

# Compilers

**Lexical Specification** 

- At least one:  $A^+ \equiv AA^*$
- Union:  $A \mid B$   $\equiv A + B$
- Option:  $\underline{A}$ ?  $\equiv \underline{A} + \underline{\varepsilon}$
- Range:  $(a' + b' + ... + z') \equiv [a-z]$
- Excluded range:

complement of 
$$[a-z] \equiv [^a-z]$$

• Last lecture: a specification for the predicate

$$S \in L(R)$$
Sets of strings

Not enough!

#### 1. Write a rexp for the lexemes of each token class

- Number = digit<sup>+</sup>
  Keyword = 'if' + 'else' + ...
- Identifier = letter (letter + digit)\*
- OpenPar = '('

2. Construct R, matching all lexemes for all tokens

$$R \neq \text{Keyword} + \text{Identifier} + \text{Number} + \dots$$
  
=  $R_1 + R_2 + \dots$ 

3. Let input be  $x_1...x_n$ For  $1 \le i \le n$  check

$$x_1...x_i \in L(R)$$

4. If success, then we know that

$$x_1...x_i \in L(R_i)$$
 for some j

Remove  $x_1...x_i$  from input and go to (3)

How much input is used?

$$x_{i} \cdot x_{i} \in \mathcal{L}(R)$$

$$x_{i} \cdot x_{j} \in \mathcal{L}(R)$$

Chase the one listed

Which token is used?

$$X_1...X_i \in L(R)$$
  $R = R_1 + ... + R_N$   
 $X_1...X_i \in L(R_i)$   
 $X_1...X_i \in L(R_k)$   
 $L(Keywords)$  Keywords = 'if  
 $L(Leywords)$  Labortifiers = left.

What if no rule matches?

Regular expressions are a concise notation for string patterns

- Use in lexical analysis requires small extensions

  - To handle errors

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- To resolve ambiguities - matches as long as possible highest priority metch
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- Good algorithms known
  - Require only single pass over the input
  - Few operations per character (table lookup)