

Variables x_1, x_2, \dots, x_n values domain say R, G, B satisfy
 but depends upon constraints.

State is defined by variables X , with values from domain D .

$$C_i = (\text{scope}, \text{rel})$$

Constraint Satisfaction Problem (CSP)

Goal is to find

CSP consists of 3 components (V, D, C)

Combination of values for variables that satisfies all constraints

V is set of variables $\{V_1, V_2, \dots, V_n\}$

D is set of domain $\{D_1, D_2, D_3, \dots, D_n\}$ one for

(natural, whole no) each variable.

C is set of constraints that specify allowable combination of values, $\{C_1, C_2, C_3, \dots\}$

Where scope is set of variables that participate in constraints.

Rel is relation that defines values that variables can take.

CSP differs from programming, bcz they don't tell how to solve a problem, they focus on what needs to be done to reach solution.

Example:

V_1

V_2

$$C_1 = (V_1, V_2), (V_1 \neq V_2)$$

A

B

(1, 2)

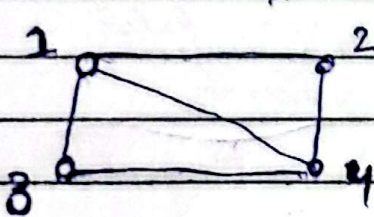
(2, 4)

$$C_1 = (V_1, V_2), (A, B)$$

$$C_1 = (V_1, V_2), (1, 2), (1, 4), (2, 4)$$

Example: intelligent.

We use backtracking method to solve CSP.



Intelligent backtracking
 as Rize by 9
 key ham exact
 aur jagga pr jae
 gie jae conflict
 hogi.

$$V = \{1, 2, 3, 4\}$$

$$D = \{\text{Red, Blue, Green}\}$$

$$C = \{1 \neq 2, 1 \neq 3, 1 \neq 4, 2 \neq 4, 3 \neq 4\}$$

→

	1	2	3	4
Initial domain	R, G, B	R, G, B	R, G, B	R, G, B
1 = R	R	GB	GB	GB
2 = G	R	G	GB	B
3 = B	R	G	B	B (Conflict)
3 = G	R	G	G	D

Verifiers of constraints

- Unary (SA ≠ green)
- Binary (1A ≠ 2A)
- Higher-order constraints

Constraint graph (nodes are variables, arcs (edges) are constraints)

Standard search formulation (incremental)

- Goal!*
- It solves CSP by starting with an empty state and incrementally filling in values that satisfy constraints.
 - initial state (empty assignment)
 - Successor function (defines how to move forward - by assigning a value to a variable that hasn't been assigned a value yet - If no valid assignment can be found, algorithm fails (backtrack to try new))
 - Goal test (when goal is achieved - every variable has been assigned with value. no constraints violated)

Same for all CSPs

- Use DFS
- Exact path doesn't matter as long as valid solution is reached. This allows to use of a Complete-State formulation - entire state of problem is considered.

• - Backtracking Search \rightarrow DFS for CSP, with single variable assignments

Commutative assignments { • In CSP's, order of variables does not matter. $NT = green$, $green = NT$.

Single variable assignment { • At each point in search, you can assign a value to one variable at a time then move to next unassigned variable.

• - Most Constrained Variable. \rightarrow No variable phy assign krns chahiye jis key pass sb my karm legal values hoo:
Variable which holds less valid options, we have to assign value to it firstly.

• Also called as minimum remaining values (MRV) \Rightarrow purpose is to handle. ~~by~~ that by handling most constrained variables firstly, we can reduce complexity of finding a solution.

• No variable select krta hain jo baqi variables per sb se zyada constraints dalta hy.

- Least Constraining Value.
- Forward Checking.