Lecture no 3 1. Storage Structure
After extraction of operational date, information is
Stored in databases.
RDB & MDB are Complementary and do not have to
exclude arch other
· In steging ones -> ROBMS can be used, but it should
To steging ones > ROBMS can be used, but it should (temporarize stores not be avaleble for user queries date) Per Performance Yeasons.
date) date)
. In presentation aver > normalized data bases (like ROBMS)
are excluded from presentation cives.
This area used multidimensional detabrie
- Multidimensional DB → . MDB are optimized for OH &
· Performance + · faster than OLAP application
advantage ROB . Crested by using input from staging over.
e Performance e designed for efficient & convenient Storage & vetvieval of large volumes of data. For gueries that collect dimensions
for queries that colled dimensions.
Jenerates Crost-to . Complexity grows quickly with no of dimension
Virws of date & position. P.g., Horm (3 dimension with lavelue) ROBMS (103 2 1000)
Lecture NOB MOB
· Obtaining same views as of . Ease of data presentation maintenance
MDB, KDB requires compret query, E. per tormance
· Complex queries like top-K . Date views are natural extent of MDB
can be tricky but are possible . Early to maintain ble date is organized
with JUL extentions. Performance of MDB can be
need for indexes. Regular toning tuning tuning.
Simplified or could ble and by decimal for modern during
date is stered in normalized of queries.

Feature)	MDB RDB
Data Structure	Multidimensional Normalized tables (vous and
	and the state of t
Optimization	OLAP, goeries OLTP.
focos	goeries)
Performance (Fast for Structured Requires tuning for complex queries queries
	querier-
Ad-HOC Guery	Dimited Plaxibilty High flexibilty
1 1 1 1 1 1 1 1	The second secon
Maintenance	Minimal High
and the first	any the minute make the first with the same to be a substant of
Drawactioned data handling	not well suited well suited for transaction.
	Por transactions data
O Cail	and a topic of the design and public 1981, and 1981.
	MDB's are inappropriate -> related to each other vesults
	In sparse representation (only)
	MDB's are appropriate - when there are highly
	intervelated detatures
	Brample + + Budgering Product Profit
Tiev	1 Ctrusture framework for overprizing roftware component
1164	Avchitecture (in pystem into different layers or tien)
DW o	wehitecture > generic two tier architecture
	(Nee- fier architection
	Web-based architecture.
	A MARINE TO THE RESERVE TO THE PARTY OF THE
_ 4_5	The second control of the second control of

Layered Architecture - Data analysis Comes in two flowers.

Thin client > Analytics executed on server

(HTTP/ ELOP) • Client jut dipploy

of the well for Internet DW access > Fat client > · Server just deliver dota (ODBK/NFS/JBDK) . Hamplythes executed on client - data is not completely current in DH . Periodic extraction. Generic Two-tier architecture emini watchevies Independent data mart - separate 5T1 for each independent dem mart e High Olate mart accen complexity

Single FIL for DN Dependent data mort -> Data marts localed from DW (Operational data store) · Simple dete access as compared to Independent. - ETL is near yeal-time Logical data mand (Active workhouse) . Det marts ove not separate databases but logical views of DW. Conexic Three-tier Architecture > · Derived data > selected, formated or aggregated for OPS support. · Reconciled data > Detailed, current date intended to