Registry > setting and options. Tashfeen Chapter no 2 (21, 2.2, 2.3, 2.4, 2.6, 2.8, 2.7) Abbesi Operating system Services. User Interface - graphical user interface windows and mice. (regular users touch interfaces on mobile phones · Command Dine Interfaces (Ch) text-based inputs Communication (TPC) - allows processes to Interplaces Communicate with each other Shaved Memory Message Parsing Processes Processes Communicate by reading receive mess ages & writing data exchange Shared region ky piper or memory So chets and executed by os. Command-line-interface CLI relies on command interpreter by Users with deep system Knowledge. Shell in UXINE examples , [. Bourne-Again Shell (Bash) · C shell (cm) · Kora shell tesh Methods to implement commands in (CLI · Command interpreter (Shell) · Interpreter - Inclustive > But Stain - commands) or copying files directly part of shell. trigger code incide interpreter Commands are programs to perform commands programs to perform commands handled by MXXXXXX sepayate program · UNIX - Ym command to delete file but executed shell finds & vun external program named these commands.

Centralized database

+ Store Configuration

Xetox Alto (1973) -> first computer to feature a Cul. Apple Macintosh in 1980. · GUZ gained populatity with evolved and adopt aqua interface in Macos. · UNIX system used CLI but now feature QUI DIKE KOE & GNOME. Touch-Skreen interface + " modern phones uses. Ma COS → · Historically GUI - focused (Aque Qui)
· Now, provides QuI & 4LI due to Unix - based tremal 2.3 System GODs. (way of a program to request services from Os.) Bridge blu softwere & operating system. (function in API).

alls) · getting files names -> opening -> verding & -> closing writing files. est of functions i) - Application Programming interface (API) - avalible to programmer, (Pet of vules and tools that allows software application · Window API · POSIX API to commonicate with each other) · Java API. Programmen access Asts vic Dibrovies provided by OS, e.s., UNIX & linux in Colbrary alled as C Dibyery or gorba posts or computers as long as those system also supports same API) · Advantages - Portability API & RTE hide implementation details from programmer easy to use complex · Astraction aptem function) API interact with OS easily rather than direct system calls. -> provides system and interface which links code with systems . - Runtime environment RTE) system calls. . - System call interface - maintains table of system call , associated number

· GUI offers a user-friendly afternative to CLI.

Boot loader > small program that initilizes system hardware and loads Os into memory When a computer starts up. Grob parameters are passed · - Parameters passing + · Registers (directly in register Parameters Block/memory Stored in homery blocking address of this black is prised in a register · Stack (Parameter pushed onto a stack by program and popped off by os. System Calls Categories · Process Control (manages executions of the process) · Device Management (manage hardware devices) Single Step mode / Trap mode · Musage Passing Model Information Maintenance · Shared-memory Model Communication Protection. · CPU mide for debugging · After every instruction is executed, 2.4 System Services System Utilities. a trap or interrupt 1) File Management 2) Status information s) File Modificatio 4) Programming alanguage support cl Program loading & execution 6) Communications a) Beckground Cervices

Compiled of	Binary interfac	ee (ABI) ->	interact at 100 Devel e.g., addressing, parameter passing
€ 2-6 Why A	pplications a	ve Os	Specific?
•		11	
· - System Ca	DOJ E APII	> . Or b	vovides Unique set of
		393ter	cation we to interact
		wit	th os.
·- Application	, to be made -) Inte	prieted languages python,
avalible on	multiple 05	(- Runs	on any OJ (ruby)
		with Slowe	on any OJ (ruby) interpreter. reversered
	;;) <u>k</u>	anguage u	with VM (Jana)
i ž	•	Compiles con	de into buterode
The state of the s	iii) sa	andard Lang	weed of ARIS (POSIX or
		Complied 1	vages or APIs (DNZX) rode for specific Os.
		Compare	yearne os
2.8 05 8	tructure		
·- Mono Oithic	Structure >:	Core function	s von in a single address space.
		combines os	Component into and leter
	Marger	executable file	in me call
Acolication	1) larger 2) faster	executable fole Ovignal Ox Linux also us	Son in a single address space. Components into one large NAX used monolithic.
Application	1) larger 2) faster (computer to the appear	executable fole Ovignal Ox	NDX used monolithic. es monolithic Kernal approach. It note address space. Used a
\uparrow	1) larger 2) faster composer to composer	executable file Ovignal Ox Linux also us on a st	NEX used monolithic. es monolithic Kernal approach. It ngte address space. Used a so > components can be located or orthogoded dynamically.
system cools, ves	Crash whole by System Crashes.	noduler desi	ngle address space. Used a 50 > Components can be located overloaded dynamically.
\uparrow	Crash, whole system crashes. 4) difficult to	vantage -> of	of the performance
system calls, VES IP (, File Sheduler,	3) If one agreement Crash, whole as System crashes. 4) difficult to add new function ality	vantage -> od	es mono within Remain approach. It ngle address space. Used a 50 > components can be located of unloaded dynamically. Plev high performance unning all Kernal operations single address operations
system colls, viss Ipc, file Sheduler, Virtual memory device onivers,	3) If one again Crash, whole System craches. 4) difficult to add new function ality 5) difficult debu	vantage -> of	es mono lithic rernex approach. It ngle address space. Used a 50 > components can be located or unloaded dynamically. Plev high performance unning all Kernal operations single address space, avoid elbormance costs.
system colls, viss Ipc, fite Sheduler, Vistual memory	3) If one again Crash, whole System craches. 4) difficult to add new function ality 5) difficult debu	vantage - of Sing. Olenges -	of monoplithic normax approach. It ngle address space. Used a go > components can be located of on loaded dynamically. Plev high performance unning all Kern-l operations single address space, avoid ellormance costs. • Complex due to lorge size
System colls, VES IP () fite Sheduler, Virtual memory device drivers, dispatcher	3) If one again Crash, whole System craches. 4) difficult to add new function ality 5) difficult debu	vantage - of Sing. Olenges -	of monoplithic normax approach. It ngle address space. Used a go > components can be located of on loaded dynamically. Plev high performance unning all Kern-l operations single address space, avoid ellormance costs. • Complex due to lorge size
system colls, viss Ipc, file Sheduler, Virtual memory device onivers,	3) If one symmet Crash, whole System crashes. 4) difficult to add new function ality Cha	vantage - of	es mono lithic rernex approach. It ngle address space. Used a 50 > components can be located or unloaded dynamically. Plev high performance unning all Kernal operations single address space, avoid elbormance costs.
System colls, VES IP () fite Sheduler, Virtual memory device drivers, dispatcher	3) If one symmet Crash, whole System crashes. 4) difficult to add new function ality Cha	vantage - of Sing. Olenges -	es mono lithic Normax approach. It ngle address space. Used a 50 > Components can be located of unloaded dynamically. Plev high performance unning all Kern-l operations single address space, avoid ellormance costs. • Complex due to lorge size • Changes might effect other parts of Kernal.
System colls, VES IP () fite Sheduler, Virtual memory device drivers, dispatcher	3) If one symmet Crash, whole System crashes. 4) difficult to add new function ality Cha	vantage - of sing. Olenges -	es mono lithic Normax approach. It ngle address space. Used a 50 > Components can be located of unloaded dynamically. Plev high performance unning all Kernal operations single address space, avoid elormance costs. • Complex due to lorge size • Changes might effect other parts of Kernal.
System colls, VES IP () fite Sheduler, Virtual memory device drivers, dispatcher	3) If one symmet Crash, whole System crashes. 4) difficult to add new function ality Cha	vantage - of	es mono lithic Normax approach. It ngle address space. Used a 50 > Components can be located of unloaded dynamically. Plev high performance unning all Kern-l operations single address space, avoid ellormance costs. • Complex due to lorge size • Changes might effect other parts of Kernal.

Each layer hider complexity of its lower Dayers from higher Bottom Payer: Hardware Higest layer: user interface. Layered Approach > · Each layer only interacts with its adjacent layer (higher layers be function no. of layers (levels) each built on top of lower layer) from lower layer) · Each Dayer is implemented using oney operations and sprvices pr-vided by comer used in This design allows Changes in one layer to impact not the entire system. Computer networks and web application (- Advantages - Easier debugging, as layors are tested and developed independently . hayers only need to know what operations do, not how they are implement. . - Challenges > · defining clear layer can be difficult.

o reduce system performance due to
multiple layers · Minimal core: handles ersential function like process and ·- Micyokernal > · Message pessing: - Communication occurs through message (minimized functionality within Kernal do not directly interact with Kernal. itself new services is easier. · Examples -> Parwin (macos, 105) -> micro Kerral ONX (embedded) -> neutrino
Rystem) -> neutrino
mica o kern-l · - Challenges -> suffer from performance issues due to overhead of message passing and context switching between processes.

