The implementation of "Walking One" algorithm at SIEMENS 1200 Hardware and Software of Control Systems

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1 Introduction

1.1 Uses of the algorithm

The "walking one" algorithm is one of the most well known algorithms used for testing the independence of operations performed by changing the state of consecutive bits.

Its primal intention is checking of coupling between two consecutive bits which ideally should be independent. If by changing the output by one two of them have the undesired state this means that the input device is not functioning properly.

1.2 Hardware description

As pointed out in the requirements the picked PLC was:

SIMANTIC S7-1200 1212c DC/DC/DC - 6ES7 212-1AE40-0xB0 - Firmware V4.1

Because of the low ammount of digital outputs the following IO module was added:

DI 16x24VDC/DQ 16xRelay - 6ES7 223-1PL30-0XB0 - Firmware V1.0

Despite offering more it was used just as a digital output device.

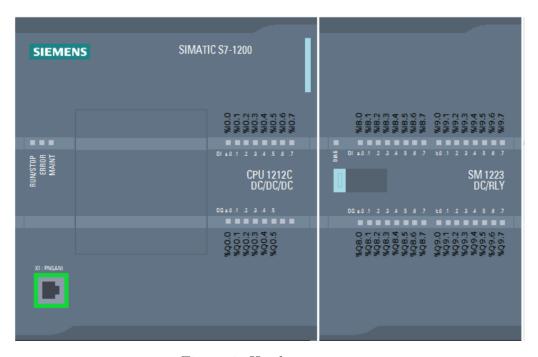


Figure 1: Hardware setup

1.3 Specification

The requirements tated for the implementation of the algorithm to be performed using TIA Portal v13 or v14 with PLCSim. The implementation had to be written using the SCL language with cyclic interrupt Organization Blocks, indirect addressing and Function Blocks.

2 Algorithm description

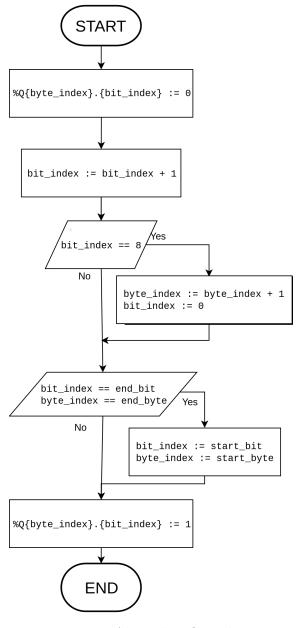


Figure 2: Algorithm flow chart

3 Implementation details

3.1 Variables and valueds

Because of indirect addressing and the technologies applied all the variables and values had to be represented by two numbers – the bitOffset and byteOffset. Both of them can be put in the output address representation like %Q{byteOffset}.{bitOffset}. An additional limitation to the bitOffset is that it can only be in range of 0 to 7. All the datatypes were picked to match the ones which were required by the POKE_BOOL function.

3.1.1 Global constants

There were four global constants – all of them were used as program constraints. They can be grouped into two parts:

- The output with the least number:
 - start_byte dint
 - start_bit int
- The output with the number after the range:
 - end_byte dint
 - end_bit int

3.1.2 Local variables of function block

The function block which was used for this algorithm consists of two variables. Both initialized at the beginning to the least allowable values. They retain their state between function block calls.

- byte_index dint
- bit_index int

These are the variables which are updated every interrupt.

3.2 Source code

Because of the use of POKE_BOOL the code became platform dependent and is not to be reused on other SIEMENS PLCs. The code consists of three separate parts:

3.2.1 Setting the previous one to zero

The following call is responsible for setting the output at %Q{byteOffset}.{bitOffset} to zero. 16#82 tells the function to look at the output memory area to change. The function changes only the bit it is pointed two with the two parameters.

3.2.2 Updating the indices

3.2.3 Setting the next one

4 Simulation

The simulated examples contain the full range of corner cases which had to be taken into consideration. Considered were cases where the limits were set either to values which are the least and highest possible or the cases just before them (one above minimum and one below maximum).

start_byte	start_bit	$\mathtt{end}_{\mathtt{-}}\mathtt{byte}$	end_bit	Experiment result
8	0	10	0	Passed
8	1	10	0	Passed
8	0	9	7	Passed
8	1	9	7	Passed

Table 1: Test sets

The only issues encountered here were that the first iteration started at the second cell instead of the first one, but the rest continued to work properly. The workaround to that is taking the last indices of the range of intrest as the starting ones – the last output would be set to zero and the first output set to one would be the first one. Working around this behaviour would require additional setup logic which should calculate the last valid index with all corner cases.

5 Summary

5.1 Results obtained

The results obtained is a fully functioning walking one implementation with a change period of one second. The first iteration warries just a bit from the expected value (it does not pick up the first ouput in the first iteration), but despite that it works well and does not have the same issue in the iterations afterwards.

5.2 Difficulties encountered

The biggest issue here was with the 21-day trial license. SIEMENS did not handle the usecase when somebody decided to use another trial version of their software on the same computer – after installing TIA Portal v14 there were issues with using TIA Portal v13 because it recognized the previous license instead of using its own. The solution to this problem was installing the software on another computer (another would be to completely reinstall the operating system).

Another issue was a bug which did not allow running PLCSim because of missing DLL – this required just installing another patch.

5.3 Lessons for the future

One of the most meaningful lessons was a selfcheck on remembering the process of programming a PLC device, but this time in another language and using Organizat Blocks instead of LAD and Function Blocks.

Another lesson was simulating PLC behaviour and the pitfalls of it when the configurations of both programmes do not match up.

It also showed a effective way of eliminating boilerplate code by using excel like sheets for variable declaration (which take less place than actual SCL code).