MIT 2.853/2.854

Introduction to Manufacturing Systems

Markov Processes and Queues

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- X(t) can be discrete or continuous, scalar or vector.

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- Or, let x(s), s ≤ t, be the history of the values of X before time t and let A be a possible value of X.
 Then

$$P\{X(t+\delta t) = A|X(s) = x(s), s \le t\} = P\{X(t+\delta t) = A|X(t) = x(t)\}$$

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- This is ONLY the definition of a class of mathematical models. It is <u>NOT</u> a statement about reality!! That is, not everything is a Markov process.

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N(t) is a Markov process. Why?

Discrete state, discrete time States and transitions

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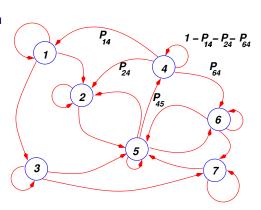
States and transitions

- States can be numbered 0, 1, 2, 3, ... (or with multiple indices if that is more convenient).
- Time can be numbered 0, 1, 2, 3, ... (or 0, Δ , 2Δ , 3Δ , ... if more convenient).
- The probability of a transition from j to i in one time unit is often written P_{ij} , where

$$P_{ij} = P\{X(t+1) = i | X(t) = j\}$$

Transition graph

Transition graph



 P_{ij} is a probability. Note that $P_{ii}=1-\sum_{m,m\neq i}P_{mi}$. This is the self-loop probability.

Transition graph



Example: H(t) is the number of Hs after t coin flips.

Transition graph

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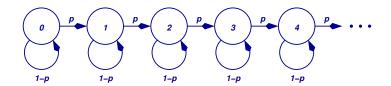
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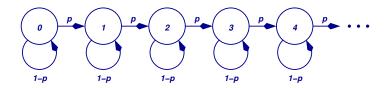
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Transition graph

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This is a system with an infinite state space.

Transition graph

Example: Coin flip bets on Slide 5.

Transition graph

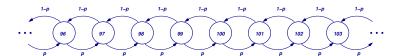
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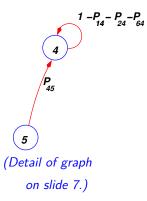
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- $\{X(t) = i\}$ is the event that random quantity X(t) has value i.
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- Define $\pi_i(t) = P\{X(t) = i\}.$
- Normalization equation: $\sum_i \pi_i(t) = 1$.

Transition equations

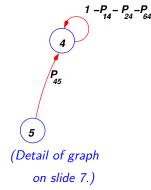
Transition equations: application of the law of total probability.



$$\pi_4(t+1) = \pi_5(t)P_{45} + \pi_4(t)(1 - P_{14} - P_{24} - P_{64})$$

Transition equations

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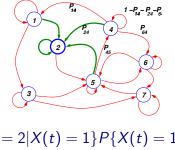
$$\pi_4(t+1) = \pi_5(t)P_{45} + \pi_4(t)(1 - P_{14} - P_{24} - P_{64})$$

(Remember that

$$P_{45} = P\{X(t+1) = 4 | X(t) = 5\},$$

 $P_{44} = P\{X(t+1) = 4 | X(t) = 4\}$
 $= 1 - P_{14} - P_{24} - P_{64}$

Transition equations



$$P\{X(t+1) = 2\}$$

$$= P\{X(t+1) = 2 | X(t) = 1\} P\{X(t) = 1\}$$

$$+ P\{X(t+1) = 2 | X(t) = 2\} P\{X(t) = 2\}$$

$$+ P\{X(t+1) = 2 | X(t) = 4\} P\{X(t) = 4\}$$

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Transition equations

• Define $P_{ij} = P\{X(t+1) = i | X(t) = j\}$

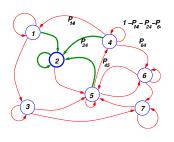
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Transition equations



Therefore, since

$$P_{ij}=P\{X(t+1)=i|X(t)=j\}$$
 and $\pi_i(t)=P\{X(t)=i\},$

we can write

$$\pi_2(t+1) = P_{21}\pi_1(t) + P_{22}\pi_2(t) + P_{24}\pi_4(t) + P_{25}\pi_5(t).$$

Note that $P_{22} = 1 - P_{52}$.

Transition equations — Matrix-Vector Form

For an *n*-state system,

$$\pi(t) = \begin{bmatrix} \pi_1(t) \\ \pi_2(t) \\ \dots \\ \pi_n(t) \end{bmatrix}, \quad P = \begin{bmatrix} P_{11} & P_{12} & \dots & P_{1n} \\ P_{21} & P_{22} & \dots & P_{2n} \\ \dots & \dots & \dots & \dots \\ P_{n1} & P_{n2} & \dots & P_{nn} \end{bmatrix}, \quad \nu = \begin{bmatrix} 1 \\ 1 \\ \dots \\ 1 \end{bmatrix}$$

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• Transition equations: $\pi(t+1) = P\pi(t)$

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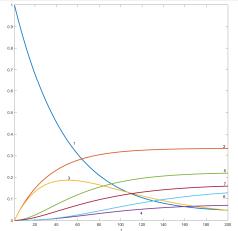
- Transition equations: $\pi(t+1) = P\pi(t)$
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- · Other facts:
 - * $\nu^T P = \nu^T$ (Each column of P sums to 1.)

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 - $\star \pi(t) = P^t \pi(0)$



State probabilities vs. t for system in Slide 7

Steady state

• Steady state: $\pi_i = \lim_{t \to \infty} \pi_i(t) = \lim_{t \to \infty} P^t \pi(0)$,

Steady state

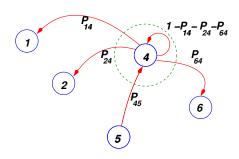
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Balance equations

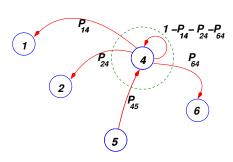


Balance equation:

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in steady state <u>only</u>.

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Intuitive meaning: The average number of transitions *into* the circle per unit time equals the average number of transitions *out of* the circle per unit time.

Steady state

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How to calculate the steady-state probability distribution π

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- Delete one transition equation and replace it with the normalization equation.
- Solve the system of *N* linear equations in *N* unknowns.

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- If a system has an infinite number of states and it has a steady state probability distribution, there are two possibilties for finding it:
 - It might be possible to solve the equations analytically. We will see an example of that.
 - Truncate the system. That is, solve a system with a large but finite subset of the states. If you understand the system, you can guess which are the highest probability states. Keep those. This provides an approximate solution.

Geometric distribution

Consider a two-state system. The system can go from 1 to 0, but not from 0 to 1.

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Let p be the conditional probability that the system is in state 0 at time t+1, given that it is in state 1 at time t. Then

$$p = P\left[\alpha(t+1) = 0 \middle| \alpha(t) = 1\right].$$

Geometric distribution — Transition equations

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Geometric distribution — Transition equations

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$$\pi(0, t+1) = P\left[\alpha(t+1) = 0 \middle| \alpha(t) = 1\right] P[\alpha(t) = 1]$$
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we have

$$\pi(0, t+1) = p\pi(1, t) + \pi(0, t),$$

 $\pi(1, t+1) = (1-p)\pi(1, t),$

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we have

$$\pi(0, t+1) = p\pi(1, t) + \pi(0, t),$$

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and the normalization equation

$$\pi(1, t) + \pi(0, t) = 1.$$

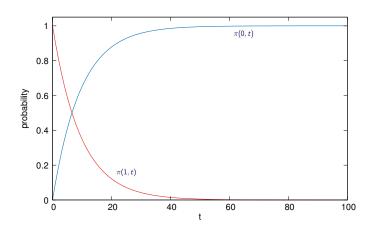
Geometric distribution — transient probability distribution

Assume that $\pi(1,0) = 1$. Then the solution is

$$\pi(0,t) = 1 - (1-p)^t,$$

 $\pi(1,t) = (1-p)^t.$

Geometric distribution — transient probability distribution



Geometric distribution

We have shown that the probability that the state goes from 1 to 0 at time t is

$$P(t) = (1-p)^{t-1}p$$

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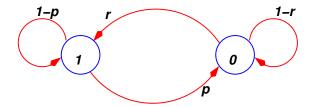
$$\bar{t} = \sum_{t=1}^{\infty} t P(t) = \sum_{t=1}^{\infty} t (1-p)^{t-1} p$$

It is not hard to show that

$$\overline{t} = \frac{1}{p}$$

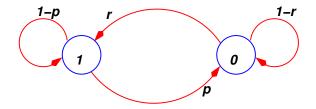
Unreliable machine

1=up; 0=down.



Unreliable machine

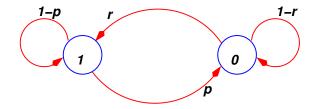
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Mean up time = Mean time to fail = MTTF = 1/p

Unreliable machine

1=up; 0=down.



Mean up time = Mean time to fail = MTTF= 1/pMean down time = Mean time to repair = MTTR= 1/r

Unreliable machine — transient probability distribution

The probability distribution satisfies

$$\pi(0, t+1) = \pi(0, t)(1-r) + \pi(1, t)p,$$

$$\pi(1, t+1) = \pi(0, t)r + \pi(1, t)(1-p).$$

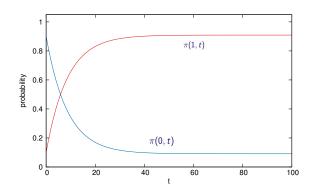
Unreliable machine — transient probability distribution

It is not hard to show that

$$\pi(0,t) = \pi(0,0)(1-p-r)^{t} + \frac{p}{r+p} [1-(1-p-r)^{t}],$$

$$\pi(1,t) = \pi(1,0)(1-p-r)^{t} + \frac{r}{r+p} [1-(1-p-r)^{t}].$$

Unreliable machine — transient probability distribution



Unreliable machine — steady-state probability distribution

As
$$t \to \infty$$
,

$$\pi(0,t) \rightarrow \frac{p}{r+p},$$

$$\pi(1,t) \rightarrow \frac{r}{r+p}$$

Unreliable machine — steady-state probability distribution

As $t \to \infty$,

$$\pi(0,t) \rightarrow \frac{p}{r+p},$$
 $\pi(1,t) \rightarrow \frac{r}{r+p}$

which is the solution of

$$\pi(0) = \pi(0)(1-r) + \pi(1)p,$$

 $\pi(1) = \pi(0)r + \pi(1)(1-p).$

Unreliable machine — efficiency

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Note that we can also write

$$e = \frac{\text{MTTF}}{\text{MTTF} + \text{MTTR}}$$

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States and transitions

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- Time is a real number, defined on $(-\infty, \infty)$ or a smaller interval.
- The probability of a transition from j to i during $[t, t + \delta t]$ is approximately $\lambda_{ij}\delta t$, where δt is small, and

$$\lambda_{ij}\delta t \approx P\{X(t+\delta t)=i|X(t)=j\}$$
 for $i\neq j$

32

More precisely,

$$\lambda_{ij}\delta t = P\{X(t+\delta t) = i|X(t) = j\} + o(\delta t)$$

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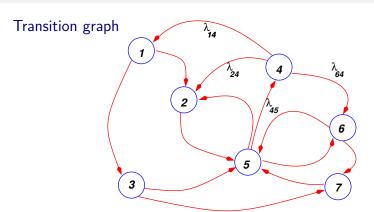
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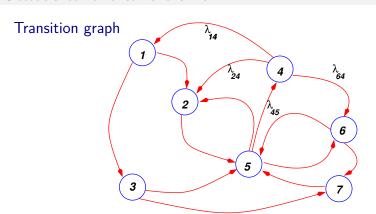
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This implies that for small δt , $o(\delta t) \ll \delta t$.



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Compare with the discrete-time graph.

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States and transitions

Or,

$$\pi_5(t+\delta t)\approx$$

$$\pi_5(t) - (\lambda_{25} + \lambda_{45} + \lambda_{65})\pi_5(t)\delta t$$

$$+(\lambda_{52}\pi_2(t)+\lambda_{53}\pi_3(t)+\lambda_{56}\pi_6(t)+\lambda_{57}\pi_7(t))\delta t$$

States and transitions

Or,

$$\lim_{\delta t o 0} rac{\pi_5(t+\delta t) - \pi_5(t)}{\delta t} =$$

$$\frac{d\pi_5}{dt}(t) = -(\lambda_{25} + \lambda_{45} + \lambda_{65})\pi_5(t)$$

$$+\lambda_{52}\pi_2(t) + \lambda_{53}\pi_3(t) + \lambda_{56}\pi_6(t) + \lambda_{57}\pi_7(t)$$

States and transitions

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- * * * Often confusing!!!

Discrete state, continuous time Transition equations — Matrix-Vector Form

• Define $\pi(t)$, ν as before *.

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Define
$$\Lambda = \begin{bmatrix} \lambda_{11} & \lambda_{12} & \dots & \lambda_{1n} \\ \lambda_{21} & \lambda_{22} & \dots & \lambda_{2n} \\ & & \dots & \\ \lambda_{n1} & \lambda_{n2} & \dots & \lambda_{nn} \end{bmatrix}$$

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 - $\star \ \nu^T P = 0$ (Each column of P sums to 0.)
 - $\star \pi(t) = e^{\Lambda t} \pi(0)$

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Discrete state, continuous time Sources of confusion in continuous time models

Never Draw self-loops in continuous time Markov process graphs.

Discrete state, continuous time Sources of confusion in continuous time models

- *Never* Draw self-loops in continuous time Markov process graphs.
 - Never write $1 \lambda_{14} \lambda_{24} \lambda_{64}$. Write $1 (\lambda_{14} + \lambda_{24} + \lambda_{64})\delta t$, or
 - $\star -(\lambda_{14} + \lambda_{24} + \lambda_{64})$

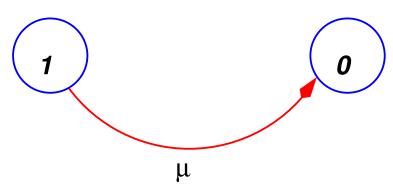
Sources of confusion in continuous time models

- *Never* Draw self-loops in continuous time Markov process graphs.
- Never write $1 \lambda_{14} \lambda_{24} \lambda_{64}$. Write $\begin{array}{c} \star & 1 (\lambda_{14} + \lambda_{24} + \lambda_{64})\delta t, \text{ or} \\ \star & -(\lambda_{14} + \lambda_{24} + \lambda_{64}) \end{array}$
- $\lambda_{ii} = -\sum_{j \neq i} \lambda_{ji}$ is **NOT** a rate and **NOT** a probability. It is **ONLY** a convenient notation.

Exponential random variable T: the time to move from state 1 to state 0.

Exponential distribution

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$$\pi(0, t + \delta t) =$$

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Transition equations
$$\left\{ \begin{array}{ll} \displaystyle \frac{d\pi(0,t)}{dt} & = & \mu\pi(1,t) \\ \\ \displaystyle \frac{d\pi(1,t)}{dt} & = -\mu\pi(1,t) \end{array} \right.$$

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The time of the transition from 1 to 0 is said to be exponentially distributed with rate μ .

Exponential distribution

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The expected transition time is $\frac{1}{\mu} = \int_0^\infty t e^{-\mu t}$.

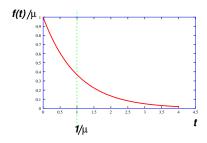
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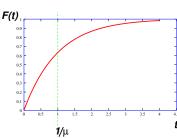
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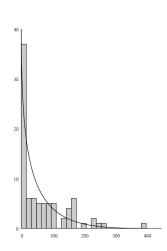




Exponential

Density function

Exponential density function and a small number of samples.



Exponential distribution: some properties

• Memorylessness:

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Exponential distribution: some properties

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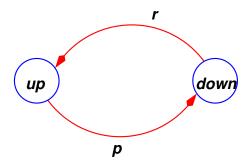
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- $T = \min(T_1, ..., T_n)$, then
- T is an exponentially distributed random variable with parameter $\mu = \mu_1 + ... + \mu_n$.
- Consequently, the time that the system stays in any state is exponentially distributed.

Continuous time unreliable machine.

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 $P(\{\text{the machine is up at time } t + \delta t\}) =$

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Unreliable machine

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and similarly for $P(\{\text{the machine is down at time } t + \delta t\})$.

Probability distribution notation and dynamics:

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\pi(1,t) = the probability that the machine is up at time t.
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Unreliable machine

Probability distribution notation and dynamics:

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$$P(ext{the machine is up at time } t + \delta t | ext{ the machine was up at time } t) = 1 - p \delta t$$

Unreliable machine

Probability distribution notation and dynamics:

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$$= r\delta t$$

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$$\pi(1, t + \delta t) = (1 - p\delta t)\pi(1, t) + r\delta t\pi(0, t) + o(\delta t)$$

Unreliable machine

Therefore

$$\pi(1, t + \delta t) = (1 - p\delta t)\pi(1, t) + r\delta t\pi(0, t) + o(\delta t)$$

Similarly,

$$\pi(0, t + \delta t) = p\delta t\pi(1, t) + (1 - r\delta t)\pi(0, t) + o(\delta t)$$

Unreliable machine

$$\pi(1, t + \delta t) - \pi(1, t) = -p\delta t\pi(1, t) + r\delta t\pi(0, t) + o(\delta t)$$

Unreliable machine

or,

$$\pi(1, t + \delta t) - \pi(1, t) = -p\delta t\pi(1, t) + r\delta t\pi(0, t) + o(\delta t)$$

$$\frac{\pi(1,t+\delta t)-\pi(1,t)}{\delta t}=-p\pi(1,t)+r\pi(0,t)+\frac{o(\delta t)}{\delta t}$$

$$\frac{d\pi(1,t)}{dt} = \pi(0,t)r - \pi(1,t)p$$

$$\frac{d\pi(1,t)}{dt} = \pi(0,t)r - \pi(1,t)p$$

$$\frac{d\pi(0,t)}{dt} = -\pi(0,t)r + \pi(1,t)p$$

Unreliable machine

Solution

$$\pi(0,t) = \frac{p}{r+p} + \left[\pi(0,0) - \frac{p}{r+p}\right] e^{-(r+p)t}$$

Unreliable machine

Solution

$$\pi(0,t) = \frac{p}{r+p} + \left[\pi(0,0) - \frac{p}{r+p}\right] e^{-(r+p)t}$$

$$\pi(1,t) = 1 - \pi(0,t).$$

Unreliable machine

Solution

$$\pi(0,t) = \frac{p}{r+p} + \left[\pi(0,0) - \frac{p}{r+p}\right] e^{-(r+p)t}$$

$$\pi(1,t) = 1 - \pi(0,t).$$

As $t \to \infty$.

$$\pi(0) \rightarrow \frac{p}{r+p},$$
 $\pi(1) \rightarrow \frac{r}{r+p}$

Unreliable machine

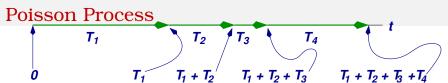
Note that MTTF=1/p; MTTR=1/r. Units are natural time units, not operation times.

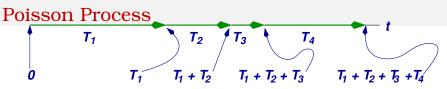
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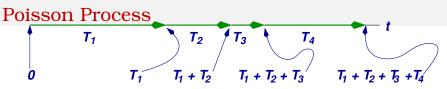
If the machine makes μ parts per time unit on the average when it is operational, the steady-state average production rate is

$$\mu\pi(1) = \mu \frac{r}{r+p} = \mu \frac{\mathsf{MTTF}}{\mathsf{MTTF} + \mathsf{MTTR}} = \mu \mathsf{e}$$

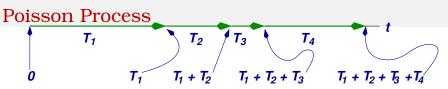




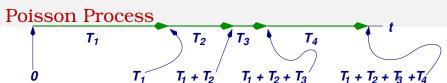
• Let T_i , i = 1, ... be a set of independent exponentially distributed random variables with parameter λ . Each random variable may represent the time between occurrences of a repeating event.

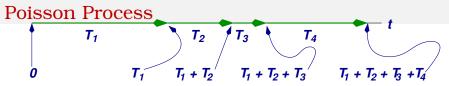


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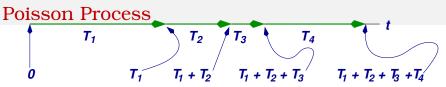


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- Then $\sum_{i=1}^{n} T_i$ is the time required for n such events.

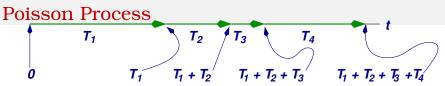




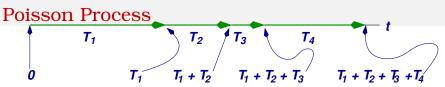
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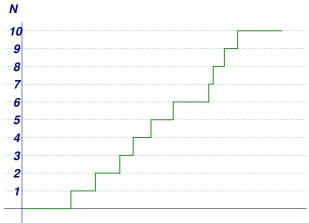
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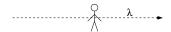
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M/M/1 Queue

Number of events N(t) during [0, t]

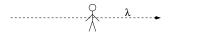


Queueing theory *M/M/1* Queue





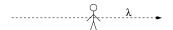
M/M/1 Queue

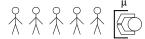




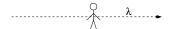
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Queueing theory *M/M/1* Queue





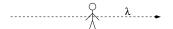
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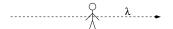
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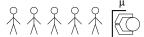
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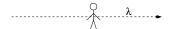


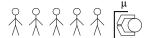
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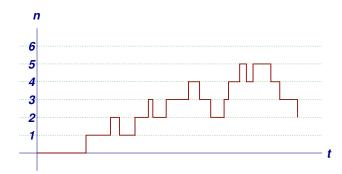




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- Define the *utilization* $\rho = \lambda/\mu$.

Queueing theory *M/M/*1 Queue

Number of customers in the system as a function of time for a M/M/1 queue.



Queueing theory D/D/1 Queue

Number of customers in the system as a function of time for a D/D/1 queue.



Queueing theory Sample path

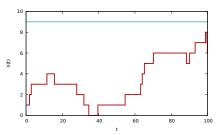
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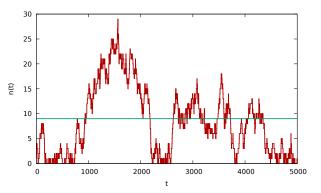
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Queue behavior over a short time interval — initial transient

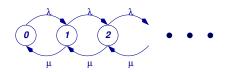
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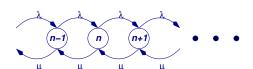


Queue behavior over a long time interval

Queueing theory *M/M/*1 Queue

State space





*M/M/*1 Queue



Let $\pi(n, t)$ be the probability that there are n parts in the system at time t. Then,

M/M/1 Queue



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$$\pi(n, t + \delta t) = \pi(n - 1, t)\lambda\delta t + \pi(n + 1, t)\mu\delta t + \pi(n, t)(1 - (\lambda\delta t + \mu\delta t)) + o(\delta t)$$

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and

$$\pi(0, t + \delta t) = \pi(1, t)\mu\delta t + \pi(0, t)(1 - \lambda\delta t) + o(\delta t).$$

M/M/1 Queue

Or,

$$\frac{d\pi(n,t)}{dt} = \pi(n-1,t)\lambda + \pi(n+1,t)\mu - \pi(n,t)(\lambda+\mu), \ n>0$$

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Why "if"?

M/M/1 Queue - Steady State

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$$\pi(n) = (1 - \rho)\rho^n, n \ge 0$$

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The average number of parts in the system is

$$\bar{n} = \sum_{n=0}^{\infty} n\pi(n) = \frac{\rho}{1-\rho} = \frac{\lambda}{\mu-\lambda}.$$

Queueing theory Little's Law

• True for most systems of practical interest (not just M/M/1).

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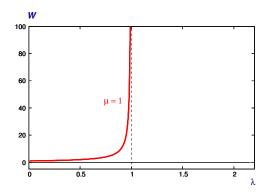
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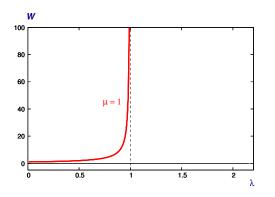
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In the M/M/1 queue, $L = \bar{n}$ and

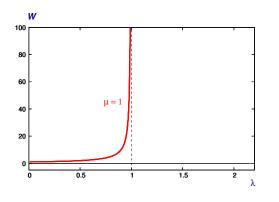
$$W = \frac{1}{\mu - \lambda}$$
.



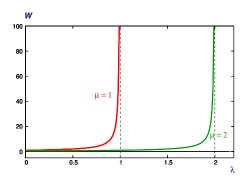
• μ is the *capacity* of the system.



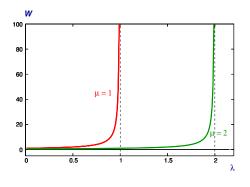
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- If $\lambda < \mu$, system is stable and waiting time remains bounded.
- If $\lambda > \mu$, waiting time grows over time.



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- To decrease delay for a given λ , increase μ .

Things get more complicated when:

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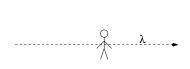
- There are multiple servers.
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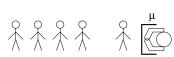
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Closed formulas and approximations exist for some, but not all, cases.

Queueing theory M/M/s Queue





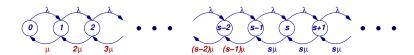




$$\bigwedge^{\mu}$$

s-Server Queue, s=3

M/M/s Queue



M/M/s Queue



• The departure rate when there are k > s customers in the system is $s\mu$ since all s servers are always busy.

M/M/s Queue



- The departure rate when there are k > s customers in the system is $s\mu$ since all s servers are always busy.
- The departure rate when there are $k \le s$ customers in the system is $k\mu$ since only k of the servers are busy.

Queueing theory *M/M/s* Queue

$$P(k) = \begin{cases} \pi(0) \frac{s^k \rho^k}{k!}, & k \leq s \end{cases}$$

Queueing theory *M/M/s* Queue

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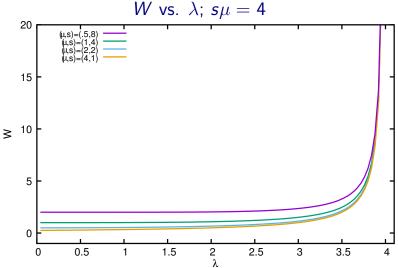
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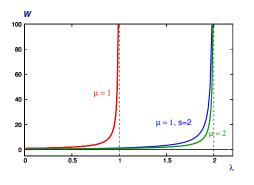
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where

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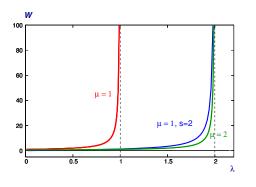
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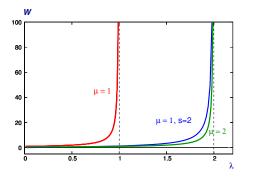
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To increase capacity or reduce delay,

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- Open network: where customers enter and leave the system. λ is known and we must find L and W.
- Closed network: where the population of the system is constant. L is known and we must find λ and W.

Examples of open networks

internet traffic

- internet traffic
- emergency room (arrive, triage, waiting room, treatment, tests, exit or hospital admission)

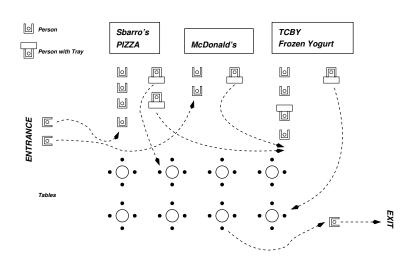
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Networks of Queues

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- airport (*arrive*, ticket counter, security, passport control, gate, board plane)
- factory with no centralized material flow control after material enters



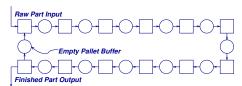
Queueing theory Networks of Queues

Examples of closed networks

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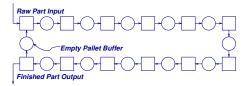
• factory with limited fixtures or pallets



Networks of Queues

Examples of closed networks

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 factory with material controlled by keeping the number of items constant (CONWIP)

Queueing networks are often modeled as Jackson networks.

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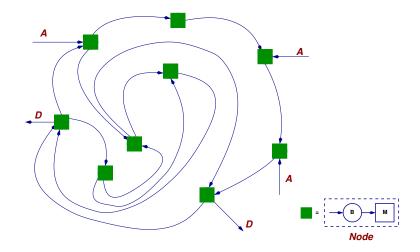
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- Valid (or good approximation) for a large class of systems ...

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* This assumption leads to bad results for systems with bottlenecks at locations other than the first station.



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- We will focus on the special case in which each node has a single server with exponential processing time. The service rate of node i is μ_i .

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- to determine the average waiting time at each node and the average time a part spends in the system.

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- $p_{ji}\lambda_j$ is the portion of the flow arriving at node j that goes to node i.

Open Jackson Networks

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• Solving for λ ,

$$\lambda = (I - \mathsf{P}^{\mathsf{T}})^{-1} \alpha$$

Probability distribution:

• If $\lambda_i < \mu_i$ for each *i*, define $\rho_i = \lambda_i/\mu_i$ and

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- If $\lambda_i \geq \mu_i$ for some i, the demand is not feasible. The system cannot handle the demand placed on it.

Solution:

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Queueing theory Open Jackson Networks

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 - * It will not work so well where blocking occurs often.

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Closed Jackson Networks

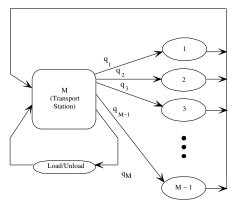
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 - ★ This means that a different solution approach is needed to analyze the system. It is used in the example that follows.

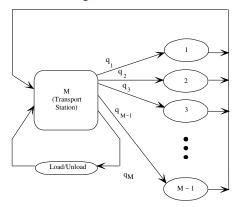
Closed Jackson Network model of an FMS

Solberg's "CANQ" model.



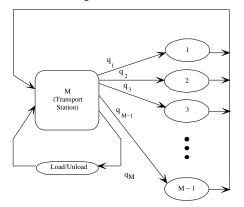
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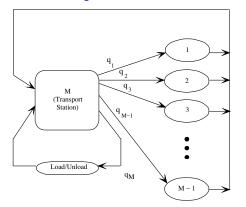
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$$p_{iM} = 1$$
 if $i \neq M$

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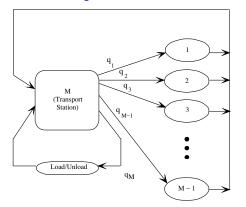


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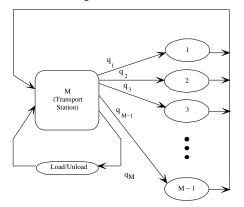
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Solberg's "CANQ" model.



$$p_{iM}=1$$
 if $i \neq M$
 $p_{Mj}=q_j$ if $j \neq M$
 $p_{ij}=0$ otherwise
Service rate at Station i is μ_i .

- Input data: $M, N, q_i, \mu_i, s_i \ (i = 1, ..., M)$
 - $\star M =$ number of stations, including transportation system
 - $\star N = \text{number of pallets}$
 - \star q_i = fraction of parts going from the transportation system to Station *i*
 - $\star \mu_i$ = processing rate of machines at Station i
 - \star s_i = number of machines at Station j

- Input data: M, N, q_j, μ_j, s_j (j = 1, ..., M)
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 - \star N = number of pallets
 - \star $q_j =$ fraction of parts going from the transportation system to Station j
 - \star μ_j = processing rate of machines at Station j
 - * s_i = number of machines at Station j
- Output data: $P, W, \rho_j \ (j = 1, ..., M)$
 - \star P = production rate
 - $\star W = \text{average time a part spends in the system}$
 - $\star \ \rho_i = \text{utilization per machine of Station } j$

Closed Jackson Network model of an FMS

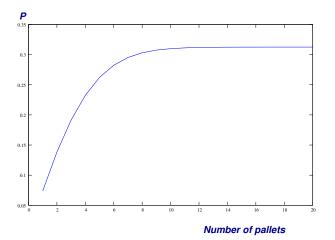
For the following graphs,

• Base input data: $M, N, q_j, \mu_j, s_j \ \ (j = 1, ..., M)$

*
$$M = 5$$

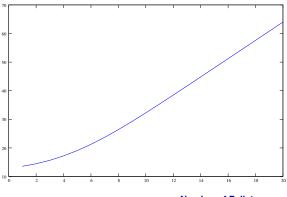
* $N = 10$
* $q_j = .1, .2, .2, .25, .25$
* $1/\mu_j = 3., 4., 3.44, 1.41, 5.$
* $s_j = 2, 1, 2, 1, 15$

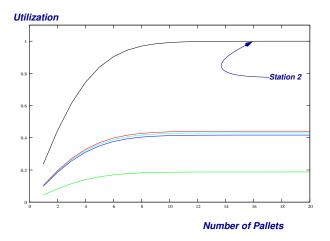
We see the effect of one of the variables on the performance measures in the following graphs.

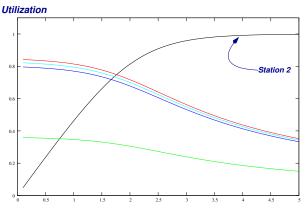


Closed Jackson Network model of an FMS

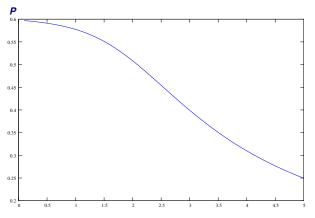
Average time in system







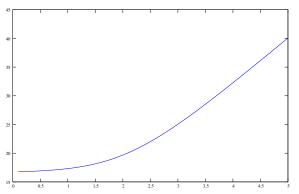
Station 2 operation time



Station 2 operation time

Closed Jackson Network model of an FMS

Average time in system



Station 2 operation time