



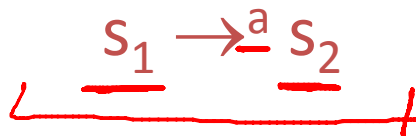
# Compilers

---

## Finite Automata

- Regular expressions = specification
- Finite automata = implementation
- A finite automaton consists of
  - An input alphabet  $\Sigma$
  - <sup>finite</sup> A set of states  $S$
  - A start state  $n$
  - A set of accepting states  $F \subseteq S$
  - A set of transitions state  $\rightarrow$  <sup>input</sup> state

- Transition



- Is read

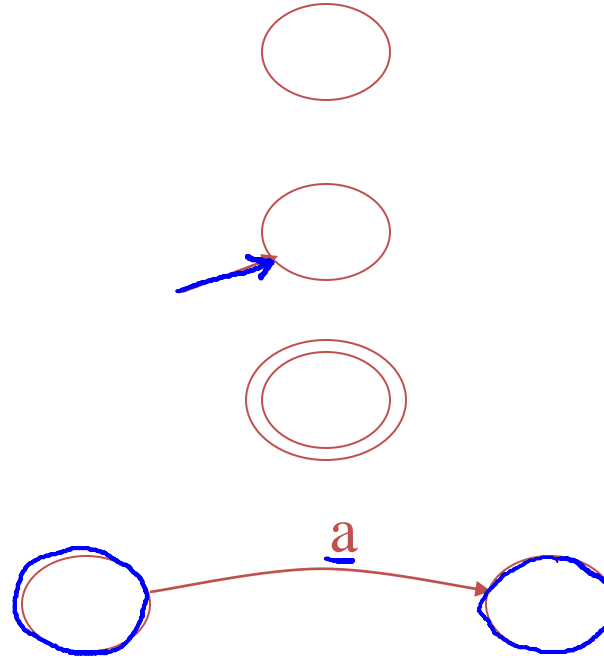
In state  $s_1$  on input  $a$  go to state  $s_2$

- If end of input and in accepting state => accept

- Otherwise => reject [ terminates in state  $s \notin F$   
gets stuck

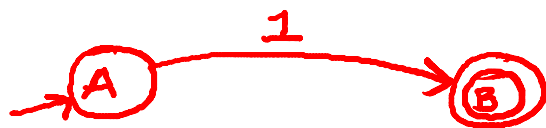
# Finite Automata

- A state
- The start state
- An accepting state
- A transition



# Finite Automata

- A finite automaton that accepts only "1"



State	Input	
<u>A</u>	<u>0</u>	<u>Reject</u>

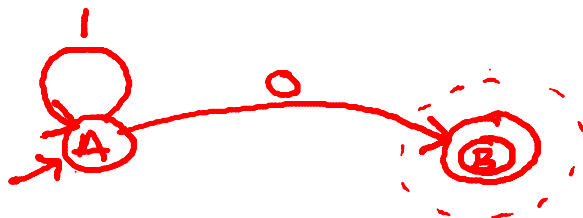
State	Input	
<u>A</u>	<u>1</u>	<u>Accept</u>
<u>B</u>	<u>1</u>	

State	Input	
<u>A</u>	<u>1</u> <u>0</u>	<u>Reject</u>
<u>B</u>	<u>1</u> <u>0</u>	

Language of a FA  $\equiv$   
set of accepted  
strings

# Finite Automata

- A finite automaton accepting any number of 1's followed by a single 0
- Alphabet:  $\{0,1\}$



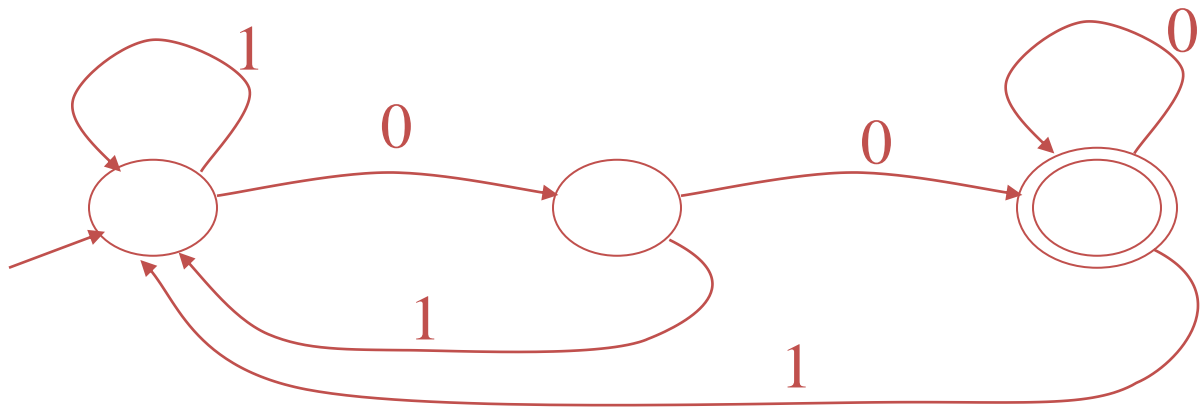
<u>State</u>	<u>Input</u>
<u>A</u>	↑100
A	1↑00
<u>B</u>	10↑0

Reject

<u>State</u>	<u>Input</u>
A	↑110
A	1↑10
A	11↑0
<u>B</u>	110↑

Accepts

Select the regular language that denotes the same language as this finite automaton



- ☐  $(0 + 1)^*$
- ☐  $(1^* + 0)(1 + 0)$
- ☐  $1^* + (01)^* + (001)^* + (000^*1)^*$
- ☐  $(0 + 1)^*00$

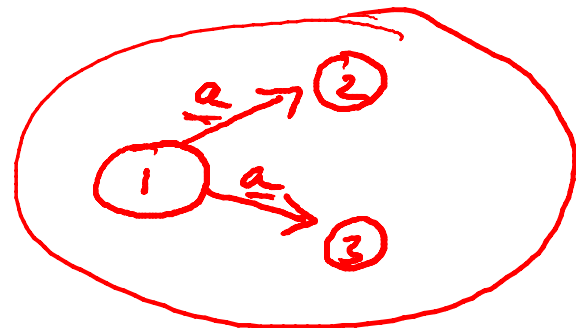
- Another kind of transition:  $\epsilon$ -moves



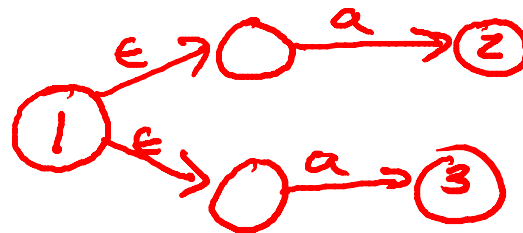
<u>State</u>	<u>Input</u>
A	$x_1$ $x_2$ $x_3$ ↑
<u>B</u>	<u><math>x_1</math></u> $x_2$ $x_3$ ↑



- Deterministic Finite Automata (DFA)
  - One transition per input per state
  - No  $\epsilon$ -moves



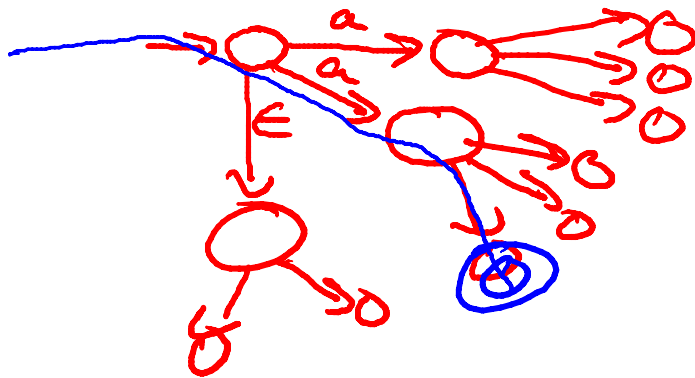
- Nondeterministic Finite Automata (NFA)
  - Can have multiple transitions for one input in a given state
  - Can have  $\epsilon$ -moves



- A DFA takes only one path through the state graph

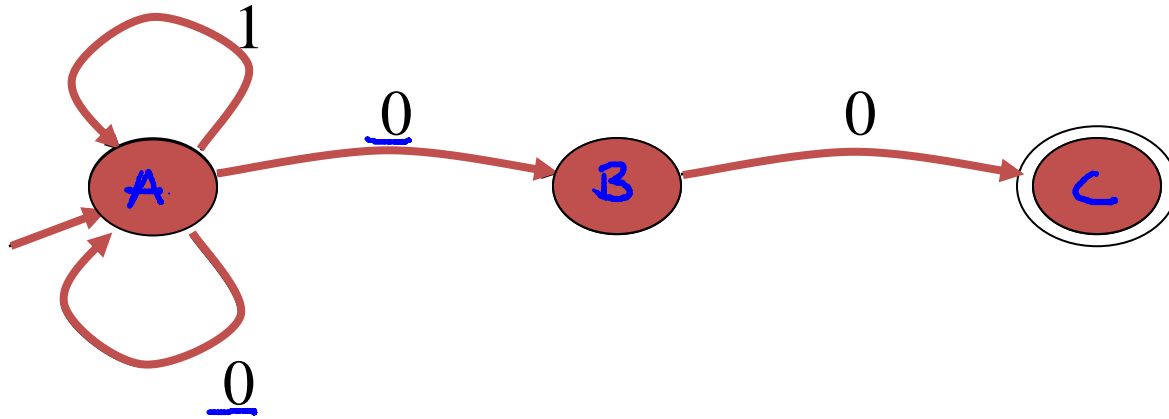


- An NFA can choose



An NFA  
accepts if  
some choices  
lead to an  
accepting state.

- An NFA can get into multiple states



- Input:

- States:

1  
0  
{A}

0  
{A,B}

0  
{A,B,C}

accepts

- NFAs and DFAs recognize the same set of languages
  - regular languages
- DFAs are faster to execute
  - There are no choices to consider
- NFAs are, in general, smaller
  - exponentially smaller