



Compilers

Self Type Checking


- SELF_TYPE's meaning depends on the enclosing class

SELF_TYPE_C


$$\underline{O, M, C} \vdash \underline{e : T}$$

An expression e occurring in the body of C has static type T given a variable type environment O and method signatures M

- The next step is to design type rules using **SELF_TYPE**
- Most of the rules remain the same
 - But use the new \leq and **lub**


$$\frac{\begin{array}{l} O(\text{Id}) = T_0 \\ O, M, C \vdash e_1 : T_0 \\ T_1 \leq T_0 \end{array}}{O, M, C \vdash \text{Id} \leftarrow e_1 : T_1}$$

- Recall the old rule for dispatch


$$\frac{\begin{array}{l} O, M, C \vdash e_0 : T_0 \\ \vdots \\ O, M, C \vdash e_n : T_n \\ M(T_0, f) = (T_1', \dots, T_n', T_{n+1}') \\ \underline{T_{n+1}' \neq \text{SELF_TYPE}} \\ T_i \leq T_i' \quad 1 \leq i \leq n \end{array}}{O, M, C \vdash e_0.f(e_1, \dots, e_n) : T_{n+1}'}$$


- If the return type of the method is **SELF_TYPE** then the type of the dispatch is the type of the dispatch expression:

$$\begin{array}{c}
 \left[\begin{array}{l}
 O, M, C \vdash \underline{e_0} : \underline{T_0} \\
 \vdots \\
 O, M, C \vdash e_n : T_n \\
 M(\underline{T_0}, f) = (\underline{T_1'}, \dots, \underline{T_n'}, \text{SELF_TYPE}) \\
 \underline{T_i} \leq \underline{T_i'} \quad 1 \leq i \leq n
 \end{array} \right] \\
 \hline
 O, M, C \vdash \underline{e_0.f(e_1, \dots, e_n)} : \underline{T_0}
 \end{array}$$

- Formal parameters cannot be SELF_TYPE
- Actual arguments can be SELF_TYPE
 - The extended \leq relation handles this case
- The type T_0 of the dispatch expression could be SELF_TYPE
 - Which class is used to find the declaration of f ?
 - Answer: it is safe to use the class where the dispatch appears

$$\frac{\begin{array}{c} e_0 : \text{SELF_TYPE } C \\ M(\underline{C}, f) = (\dots) \end{array}}{C, M, C \vdash e_0.f(e_1)}$$

- Recall the original rule for static dispatch


$$\frac{\begin{array}{l} O, M, C \vdash e_0 : T_0 \\ \vdots \\ O, M, C \vdash e_n : T_n \\ T_0 \leq T \\ M(T, f) = (T_1', \dots, T_n', T_{n+1}') \\ T_{n+1}' \neq \text{SELF_TYPE} \\ T_i \leq T_i' \quad 1 \leq i \leq n \end{array}}{O, M, C \vdash e_0 @ T.f(e_1, \dots, e_n) : T_{n+1}'}$$

- If the return type of the method is **SELF_TYPE** we have:

$$\begin{array}{l}
 \left[\begin{array}{l}
 O, M, C \vdash e_0 : T_0 \\
 \vdots \\
 O, M, C \vdash e_n : T_n
 \end{array} \right. \\
 \rightarrow T_0 \leq T \\
 \rightarrow M(\underline{T}, f) = (T_1', \dots, T_n', \underline{\text{SELF_TYPE}}) \\
 \underline{T_i \leq T_i'} \quad 1 \leq i \leq n \\
 \hline
 \underline{O, M, C} \vdash e_0 @ \underline{T}. f(e_1, \dots, e_n) : \underline{\underline{T_0}}
 \end{array}$$

- Why is this rule correct?
- If we dispatch a method returning `SELF_TYPE` in class `T`, don't we get back a `T`?

$$\underline{T_0} \leq \underline{T}$$

- No. `SELF_TYPE` is the type of the self parameter, which may be a subtype of the class in which the method appears

- There are two new rules using **SELF_TYPE**

$$O, M, \underline{\underline{C}} \vdash \text{self} : \text{SELF_TYPE}_{\underline{\underline{C}}}$$

$$O, M, C \vdash \underline{\text{new SELF_TYPE}} : \text{SELF_TYPE}_{\underline{\underline{C}}}$$

Choose the static/dynamic type pairs that are correct. For dynamic type, assume execution has halted at line 15.

	<u>Var</u>	<u>Static Type</u>	<u>Dynamic Type</u>
<input type="checkbox"/>	w	Animal	Pet
<input type="checkbox"/>	x	Animal	Lion
<input type="checkbox"/>	y	Pet	Pet
<input type="checkbox"/>	z	Animal	Dog

Self Type Checking

```
1  class Animal {  
2      clone() : SELF_TYPE { new SELF_TYPE }  
3  }  
4      class Pet inherits Animal {  
5          clone() : Pet { new SELF_TYPE }  
6      }  
7      class Cat inherits Pet { ... }  
8      class Dog inherits Pet { ... }  
9      class Lion inherits Animal { ... }  
10 class Main {  
11     w:Animal <- (new Animal).clone();  
12     x:Animal <- (new Lion).clone();  
13     y:Pet <- (new Cat).clone();  
14     z:Animal <- (new Dog).clone();  
→ 15     ...  
16 }
```

- The extended \leq and **lub** operations can do a lot of the work.
- **SELF_TYPE** can be used only in a few places. Be sure it isn't used anywhere else.
- A use of **SELF_TYPE** always refers to any subtype of the current class
 - The exception is the type checking of dispatch. The method return type of **SELF_TYPE** might have nothing to do with the current class

$$M(\underline{c}, f) = (\dots, \text{SELF_TYPE})$$

- SELF_TYPE is a research idea
 - It adds more expressiveness to the type system
- SELF_TYPE is itself not so important
 - except for the project
- Rather, SELF_TYPE is meant to illustrate that type checking can be quite subtle
- In practice, there should be a balance between the complexity of the type system and its expressiveness