



Compilers

Reference Counting

- Rather than wait for memory to be exhausted, try to collect an object when there are no more pointers to it
- Store in each object the number of pointers to that object
 - this is the reference count
- Each assignment operation manipulates the reference count



Reference Counting

- new returns an object with reference count 1
- Let rc(x) be the reference count of x

- Assume x, y point to objects o, p

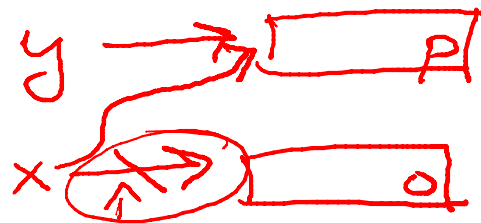
- Every assignment $x \leftarrow y$ becomes:

$rc(p) \leftarrow rc(p) + 1$

$rc(o) \leftarrow rc(o) - 1$

if($rc(o) == 0$) then free o

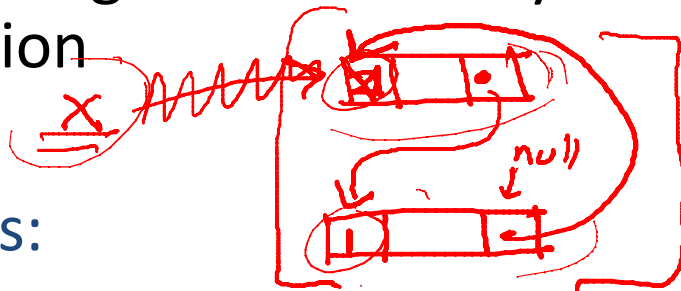
$\rightarrow x \leftarrow y$



- Advantages:

- – easy to implement
- – collects garbage incrementally without large pauses in the execution

$x \leftarrow \text{null}$



unreachable

```
ref() ←  
ref() ←  
if —  
  x ← y
```

- Disadvantages:

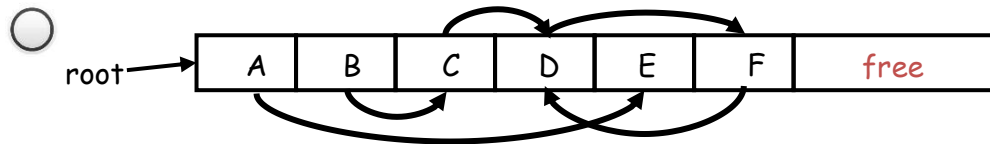
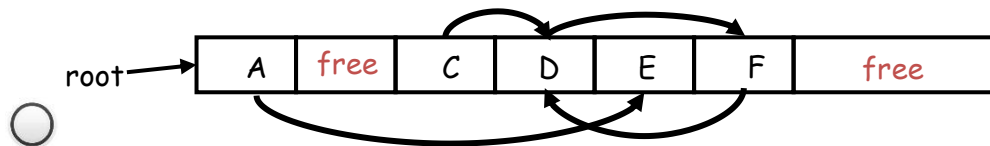
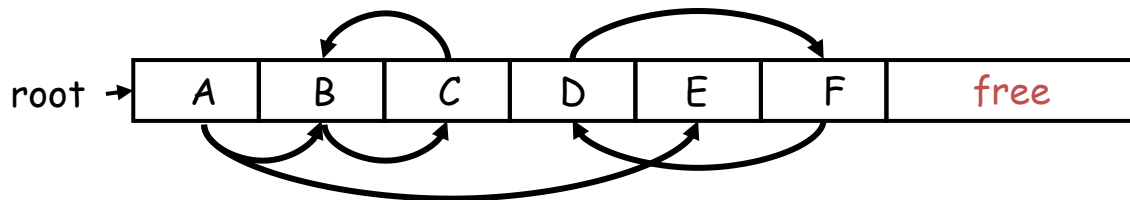
- – cannot collect circular structures
- – manipulating reference counts at each assignment is very slow

Reference Counting

Choose the final heap after executing the following two assignments and updating reference counts:

$C.\text{ptrToB} = D$

$A.\text{ptrToB} = \text{NULL}$



- Automatic memory management prevents serious storage bugs

- But reduces programmer control
 - e.g., layout of data in memory
 - e.g., when is memory deallocated

- • Pauses problematic in real-time applications
- • Memory leaks possible (even likely)

$x = \text{Null}$

~~x~~

AST

- Garbage collection is very important
- There are more advanced garbage collection algorithms:
 - concurrent: allow the program to run while the collection is happening
 - generational: do not scan long-lived objects at every collection
 - real time: bound the length of pauses
 - parallel: several collectors working at once