

## **WIA2004 Operating Systems**

**Lab 7: Deadlock Management** 

Group: F5U

Occ: 10

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## 1.0 Introduction

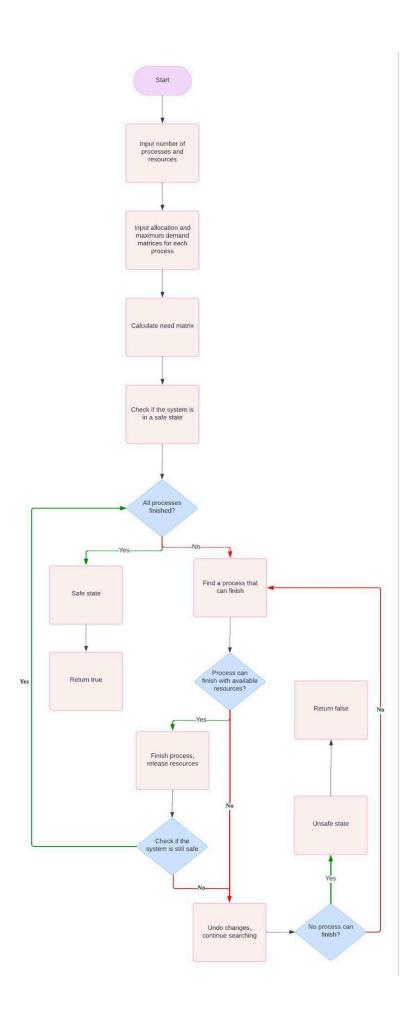
Deadlock is a situation in computing where a set of processes become stuck in a cycle of waiting for each other to release resources. Specifically, each process in the set is waiting for a resource that is currently held by another process in the same set, creating a scenario where none of the processes can proceed. The necessary conditions needed for deadlock include mutual exclusion, hold and wait, no preemption, and circular wait.

The Banker's Algorithm is a resource allocation and deadlock avoidance technique used in operating systems to ensure efficient sharing of resources among multiple processes. It operates by simulating the allocation of resources to processes while checking for safety at each step, ensuring that granting resource requests will not lead to deadlock. Named after banking practices, where loans are granted based on the availability of funds to ensure that all customers' needs can be met, the algorithm similarly ensures that the system remains in a "safe state" where processes can continue execution without the risk of deadlock.

Each process informs the system about its resource needs and how long it will use them. This helps the system allocate resources efficiently. The algorithm tracks the maximum resources each process can use, preventing deadlock. However, it has limitations: it works only with a fixed number of processes, doesn't allow processes to change their needs during execution, and requires processes to declare their maximum needs upfront. Additionally, resource allocation must occur within a set time limit, typically one year.

# 2.0 Methodology

Flowchart



### Pseudocode

## Class BankersAlgorithm: Integer numberOfProcesses Integer numberOfResources Array available Array maximum Array allocation Array need Constructor BankersAlgorithm(numberOfProcesses, numberOfResources): Initialize numberOfProcesses, numberOfResources Initialize available, maximum, allocation, need arrays Method inputMatrices(scanner): Print "Enter details for each process:" For each process i from 0 to numberOfProcesses: Print "P" + i + " Allocation: " Read allocation for process i from scanner For each resource j from 0 to numberOfResources: Set allocation[i][j] to allocation value for resource j Print "P" + i + " Max: " Read maximum claim for process i from scanner For each resource j from 0 to numberOfResources: Set maximum[i][j] to maximum claim value for resource j Print "Enter the Available Resources:" Read available resources from scanner For each resource i from 0 to numberOfResources: Set available[i] to available resource value for resource i Call calculateNeed() Method calculateNeed(): For each process i from 0 to numberOfProcesses: For each resource j from 0 to numberOfResources: Set need[i][j] to maximum[i][j] - allocation[i][j] Method isSafeState(): Initialize work as a copy of available Initialize finish as an array of boolean values of size numberOfProcesses, all set to false Initialize safeSequence as an array of integers of size numberOfProcesses Initialize count to 0

```
While count < numberOfProcesses:
     Set foundProcess to false
     For each process i from 0 to numberOfProcesses:
       If process i is not finished (finish[i] is false):
          Set canProceed to true
          For each resource j from 0 to numberOfResources:
            If need[i][j] > work[j]:
               Set canProceed to false
               Break out of inner loop
          If canProceed is true:
             For each resource k from 0 to numberOfResources:
               Increment work[k] by allocation[i][k]
            Add process i to safeSequence at position count
            Increment count by 1
            Set finish[i] to true
            Set foundProcess to true
            Print "P" + i + " is visited: ("
            For each resource k from 0 to numberOfResources:
               Print work[k]
               If k is not the last resource:
                 Print "."
            Print ")"
     If no process is found:
       Print "Deadlock! The system is in an unsafe state."
       Return false
  Print "\nThe system is in a safe state."
  Print "The safe sequence is: "
  For each process i from 0 to numberOfProcesses:
     Print "P" + safeSequence[i]
     If i is not the last process:
       Print " -> "
  Print ""
  Return true
Method requestResources(processID, request):
  For each resource i from 0 to numberOfResources:
     If request[i] > need[processID][i]:
       Print "Error: Process has exceeded its maximum claim."
       Return false
     If request[i] > available[i]:
       Print "Resources are not available."
       Return false
  For each resource i from 0 to numberOfResources:
     Decrement available[i] by request[i]
```

Increment allocation[processID][i] by request[i] Decrement need[processID][i] by request[i]

If isSafeState() returns false:

Print "Request cannot be granted. Rolling back."
For each resource i from 0 to numberOfResources:
Increment available[i] by request[i]
Decrement allocation[processID][i] by request[i]
Increment need[processID][i] by request[i]
Return false

Return true

Method printState():
 Print "Process\tAllocation\tMax\t\tNeed"

## 3.0 Implementation

## 3.1 Coding

```
public class Lab_7 {
    private int numberOfProcesses;
    private int[] available;
    private int[][] maximum;
    private int[][] allocation;
    private int[][] need;

public Lab_7(int numberOfProcesses, int numberOfResources) {
        this.numberOfProcesses = numberOfProcesses;
        this.numberOfResources = numberOfResources;
        this.available = new int[numberOfProcesses][numberOfResources];
        this.allocation = new int[numberOfProcesses][numberOfResources];
        this.need = new int[numberOfProcesses][numberOfResources];
    }
}
```

```
public void inputMatrices(Scanner scanner) {
   System. out. println("Enter details for each process:");
        System. out. println("P" + i + " Allocation: ");
        String[] alloc = scanner.nextLine().split(",");
            allocation[i][j] = Integer. parseInt(alloc[j]);
        System. out. println("P" + i + " Max: \t');
        String[] max = scanner.nextLine().split(",");
        for (int j = 0; j < numberOfResources; <math>j++) {
            maximum[i][j] = Integer. parseInt(max[j]);
   System. out. println("Enter the Available Resources:");
    String[] avail = scanner.nextLine().split(",");
    for (int i = 0; i < numberOfResources; i++) {</pre>
        available[i] = Integer. parseInt(avail[i]);
    calculateNeed();
```

```
//need matrix = max matrix - allocation matrix
private void calculateNeed() {
   for (int i = 0; i < numberOfProcesses; i++) {
      for (int j = 0; j < numberOfResources; j++) {
         need[i][j] = maximum[i][j] - allocation[i][j];
      }
}</pre>
```

```
public boolean isSafeState() {
    int[] safeSequence = new int[numberOfProcesses];
        boolean foundProcess = false;
            if (!finish[i]) {
                boolean canProceed = true;
                    if (need[i][j] > work[j]) {
                        canProceed = false;
                if (canProceed) {
                        work[k] += allocation[i][k];
                    safeSequence[count++] = i;
                    finish[i] = true;
                    foundProcess = true;
                    System. out. print("P" + i + " is visited:\t(");
                        System. out. print(work[k]);
                    System. out. println(")");
```

```
System. out. println(")");
}

if (!foundProcess) {
    System. out. println("Deadlock! The system is in an unsafe state.");
    return false;
}

System. out. println("\nThe system is in a safe state.");
System. out. print("The safe sequence is: ");
for (int i = 0; i < numberOfProcesses; i++) {
    System. out. print("P" + safeSequence[i]);
    if (i != numberOfProcesses - 1) System. out. print(" -> ");
}
System. out. println();
return true;
```

```
public boolean requestResources(int processID, int[] request) {
            System. out. println("Resources are not available.");
        allocation[processID][i] += request[i];
        need[processID][i] -= request[i];
    if (!isSafeState()) {
        System. out. println("Request cannot be granted. Rolling back.");
            allocation[processID][i] -= request[i];
            need[processID][i] += request[i];
```

```
public void printState() {
    System. out. println("Process\tAllocation\tMax\t\tNeed");
    for (int i = 0; i < numberOfProcesses; i++) {
        System. out. print("P" + i + "\t");
        for (int j = 0; j < numberOfResources; j++) {
                System. out. print(allocation[i][j] + (j == numberOfResources - 1 ? "" : ","));
        }
        System. out. print("\t\t");
        for (int j = 0; j < numberOfResources; j++) {
                System. out. print(maximum[i][j] + (j == numberOfResources - 1 ? "" : ","));
        }
        System. out. print("\t\t");
        for (int j = 0; j < numberOfResources; j++) {
                System. out. print(need[i][j] + (j == numberOfResources - 1 ? "" : ","));
        }
        System. out. println();
    }
}</pre>
```

```
Scanner scanner = new Scanner (System. in);
Lab_7 bankersAlgorithm = new Lab_7 (numberOfProcesses, numberOfResources);
System. out. println("Enter the requested resources:");
String[] requestInput = scanner.nextLine().split(",");
} catch (Exception e) {
    e. printStackTrace();
```

```
if (bankersAlgorithm.requestResources(processID, request)) {
    System.out.println("The request has been granted.");
} else {
    System.out.println("The request has been denied.");
}

bankersAlgorithm.printState();
scanner.close();
}
```

# 3.2 Sample output

## 3.2.1 No deadlock

#### Allocation matrix:

P0: [0, 1, 0]

P1: [2, 0, 0]

P2: [3, 0, 2]

P3: [2, 1, 1]

P4: [0, 0, 2]

### Max matrix:

P0: [7, 5, 3]

P1: [3, 2, 2]

P2: [9, 0, 2]

P3: [2, 2, 2]

P4: [4, 3, 3]

### Need matrix:

P0: [7, 4, 3]

P1: [1, 2, 2]

P2: [6, 0, 0]

P3: [0, 1, 1]

P4: [4, 3, 1]

### User input:

```
Enter number of processes:
Enter number of resources:
PO Allocation:
PO Max:
P1 Allocation:
P1 Max:
P2 Allocation:
P2 Max:
P3 Allocation:
P3 Max:
P4 Allocation:
P4 Max:
Enter the Available Resources:
Enter the process number(PID) making the request:
Enter the requested resources:
```

#### Step 1: Check the Request

- Process (PID): 1
- Request: (1, 0, 2)

#### Step 2: Allocate Resources

- Update Available Resources: [3-1, 3-0, 2-2] = [2, 3, 0]
- Update Allocation for P1(Allocation[P1] + Request): [2+1, 0+0, 0+2] = [3, 0, 2]
- Update Need for P1(Need[P1] Request): [1-1, 2-0, 2-2] = [0, 2, 0]

#### Step 3: Check for Safety

#### Initialize:

- Work Array: [2, 3, 0]
- Finish Array: [false, false, false, false, false]
- Safe Sequence: []

Iterate through each process to check if it can proceed:

- 1. Check Process P0:
  - Need[P0] is [7, 4, 3], but Work is [2, 3, 0], so P0 cannot proceed.
  - Break; i++
- 2. Check Process P1:
  - Need[P1] is [0, 2, 0], and Work is [2, 3, 0], so P1 can proceed.
  - work[] =work[] + allocation[1] = [2,3,0] + [3,0,2] = [5,3,2]
  - i++
- Check Process P2:
  - Need[P2] is [6, 0, 0], but Work is [5, 3, 2], so P2 cannot proceed.
  - Break; i++
- 4. Check Process P3:
  - Need[P3] is [0, 1, 1], and Work is [5, 3, 2], so P3 can proceed.
  - work[] =work[] + allocation[3] = [5,3,2] + [2,1,1] = [7,4,3]
  - i++
- 5. Check Process P4:
  - Need[P4] is [4, 3, 1], and Work is [2, 3, 0], so P4 can proceed.
  - work[] =work[] + allocation[4] = [7,4,3] + [0,0,2] = [7,4,5]
  - i++

- 6. Finish Array: [false, true, false, true, true]
  - ->not all element are true(count<numberOfProcess)
  - ->while()loop keep executing
- 7. Check Process P0:
  - Need[P0] is [7, 4, 3], and Work is [7, 4, 5], so P0 can proceed.
  - work[] =work[] + allocation[0] = [7,4,5] + [0,1,0] = [7,5,5]
  - Break; i++
- 8. Check Process P2:
  - Need[P2] is [6, 0, 0], and Work is [7, 5, 5], so P2 can proceed.
  - work[] =work[] + allocation[2] = [7,5,5] + [3,0,2] = [10,5,7]
  - Break; i++
- 9. Finish Array: [true, true, true, true, true]
  - ->all element are true(count !< numberOfProcess)
  - ->while()loop step executing
- 10. Safe state find, no deadlock

## Sample output:

```
P1 is visited: (5, 3, 2)
P4 is visited: (7, 4, 5)
P2 is visited: (10, 5, 7)
The safe sequence is: P1 \rightarrow P3 \rightarrow P4 \rightarrow P0 \rightarrow P2
The request has been granted.
P0
        0, 1, 0
                         7, 5, 3
P1
                         3, 2, 2
P2
       3, 0, 2
                         9, 0, 2
                                           6, 0, 0
P3
        2, 1, 1
                         2, 2, 2
                                           0, 1, 1
P4
      0, 0, 2
                                           4, 3, 1
BUILD SUCCESS
Total time: 47.649 s
Finished at: 2024-05-14T15:32:47+08:00
```

### 3.2.2 Exist deadlock:

#### Allocation matrix:

$\mathbf{A}$	В	$\mathbf{C}$	
P0	0	1	0
P1	2	0	0
P2	3	0	2
P3	2	1	1
P4	0	0	2

Max matrix:

$\mathbf{A}$	В	$\mathbf{C}$	
P0	7	5	3
P1	3	2	2
P2	9	0	2
P3	2	2	2
P4	4	3	3

Need matrix:

$\mathbf{A}$	В	$\mathbf{C}$	
P0	7	4	3
P1	1	2	2
P2	6	0	0
P3	0	1	1
P4	4	3	1

#### User input:

```
Enter number of resources:
PO Allocation:
PO Max:
P1 Allocation:
P2 Allocation:
P2 Max:
P3 Allocation:
P3 Max:
P4 Allocation:
P4 Max:
Enter the Available Resources:
Enter the process number(PID) making the request:
Enter the requested resources:
```

#### Step 1: Check the Request

Process (PID): 1Request: (1, 0, 0)

#### Step 2: Allocate Resources

• Update Available Resources: [1-1, 1-0, 0-0] = [0, 1, 0]

• Update Allocation for P1: [2+1, 0+0, 0+0] = [3, 0, 0]

• Update Need for P1: [1-1, 2-0, 2-0] = [0, 2, 2]

#### Step 3: Check for Safety

#### Initialize:

• Work Array: [0, 1, 0]

• Finish Array: [false, false, false, false, false]

• Safe Sequence: []

Iterate through each process to check if it can proceed:

- 11. Check Process P0:
  - Need[P0] is [7, 4, 3], but Work is [0, 1, 0], so P0 cannot proceed.
- 12. Check Process P1:
  - Need[P1] is [1, 2, 2], but Work is [0, 1, 0], so P1 cannot proceed.
- 13. Check Process P2:
  - Need[P2] is [6, 0, 0], but Work is [0, 1, 0], so P2 cannot proceed.
- 14. Check Process P3:
  - Need[P3] is [0, 1, 1], but Work is [0, 1, 0], so P3 cannot proceed.
- 15. Check Process P4:
  - Need[P4] is [4, 3, 1], but Work is [0, 1, 0], so P4 cannot proceed.

#### Deadlock Result

Since none of the processes can proceed, the system is in a deadlock state. The system will output that no safe sequence was found and declare the system to be in an unsafe state:

## Sample output:

Deadloc	Deadlock! The system is in an unsafe state.			
Request	Request cannot be granted. Rolling back.			
The request has been denied.				
Process	Allocation	Max	Need	
P0	0, 1, 0	7, 5, 3	7, 4, 3	
P1	2, 0, 0	3, 2, 2	1, 2, 2	
P2	3, 0, 2	9, 0, 2	6, 0, 0	
P3	2, 1, 1	2, 2, 2	0, 1, 1	
l P4	0, 0, 2	4, 3, 3	4, 3, 1	

#### BUILD SUCCESS

\_\_\_\_\_

Total time: 53.513 s

Finished at: 2024-05-14T21:39:09+08:00

------

## 4.0 Conclusion

The Banker's algorithm is a valuable tool for ensuring safe resource allocation in systems with multiple competing requests. It prevents deadlocks by simulating resource allocation scenarios and guaranteeing that the system can reach a state where all processes can finish their tasks. In essence, it prioritizes system stability by ensuring sufficient resources are always available to meet future demands.

However, the Banker's algorithm does have limitations. It requires advanced knowledge of the maximum resource needs for each process, which may not always be practical. Additionally, it doesn't allow for dynamic changes in resource requirements or the addition of new processes mid-execution.

Despite these limitations, the Banker's algorithm remains a crucial concept for understanding resource allocation and deadlock avoidance in various systems, including operating systems and banking environments.

## 5.0 References

- Banker's Algorithm in Operating System javatpoint. (n.d.). Www.javatpoint.com. https://www.javatpoint.com/bankers-algorithm-in-operating-system
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