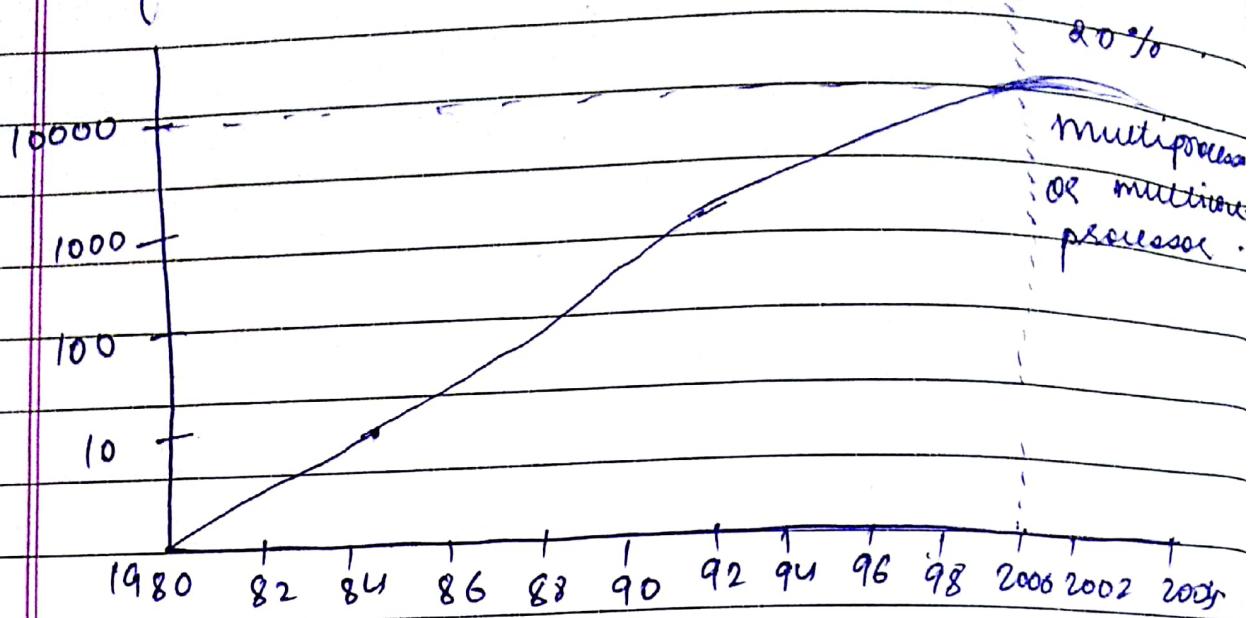


growth in uniprocessor performance



Trend in power in integrated circuit.

CMOS chip.

(i) dynamic power = $\frac{1}{2} C V^2 f$.

C = capacitive load

5 volt \rightarrow 1 volt

voltage scaling

clock scaling

(ii) static power = $I_{\text{leakage}} \times \text{Voltage}$
a 5% of total power.

Ex:- Some microprocessors today are designed to have adjustable voltage so that if 15% reduction in voltage may result in a 15% reduction in frequency. What would be impact on

dynamic power

$$P_{\text{old}} = \frac{1}{2} CV^2 R$$

$$P_{\text{new}} = \frac{1}{2} C (0.85V)^2 \times (0.85R)$$

$$\frac{P_{\text{old}}}{P_{\text{new}}} = \frac{1}{(0.85)^3} = 1.628$$

Multiprocessors or multicore processors

A summary of 3 Main stream computing classes (1980s onwards).

Feature	Desktop	Server	Embedded
price of system	\$500 - \$5000	\$5000 - \$5,000,000	\$10 - \$10,000
price of microprocessor module	\$50 - \$500	\$200 - \$10,000	\$0.01 - \$100
critical system design issues	price-performance graphics, performance	throughput, availability, power, scalability	cost, "apple" specific performance

time	decimal value	decimal value	binary value	decimal value
sec	$10^3 B$	$10^3 B$	2^{10}	1
millise	MB	10^6	2^{20}	10^{-3}
microse	QB	10^9	2^{30}	10^{-6}
nanose	TB	10^{12}	2^{40}	10^{-9}
picose	PB	10^{15}	2^{50}	10^{-12}
sesto	E B	10^{18}	2^{60}	10^{-15}
	Z B	10^{21}	2^{70}	
	Y B	10^{24}	2^{80}	

Measuring and comparing computer performance.

$$\text{performance} \propto \frac{1}{\text{execution time}}$$

$$\text{Execution time} = I_c \times CPI \times T$$

T = clock period, CPI \rightarrow cycles per instruction.
 I_c = Instruction count.

A program runs in 10s on comp A with 400 MHz. We want a comp B that could run the program in 6 sec. Substantial rise in clock speed is possible, 10 times or many

clock cycle op computer A. what would be the clock rate of computer B.

given

Comp. A

Comp B

Ex. Time of op 10 sec. 6 sec.

Clock Rate 400 MHz

CPI = 1.2 n

$$6 \text{ sec} = \frac{IC \times 1.2 n}{400 \times 10^6}$$

$$10 = \frac{IC \times n}{400 \times 10^6}$$

$$IC \times n = 400 \times 10^7$$

$$6 = \frac{400 \times 10^7 \times 1.2}{f}$$

$$f = 8 \times 10^8$$

$$= 800 \text{ MHz}$$

measuring and comparing performance

Execution Time $\propto IC \times CPI \times n \times T$.

MIPS (Million Instruction Per sec.)

$$\text{Ex. Time} = \frac{IC}{n \times 10^6}$$

1 sec = $\frac{IC}{Ex. Time \times 10^6}$ (MIPS).

$$= \frac{IC}{Ex. Time \times CPI \times 10^6} = \frac{f}{CPI \times 10^6} \text{ (MIPS).}$$

- Q1:- A 400 MHz Computer is running a program of 100,000 instruction. The program mix and their CPI is given in the below table. Calculate the execution time of the program and its MIPS rating.

<u>Instruction type</u>	<u>CPI</u>	<u>% of the prog.</u>
Integer arithmetic	1	60
data transfer	4	20
control instruction	2	20

Average CPI

$$= 1 \times 0.6 + 4 \times 0.8 + 2 \times 0.2$$

$$= 1.8$$

$$\frac{1 \times 60 \times 100,000 + 4 \times 0.2 \times 100,000 + 2 \times 0.2 \times 100,000}{100,000}$$

$$\text{Execution Time} = DC \times CPI \times IN = \frac{100,000 \times 1.8}{400 \times 10^6}$$

$$= 0.45 \text{ usec}$$

$$\text{MIPS Rating} = \frac{400 \times 10^6}{1.8 \times 10^6} = 222.222$$

Execution now when the task is performed in a 4 processor system due to synchronisation 8000 instructions are added per processor. Now find the avg. CPI.

$$1 \times 0.6 + 4 \times 0.2 + 8 \times 0.2 = 1.8.$$

Determine the corresponding MIPS rate.

$$\frac{400 \times 10^6}{1.8 \times 10^6} = 222.22.$$

$$\text{Execution time} = \frac{7000 \times 1.8}{400 \times 10^6} = 181.5 \text{ usec.}$$

$$\text{MIPS rating of comp} = 222.22 \times 4.$$

O:- Suppose we have made the following frequency of FP operations = 0.5%. Average CPI of FP operation = 4.0. Average CPI of other instruction = 1.33. Frequency of FPSQR = 0.1%, CPI of FPSQR = 0.0.

Assume that the two design alternation are to decrease the CPI of FPSQR to 0.0 or to decrease the avg. CPI of all FP operation to 0.5. Compare these two design alternatives using processor performance instruction equation

(all FP operation). (only FPSQR operation).

Original

design - 1

design - 2:

$$\begin{array}{lll}
 \text{(1) avg: } 0.25 \times 4 + & 0.25 \times 0.5 + & (0.0 \times 0.2) (20-2) \\
 0.75 \times 1.33 & 0.75 \times 1.33 & = 1.64 \\
 = 2 & = 1.624 &
 \end{array}$$

design - 1 is better all FP operation (common case fast)

Complete performance
Execution Time = $IC \times CP / \text{avg} \times T$

$$\text{MIPS Rating} = \frac{f}{CPI_{\text{avg}} \times 10^6}$$

a) we have two implementation of the same ISA.

Machine A has clock cycle time 1.2 and 1.5 of μ for same program

Machine B has clock cycle time 1.2 and 1.5 of μ for same program

which machine is faster for this program and by how much

Soln Ex. Time A = $IC \times Q \times 10^{-9}$

Ex. Time B = $IC \times 1.2 \times 10^{-9}$

Ex. Time 'B' = $IC \times 1.2 \times 10^{-9}$

Ex. Time 'A' = $IC \times Q \times 10^{-9}$

ratio of time = 1.2 times

Ex: Consider the machine with the following 3 instruction classes and CPI

Instruction	Instruction class	CPI for instruction
A	A	1
B	B	2
C	C	3

Suppose we measure the code of the same program from two compilers

Code from compiler - 1	A	B	C
" A + B + C "	5	1	1
" A * B + C "	10	1	1
" A + B * C "	0	1	1
" A * B * C "	0.5	1	1

Assume that the machine's clock rate is 500 MHz. which code sequence will execute faster acc. to execution time? Acc. to MIPS?

$$\text{Execution Time compiler - 1} = 5 \times 1 + 1 \times 2 + 1 \times 3 \\ = 10 \text{ billion cycles.}$$

$$\text{Ex. Time compiler - 2} = 10 \times 1 + 1 \times 2 + 1 \times 3 \\ = 15 \text{ billion cycles.}$$

$$\text{MIPS compiler - 1} = \frac{f}{\text{CPI avg} \times 10^6}$$

$$\text{CPI avg (compiler - 1)} = \frac{5 \times 1 + 1 \times 2 + 1 \times 3}{7} = 10 = 1.42$$

$$\text{MIPS compiler - 2} = \frac{f}{\text{CPI avg} \times 10^6} = \frac{15}{12} = 1.25$$

$$\text{CPI avg (compiler - 2)} = \frac{10 \times 1 + 1 \times 2 + 1 \times 3}{12} = 15 = 1.25$$

$$MIPS_{\text{Compiler-A}} = f \cdot 1.25 \times 10^6$$

Performance Summary

	Compiler-A	Compiler-B	Y-time
program-1(sec)	1	10	50%
program-2(sec)	1000	100	20%
total	1001	110	

$$\text{Time A.M} = 600.5$$

$$\text{Arithmetic mean} = \frac{1}{n} \sum_{i=1}^n \text{Ex-Time}_i$$

Weighted Arithmetic mean (WAM)

$$= \sum_{i=1}^n w_i \times \text{Ex-Time}_i$$

$$WAM_{\text{Comp-A}} = 0.8 \times 1 + 0.2 \times 1000 = 200.8$$

$$WAM_{\text{Comp-B}} = 0.8 \times 10 + 0.2 \times 100 = 28$$

System Performance Evaluation

Cooperative (SPEC) → 1980

	Benchmark program	Instruction count (10^9)	CPI
1.	Interpreted string processing	2252	0.60
2.	Block sorting compressor	2390	0.36
3.	GNU C compiler	794	0.376
4.	Combinational optimization	221	0.66
5.	Gro Kram (AI)	250	0.60

- 8. Search gene sequence
- 7. Video compression
- 8. XML parsing

2216	1.70
3793	1.5
1045	0.7

$$\text{SPEC Ratio}_i = \frac{T_{ref-i}}{T_{compA-i}}$$

$$\text{Geometric Mean} = \sqrt[n]{\prod_{i=1}^{20} \text{Execution Time Ratio}_i}$$

SPEC (System Performance Evaluation cooperative).

SPEC 1989

SPEC 2000

SPEC 2006

$$\text{SPEC Ratio} = \frac{T_{ref-i}}{T_{test-i}}$$

$$\text{SPEC Rating} = \sqrt[n]{\prod_{i=1}^n (\text{SPEC Ratio})_i}$$

$$= \sqrt[n]{(\text{SPEC-Ratio})_1 \times (\text{SPEC-Ratio})_2 \times \dots \times (\text{SPEC-Ratio})_n}$$

Ex:- There are two machines A and B and a sequence machine. These are two test T₁ and T₂ and we obtain the score for the machine.

Machine	Test T ₁	Test T ₂
Machine - A	10 sec	100 sec
" - B	1 sec	500
" - C	1 sec	100

what is the SPEC Rating for Machine A & B?

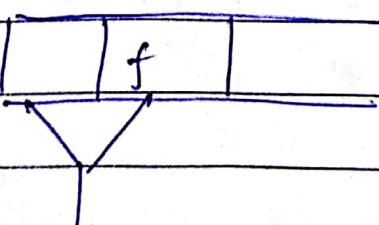
$$\text{SPEC Rating comp-A} = 2 \sqrt{\frac{1 \times 100}{10 \times 100}} = 0.311$$

$$\text{SPEC Rating comp-B} = 2 \sqrt{\frac{1 \times 100}{1 \times 500}} = 0.4$$

Speed-up.

Chipprocessor \rightarrow Multi processor.

Andahl's law



$(1-f)$ = function of code that can't be parallelized.
 $N = \text{No. of processes}$

T_{old} = Time taken to run the code
on a uniprocessor

$$T_{new} = \frac{(1-f) T_{old} + f T_{old}}{N}$$

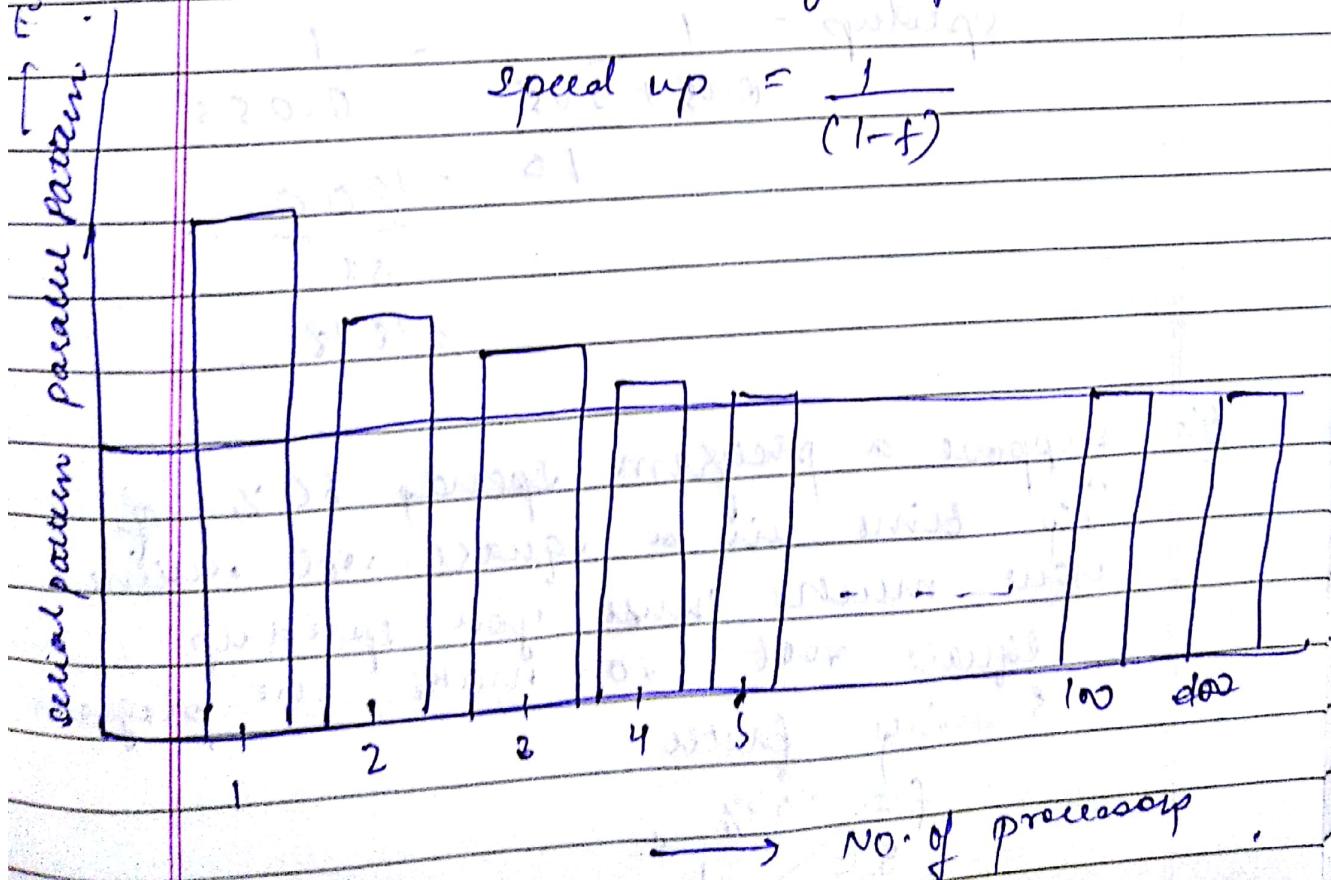
$$\text{Speedup} = \frac{T_{old}}{T_{new}}$$

$$= \frac{T_{old}}{(1-f) T_{old} + f \cdot T_{old}}$$

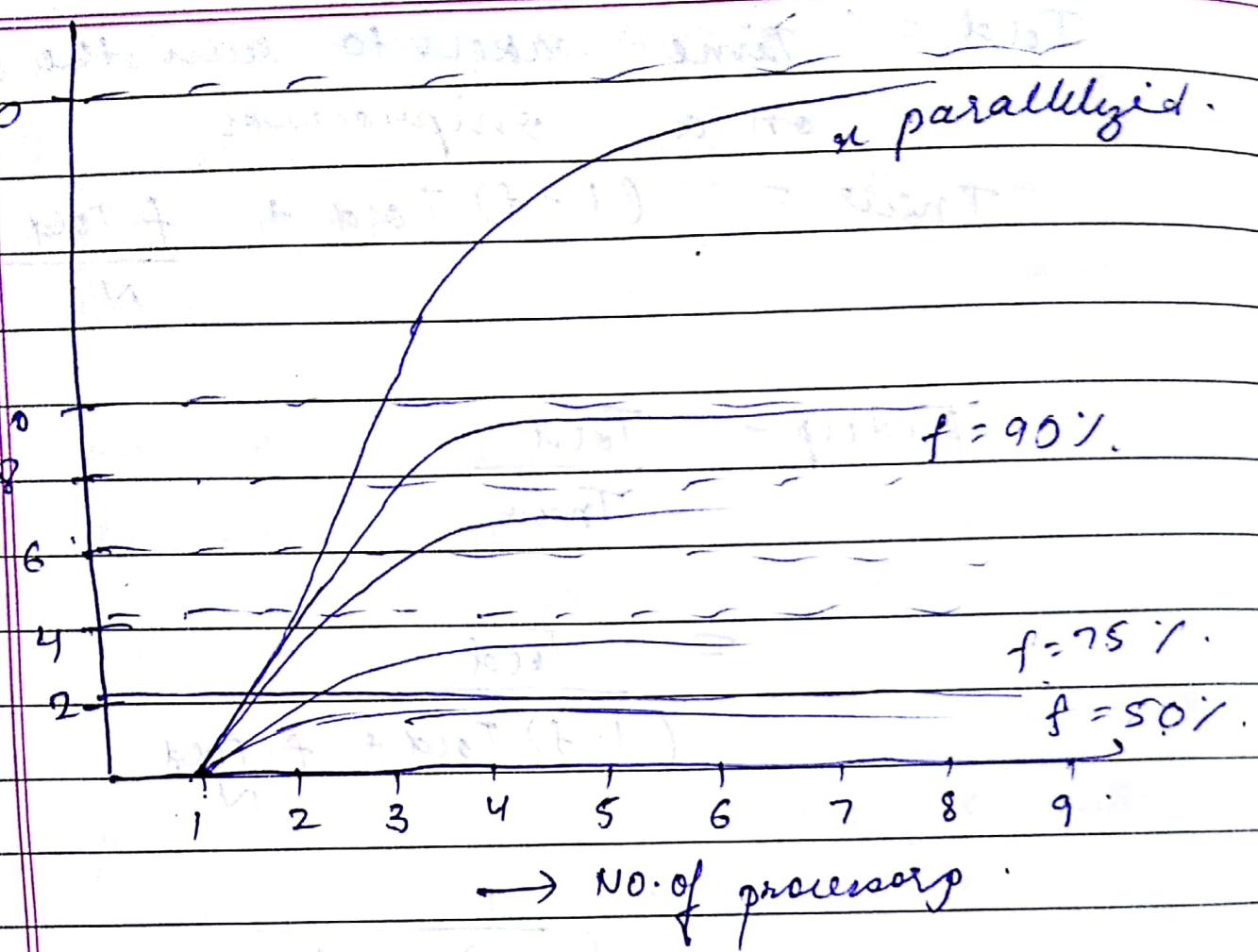
$$\text{Speedup} = \frac{1}{(1-f) + \frac{f}{N}}$$

When N is infinity then :

$$\text{Speedup} = \frac{1}{(1-f)}$$



20



speedup =

$$f = 95\% : N = 10$$

$$\text{speedup} = \frac{1}{0.05 + \frac{0.05}{N}} = \frac{1}{0.055}$$

$$10 = \frac{1000}{55}$$

$$= 18.18$$

Q:- Suppose a program spends 80% of its time in a square-root routine. How much must you speed up the square root to make the program 5 times faster?

Charlie

Date

Page No.

$$5 = \frac{1}{(-0.8) + \frac{0.8}{N}}$$

$$1 = \cancel{0.8} + \frac{4}{N} \quad (\text{Not possible})$$

$$\cancel{+} \cancel{= 4}$$