



Compilers

Dataflow Analysis

Recall the simple basic-block optimizations

- Constant propagation
- Dead code elimination

$X := 3$
 $Y := Z * W$
 $Q := X + Y$

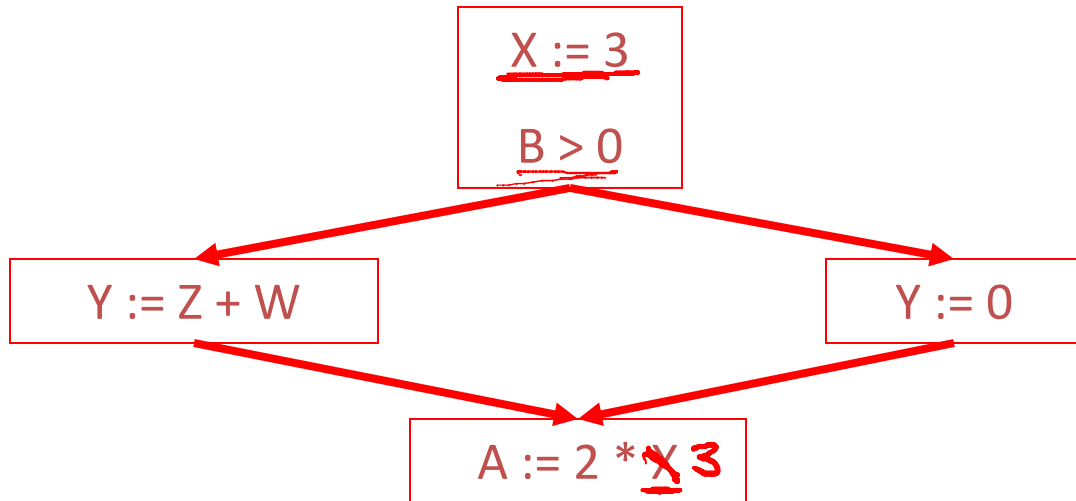


$X := 3$
 $Y := Z * W$
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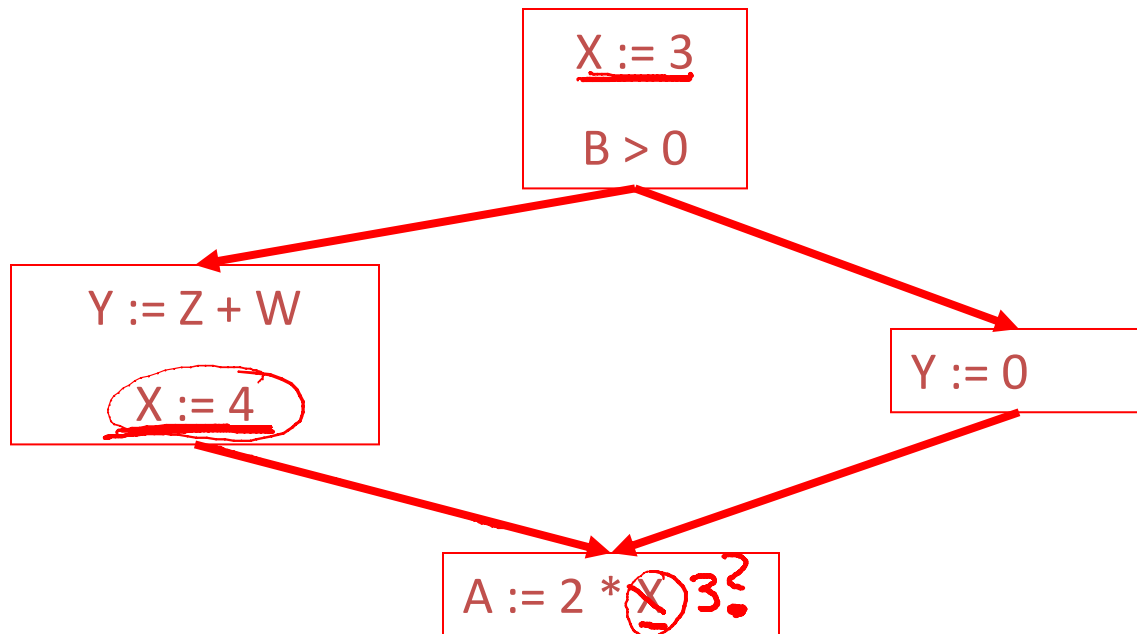


$Y := Z * W$
 $Q := 3 + Y$

These optimizations can be extended to an entire control-flow graph



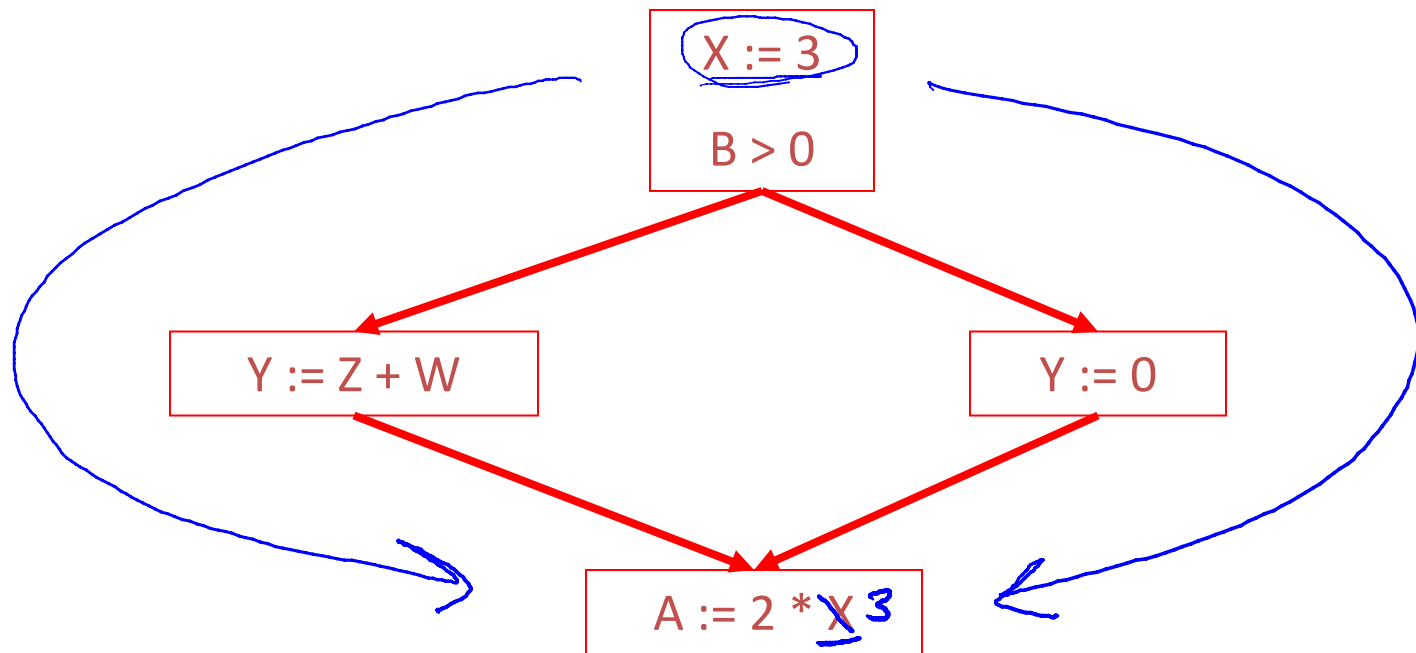
- How do we know it is OK to globally propagate constants?



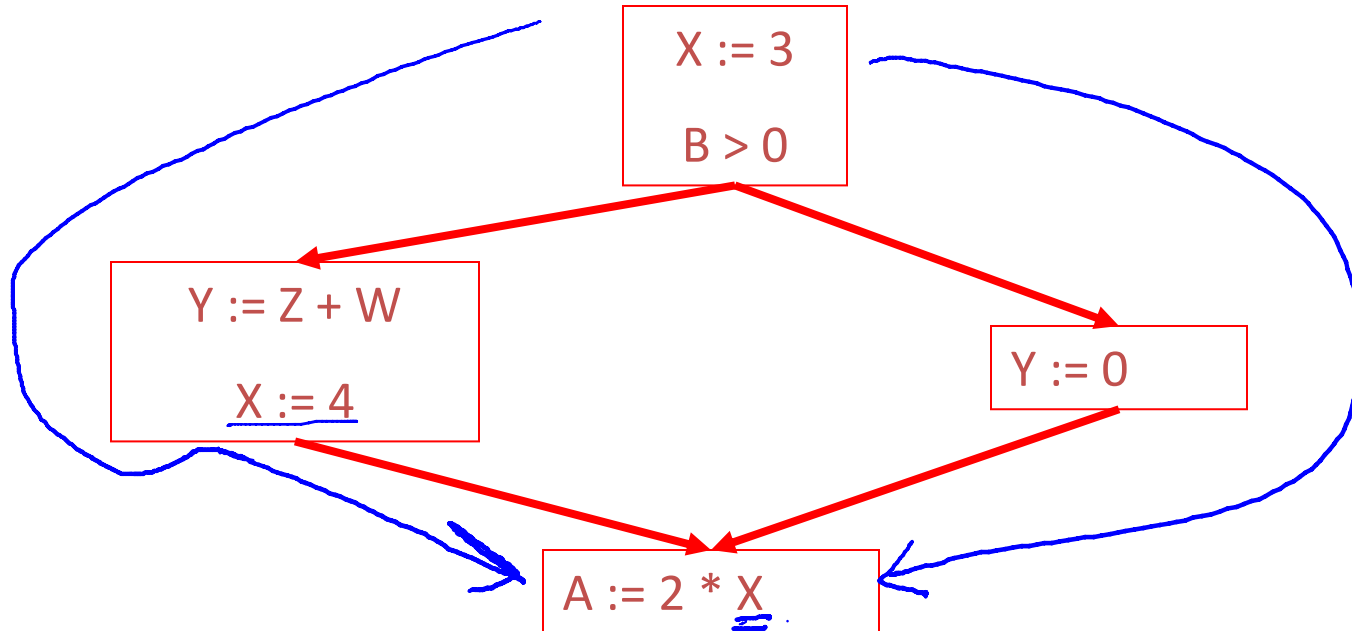
To replace a use of x by a constant k we must know:

On every path to the use of x, the last assignment to x is
 $x := k$

Dataflow Analysis



Dataflow Analysis



- The correctness condition is not trivial to check
- “All paths” includes paths around loops and through branches of conditionals
- Checking the condition requires global dataflow analysis
 - An analysis of the entire control-flow graph

Global optimization tasks share several traits:

- The optimization depends on knowing a property X at a particular point in program execution
- Proving X at any point requires knowledge of the entire program
- It is OK to be conservative. If the optimization requires X to be true, then want to know either
 - X is definitely true
 - Don't know if X is true
 - It is always safe to say “don't know”

- *Global dataflow analysis* is a standard technique for solving problems with these characteristics
- Global constant propagation is one example of an optimization that requires global dataflow analysis