An Illustrated Guide to LLVM

or: an introduction to building simple and not-so-simple compilers with Rust and LLVM

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What?

A library for building compilers

Formerly "Low level virtual machine"

- Compiler backends
 - Code generation, optimization
 - Not lexing, parsing

Supports both ahead-of-time and just-in-time

Industrial-grade

- Used in industry
 - Apple
 - Google
 - Others

- Mature
 - First release in 2003
 - ~5 million LOC

Portable

Supports many systems:

- High-performance
 - x86
 - PowerPC
 - SPARC
- Embedded
 - ARM
 - PowerPC
 - SPARC

- GPUs
 - AMD GCN
 - Nvidia PTX
- Exotics
 - BPF
 - Hexagon

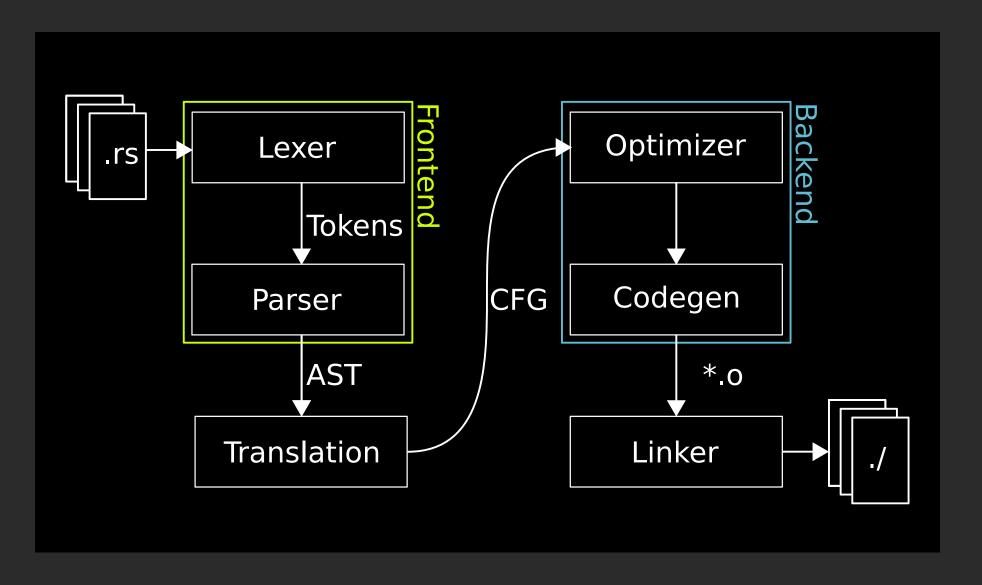
Numerous frontends

- Clang (C)
- GHC (Haskell)
- LDC (D)
- OpenJDK (Java)

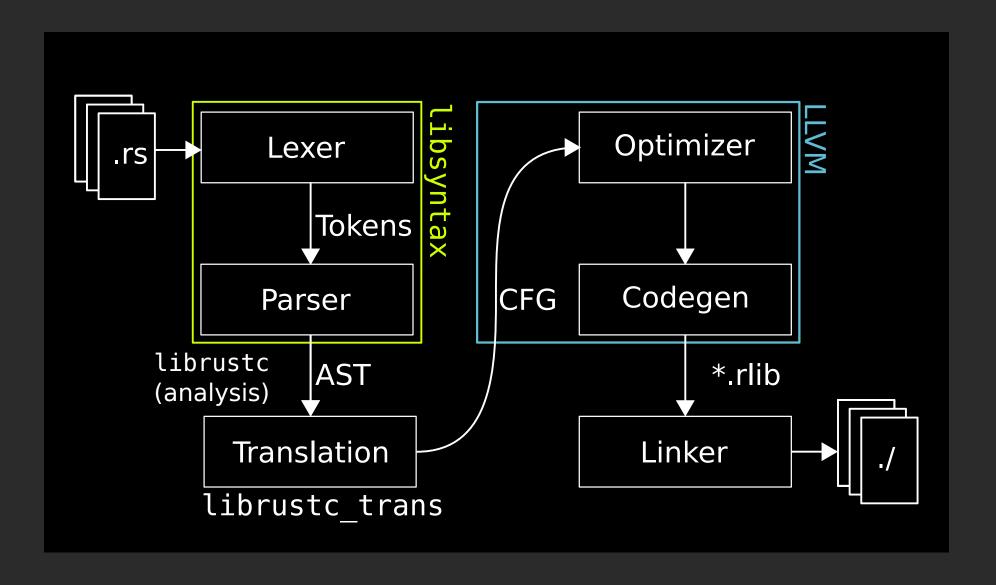
.. and of course rustc.

How?

Compiler structure



rustc



Talking to LLVM

IR goes in..

```
fn add(x: i32, y: i32) -> i33 {
    (x as i33) + (y as i33)
}

target triple = "x86_64-unknown-linux"

define external i33 @add(i32 %x, i32 %y) nounwind readnone {
    %xe = sext i32 %x to i33
    %ye = sext i32 %y to i33
    %result = add i33 %x, %y
    ret i33 %result
}
```

..code comes out

```
.text
.globl add
add:

movsxd rcx, edi
movsxd rax, esi
add rax, rcx
ret
```

More complex

Testing the Collatz conjecture:

```
fn collatz(x: u32) -> bool {
    if x == 1 {
        return true;
    }

    let next = if x % 2 == 0 {
        x / 2
    } else {
        (3 * x) + 1
    }
    collatz(next)
}
```

collatz(u32) in IR

```
define "fastcc" i1 @collatz(i32 %x) {
    %finished = icmp eq i32 %x, 1
    br i1 %finished, label %Base, label %Continue

Base:
    ret i1 1

Continue:
    %next = alloca i32
    %odd = urem i32 %x, 2
    %odd1 = trunc i32 %odd to i1
    br i1 %odd1, label %Odd, label %Even
```

collatz(u32) continued

```
Odd:
    %halved = udiv i32 %x, 2
    store i32 %halved, i32* %next
    br label %Recurse
Even:
    %larger = mul i32 %x, 3
    %larger1 = add i32 %larger, 1
    store i32 %larger1, i32* %next
    br label %Recurse
Recurse:
    %nextval = load i32, i32* %next
    %result = musttail call i1 @collatz(i32 %nextval)
   ret il %result
```

SSA

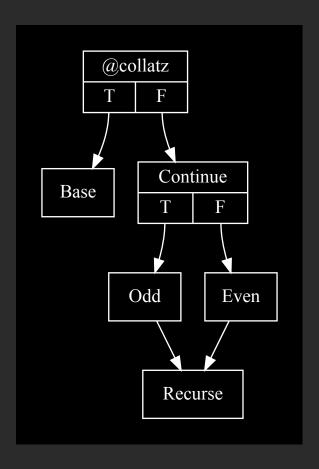
Why not this?

```
%nextval = %halved
br label %Recurse
Even:
    // ...
    %nextval = %larger1
Recurse:
    // use nextval
```

Single static assignment

- Every value has exactly one assignment
- Allows the system to work with a true CFG

Control flow graph



Native format for optimization.

Speaking Merthese

Token	Action
m	Print "merth"
е	Print "\n"
r	Print " "
t	Print random [a-z] [0, 13.4) times
h	Jump to after the next 'h'
_	Do nothing

Planning

Primitive operations

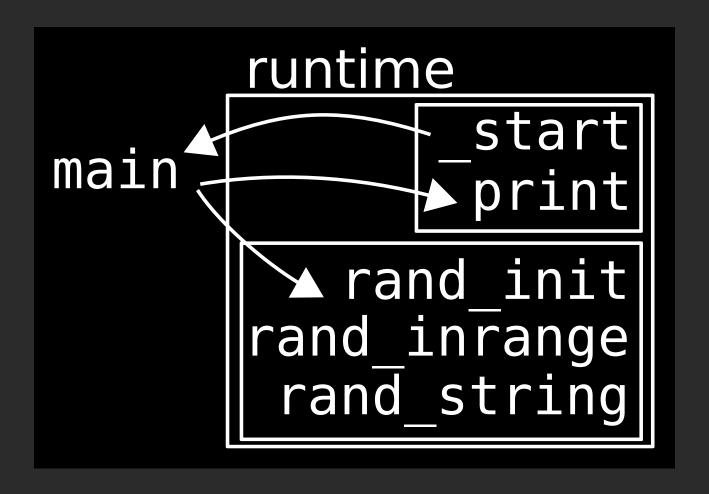
- Print string
 - fn print(s: *mut u8, len: u8)
- Random integer
 - fn rand_inrange(max: u8) -> u8

These are not supported by any CPU...

Runtime library

- Statically linked
 - Better optimization (inlining!)
 - Self-contained
- Written in IR
 - Because we can
 - More portable

Program structure



start

- Initialize RNG
 - Open/dev/urandom
 - Read bytes
 - Close file
- Call main
- exit(0)

We need to do syscalls

extern fn syscall(nr: i64, ...) -> i64;

- RAX: nr
- RDI: p1
- RSI: p2
- Result: RAX


```
@__NR_open = private constant i64 2

define private i32 @open(i8 *%path0, i32 %flags0) {
    %nr = load i64, i64* @__NR_open
    %path = ptrtoint i8* %path0 to i64
    %flags = zext i32 %flags0 to i64
    %out0 = call i64 @syscall2(i64 %nr, i64 %path, i64 %flags)
    %out = trunc i64 %out0 to i32
    ret i32 %out
}
```

Turns out syscalls are boring. Back to start.

```
declare void @main()
define void @exit(i32 %code) noreturn {}

define void @_start() noreturn {
    // initialize RNG
    call void @main()
    call void @exit(i32 0)
    unreachable
}
```

Feels like C: declare external functions, glue them together.

Writing some Rust

Skeleton

```
extern crate llvm_sys as llvm;

fn main() {
    unsafe {
        LLVM_InitializeNativeTarget();
        LLVM_InitializeNativeAsmPrinter();
        LLVM_InitializeNativeAsmParser();

        let ctxt = LLVMContextCreate();
        /* Use ctxt */
        LLVMContextDispose(ctxt);
    }
}
```

Create main

declare void @main()

Emitting IR

Could we use io::Write instead?

Dropping to C++

```
extern "C" typedef int (*cb t)(const void *, size t, void *);
class raw callback ostream : public llvm::raw ostream {
    cb t callback;
    void *callback data;
public:
    raw callback ostream(cb t cb, void *cb data) { /* ... */ }
private:
    void write impl(const char *p, size t sz) override {
        callback(p, sz, callback data);
        offset += sz;
};
```

A C function to expose it:

Rust adapter from Write to callbacks

Safe wrapper

Emit some IR



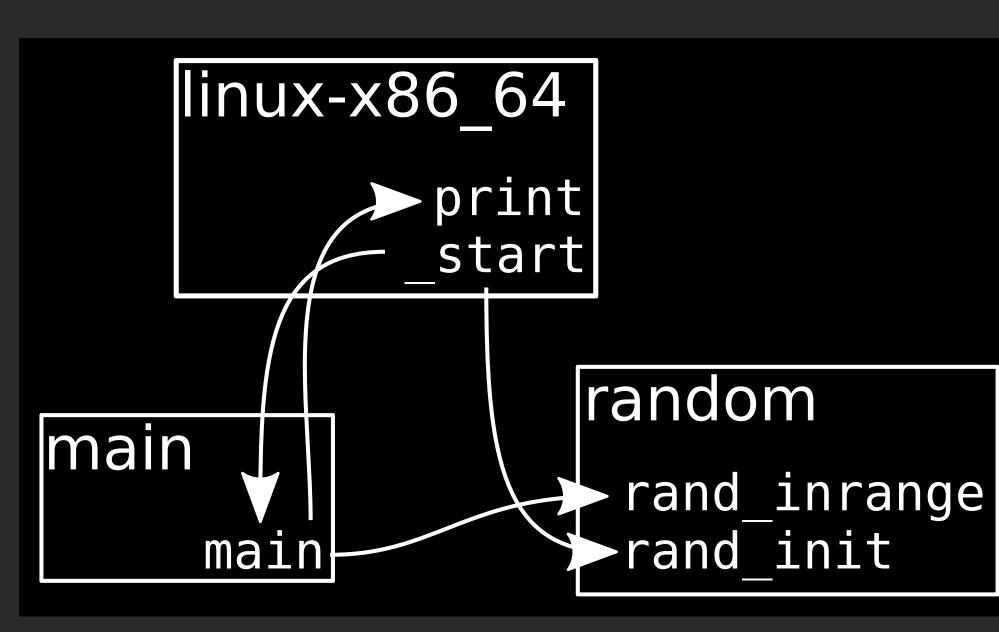
Using Builders



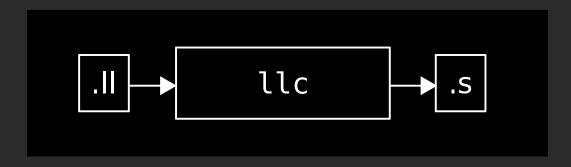
Manual testing

The pieces so far

```
linux-x86_64.ll
   _start and print functions for x86_64 Linux
random.ll
   rand_inrange function and RNG impl
main.ll
   Generated main function
```



Code generation







```
llc linux-x86_64.ll
as -o linux-x86_64.o linux-x86_64.s

llc random.ll
as -o random.o random.s

llc main.ll
as -o main.o main.s

ld main.o linux-x86_64.o random.o
./a.out
```

Optimization

The generated code is inefficient!

- Less so if optimization is turned on
 - llc -02
- Better: LTO
 - llvm-link *.ll -o a.bc
 - llc -02 a.bc, etc

We can do better (but not right now)

More Rust

Using the runtime

```
let ty void = LLVMVoidTypeInContext(ctxt);
let ty i8 = LLVMIntTypeInContext(ctxt, 8);
let ty i8p = LLVMPointerType(ty i8, 0);
let param types = [ty i8p, ty i8];
let ty rt print = LLVMFunctionType(
        ty void,
        param types.as ptr() as *mut ,
       param types.len(),
        0);
let rt print = LLVMAddFunction(
        llmod,
        b"print\sqrt{0}".as ptr() as *const _,
        ty rt print);
```

declare void print(i8*, i8)

Other functions left as an exercise for the reader.

Being wrong

LLVMVoidType ↔ LLVMVoidTypeInContext

Implicit global context ↔ explicit

- Mixing contexts is wrong and can cause miscompilation
- Tools to help prevent bugs?

Bug swatting

- LLVMVerifyModule
 - Print message or abort
- Debug assertions
 - LLVM_ENABLE_ASSERTIONS
 - Usually not enabled in binary releases
- Manual inspection
 - As done earlier

Filling in main

Parsing code

```
let code: Iterator<char>;
while let Some(c) = code.next() {
    match c {
        'm' => { /* Print "merth" */ },
        'e' \Rightarrow { /* Print newline */ },
        'r' => { /* Print space */ },
        't' \Rightarrow { /* Random string */ },
        'h' => { loop { match code.next() {
             Some('h') | None => break,
             => continue,
        } } }
        _ => { /* Do nothing */ },
```

m is for merth

 $m \rightarrow print("merth", 5)$

e is for newline

 $e \rightarrow print("\n", 1)$

ptr to const

- Save a byte vs GlobalStringPtr
- GEP: GetElementPointer
 - Inbounds: must not be out of bounds

r is for space

This space intentionally left blank.



t is for randomness

- [0, 13.4) times [a-z]
- Runtime: rand_inrange + rand_string

```
let len = rand_inrange(13.4 as i8 + 1);
rand_string(len);
```

```
let v len = LLVMBuildCall(
        b,
        rt rand inrange,
        [LLVMConstAdd(
             LLVMConstFPToUI(
                 LLVMConstReal(LLVMFloatTypeInContext(ctxt),
                               13.4),
                 ty i8),
             v 1i8)
        ].as ptr() as *mut ,
        1,
        b"\0".as ptr() as *const );
LLVMBuildCall(rt rand string, [v len]);
```

FP is slow, do it all as const for speed.

Codegen & optimization

Load runtime

```
static RT SOURCES: &'static [&'static [u8]] = &[
    include bytes!("../runtime/random.ll")
];
let mbuf = LLVMCreateMemoryBufferWithMemoryRange(
        code.as_ptr() as *const ,
        code.len() - 1 as libc::size t,
        b"\0".as ptr() as *const ,
        /* RequiresNullTerminator */ 1);
let module: LLVMModuleRef;
let err msg: *mut i8;
LLVMParseIRInContext(ctxt, mbuf, &mut module, &mut err msg);
```

Platform runtime

```
static RT_TARGET_SOURCE: phf::Map<
    &'static str,
    &'static [u8]
> = ...;
```

Use phf for excessively efficient lookup tables (built at compile-time)

```
let target = LLVMGetDefaultTargetTriple();
RT_TARGET_SOURCES.get(target);
```

Linking

```
let main_module: LLVMModuleRef;

for module in rt_modules {
    // Destroys module
    LLVMLinkModules2(main_module, module);
}
```

Easy as llvm-link

Codegen

Get a target, make a target machine

Emit code to memory, write to a file.

Not pictured: linker invocation (as subprocess)

Optimization

```
let fpm = LLVMCreateFunctionPassManagerForModule(llmod);
let mpm = LLVMCreatePassManager();

let pmb = LLVMPassManagerBuilderCreate();
LLVMPassManagerBuilderSetOptLevel(pmb, 2);
LLVMPassManagerBuilderUseInlinerWithThreshold(pmb, 225);
LLVMPassManagerBuilderPopulateModulePassManager(pmb, mpm);
LLVMPassManagerBuilderPopulateFunctionPassManager(pmb, fpm);
LLVMPassManagerBuilderDispose(pmb);
```

Pass manager wrangles optimizer passes

Including: DCE, GVN, constant propagation, LICM, loop unrolling, inlining...

Running passes

```
LLVMInitializeFunctionPassManager(fpm);

let mut func = LLVMGetFirstFunction(llmod);
while func != ptr::null_mut() {
    LLVMRunFunctionPassManager(fpm, func);
    func = LLVMGetNextFunction(func);
}
LLVMFinalizeFunctionPassManager(fpm);
```

Iterate over functions, optimizing each

```
LLVMRunPassManager(mpm, llmod);
```

LTO

- Link modules together, then optimize
 - ..retaining external symbols
- Internalize symbols first

```
let mut func = LLVMGetFirstFunction(llmod);
while func != ptr::null_mut() {
    let func_name = CStr::from_ptr(LLVMGetValueName(func));
    if func_name != "_start" {
        LLVMSetLinkage(func, LLVMLinkage::LLVMPrivateLinkage);
    }
    func = LLVMGetNextFunction(func);
}
```

References & advertising

LLVM IIvm.org Ilvm-sys Rust → L

Rust → LLVM C library bindings crates.io:llvm-sys

Reference merthc bitbucket.org/tari/merthc

Me



• @tari

Incidental foxes

きつねさんでもわかるLLV M

Ferris the Rustacean

Karen Rustad Tölva (public domain)

Ask me questions now.

