



Compilers

Recursive Descent Algorithm

- Let TOKEN be the type of tokens
 - Special tokens INT, OPEN, CLOSE, PLUS, TIMES
- Let the global next point to the next input token



- Define boolean functions that check for a match of:
 - A given token terminal

```
bool term(TOKEN tok) { return *next++ == tok; }
```

Handwritten notes: "true/false" above the code, an arrow pointing down to the "++" operator, and red underlines under "TOKEN", "tok", "next", and "tok".
 - The n th production of S :

```
→ bool  $S_n$ () { ... }
```

Handwritten notes: red arrow pointing to the code, and red circles around the ellipsis "..." and the subscript n .
 - Try all productions of S :

```
→ bool  $S$ () { ... }
```

Handwritten notes: red arrow pointing to the code, and a red underline under the S in the function name.

- For production $E \rightarrow T$
`bool E1() { return T(); }`
- For production $E \rightarrow T + E$
`bool E2() { return T() && term(PLUS) && E(); }`
- For all productions of E (with backtracking)

```
bool E() {
  TOKEN *save = next;
  return (next = save, E1())
  || (next = save, E2()); }
```

- Functions for non-terminal T

bool T₁() { return term(INT); } $T \rightarrow int$

bool T₂() { return term(INT) && term(TIMES) && T(); } $T \rightarrow int * T$

bool T₃() { return term(OPEN) && E() && term(CLOSE); } $T \rightarrow (E)$

```
bool T() {
    TOKEN *save = next;
    return  (next = save, T1())
           || (next = save, T2())
           || (next = save, T3()); }
```

- To start the parser
 - Initialize next to point to first token
 - Invoke E()
- Easy to implement by hand

RD Algorithm

$E \rightarrow T \mid T + E$

$T \rightarrow \text{int} \mid \text{int} * T \mid (E)$

(int)
✗ ✗ ✗ ✗ ✗

```
bool term(TOKEN tok) { return *next++ == tok; }
```

```
bool E1() { return T(); }
```

```
bool E2() { return T() && term(PLUS) && E(); }
```

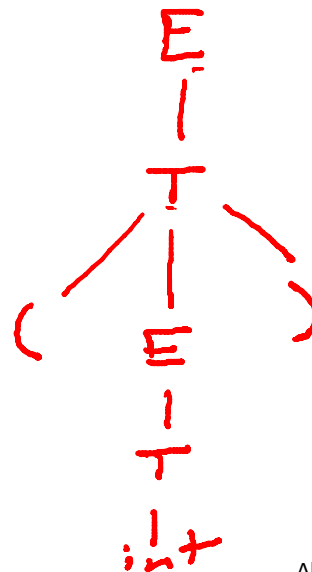
```
bool E() { TOKEN *save = next; return (next = save, E1())  
|| (next = save, E2()); }
```

```
bool T1() { return term(INT); }
```

```
bool T2() { return term(INT) && term(TIMES) && T(); }
```

```
bool T3() { return term(OPEN) && E() && term(CLOSE); }
```

```
bool T() { TOKEN *save = next; return (next = save, T1())  
|| (next = save, T2())  
|| (next = save, T3()); }
```



Which lines are incorrect in the recursive descent implementation of this grammar?

$$E \rightarrow E' \mid E' + id$$
$$E' \rightarrow -E' \mid id \mid (E)$$

☐ Line 3

☐ Line 5

☐ Line 6

☐ Line 12

RD Algorithm

```
1  bool term(TOKEN tok) { return *next++ == tok; }
2  bool E1() { return E'(); }
3  bool E2() { return E'() && term(PLUS) && term(ID); }
4  bool E() {
5      TOKEN *save = next;
6      return (next = save, E1()) && (next = save, E2());
7  }

8  bool E'1() { return term(MINUS) && E'(); }
9  bool E'2() { return term(ID); }
10 bool E'3() { return term(OPEN) && E() && term(CLOSE); }
11 bool E'() {
12     TOKEN *next = save; return (next = save, T1())
13                             || (next = save, T2())
14                             || (next = save, T3());
15 }
```