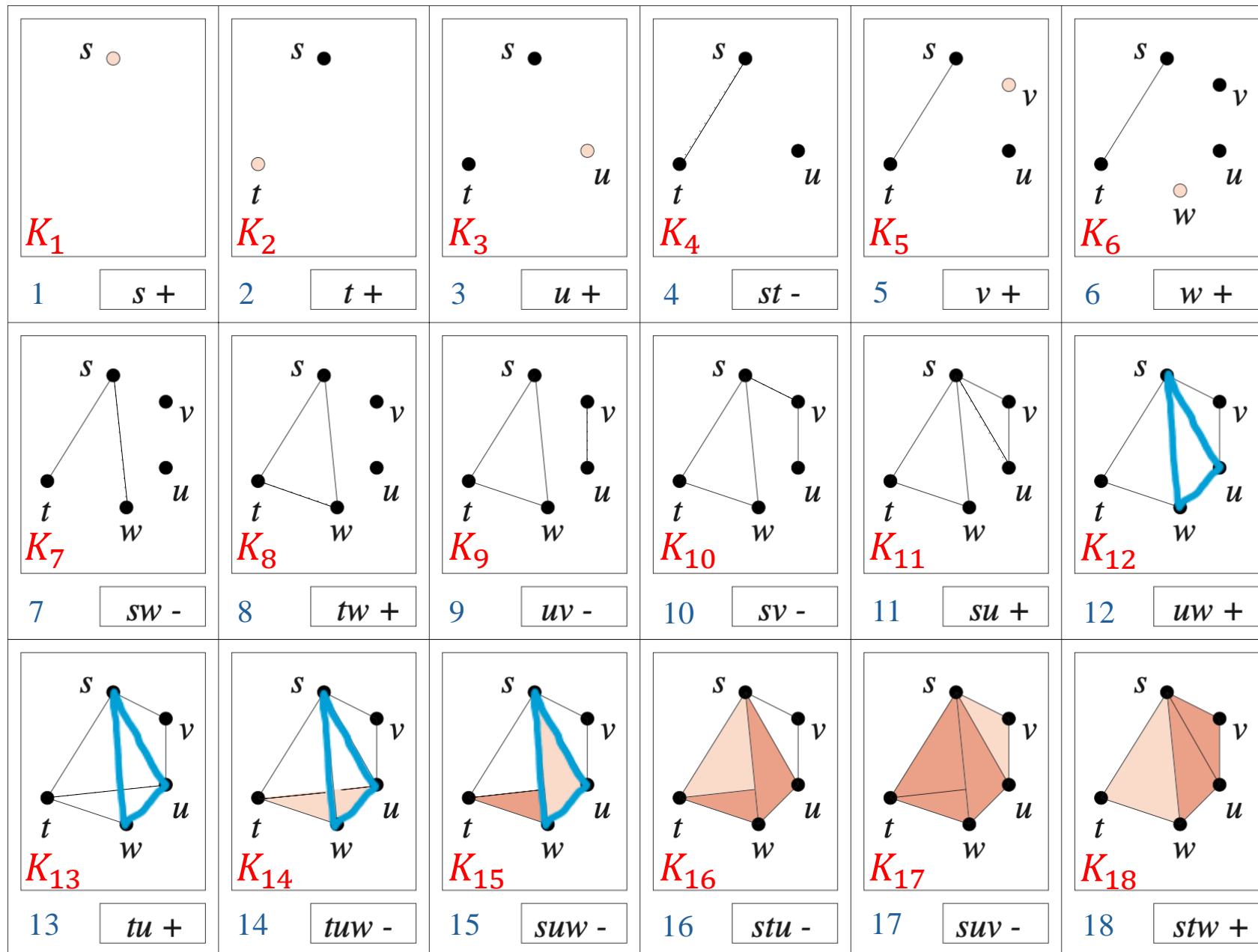


- We also notice that the cycle recorded in the “ ζ table” indeed captures the homology hole born and died with a birth-death interval in the barcode (point in the PD)
- i.e., for an interval $[b, d)$, $\zeta[\sigma_b]$ represents the homology feature born at the index b and dying at index d .

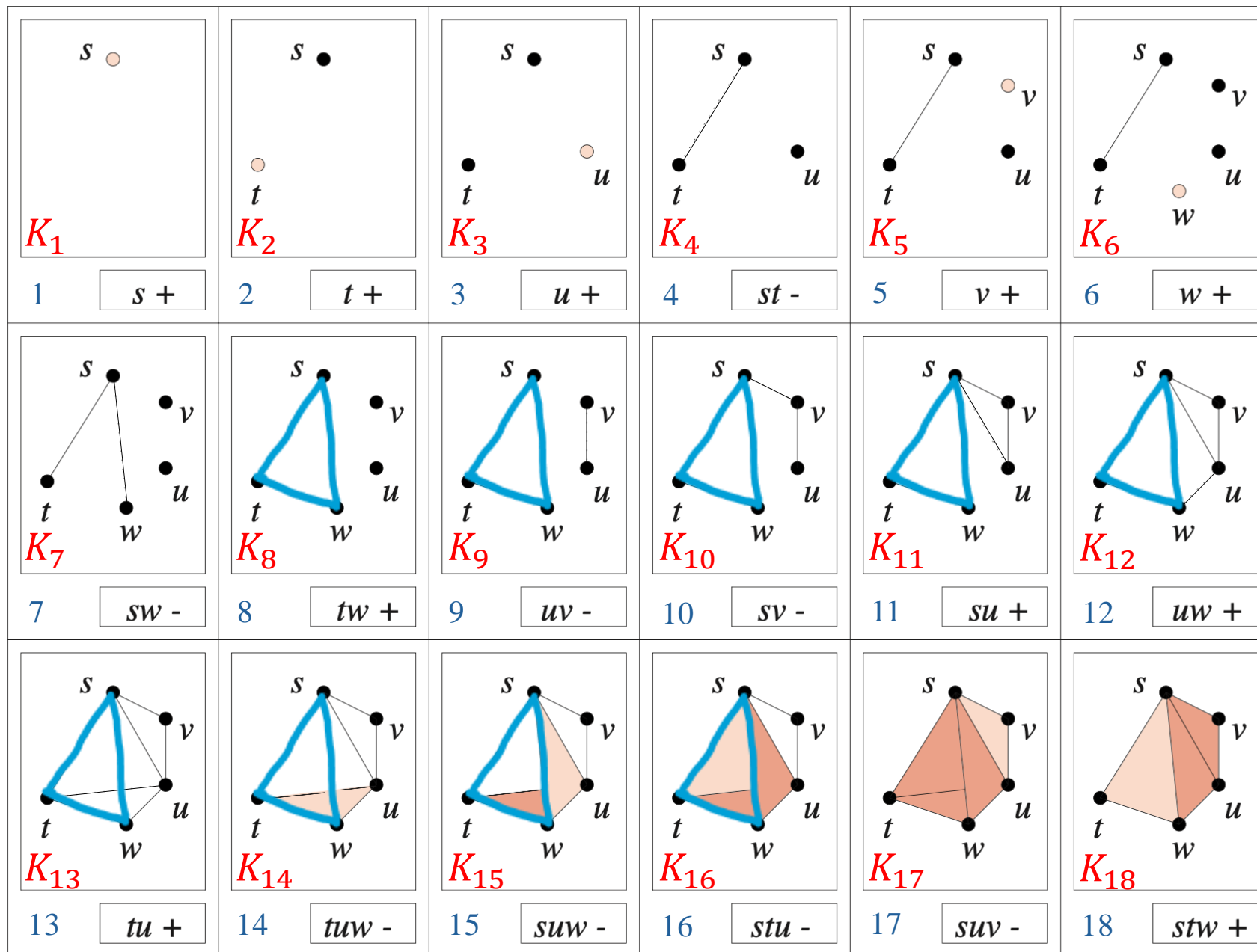
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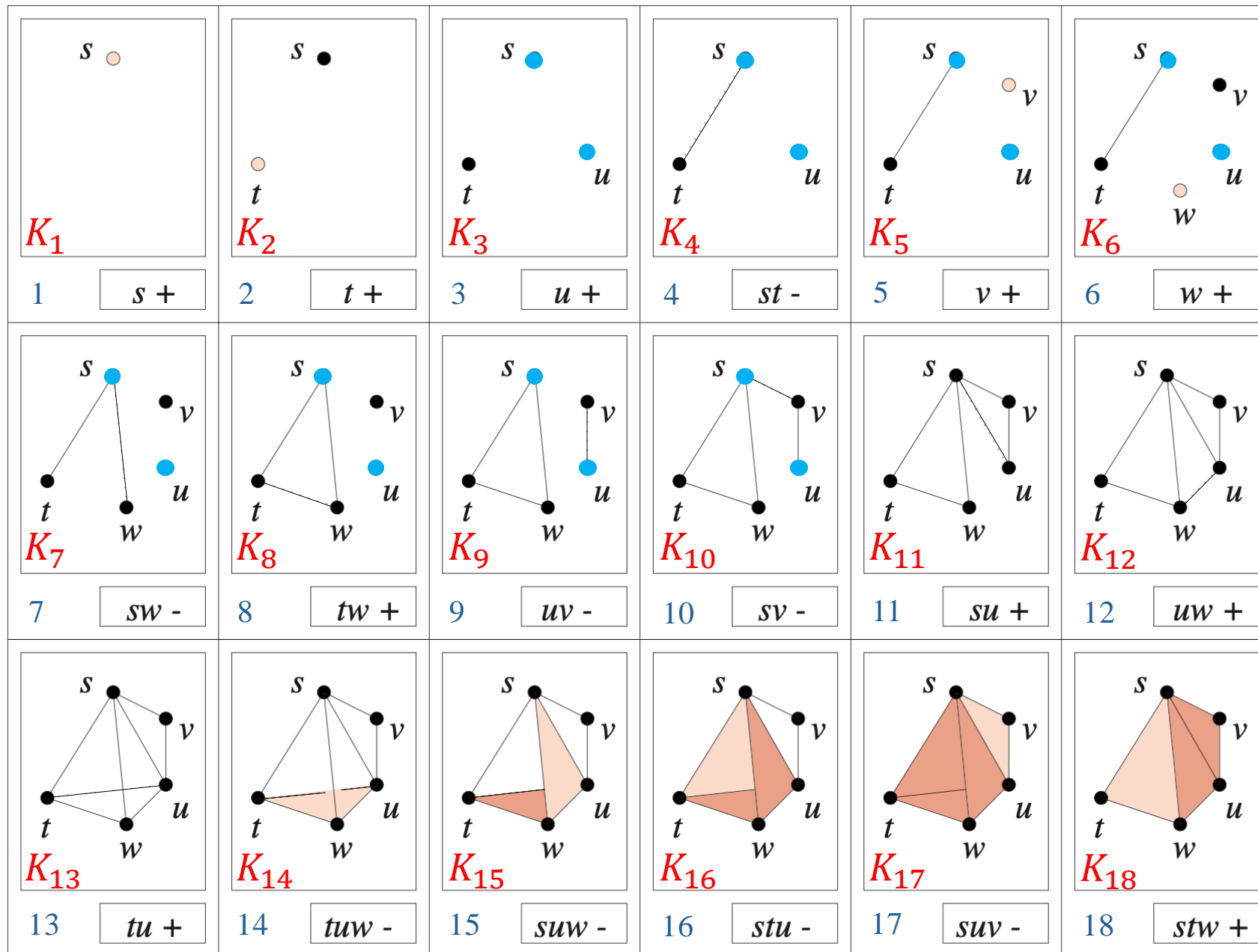
- We also notice that the cycle recorded in the “ ζ table” indeed captures the homology hole born and died with a birth-death interval in the barcode (point in the PD)
- i.e., for an interval $[b, d)$, $\zeta[\sigma_b]$ represents the homology feature born at the index b and dying at index d .
- This $\zeta[\sigma_b]$ is also called the **representative** for the interval $[b, d)$.



1d hole captured by
interval $[12,15) \in PD_1$



1d hole captured by
interval $[8,16) \in PD_1$



- 0d hole captured by interval
 $[3,10) \in PD_0$,
 which is the gap
 between s and u .
- The gap disappears
 when the two points
 become connected

More interpretations of the algorithm:

- When processing each σ_i , if the while loop ends with $z = 0$, then the simplex σ_i is called **positive**

Input: a filtration \mathcal{F} as a sequence of simplices $\sigma_1, \sigma_2, \dots, \sigma_m$

Output: p -th PD of \mathcal{F} , $\text{PD}_p(\mathcal{F})$, for each dimension p

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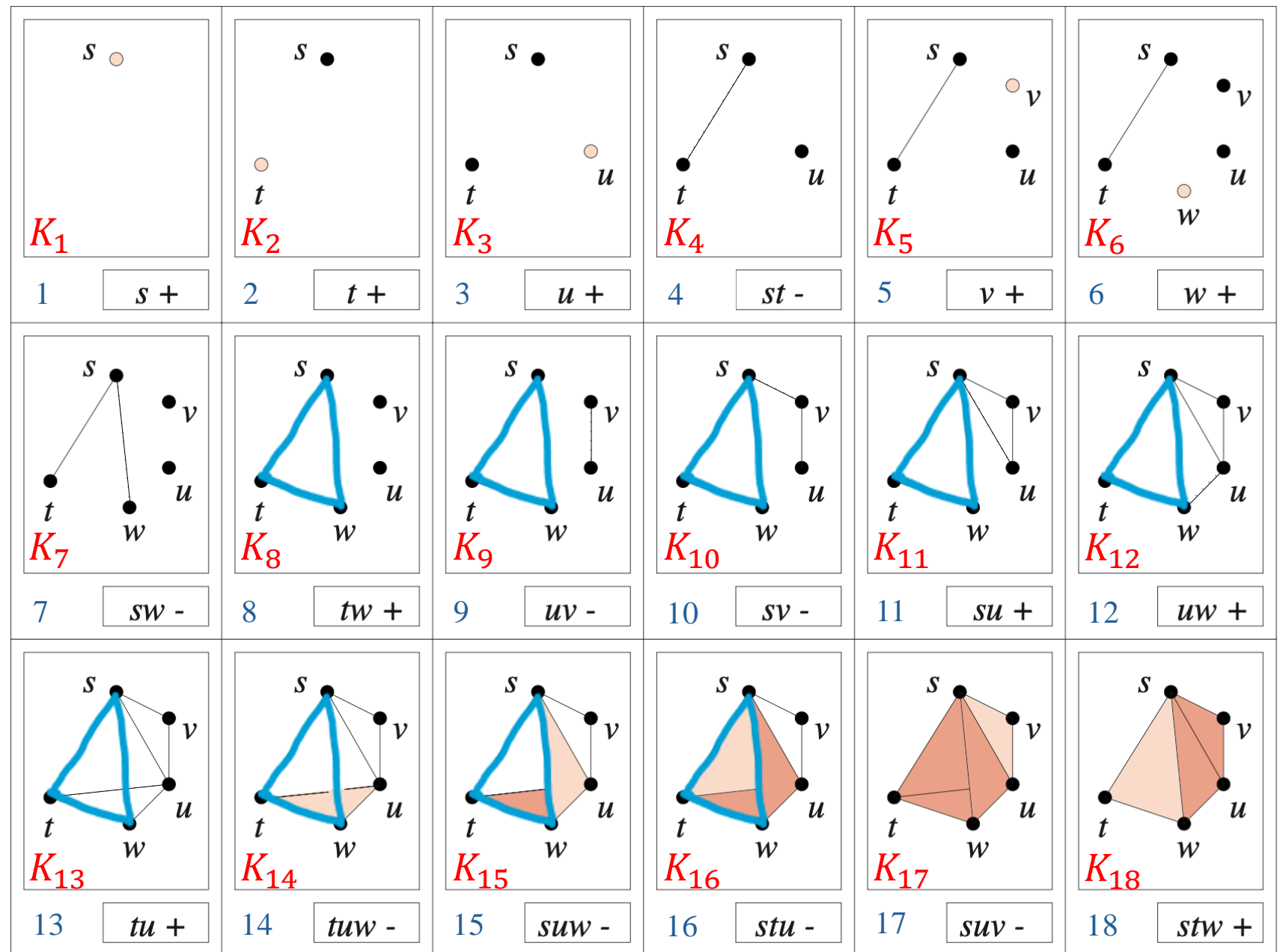
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- It means that inserting σ_i creates a new homology hole

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- E.g., inserting $\sigma_8 = tw$ creates the blue 1d hole



- If the while loop ends with $z \neq 0$, then the simplex σ_i is called **negative**

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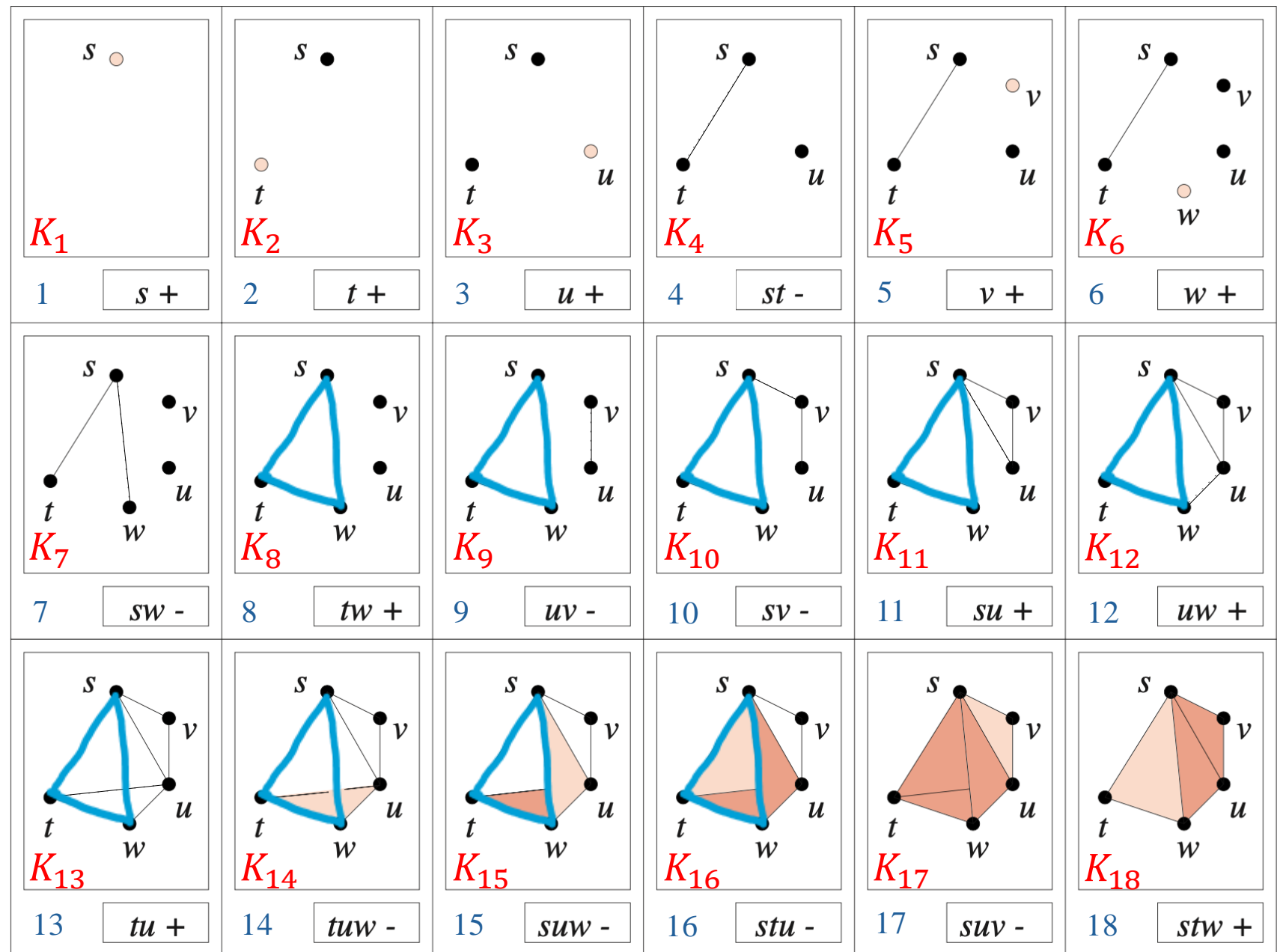
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- E.g., inserting $\sigma_{16} = stu$ kills the blue 1d hole



We have that line 10 in the algorithm is always pairing

- a positive simplex σ_j
- with
- a negative simplex σ_i

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In the worst case, both inner and outer loop iterates $O(m)$ time, and hence $O(m^3)$ overall

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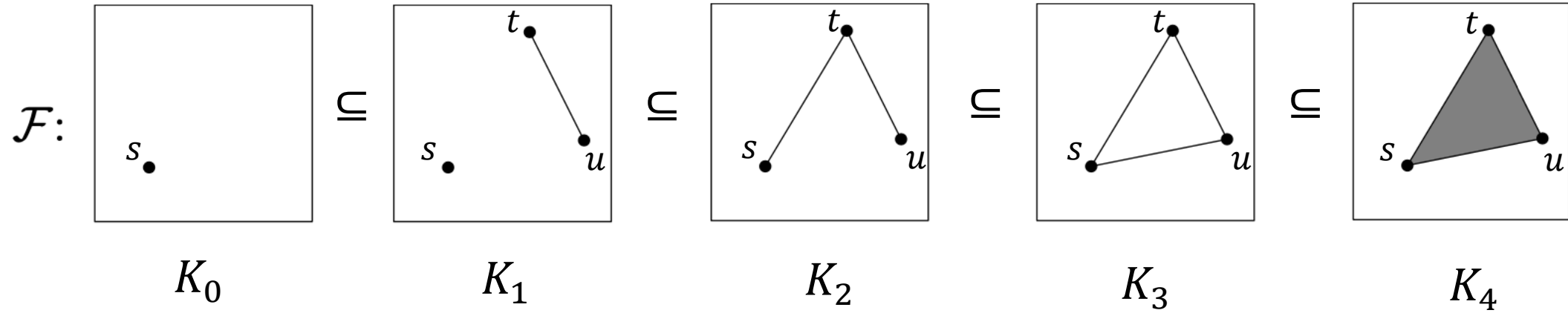
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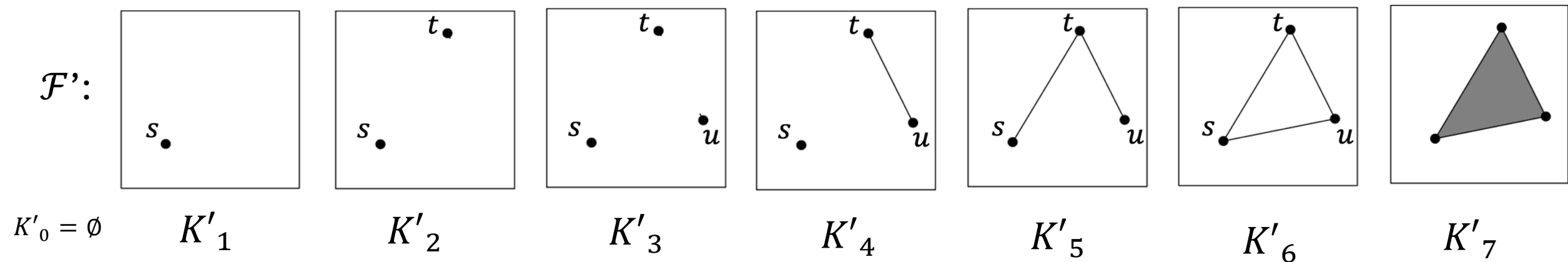
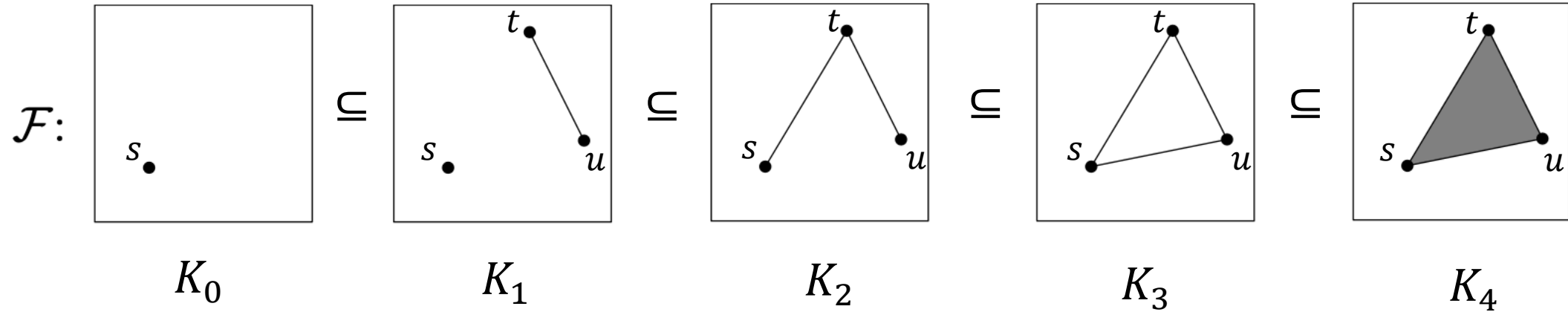
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 4. During the contraction, some intervals in $\text{PD}(\mathcal{F}')$ may disappear (birth and death coincide)

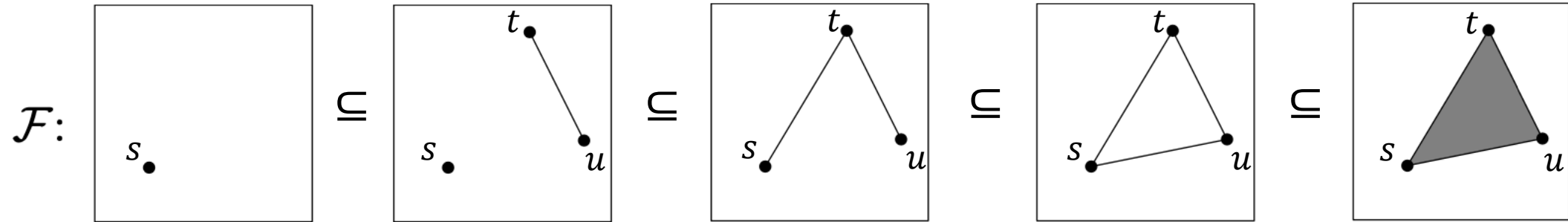
Expand general filtration to simplex wise



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K_0

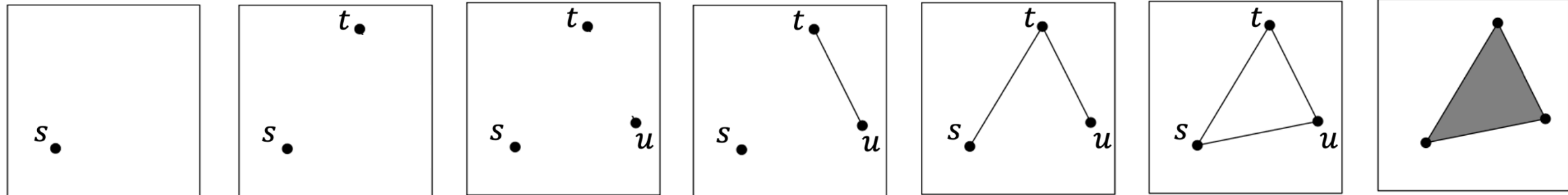
K_1

K_2

K_3

K_4

\mathcal{F}' :



$K'_0 = \emptyset$

K'_1

K'_2

K'_3

K'_4

K'_5

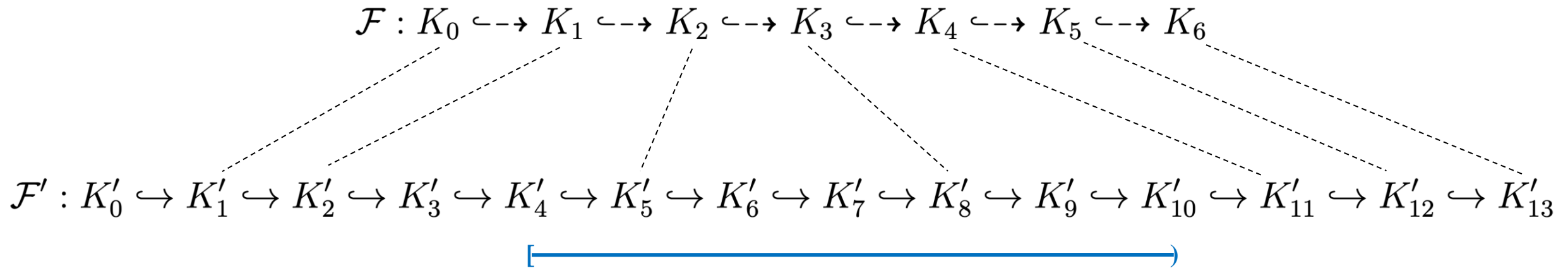
K'_6

K'_7

Expand general filtration to simplex wise

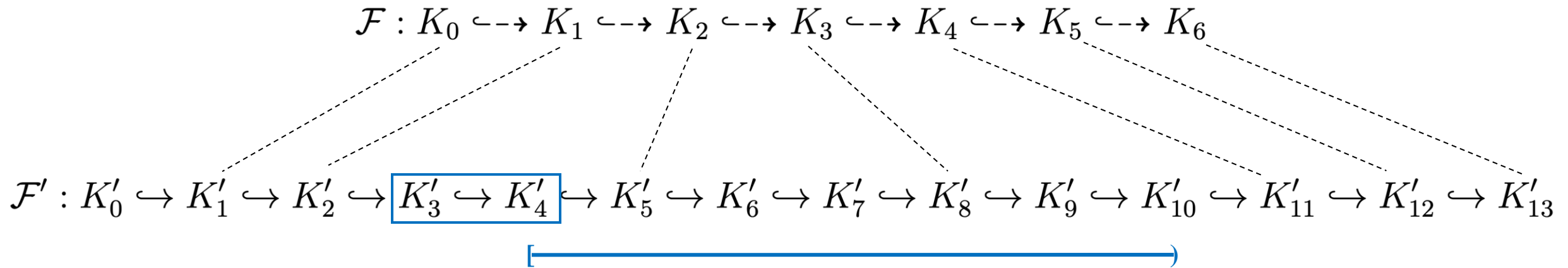
- Another interactive example for correspondence between a general filtration and its simplex-wise version: <https://iuricichf.github.io/ICT/algorithm.html>

PD for General Filtration



“Contracting” $[4,10) \in PD(\mathcal{F}')$ into one for $PD(\mathcal{F})$:

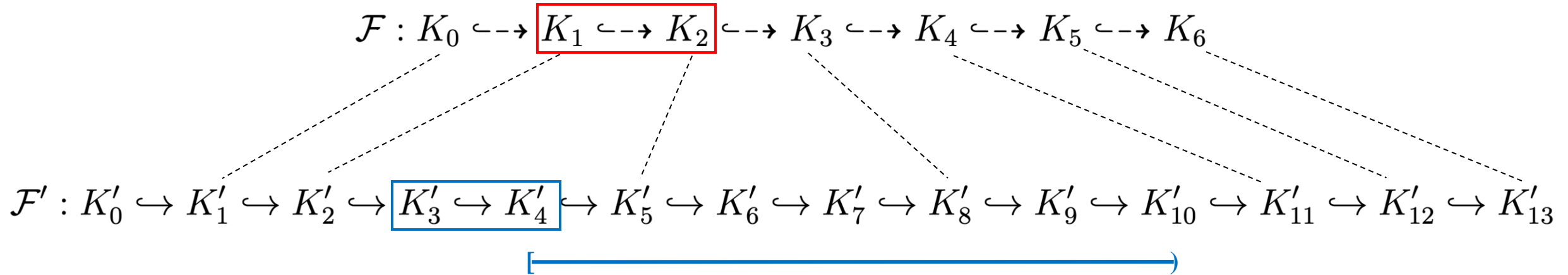
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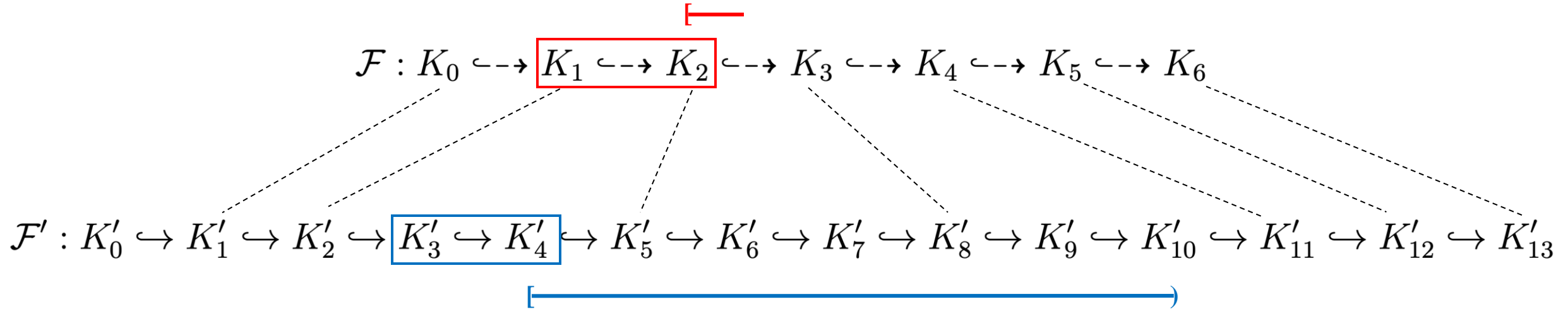
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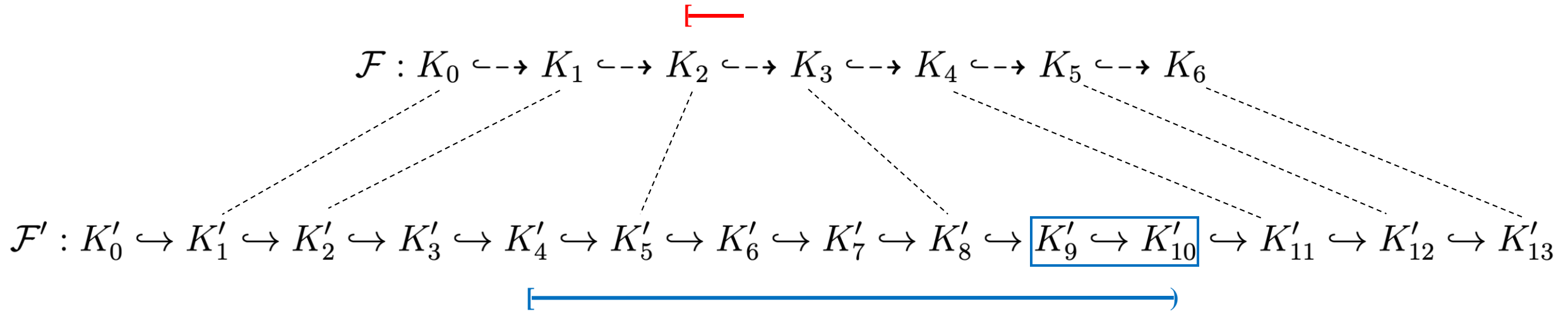
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- So the birth of the corresponding interval in $PD(\mathcal{F})$ is 2

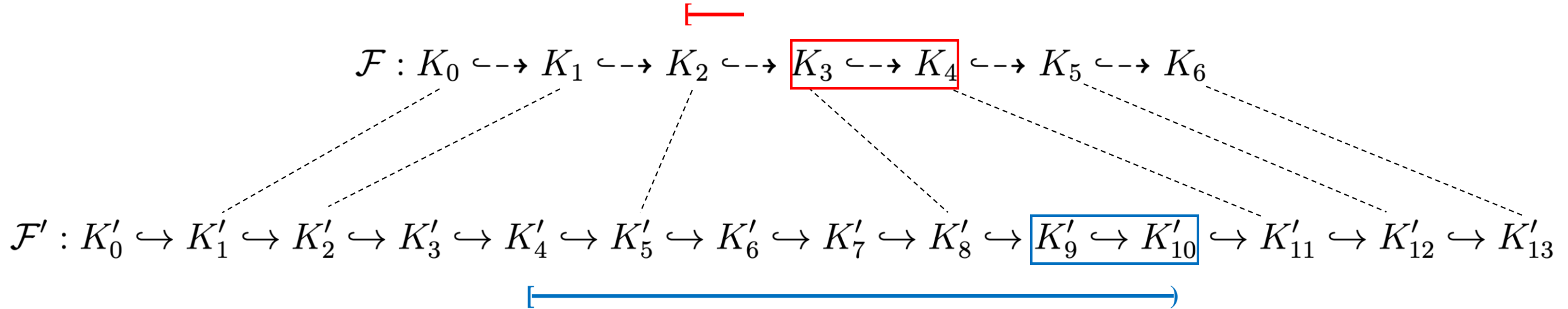
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“Contracting” $[4,10) \in PD(\mathcal{F}')$ into one for $PD(\mathcal{F})$:

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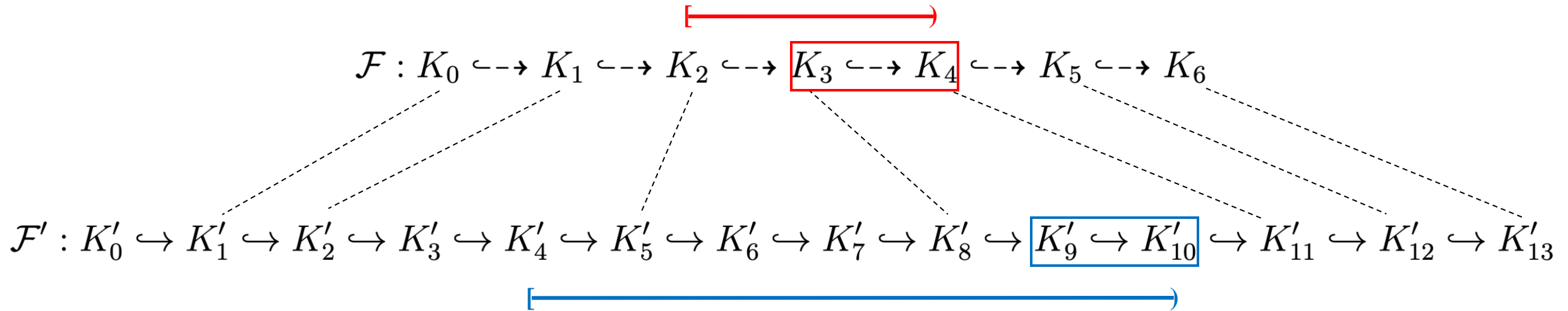
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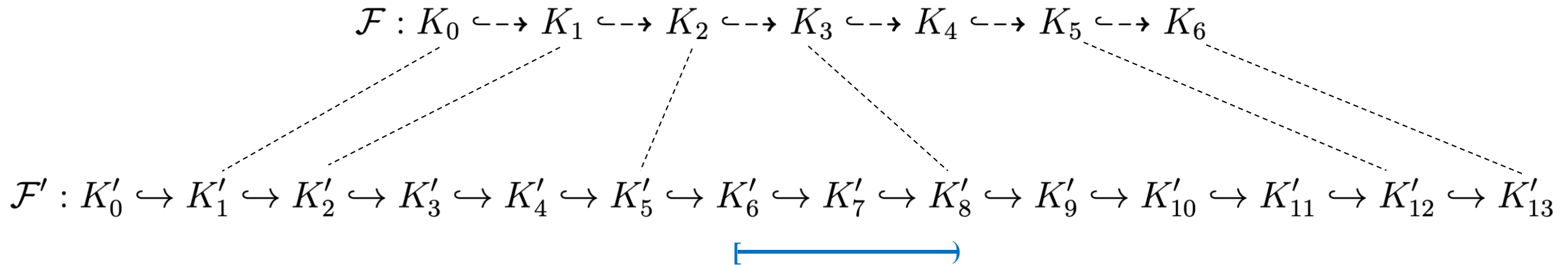
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- In \mathcal{F} , the homology feature dies when we go from K_3 to K_4 , aka in K_4
- So the corresponding interval in $PD(\mathcal{F})$ is $[2, 4)$

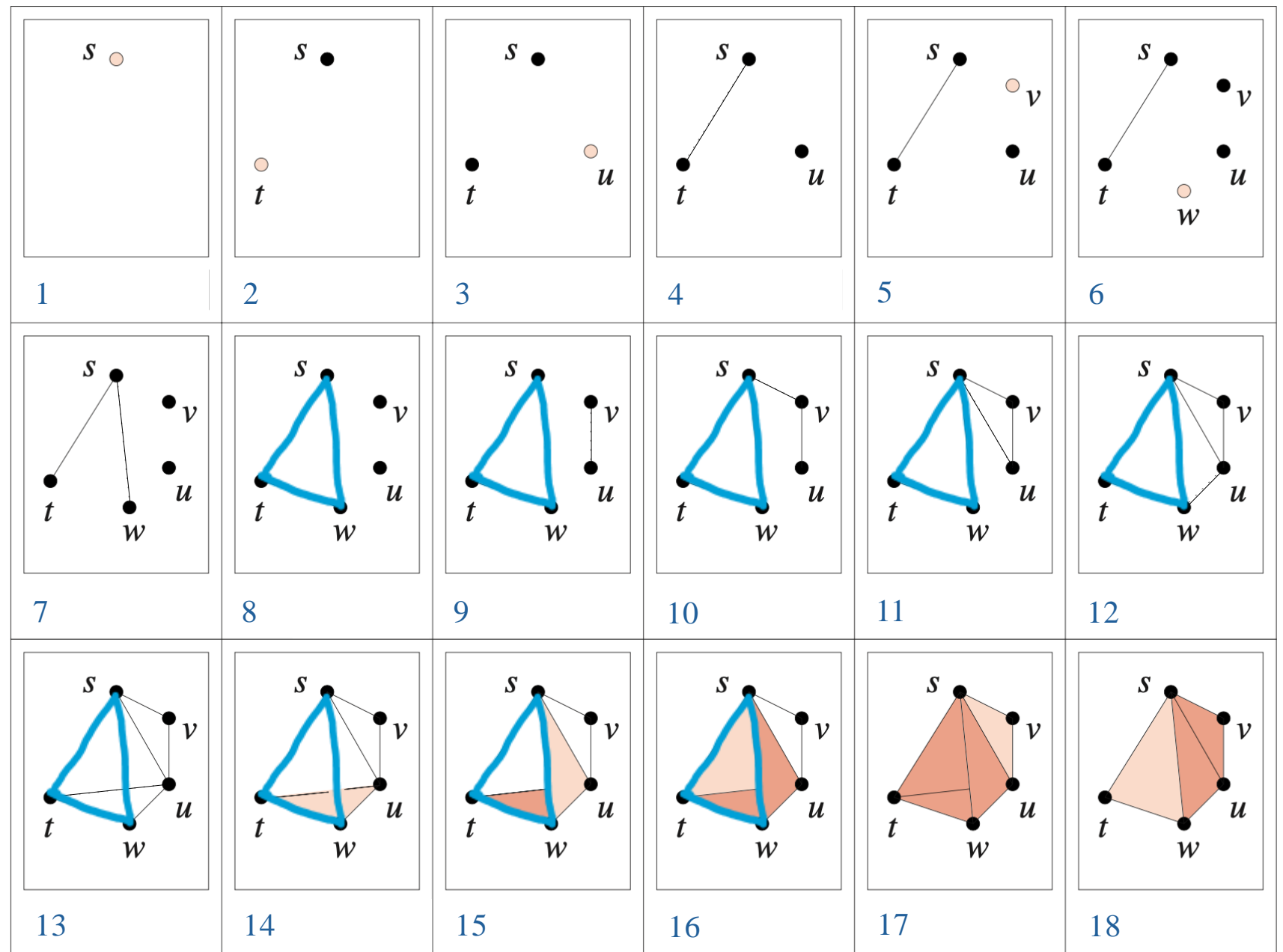
PD for General Filtration



$[5,8) \in PD(\mathcal{F}')$ does not correspond to any interval in $PD(\mathcal{F})$:

- In \mathcal{F} , the homology feature is born in K_3 and dies also K_3 (so it's ephemeral)

- For the previous simplex-wise filtration, we can skip some complexes and renumber them



- For the previous simplex-wise filtration, we can skip some complexes and renumber them
- Then [8,16] in the simplex-wise filtration becomes [5,10) in the non-simplex-wise
- But they are essential “same” interval (representatives are the same)

