



# Compilers

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Activations

- Two goals:
  - Correctness ←
  - Speed ←
- Complications in code generation come from trying to be fast as well as correct

Two assumptions:

1. Execution is sequential; control moves from one point in a program to another in a well-defined order
2. When a procedure is called, control always returns to the point immediately after the call

concurrency

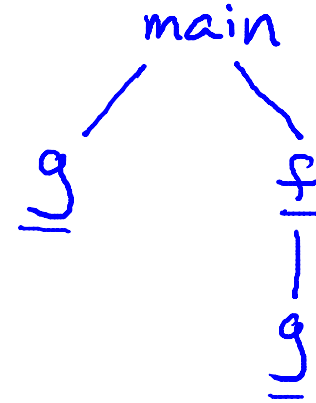
exceptions call/cc

- An invocation of procedure **P** is an activation of **P**
- The lifetime of an activation of **P** is
  - All the steps to execute **P**
  - Including all the steps in procedures **P** calls

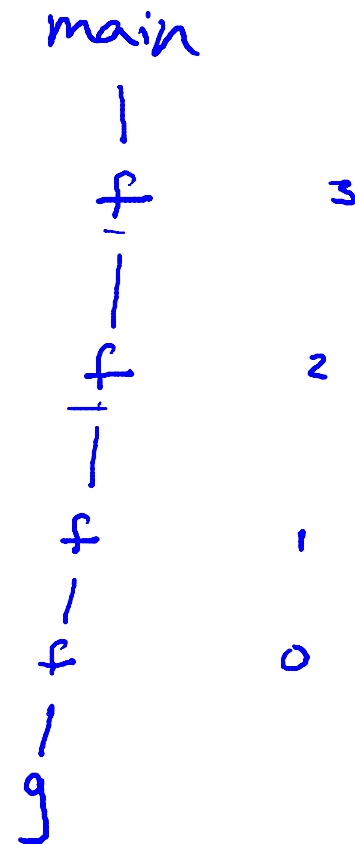
- The lifetime of a variable x is the portion of execution in which x is defined
- Note that
  - Lifetime is a dynamic (run-time) concept
  - Scope is a static concept

- Observation
  - When P calls Q, then Q returns before P returns
- Lifetimes of procedure activations are properly nested
- Activation lifetimes can be depicted as a tree

```
Class Main {  
    g() : Int { 1 };  
    f(): Int { g() };  
    main(): Int {{ g(); f(); }};  
}
```



```
Class Main {  
    g() : Int { 1 };  
    f(x:Int): Int { if x = 0 then g() else f(x - 1) fi};  
    main(): Int {{f(3); }};  
}
```

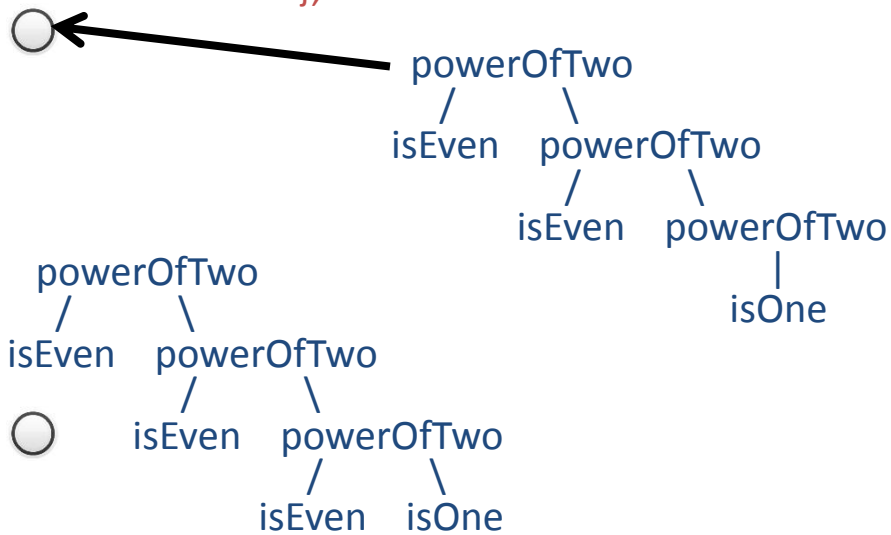
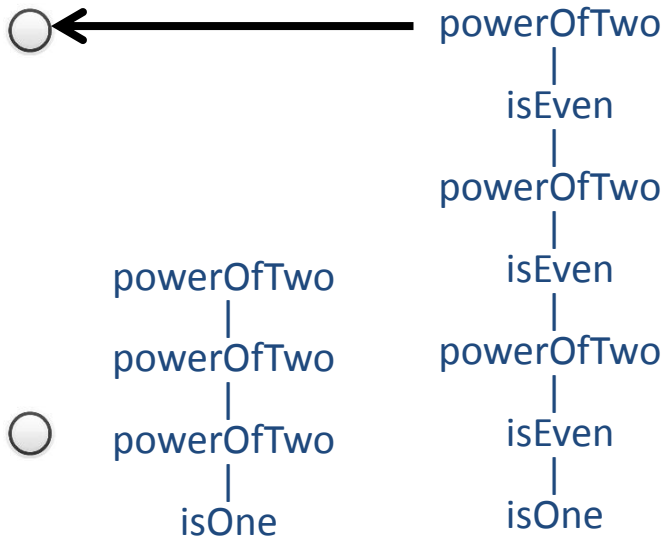




The `powerOfTwo()` function, shown to the right, returns true if its argument is a power of two, false otherwise. What is the activation tree for `powerOfTwo(4)`?

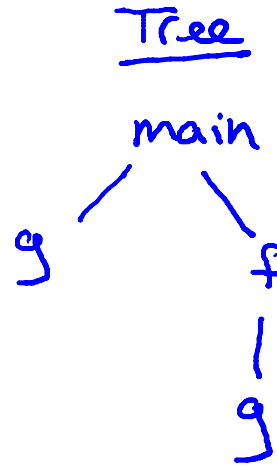
# Activations

```
isEven(x:Int) : Bool { x % 2 == 0 };
isOne(x:Int) : Bool { x == 1 };
powerOfTwo(x:Int) : Bool {
  if isEven(x) then powerOfTwo(x / 2)
  else isOne(x)
};
```



- The activation tree depends on run-time behavior
- The activation tree may be different for every program input
- Since activations are properly nested, a stack can track currently active procedures

```
Class Main {  
  g() : Int { 1 };  
  f(): Int { g() };  
  main(): Int {{ g(); f(); }};  
}
```



Stack

**Memory**

