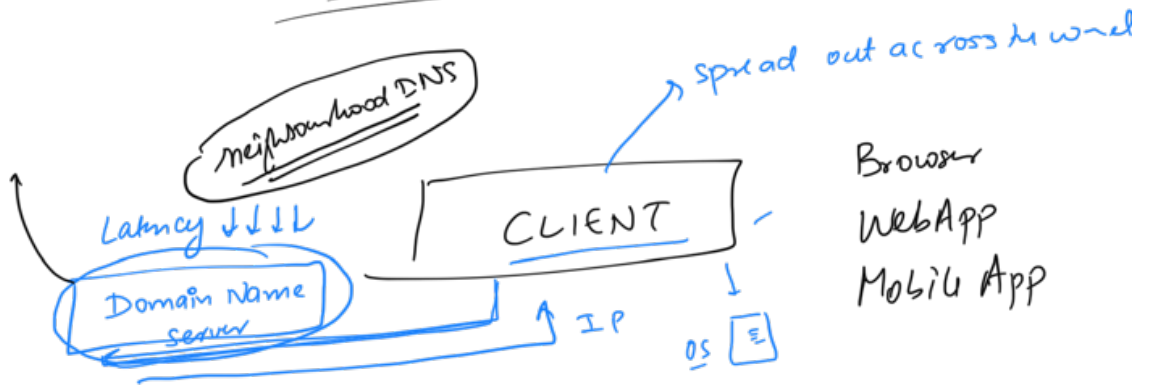
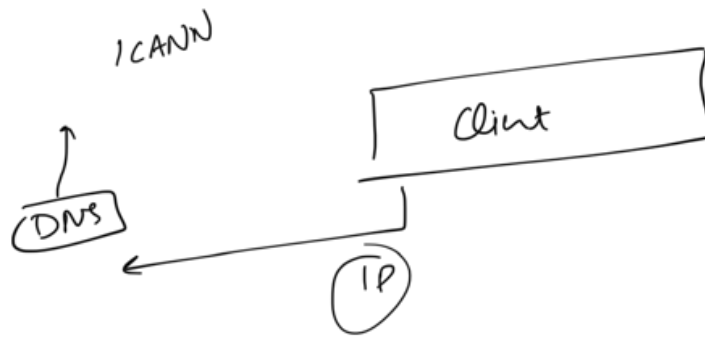
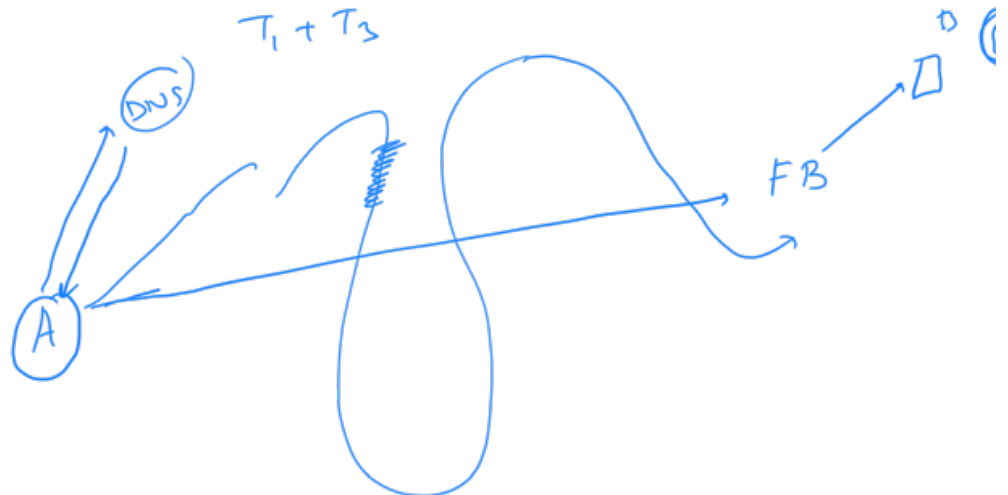
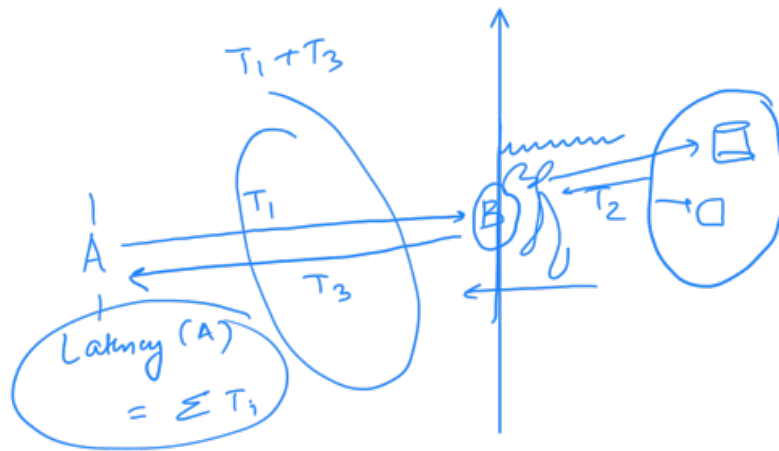


# High Level Design



1

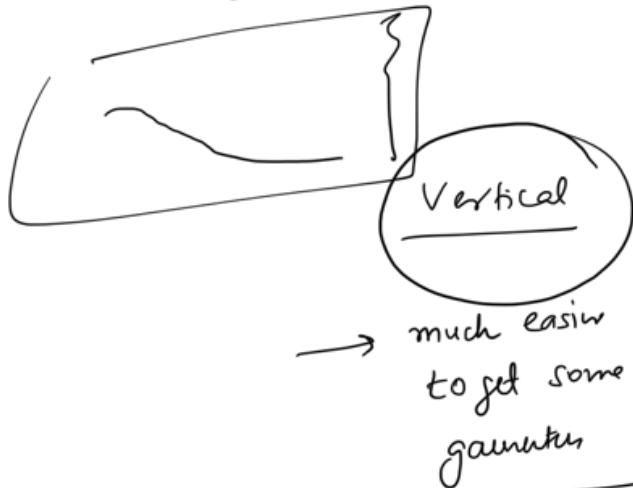
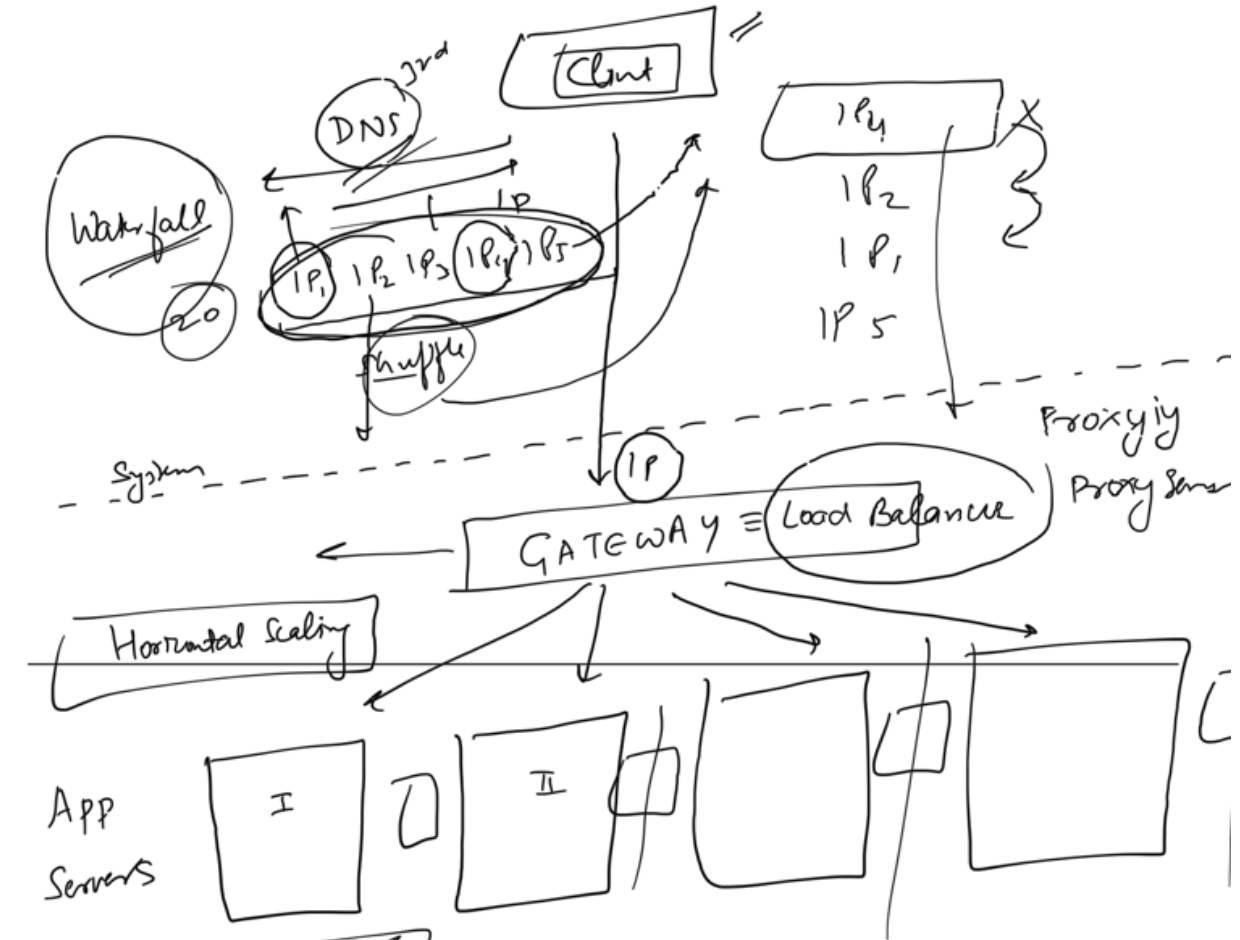


- ISPs
- Google
- FB

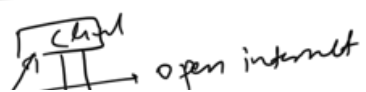


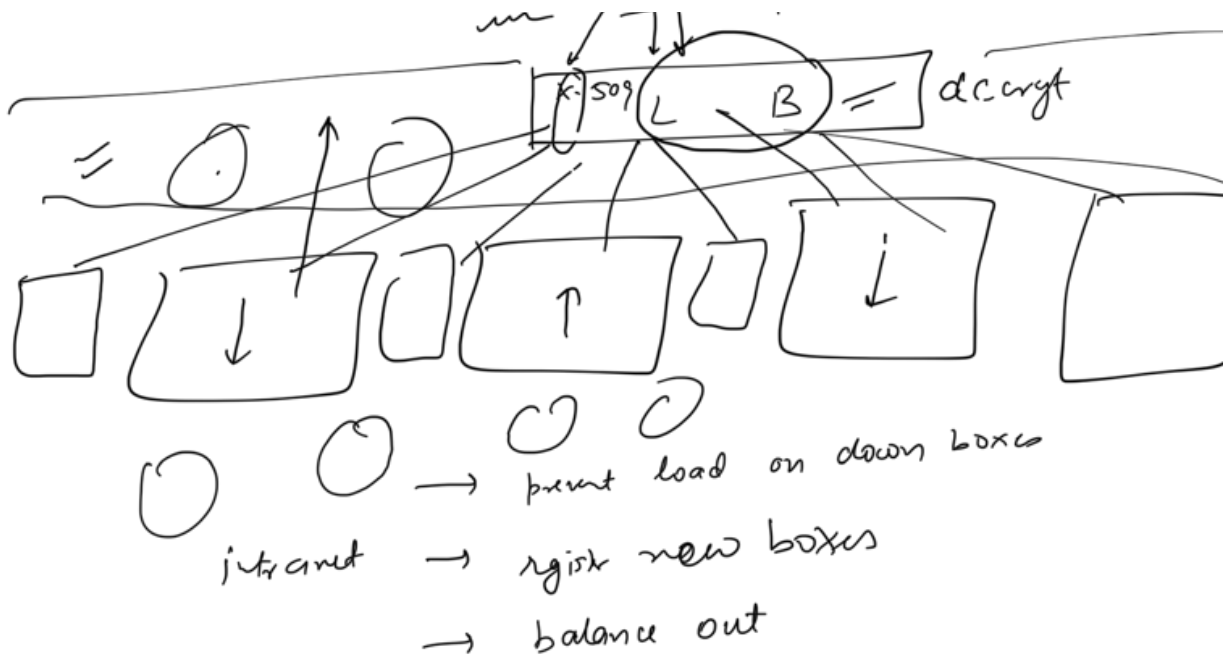
FB.com  
newFB.com

hdfc.com

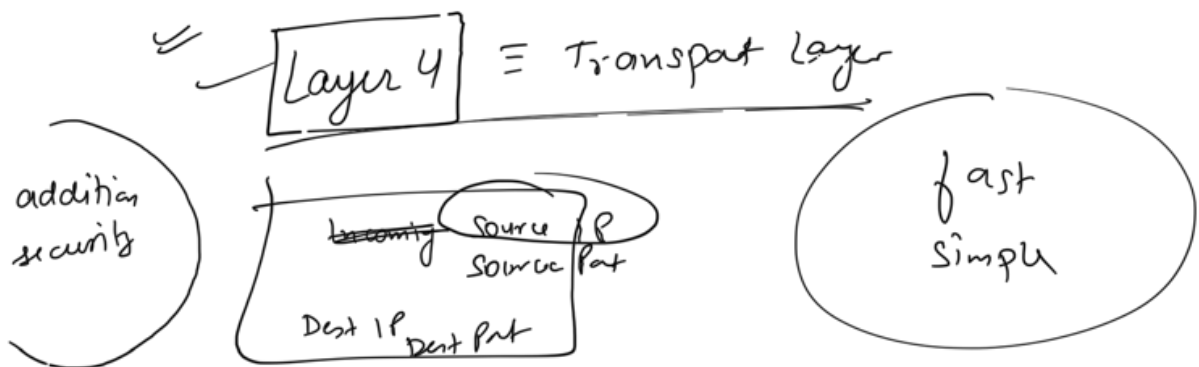


CAP Theorem





— X.509 —

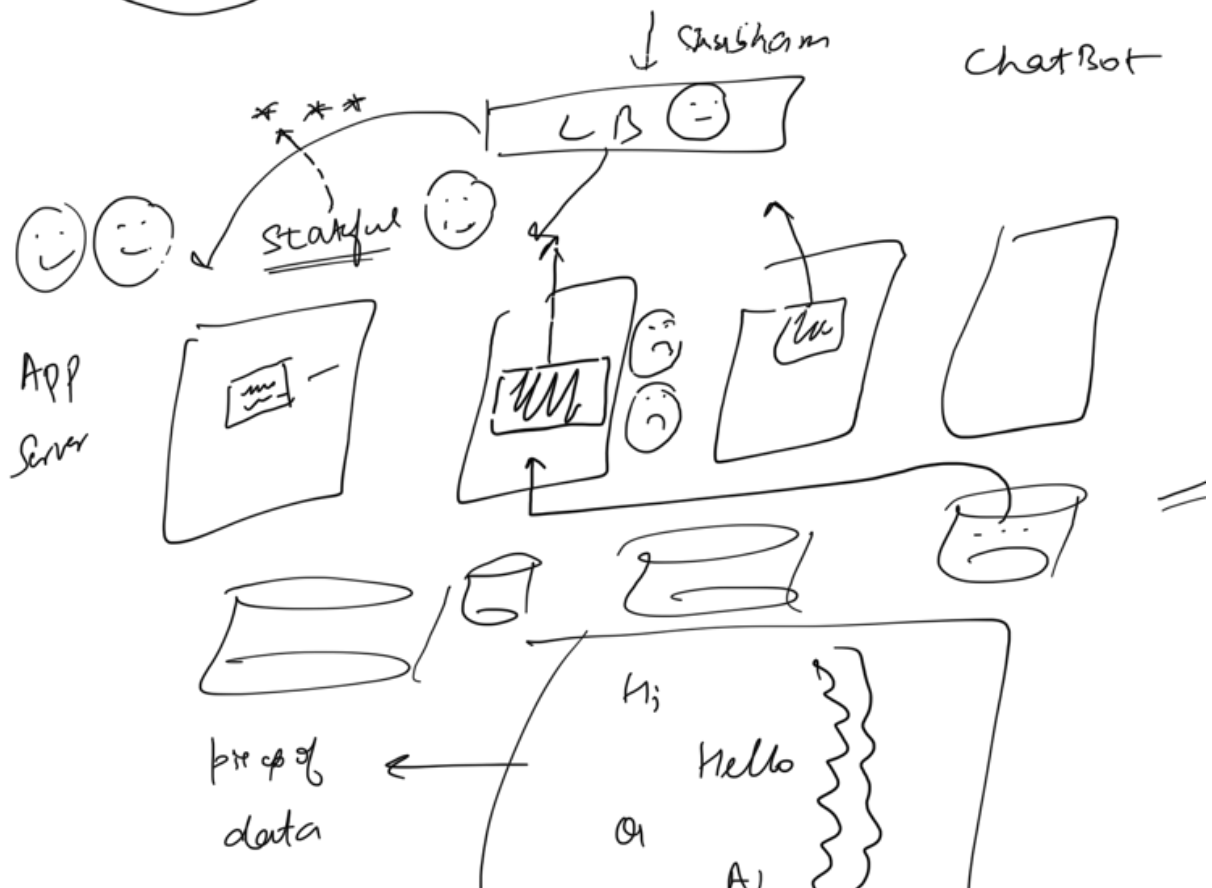
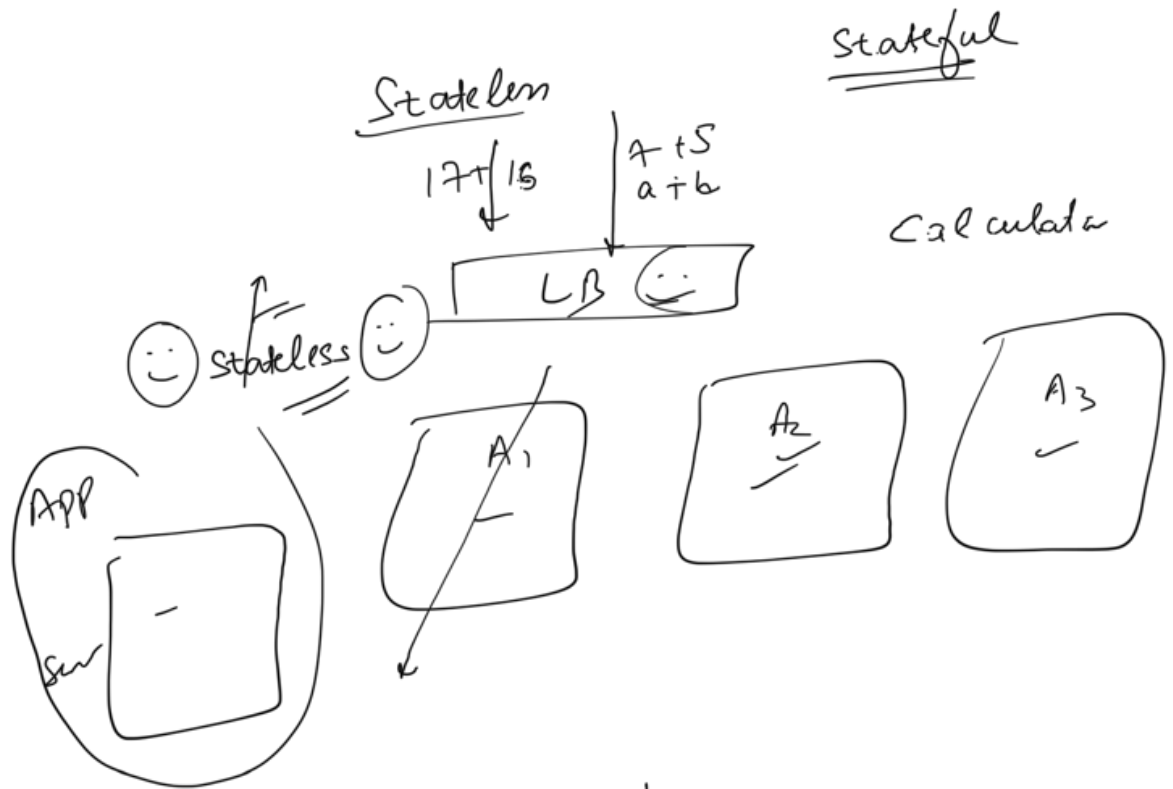


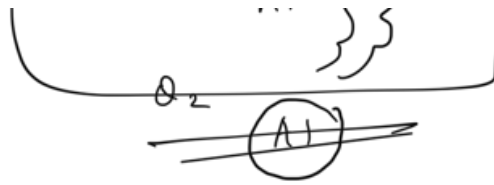
✓ Layer 7 LB = Application

user id = 707  
vid = 107

✓ much better flexibility

Slower

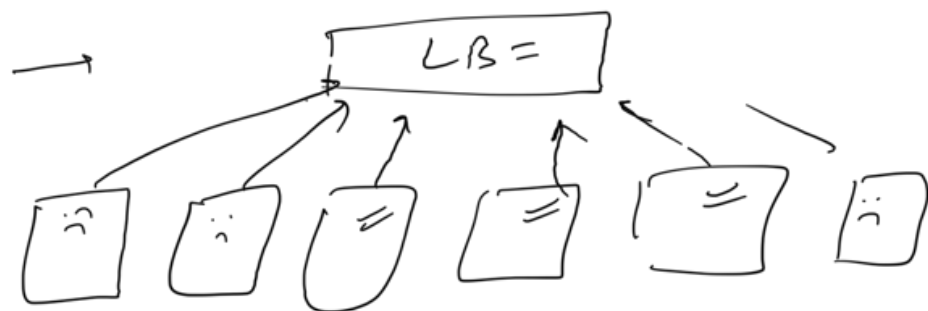
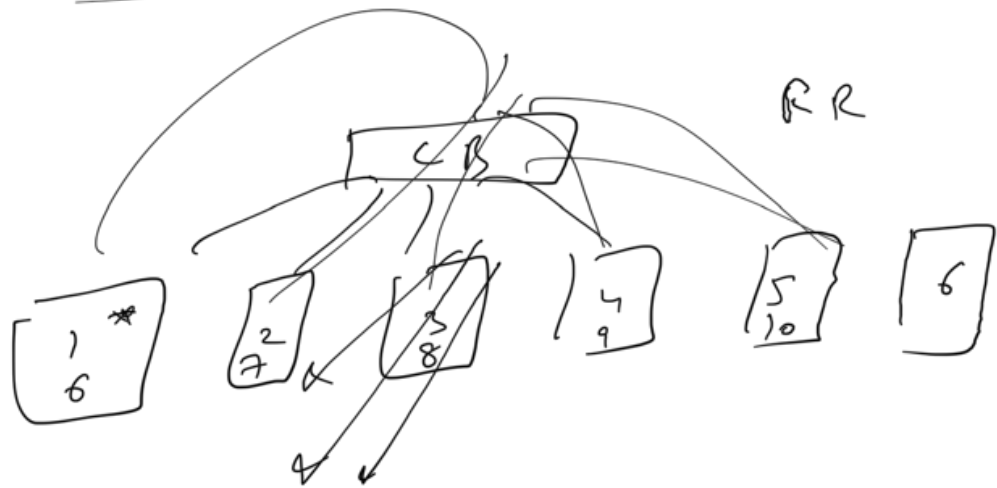




## Load Balancing //

statiken

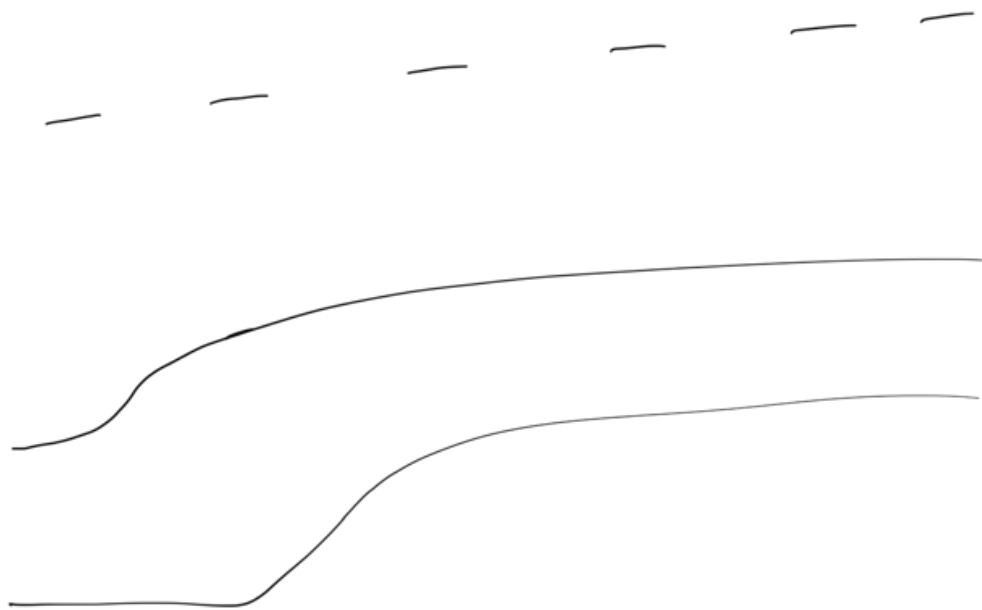
→ Round Robin



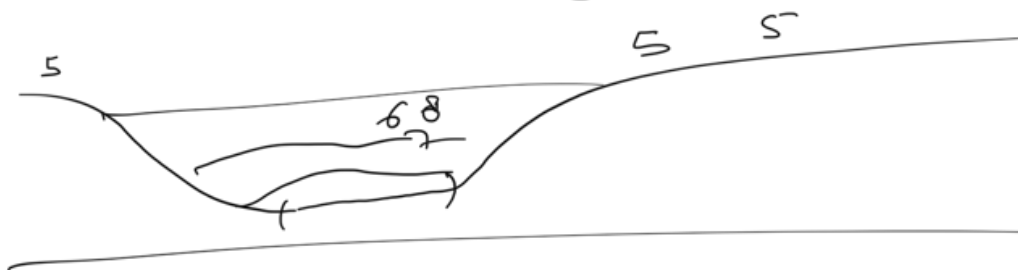
WRR

→ (Least Connection first)

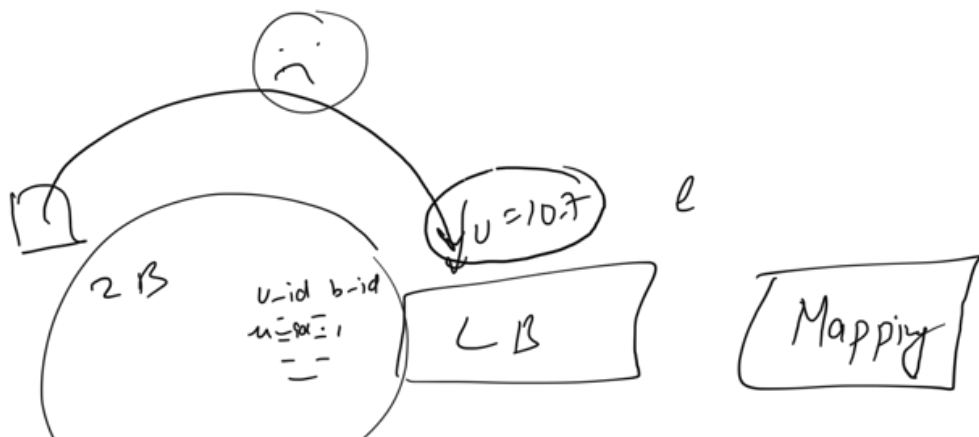
→ WRR ARR



equib



Load Balancing in Statful







✓✓✓ Sol ① Mainday Map **BAD**

map-size = order of key-size

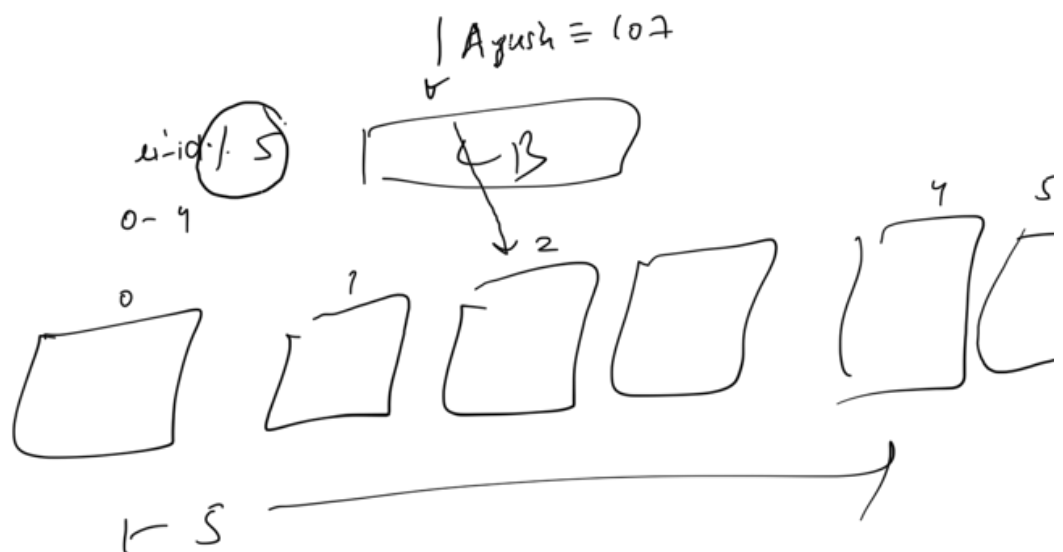
Load Balancing func  
will slow down

Sol ② Hash

**BAD!!!**

10 boxes  $\equiv$  10 App Sess

$$uid \% 10 = \overset{0-9}{\square}$$



→ amount of stickiness ... hashing

→ L.B. (smiley face)

Download  
on server in c/dic  
you will have 1.6  
to move state of all boxes  
internally

1	1.5	→	1	1.6	1
2		→	2		2
3		→	3		3
4		→	4		4
5		→	5		5
6		→	6		6
7		→	7		7
8		→	8		8
9		→	9		9

↑

③

Ranges

in L.B.

1 - 10,000

1 - 1000

1001 - 5000

5001 - 10000

↓  
I  
II  
III  
→ 10

5001 - 10000

Bottleneck: Add/Removal  
of boxes is  
inefficient



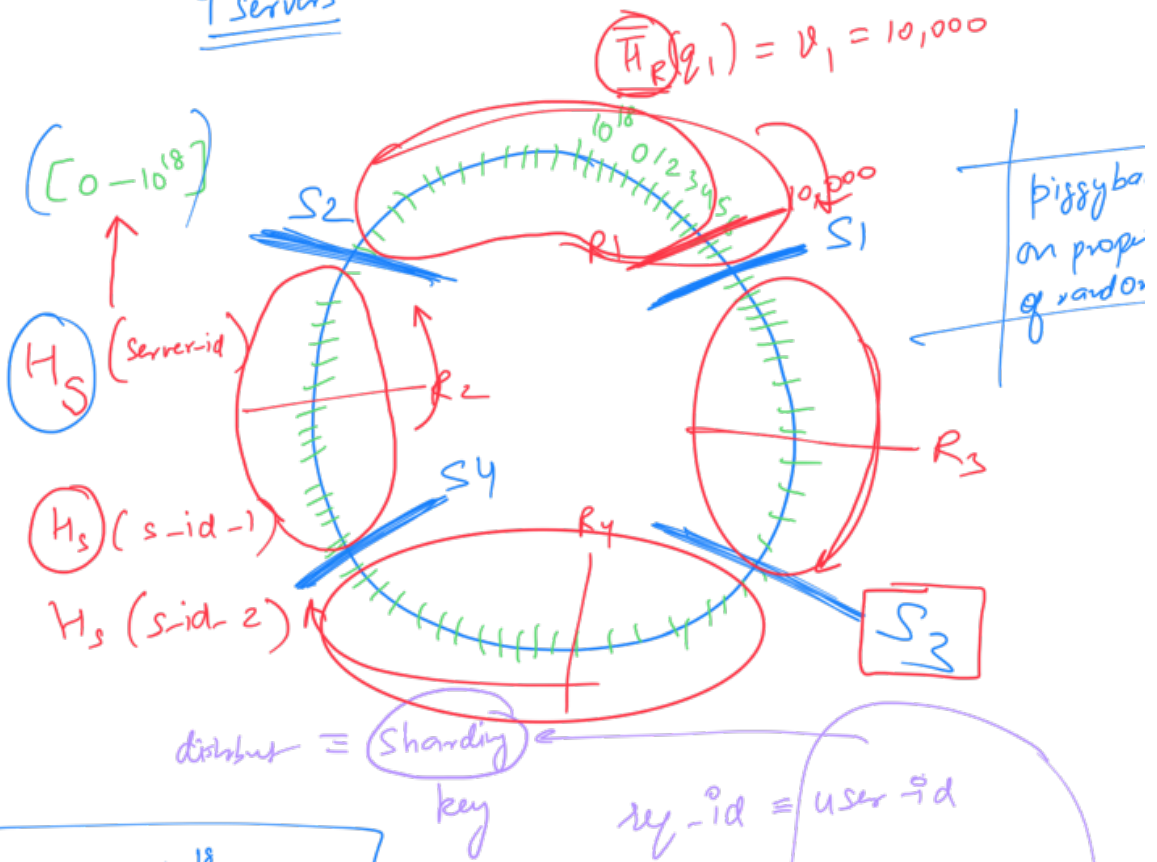
4

## Consistent Hashing

LB for Stateful Application

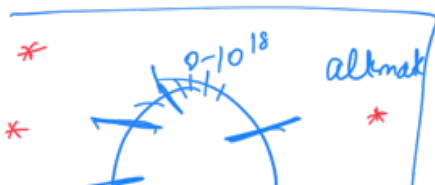
const addition/removal  
of boxes

4 servers



distributed  $\equiv$  Sharding  
key

$ry\_id \equiv \text{user id}$   
 $ry\_id \equiv \text{location id}$





Lokesh

✓ (I)

Amount of state transfer that happens  
has been reduced a lot

😊😊

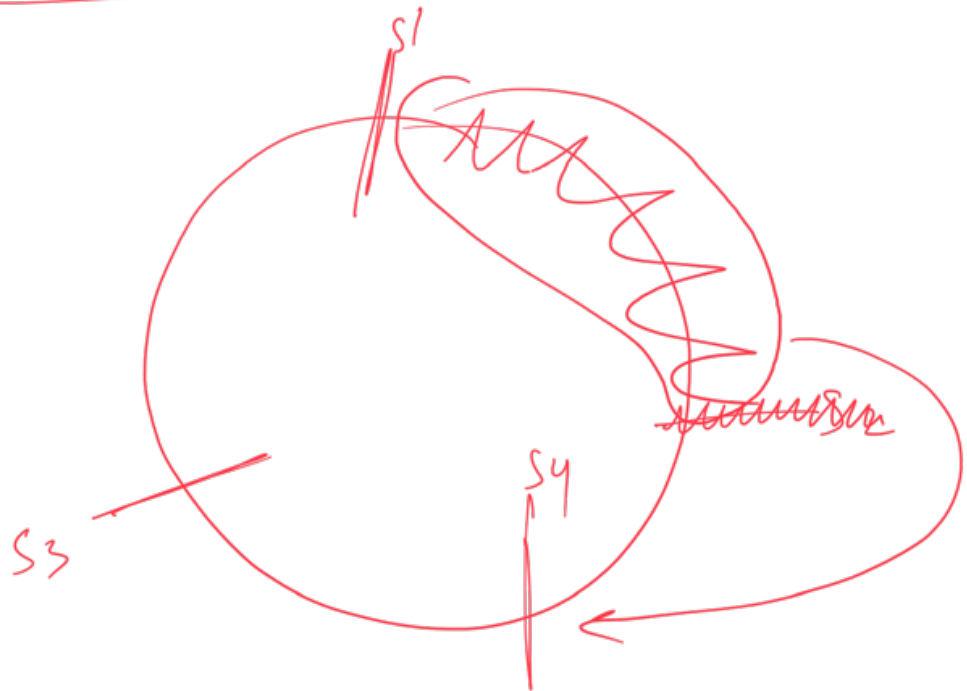
😊😊 by a factor of # servers 😊

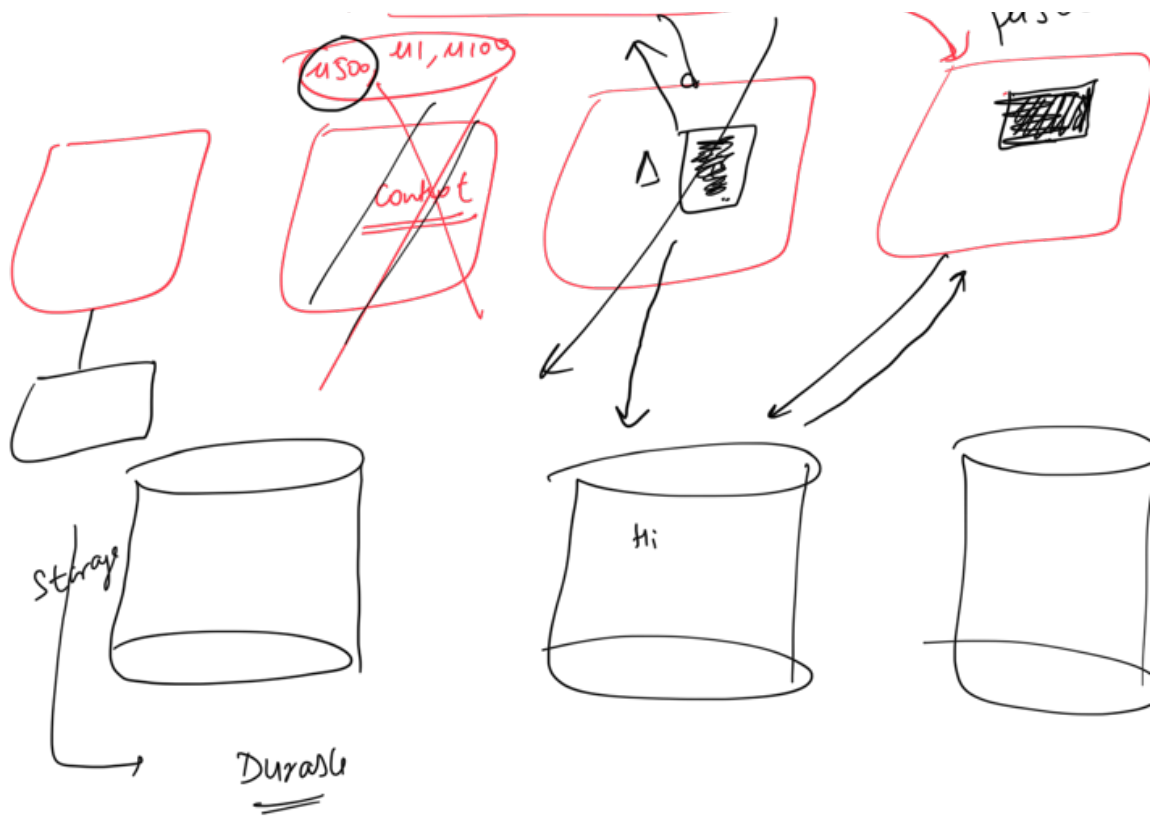
✓ (II)

downtime will also be low.....

✓ (III)

you don't have downtime for every server. Small set of  
servers get impo.





→ we are not losing the ultimate source of truth

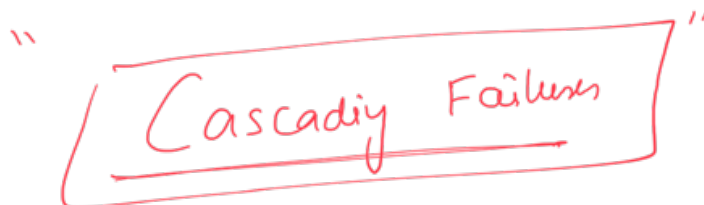
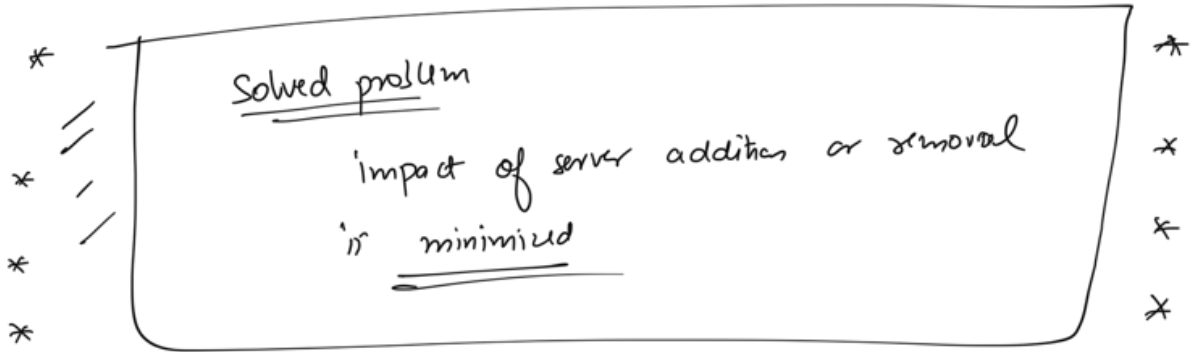
(A) cold start  $\equiv$  requests latency inc for some time  
 Hi

Model  
 01  
 A1  
 02  
 A2

Online Learning  
 Algo's  
 model

(B) pre-emptive copying

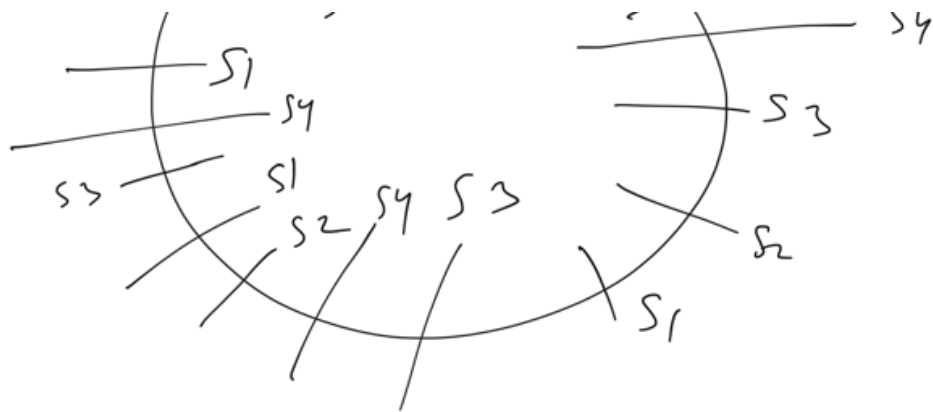
n ...



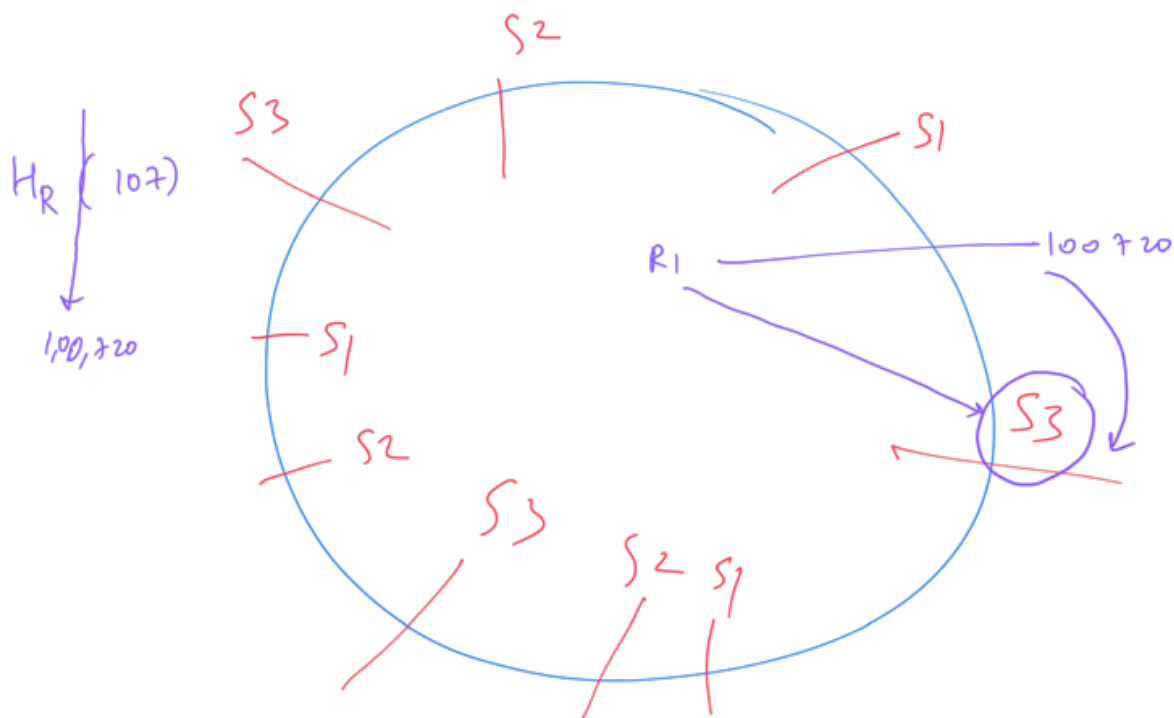
1 will go down







Solved Cascading Failures  
 [Now, the load will get divided almost equally]

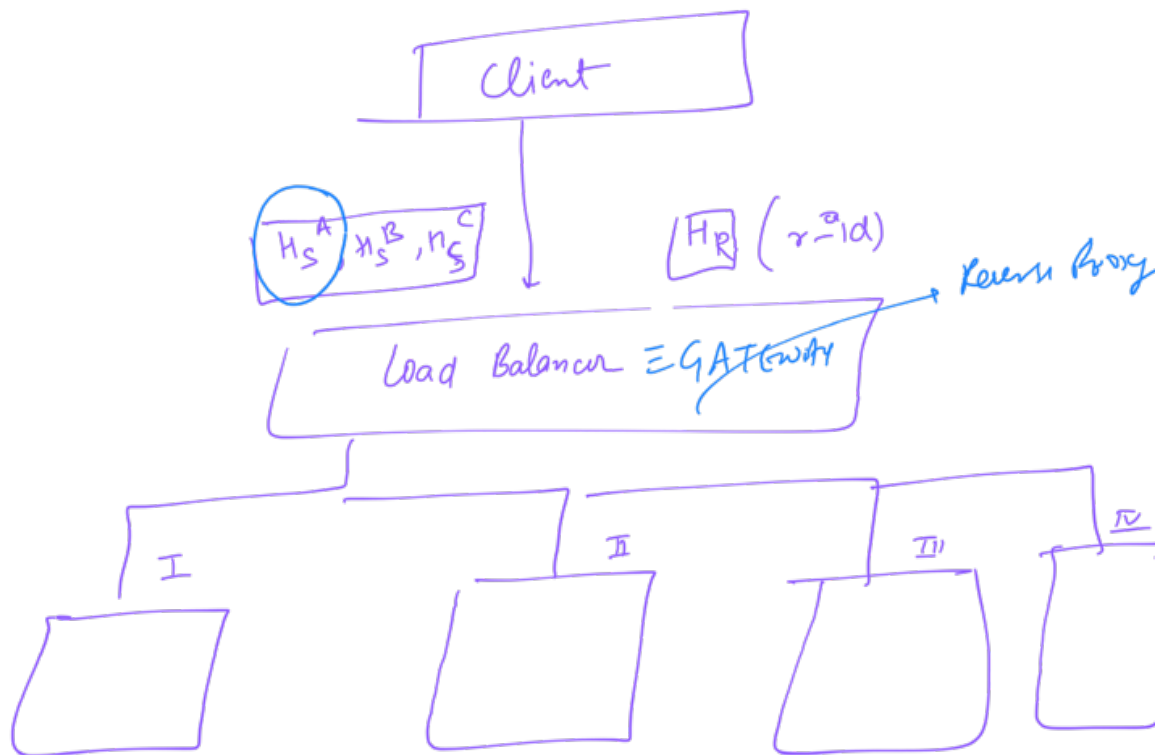


$|H_S| \uparrow \uparrow \equiv \text{Randomness} \uparrow \uparrow$   
 $\downarrow$   
 Equitable Distribution  $\uparrow \uparrow$

✓ L

$$\left[ (H_S) TTT \equiv \text{Time for Resolution} / TTT \right]$$

ASG



10 servers

ASG (Application Security Gateway)

$$h_s^n(1-10)$$

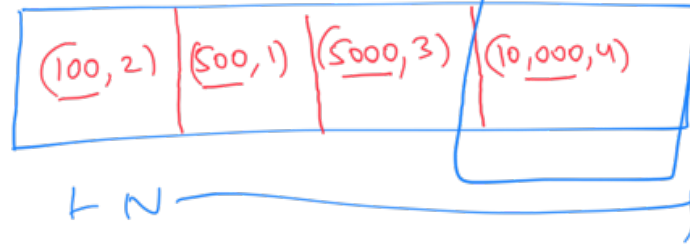
$[0, 10^{18}]$

$\sigma_{id} = 107$

$$H_R(107) = 6220$$

$H_S^A$

\*  
 $O(\log N)$   
 Map a key to  
 a server  
 \*

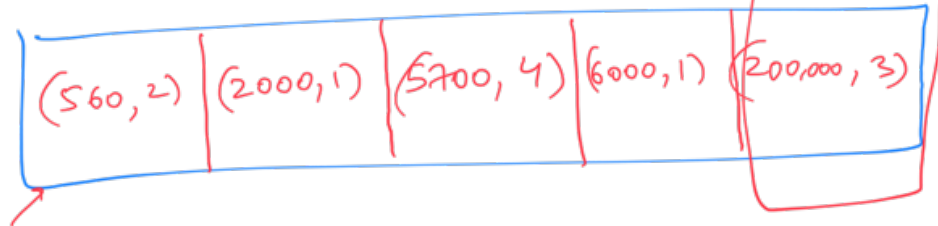


$H_S^A$

$H_S^B$

$H_S^C$

$4 \times 3$



$$H_R(107) = 6220$$

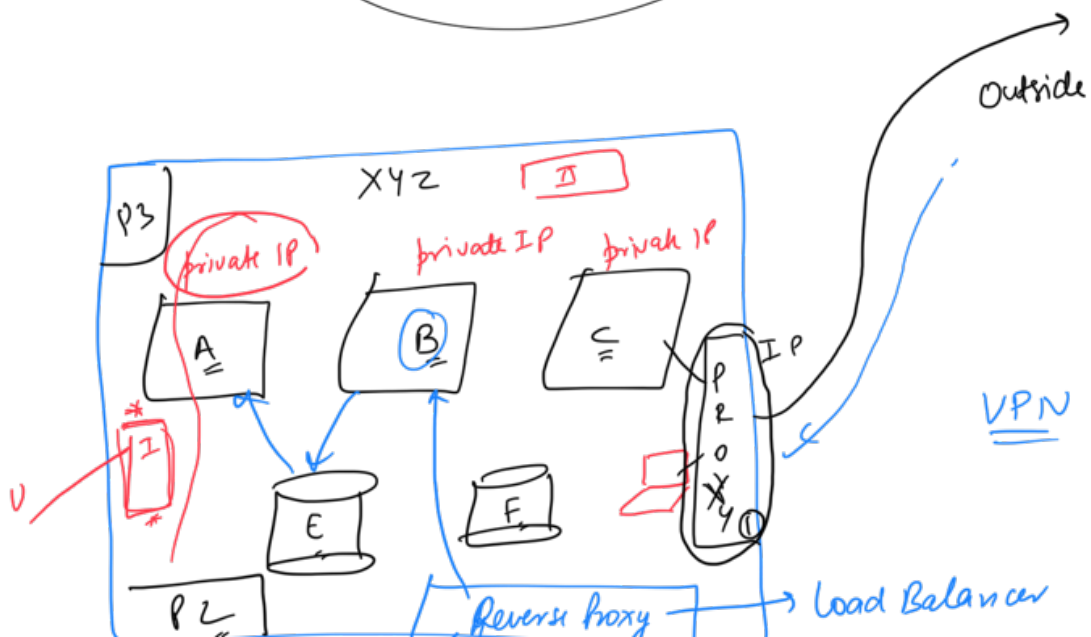
3

High Level Design

OSI Layers

The diagram illustrates a network topology with three main components:

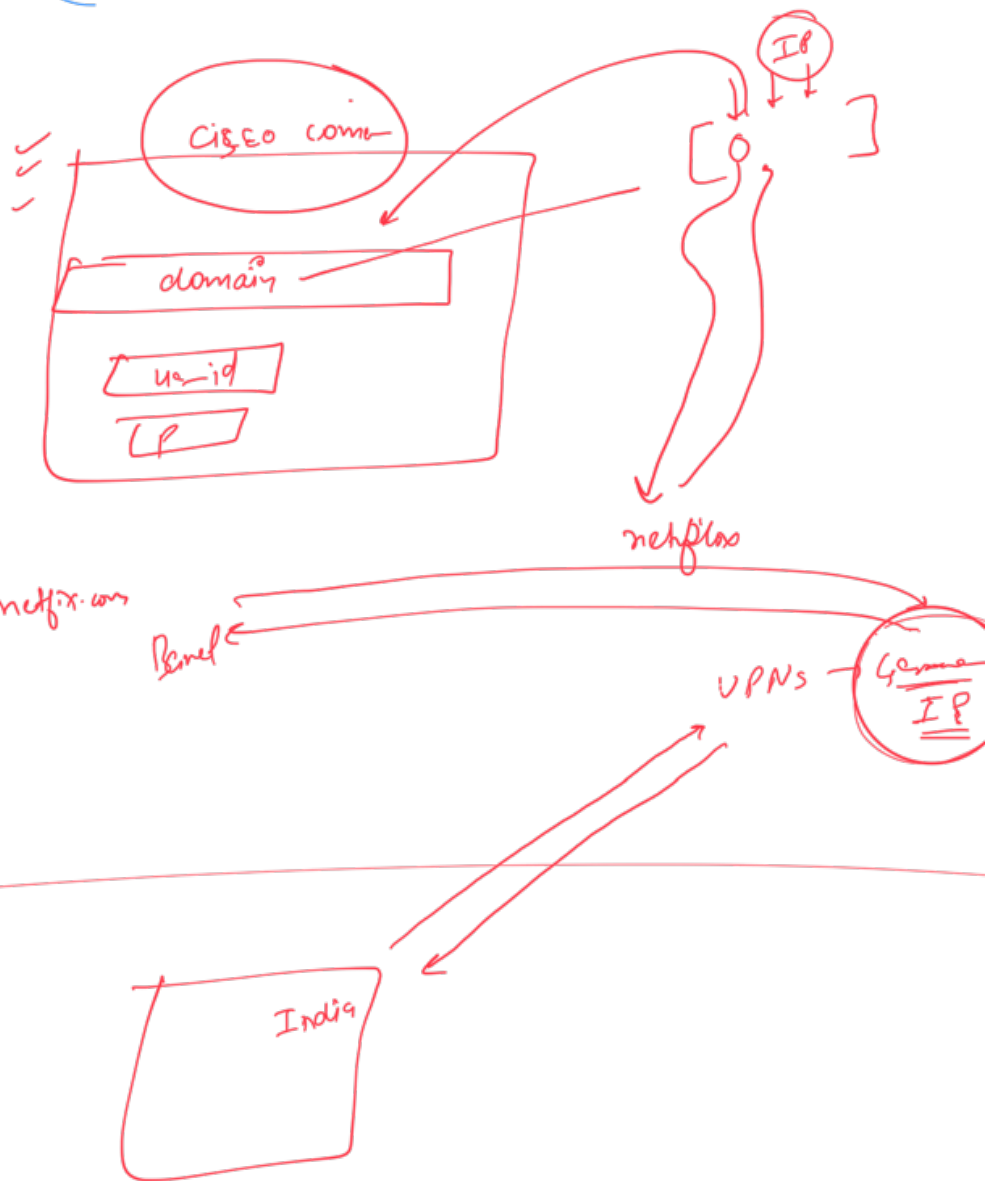
- Intranet (Top Left):** A large circle containing several square nodes. One node is labeled Intranet. A GATEWAY node is highlighted with a blue oval and a blue rectangle. Blue lines connect the gateway to several other nodes within the Intranet.
- Internet (Top Right):** Labeled INTERNET. A blue arrow points from the GATEWAY node to this label.
- Intranet (Bottom Right):** A circle labeled Intranet. A line connects it to the GATEWAY node.
- Unlabeled Intranet (Bottom):** A large empty circle at the bottom, connected to the GATEWAY node by a vertical line.



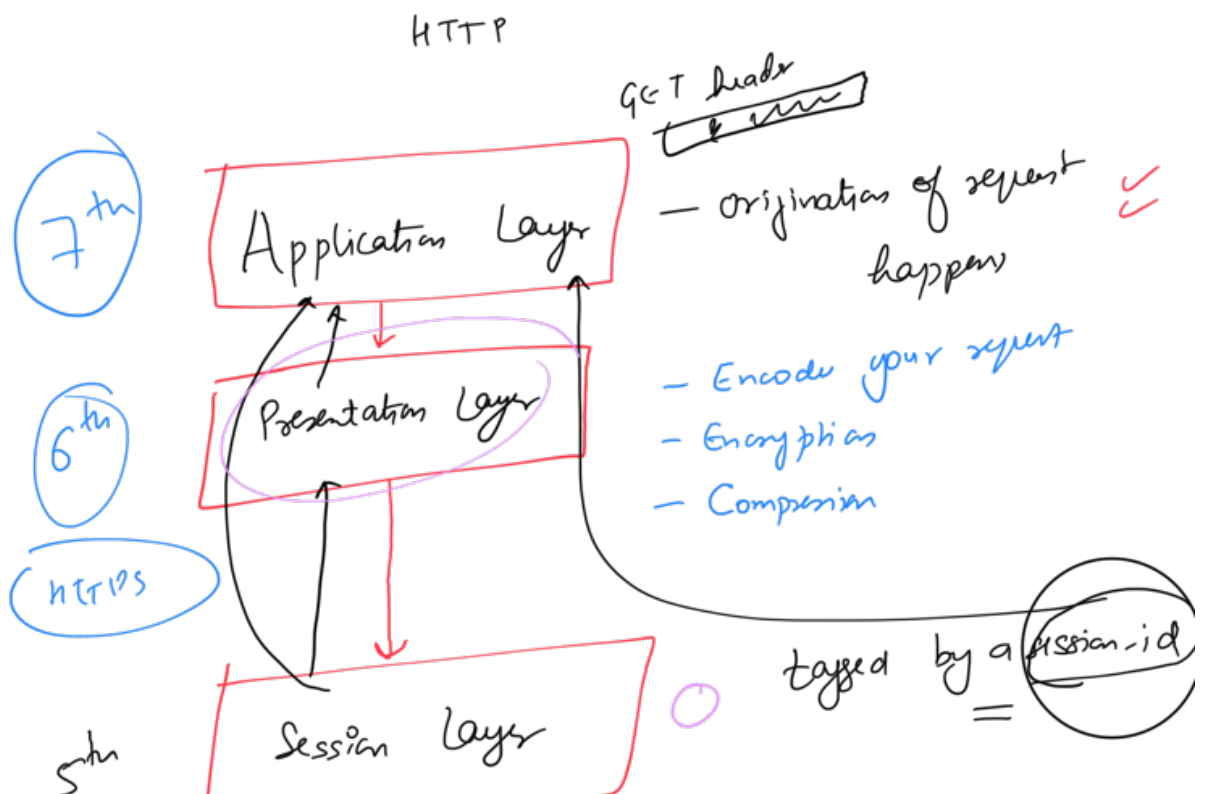
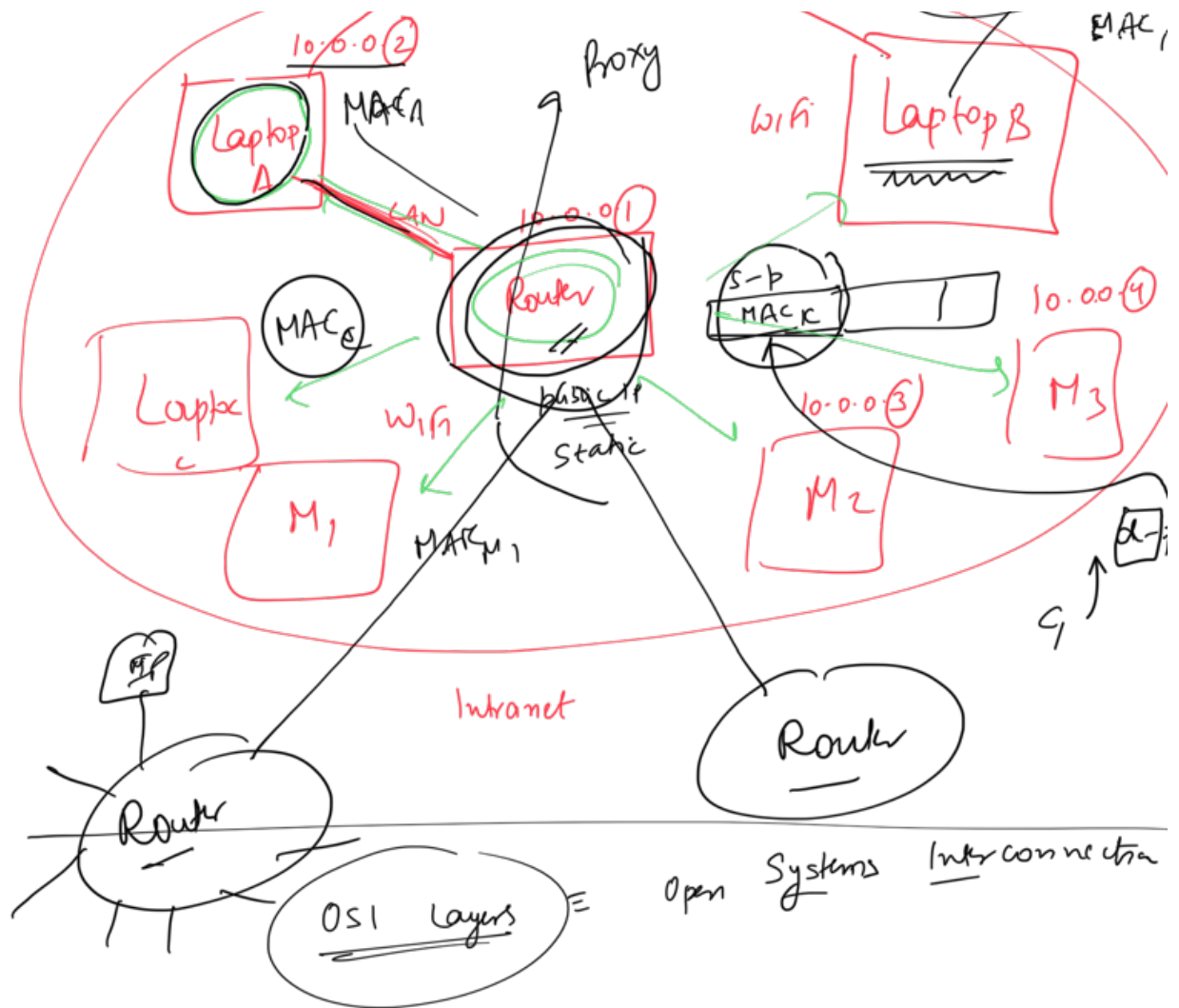
Gateway

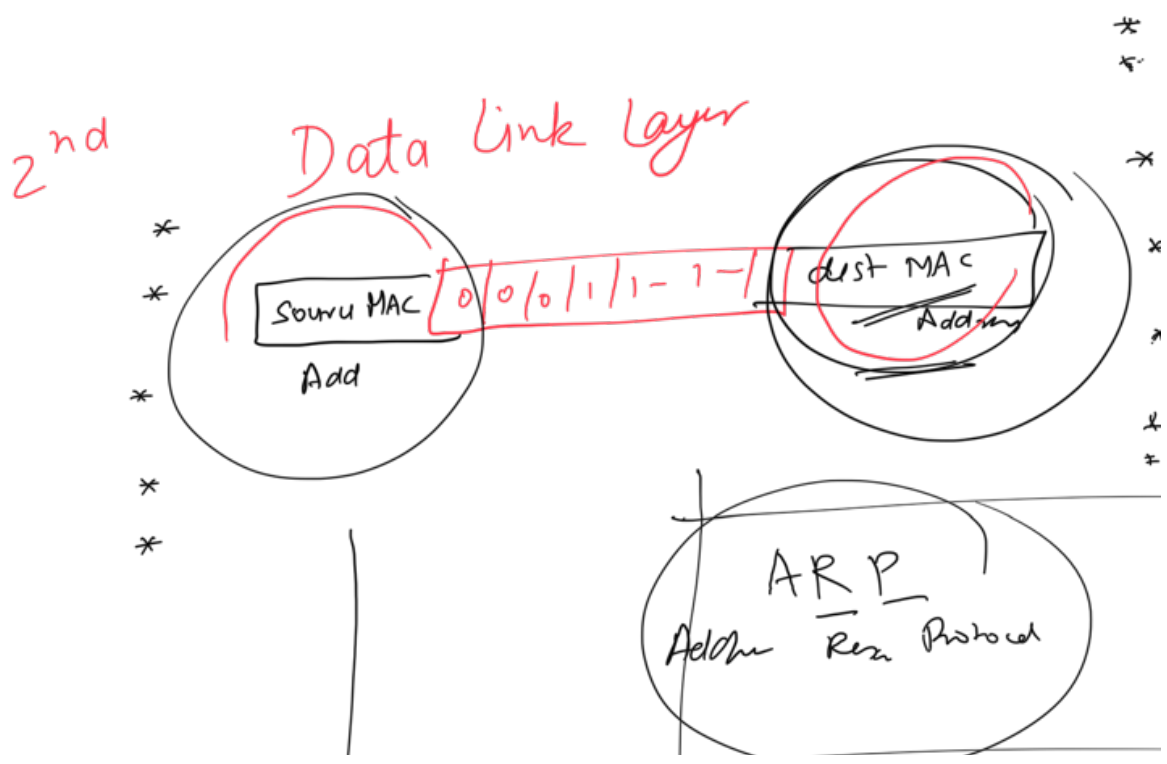
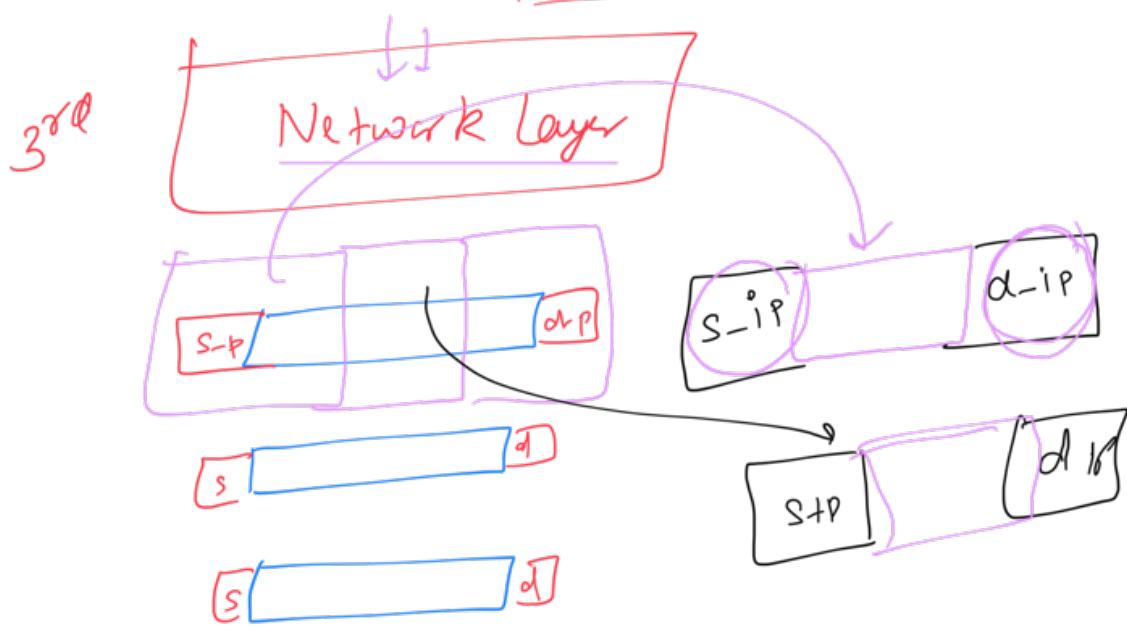
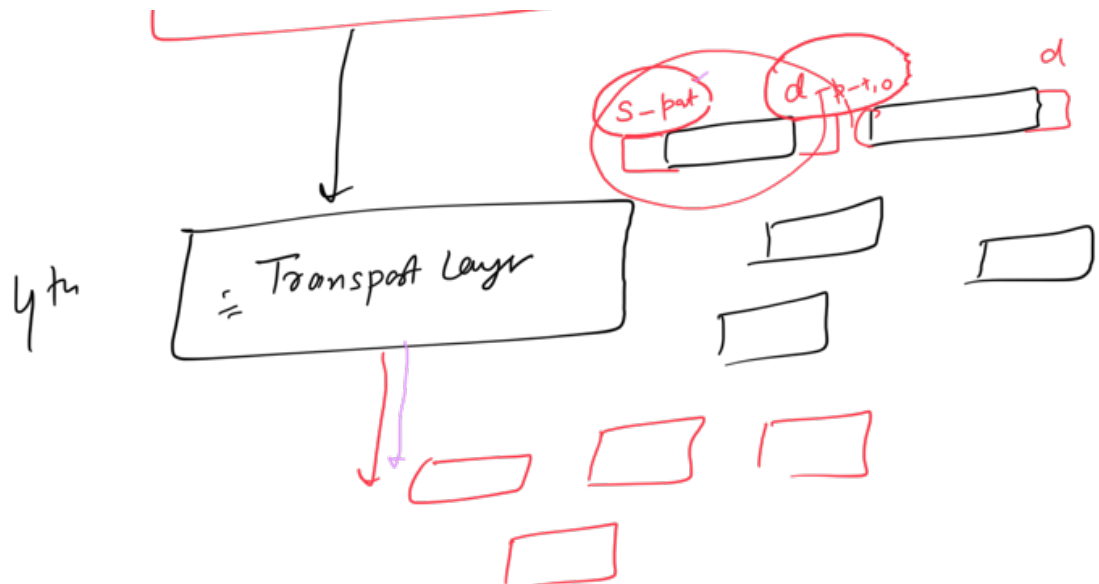
Proxy  $\equiv$  Gateway to the outside world

Reverse proxy  $\equiv$  Gateway to the inside world

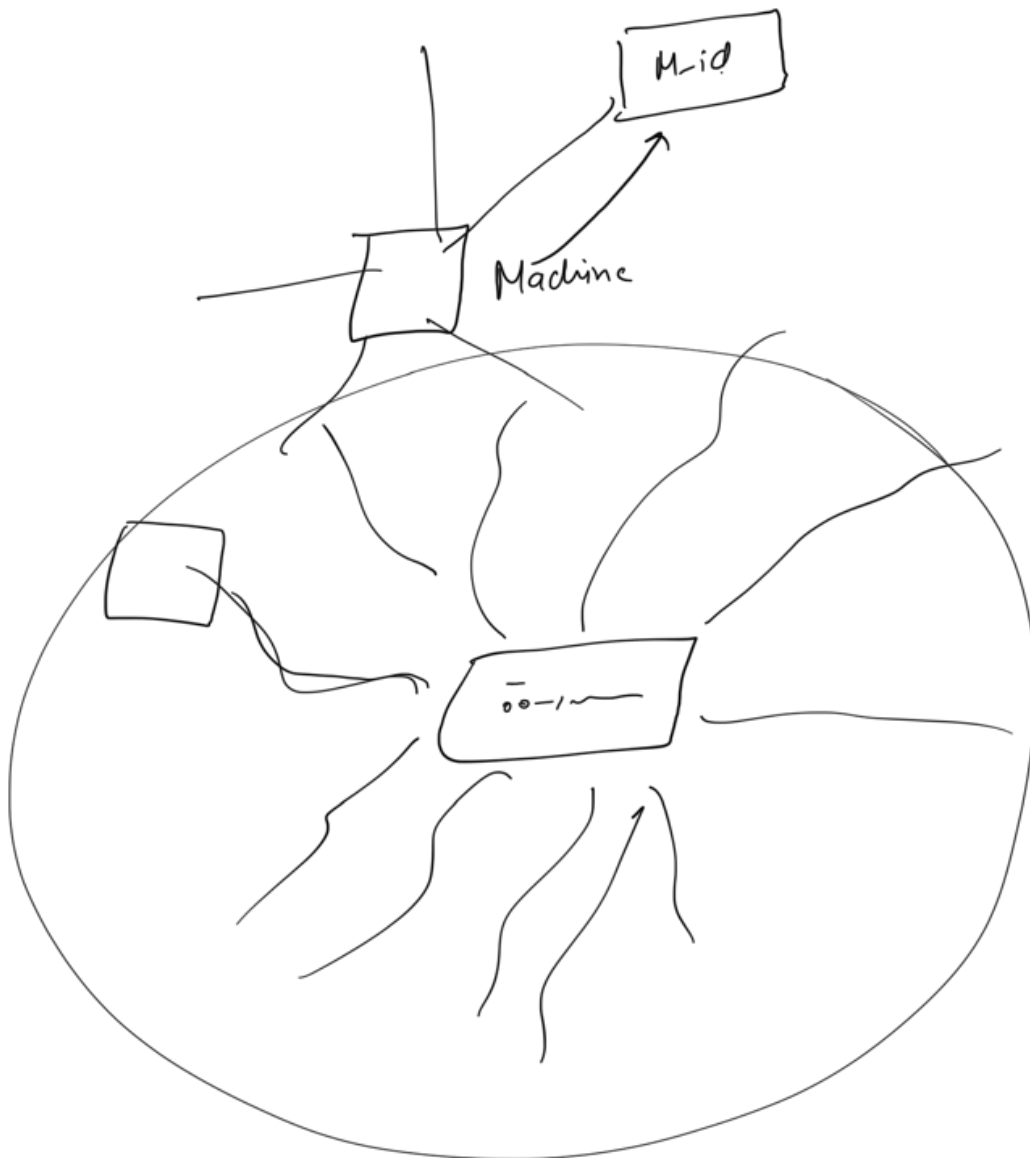
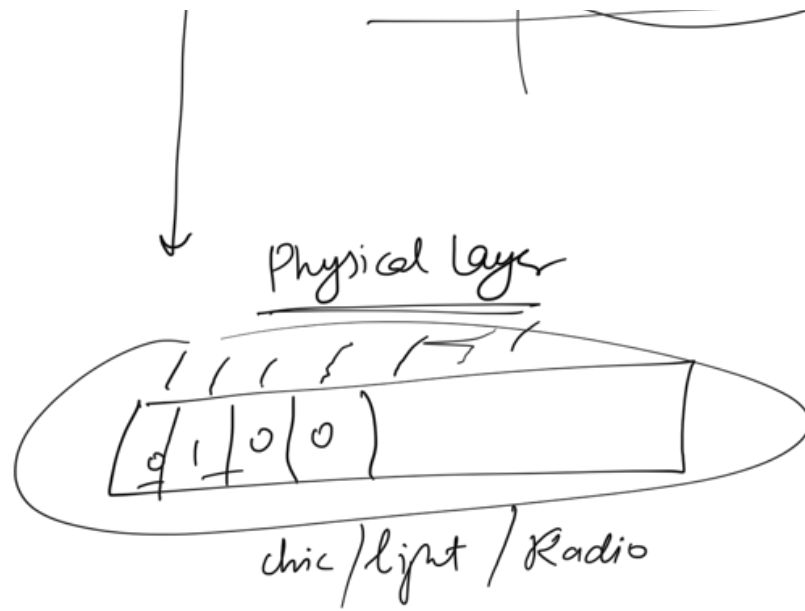


hashta service

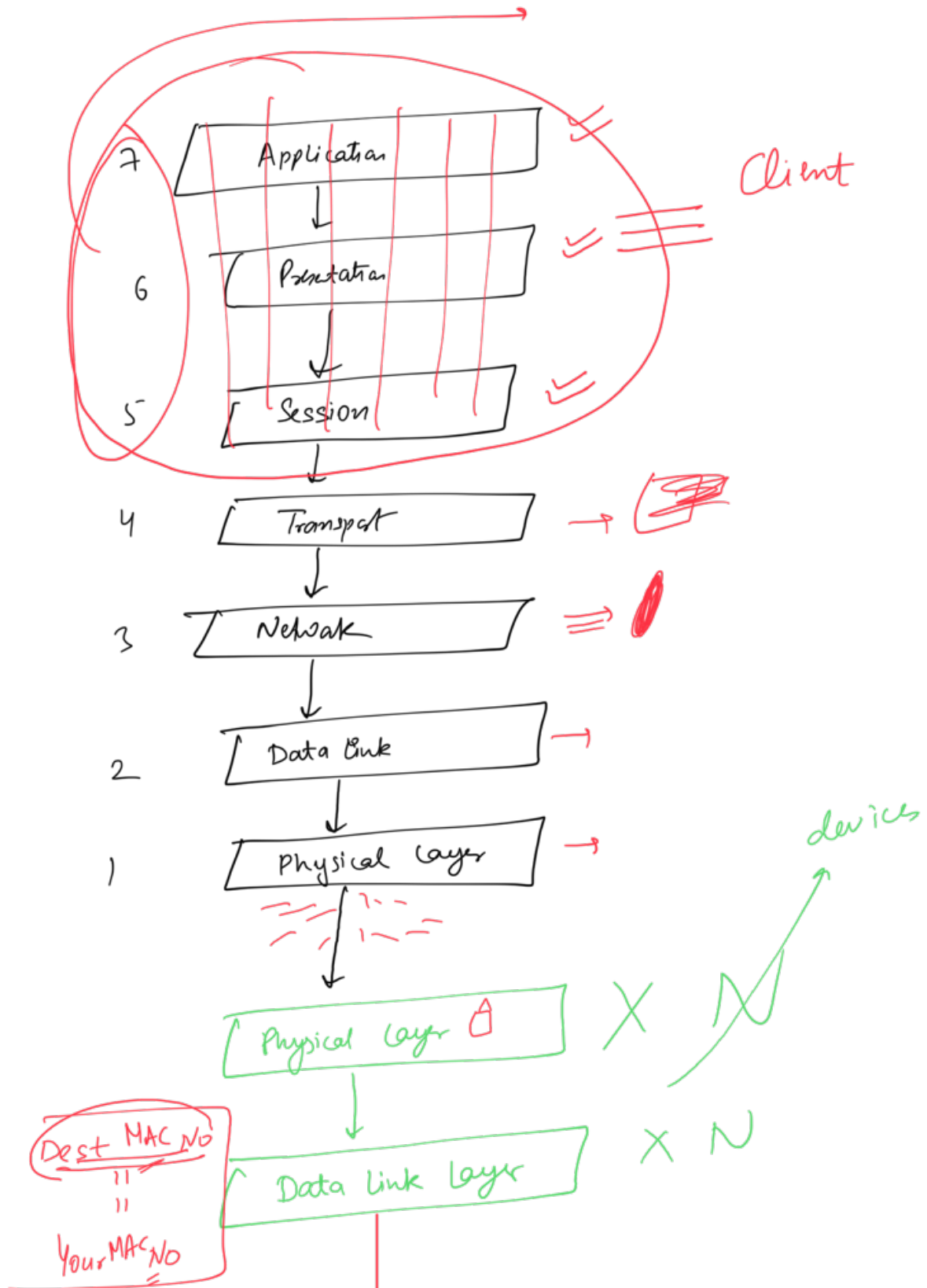
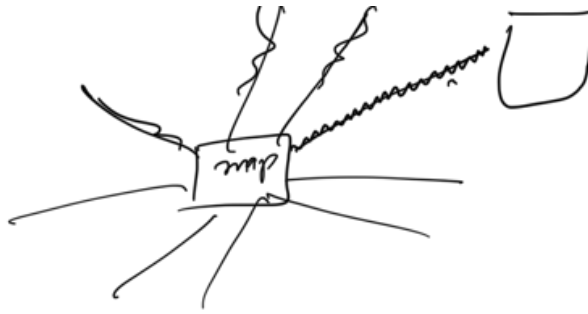


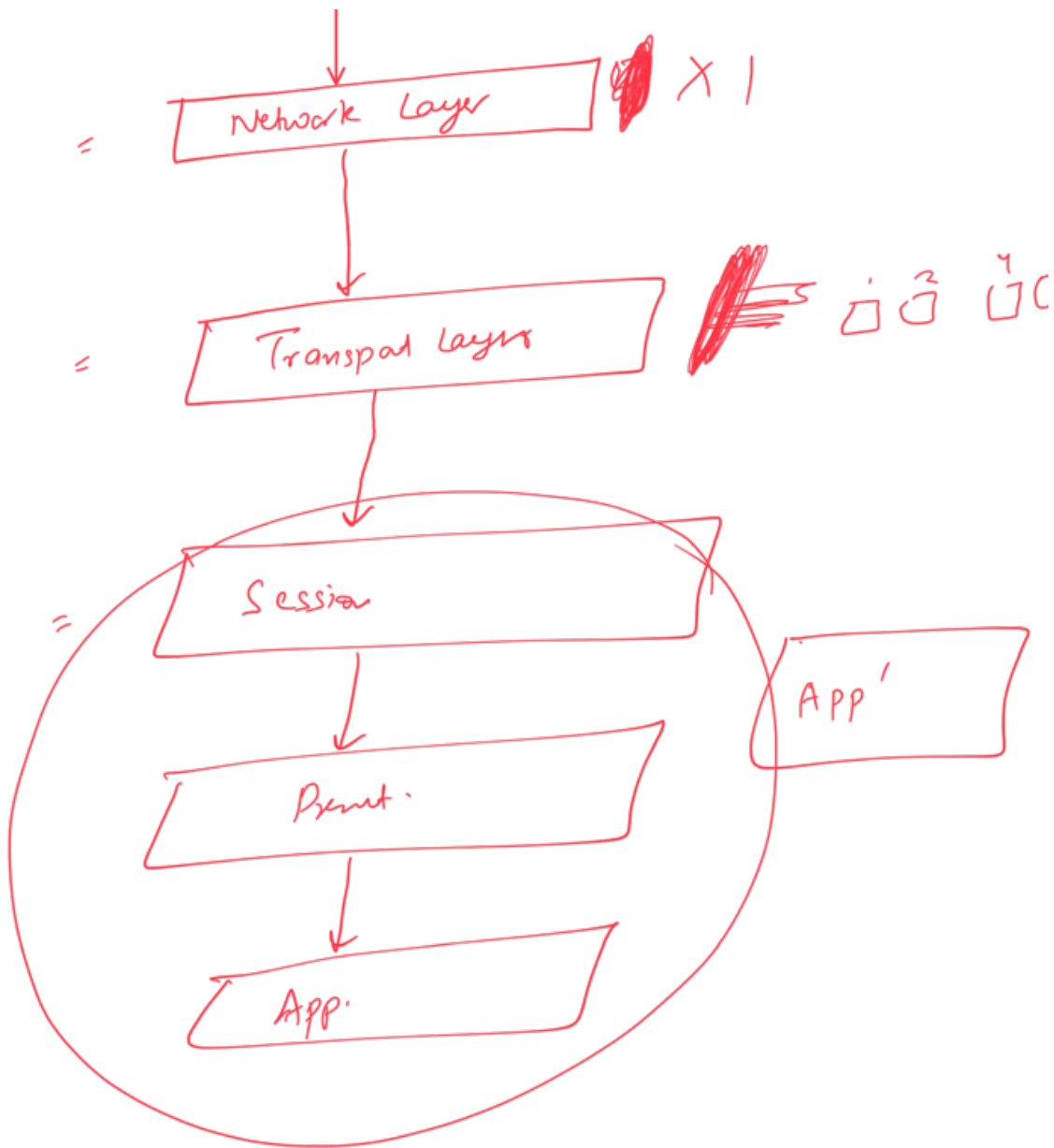


1st











~~110~~ 110 - -  
210 110 - -