Parser Design

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Section 1. Transformed Grammar Into LL(1)

Steps to correcting first and follow set conflicts

Once the transformed grammar without EBNF optional/repeating constructs were removed along with left-recursion using Dr. Paquet's grammartool.jar, the resulting grammar output was entered in the UCalgary Grammar Tool. The result of the analysis of the grammar is shown in the following code block:

```
FUNCHEAD has a first set conflict.

FUNCHEAD_1 is nullable with clashing first and follow sets.

MEMBERDECL has a first set conflict.

EXPR has a first set conflict.

FACTOR has a first set conflict.

FUNCTIONCALL_1 is nullable with clashing first and follow sets.

IDNEST has a first set conflict.

FUNCDECL has a first set conflict.

STATEMENT has a first set conflict.

VARIABLE_1 is nullable with clashing first and follow sets.

ARRAYSIZE has a first set conflict.
```

FUNCHEAD | FUNCHEAD 1

Left factored the FUNCHEAD_1 id lpar FPARAMS rpar colon ... portion and changed the FUNCHEAD_1 to take into that they both can start by id with optional sr id before the lpar terminal.

Before:

```
FUNCHEAD -> FUNCHEAD_1 id lpar FPARAMS rpar colon TYPE .

FUNCHEAD -> FUNCHEAD_1 id lpar FPARAMS rpar colon void .

FUNCHEAD_1 -> id sr .

FUNCHEAD_1 -> .
```

After:

```
FUNCHEAD -> id FUNCHEAD_1 lpar FPARAMS rpar colon TYPEORVOID .

FUNCHEAD_1 -> sr id .

FUNCHEAD_1 -> .

TYPEORVOID -> TYPE .

TYPEORVOID -> void .
```

FUNCDECL

Left factored the id lpar FARAMS rpar colon portion and replaced the rule to account for TYPEORVOID:

Before:

```
FUNCDECL -> id lpar FPARAMS rpar colon TYPE semi .
FUNCDECL -> id lpar FPARAMS rpar colon void semi .
```

After:

```
\label{eq:function} \text{FUNCDECL} \; -\!\!\!\!> \; \text{id} \;\; \text{lpar} \;\; \text{FPARAMS} \;\; \text{rpar} \;\; \text{colon} \;\; \text{TYPEORVOID} \;\; \text{semi} \;\; .
```

ARRAYSIZE

Factored the possibly nullable second intnum terminal.

Before:

```
ARRAYSIZE -> lsqbr intnum rsqbr .
ARRAYSIZE -> lsqbr rsqbr .
```

After:

```
ARRAYSIZE -> lsqbr MAYBEININUM rsqbr .

MAYBEININUM -> intnum .

MAYBEININUM -> .
```

IDNEST

Left factored the id terminal.

Before:

```
IDNEST -> id IDNEST_1 dot .
IDNEST -> id lpar APARAMS rpar dot .

IDNEST_1 -> INDEX IDNEST_1 .
IDNEST_1 -> .
```

After:

```
IDNEST_1 -> id IDNEST_1 .

IDNEST_1 -> IDNEST_2 dot .

IDNEST_1 -> lpar APARAMS rpar dot .

IDNEST_2 -> INDEX IDNEST_2 .

IDNEST_2 -> .
```

EXPR

Left factored the common ARITHEXPR in both RELEXPR and ARITHEXPR

Before:

```
EXPR -> ARITHEXPR .
EXPR -> RELEXPR .

ARITHEXPR -> TERM ARITHEXPR_1 .

ARITHEXPR_1 -> ADDOP TERM ARITHEXPR_1 .

ARITHEXPR_1 -> .

RELEXPR -> ARITHEXPR RELOP ARITHEXPR .
```

After:

```
EXPR_1 -> RELOP ARITHEXPR .

EXPR_1 -> .

ARITHEXPR -> TERM ARITHEXPR_1 .
```

```
ARITHEXPR_1 -> ADDOP TERM ARITHEXPR_1 .
ARITHEXPR_1 -> .

RELEXPR -> ARITHEXPR RELOP ARITHEXPR .
```

MEMBERDECL

The first step was to factor the FUNCDECL to take care of TYPE or void. This was easily done by eliminating the two productions into one by using a seperate production TYPEORVOID. Afterwards, the id shared by both FUNCDECL and possibly VARDECL was slightly trickier because only when the TYPE of VARDECL uses the id terminal this would clash. To eliminate the clash, the id needed to be left factored out of both FUNCDECL and VARDECL. This was done with the help of INTEGERORFLOAT to take into account the case when id isn't the first terminal of VARDECL. This left FUNCDECL unsued now, so it was removed and the new production was renamed to FUNCDECLIDLESS to signific that it used to be the old definition of FUNCDECL but is missing the first id token that was prepended.

After:

```
MEMBERDECL -> FUNCDECL .

MEMBERDECL -> VARDECL .

FUNCDECL -> id lpar FPARAMS rpar colon TYPE semi .

FUNCDECL -> id lpar FPARAMS rpar colon void semi .

VARDECL -> TYPE id VARDECL_1 semi .

VARDECL_1 -> ARRAYSIZE VARDECL_1 .

VARDECL_1 -> .
```

Before:

```
MEMBERDECL -> id MEMBERDECL_1 .

MEMBERDECL_1 -> VARDECL_1 .

MEMBERDECL_1 -> VARDECL_1 .

MEMBERDECL_1 -> FUNCDECLIDLESS .

FUNCDECLIDLESS -> lpar FPARAMS rpar colon TYPEORVOID semi .

VARDECL -> INTEGERORFLOAT VARDECL_1 .

VARDECL_> id VARDECL_1 .

VARDECL_1 -> id VARDECL_2 semi .

VARDECL_2 -> ARRAYSIZE VARDECL_2 .

VARDECL_2 -> .
```

FUNCTIONCALL_1

The removal of the first and follow set clash for this one was tricky. At first, the clash was with respect to the id terminal. This was fixed by left factoring it. However, another first and follow set clash remained afterwards with respect to the lpar terminal. This was corrected by factoring the IDNEST components that were causing the issue by further decomposing it into IDNEST_2 and IDNEST_3. This way, the conflicting lpar would not occur from a possibly null terminal.

Before:

```
FUNCTIONCALL_1 -> FUNCTIONCALL_1 id lpar APARAMS rpar .

FUNCTIONCALL_1 -> IDNEST FUNCTIONCALL_1 .

FUNCTIONCALL_1 -> .

IDNEST -> id IDNEST_1 dot .
```

```
IDNEST -> id lpar APARAMS rpar dot .

IDNEST_1 -> INDEX IDNEST_1 .

IDNEST_1 -> .
```

After:

```
FUNCTIONCALL_1 -> id FUNCTIONCALL_1 .

FUNCTIONCALL_1 -> IDNEST_1 FUNCTIONCALL_2 .

FUNCTIONCALL_2 -> dot id IDNEST_3 .

FUNCTIONCALL_2 -> .

IDNEST -> id IDNEST_1 dot .

IDNEST_1 -> IDNEST_2 .

IDNEST_1 -> IDNEST_3 .

IDNEST_2 -> INDEX IDNEST_2 .

IDNEST_2 -> INDEX IDNEST_2 .

IDNEST_2 -> .

IDNEST_3 -> lpar APARAMS rpar .
```

VARIABLE_1

This first and follow set conflict required to factor the id terminal in VARIABLE that occured inside IDNEST (which was replaced by it's IDNEST_1 id form).

Before:

```
VARIABLE_1 -> VARIABLE_1 id VARIABLE_2 .

VARIABLE_1 -> IDNEST VARIABLE_1 .

VARIABLE_1 -> .

VARIABLE_2 -> INDEX VARIABLE_2 .

VARIABLE_2 -> .
```

After:

```
VARIABLE_1 -> id VARIABLE_1 VARIABLE_2 .

VARIABLE_1 -> IDNEST_1 dot VARIABLE_3 .

VARIABLE_2 -> INDEX VARIABLE_2 .

VARIABLE_2 -> .

VARIABLE_3 -> IDNEST VARIABLE_3 .

VARIABLE_3 -> .
```

FACTOR

The first conflict for this was easy to fix with a simple left factorization of the id terminal shared between VARIABLE and FUNCTIONCALL. However, the new resulting production of this factorization introduced yet another first set conflict with respect to the lsqbr terminal. After fixing this one, yet another came along with regards to id, and another one after with regards to IDNEST_3, all of which were fixed using left factorization.

Before:

```
FACTOR -> VARIABLE .
FACTOR -> FUNCTIONCALL .
FACTOR -> intnum .
```

```
FACTOR -> floatnum .
FACTOR -> lpar ARITHEXPR rpar .
FACTOR -> not FACTOR .
FACTOR -> sign Factor .

VARIABLE -> id VARIABLE_1 VARIABLE_2 .

VARIABLE_1 -> IDNEST_1 dot VARIABLE_3 .

VARIABLE_2 -> INDEX VARIABLE_2 .

VARIABLE_2 -> .

VARIABLE_3 -> IDNEST VARIABLE_3 .

VARIABLE_3 -> id FUNCTIONCALL_1 .

FUNCTIONCALL -> id FUNCTIONCALL_1 .

FUNCTIONCALL_1 -> IDNEST_1 FUNCTIONCALL_2 .

FUNCTIONCALL_2 -> dot id IDNEST_3 .

FUNCTIONCALL_2 -> dot id IDNEST_3 .

FUNCTIONCALL_2 -> .
```

After:

```
FACTOR \rightarrow id FACTOR_1.
FACTOR \rightarrow intnum.
FACTOR -> floatnum
FACTOR \rightarrow lpar ARITHEXPR rpar.
FACTOR \rightarrow not FACTOR.
FACTOR \rightarrow SIGN FACTOR.
\label{eq:factor_1} \text{FACTOR}\_1 \ -\!\!\!> \ \text{IDNEST}\_1 \ \text{FACTOR}\_2 \ .
FACTOR_2 \rightarrow dot FACTOR_3.
FACTOR_2 -> .
FACTOR_3 \rightarrow id FACTOR_4.
FACTOR_3 \rightarrow.
FACTOR\_4 \rightarrow IDNEST\_2 dot VARIABLE\_3.
FACTOR\_4 \rightarrow IDNEST\_3 FACTOR\_5.
FACTOR_5 \rightarrow dot VARIABLE_3.
FACTOR_5 \rightarrow.
VARIABLE -> id VARIABLE_1 VARIABLE_2 .
VARIABLE_1 \rightarrow IDNEST_1 VARIABLE_4.
VARIABLE 2 \rightarrow INDEX VARIABLE 2.
VARIABLE_2 \rightarrow.
VARIABLE 3 -> IDNEST VARIABLE 3 .
VARIABLE_3 -> .
VARIABLE_4 \rightarrow dot VARIABLE_3.
FUNCTIONCALL \rightarrow id FUNCTIONCALL 1.
FUNCTIONCALL_1 \rightarrow IDNEST_1 FUNCTIONCALL_2.
FUNCTIONCALL\_2 -\!\!\!> \ dot \ id \ IDNEST\_3 \ .
FUNCTIONCALL_2 \rightarrow.
```

STATEMENT

The initial conflict with this rule was id (corrected after left factoring and creating new production STATEMENT_1), which was factored out that involved the ASSIGNSTAT and FUNCTIONCALL variables. However, multiple new first set

conflicts arose, such as: IDNEST_1, dot, and id, and IDNEST_3 as made evident by the new productions STATEMENT_2, STATEMENT_3, STATEMENT_4, and STATEMENT_5, respectively. Because of the need for decomposing the rule so much, it ended up that the FUNCTIONCALL, FUNCTIONCALL_1, and FUNCTIONCALL_2 productions were not needed anymore.

Before:

```
STATEMENT -> ASSIGNSTAT semi .

STATEMENT -> if lpar RELEXPR rpar then STATBLOCK else STATBLOCK semi .

STATEMENT -> while lpar RELEXPR rpar STATBLOCK semi .

STATEMENT -> read lpar VARIABLE rpar semi .

STATEMENT -> write lpar EXPR rpar semi .

STATEMENT -> return lpar EXPR rpar semi .

STATEMENT -> FUNCTIONCALL semi .

ASSIGNSTAT -> VARIABLE ASSIGNOP EXPR .
```

After:

```
STATEMENT -> id STATEMENT 1 .
{
m STATEMENT} \rightarrow {
m if} \ {
m lpar} \ {
m RELEXPR} \ {
m rpar} \ {
m then} \ {
m STATBLOCK} \ {
m else} \ {
m STATBLOCK} \ {
m semi} \ .
{\it STATEMENT} \longrightarrow {\it while lpar RELEXPR rpar STATBLOCK semi} .
STATEMENT -> read lpar VARIABLE rpar semi .
STATEMENT -> write lpar EXPR rpar semi
STATEMENT \rightarrow return lpar EXPR rpar semi.
STATEMENT 1 \rightarrow IDNEST 1 STATEMENT 2.
STATEMENT_2 \rightarrow dot STATEMENT_3.
STATEMENT 2 \rightarrow semi.
STATEMENT 3 -> id STATEMENT 4
STATEMENT\_3 -> VARIABLE\_2 \ ASSIGNSTAT\_1 \ semi \ .
STATEMENT 4 \rightarrow IDNEST 2 dot VARIABLE 3 VARIABLE 2 ASSIGNSTAT 1 semi .
STATEMENT_4 \rightarrow IDNEST_3 STATEMENT_5.
STATEMENT_5 -> dot VARIABLE_3 VARIABLE_2 ASSIGNSTAT_1 semi .
STATEMENT 5 \rightarrow \text{semi}.
ASSIGNSTAT \longrightarrow VARIABLE \ ASSIGNSTAT\_1 \ .
ASSIGNSTAT_1 \rightarrow ASSIGNOP EXPR.
```

Transformed grammar after removal of EBNF optional/repeating notation and left-recursion

```
START -> PROG .

PROG -> PROG_1 PROG_2 main FUNCBODY .

PROG_1 -> CLASSDECL PROG_1 .

PROG_1 -> .

PROG_2 -> FUNCDEF PROG_2 .

PROG_2 -> .

CLASSDECL -> class id CLASSDECL_1 lcurbr CLASSDECL_2 rcurbr semi .

CLASSDECL_1 -> inherits id CLASSDECL_3 .

CLASSDECL_1 -> .

CLASSDECL_2 -> VISIBILITY MEMBERDECL CLASSDECL_2 .

CLASSDECL_3 -> comma id CLASSDECL_3 .

CLASSDECL_3 -> comma id CLASSDECL_3 .

CLASSDECL_3 -> .
```

```
\label{eq:function} \text{FUNCDEF} \to \text{FUNCHEAD FUNCBODY semi} \quad .
FUNCHEAD \rightarrow FUNCHEAD_1 id lpar FPARAMS rpar colon TYPE .
FUNCHEAD -> FUNCHEAD_1 id lpar FPARAMS rpar colon void .
FUNCHEAD\_1 -> id sr .
FUNCHEAD 1 \rightarrow.
\label{eq:memberdecl} \mbox{MEMBERDECL} \ -> \ \mbox{FUNCDECL} \ \ .
MEMBERDECL \rightarrow VARDECL.
FUNCDECL \rightarrow id lpar FPARAMS rpar colon TYPE semi .
FUNCDECL \rightarrow id lpar FPARAMS rpar colon void semi.
VARDECL -> TYPE id VARDECL 1 semi .
VARDECL_1 \rightarrow ARRAYSIZE VARDECL_1.
VARDECL_1 \rightarrow.
EXPR \rightarrow ARITHEXPR.
EXPR -> RELEXPR
ARITHEXPR \rightarrow TERM ARITHEXPR_1.
ARITHEXPR_1 \rightarrow ADDOP TERM ARITHEXPR_1.
ARITHEXPR_1 \rightarrow.
RELEXPR \rightarrow ARITHEXPR RELOP ARITHEXPR .
FACTOR -> VARIABLE
FACTOR \rightarrow FUNCTIONCALL.
FACTOR \rightarrow intnum.
FACTOR -> floatnum
FACTOR -> lpar ARITHEXPR rpar .
FACTOR \rightarrow not FACTOR .
FACTOR -> SIGN FACTOR .
\label{eq:funcbody} \hbox{\it FUNCBODY}\_1 \ \ \hbox{\it do} \ \ \hbox{\it FUNCBODY}\_2 \ \ \hbox{\it end} \ \ .
FUNCBODY_1 \rightarrow local FUNCBODY_3.
FUNCBODY_1 \rightarrow .
FUNCBODY_2 \longrightarrow STATEMENT FUNCBODY_2.
FUNCBODY 2 \rightarrow.
FPARAMS \rightarrow TYPE id FPARAMS_1 FPARAMS_2 .
FPARAMS \rightarrow.
FPARAMSTAIL \rightarrow comma TYPE id FPARAMSTAIL_1.
\label{eq:fparamstail} \text{FPARAMSTAIL}\_1 \ -> \ \text{ARRAYSIZE} \ \text{FPARAMSTAIL}\_1 \ .
FPARAMSTAIL_1 \rightarrow.
FPARAMS_1 \rightarrow ARRAYSIZE FPARAMS_1.
FPARAMS 1 \rightarrow .
FPARAMS_2 \rightarrow FPARAMSTAIL FPARAMS_2.
FPARAMS 2 \rightarrow .
FUNCBODY_3 \rightarrow VARDECL FUNCBODY_3.
FUNCBODY_3 \rightarrow.
FUNCTIONCALL -> FUNCTIONCALL 1 id lpar APARAMS rpar .
FUNCTIONCALL 1 -> IDNEST FUNCTIONCALL 1 .
FUNCTIONCALL_1 \rightarrow.
APARAMS \rightarrow EXPR APARAMS 1.
APARAMS \rightarrow.
```

```
APARAMS 1 \rightarrow APARAMSTAIL APARAMS 1.
APARAMS_1 \rightarrow ...
APARAMSTAIL \rightarrow comma EXPR.
IDNEST \rightarrow id IDNEST_1 dot.
IDNEST \rightarrow id lpar APARAMS rpar dot .
IDNEST 1 \rightarrow INDEX IDNEST 1.
IDNEST_1 \rightarrow.
INDEX -> lsqbr ARITHEXPR rsqbr .
STATBLOCK \rightarrow do STATBLOCK_1 end.
STATBLOCK \rightarrow STATEMENT.
STATBLOCK \rightarrow.
STATBLOCK_1 \rightarrow STATEMENT STATBLOCK_1.
STATBLOCK_1 \rightarrow.
STATEMENT \longrightarrow ASSIGNSTAT semi.
{\tt STATEMENT} {\tt ->} {\tt if lpar RELEXPR rpar then STATBLOCK else STATBLOCK semi }.
STATEMENT -> while lpar RELEXPR rpar STATBLOCK semi .
STATEMENT \rightarrow read lpar VARIABLE rpar semi.
{\tt STATEMENT} {\tt ->} {\tt write } {\tt lpar EXPR } {\tt rpar } {\tt semi } .
STATEMENT -> return lpar EXPR rpar semi .
STATEMENT -> FUNCTIONCALL semi .
ASSIGNSTAT \longrightarrow VARIABLE \ ASSIGNOP \ EXPR \ .
TERM \rightarrow FACTOR TERM_1.
TERM_1 \rightarrow MULTOP FACTOR TERM_1.
TERM_1 \rightarrow .
VARIABLE 1 \rightarrow IDNEST VARIABLE 1.
VARIABLE_1 -> .
VARIABLE_2 \rightarrow INDEX VARIABLE_2.
VARIABLE_2 \rightarrow.
ASSIGNOP \rightarrow equal.
VISIBILITY -> public .
VISIBILITY -> private .
ADDOP -> plus .
ADDOP \rightarrow minus .
ADDOP \rightarrow or.
TYPE \rightarrow integer .
TYPE \rightarrow float.
TYPE \rightarrow id.
RELOP \rightarrow eq.
RELOP \rightarrow neq .
RELOP -> lt .
RELOP \rightarrow gt.
RELOP -> leq .
\label{eq:reloop} \text{RELOP} \to \text{geq} \ .
ARRAYSIZE -> lsqbr intnum rsqbr .
ARRAYSIZE \rightarrow lsqbr rsqbr.
SIGN -> plus .
SIGN \rightarrow minus .
MULTOP \rightarrow mult.
MULTOP \rightarrow div.
```

Ambiguity-free grammar

```
START \rightarrow PROG.
PROG \rightarrow PROG 1 PROG 2 main FUNCBODY.
PROG_1 \rightarrow CLASSDECL PROG_1.
PROG_1 \rightarrow .
PROG_2 \rightarrow FUNCDEF PROG_2.
PROG_2 \rightarrow .
CLASSDECL -> class id CLASSDECL 1 lcurbr CLASSDECL 2 rcurbr semi .
CLASSDECL 1 -> inherits id CLASSDECL 3 .
CLASSDECL_1 \rightarrow.
{\tt CLASSDECL\_2} \,\to\, {\tt VISIBILITY} \,\, {\tt MEMBERDECL} \,\, {\tt CLASSDECL\_2} \,\,\, .
CLASSDECL_2 \rightarrow.
\label{eq:classdecl} {\it CLASSDECL\_3} \ -\!\!\!\!> \ {\it comma} \ \ {\it id} \ \ {\it CLASSDECL\_3} \ \ .
CLASSDECL_3 \rightarrow.
FUNCDEF \longrightarrow FUNCHEAD FUNCBODY semi.
FUNCHEAD \rightarrow id FUNCHEAD_1 lpar FPARAMS rpar colon TYPEORVOID.
FUNCHEAD\_1 -> \ sr \ id \ .
FUNCHEAD_1 \rightarrow .
{\tt TYPEORVOID} \, -\!\!\!> \, {\tt TYPE} \  \  \, .
TYPEORVOID \rightarrow void.
MEMBERDECL -> id MEMBERDECL 1
{\tt MEMBERDECL} {\tt ->} {\tt INTEGERORFLOAT} {\tt VARDECL\_2} \ .
MEMBERDECL 1 \rightarrow VARDECL 1.
MEMBERDECL_1 \rightarrow FUNCDECLIDLESS.
VARDECL \rightarrow INTEGERORFLOAT VARDECL_1.
VARDECL \rightarrow id VARDECL_1.
VARDECL\_1 -\!\!\!> id\ VARDECL\_2\ semi\ .
VARDECL_2 \rightarrow ARRAYSIZE VARDECL_2.
VARDECL_2 \rightarrow.
\label{eq:function} \mbox{FUNCDECLIDLESS} \rightarrow \mbox{lpar FPARAMS rpar colon TYPEORVOID semi }.
EXPR \rightarrow ARITHEXPR EXPR_1.
EXPR 1 \rightarrow RELOP ARITHEXPR.
EXPR_1 \rightarrow .
ARITHEXPR \rightarrow TERM ARITHEXPR_1.
\label{eq:arithexpr_1} \text{ARITHEXPR}\_1 \ -\!\!\!> \text{ADDOP TERM ARITHEXPR}\_1 \ .
ARITHEXPR 1 \rightarrow.
RELEXPR \rightarrow ARITHEXPR RELOP ARITHEXPR .
FACTOR \rightarrow id FACTOR_1.
FACTOR \rightarrow intnum.
FACTOR -> floatnum
FACTOR \rightarrow lpar ARITHEXPR rpar .
FACTOR \rightarrow not FACTOR.
FACTOR -> SIGN FACTOR .
```

```
FACTOR_1 \rightarrow IDNEST_1 FACTOR_2.
FACTOR_2 \rightarrow dot FACTOR_3.
FACTOR_2 \rightarrow .
FACTOR_3 \rightarrow id FACTOR_4.
FACTOR_3 \rightarrow.
FACTOR\_4 -> IDNEST\_2 \ dot \ VARIABLE\_3 \ .
FACTOR\_4 \rightarrow IDNEST\_3 FACTOR\_5.
FACTOR\_5 -\!\!\!> \ dot \ VARIABLE\_3 \ .
FACTOR\_5 \rightarrow .
FUNCBODY \rightarrow FUNCBODY 1 do FUNCBODY 2 end .
FUNCBODY_1 \rightarrow local FUNCBODY_3.
FUNCBODY_1 \rightarrow .
FUNCBODY_2 \longrightarrow STATEMENT FUNCBODY_2.
FUNCBODY_2 \rightarrow.
FPARAMS \rightarrow TYPE id FPARAMS_1 FPARAMS_2.
FPARAMS \rightarrow.
FPARAMSTAIL \rightarrow comma TYPE id FPARAMSTAIL 1.
\label{eq:fparamstail} \text{FPARAMSTAIL}\_1 \ -> \ \text{ARRAYSIZE} \ \text{FPARAMSTAIL}\_1 \ .
FPARAMSTAIL_1 \rightarrow.
FPARAMS_1 \rightarrow ARRAYSIZE FPARAMS_1.
FPARAMS_1 \rightarrow .
FPARAMS 2 \rightarrow FPARAMSTAIL FPARAMS 2.
FPARAMS_2 -> .
FUNCBODY 3 \rightarrow VARDECL FUNCBODY 3.
FUNCBODY_3 \rightarrow.
APARAMS \rightarrow EXPR APARAMS_1.
APARAMS \rightarrow.
APARAMS_1 \rightarrow APARAMSTAIL APARAMS_1.
APARAMS 1 \rightarrow .
APARAMSTAIL \rightarrow comma EXPR.
INDEX -> lsqbr ARITHEXPR rsqbr .
STATBLOCK \rightarrow do STATBLOCK_1 end.
STATBLOCK \longrightarrow STATEMENT.
STATBLOCK \rightarrow .
STATBLOCK_1 \rightarrow STATEMENT STATBLOCK_1.
STATBLOCK_1 \rightarrow.
STATEMENT \rightarrow id STATEMENT_1.
STATEMENT -> if lpar RELEXPR rpar then STATBLOCK else STATBLOCK semi .
{\it STATEMENT} \longrightarrow {\it while lpar RELEXPR rpar STATBLOCK semi} .
{\tt STATEMENT} \mathrel{->} {\tt read} {\tt lpar} {\tt VARIABLE} {\tt rpar} {\tt semi} \ .
STATEMENT -> write lpar EXPR rpar semi
STATEMENT -> return lpar EXPR rpar semi .
STATEMENT_1 \rightarrow IDNEST_1 STATEMENT_2.
STATEMENT_2 \rightarrow dot STATEMENT_3.
STATEMENT_2 \rightarrow semi.
STATEMENT_3 -> id STATEMENT_4
{\tt STATEMENT\_3} \to {\tt VARIABLE\_2} \ {\tt ASSIGNSTAT} \ {\tt semi} \ .
```

```
{\tt STATEMENT\_4 -> IDNEST\_2 \ dot \ VARIABLE\_3 \ VARIABLE\_2 \ ASSIGNSTAT \ semi \ .}
STATEMENT_4 \rightarrow IDNEST_3 STATEMENT_5.
STATEMENT\_5 -> \ dot \ VARIABLE\_3 \ VARIABLE\_2 \ ASSIGNSTAT \ semi \ .
STATEMENT_5 \rightarrow semi.
ASSIGNSTAT \rightarrow ASSIGNOP EXPR.
TERM \rightarrow FACTOR TERM_1 .
TERM_1 \rightarrow MULTOP FACTOR TERM_1.
TERM_1 \rightarrow .
VARIABLE -> id VARIABLE_1 VARIABLE_2 .
VARIABLE_1 \rightarrow IDNEST_1 VARIABLE_4.
VARIABLE_2 \rightarrow INDEX VARIABLE_2.
VARIABLE_2 \rightarrow.
VARIABLE_3 \rightarrow IDNEST VARIABLE_3.
VARIABLE_3 \rightarrow.
VARIABLE_4 \rightarrow dot VARIABLE_3.
IDNEST \rightarrow id IDNEST_1 dot.
IDNEST_1 \rightarrow IDNEST_2.
IDNEST_1 \rightarrow IDNEST_3.
IDNEST_2 \rightarrow INDEX IDNEST_2.
IDNEST_2 \rightarrow .
IDNEST_3 \rightarrow lpar APARAMS rpar.
ASSIGNOP -> equal .
VISIBILITY \rightarrow public .
VISIBILITY -> private .
ADDOP -> plus .
ADDOP \rightarrow minus .
ADDOP \rightarrow or .
{\tt INTEGERORFLOAT} \,\, -\! > \,\, {\tt integer} \quad .
INTEGERORFLOAT \rightarrow float.
TYPE \rightarrow INTEGERORFLOAT .
TYPE \rightarrow id .
RELOP \rightarrow eq .
\label{eq:reloop} \text{RELOP} \to \text{neq} \ .
RELOP \rightarrow lt .
\begin{array}{c} \text{RELOP} \to \text{gt} \quad . \\ \text{RELOP} \to \text{leq} \quad . \end{array}
RELOP \longrightarrow geq .
ARRAYSIZE -> lsqbr MAYBEINTNUM rsqbr .
MAYBEINTNUM \rightarrow intnum.
MAYBEINTNUM \rightarrow.
SIGN -> plus .
\mathrm{SIGN} \ -\!\!\!> \ \mathrm{minus} .
MULTOP \rightarrow mult .
MULTOP \rightarrow div .
MULTOP \rightarrow and.
```

Section 2. First and Follow Sets

The first sets generated are shown in the following code block:

```
Start: { id class main }
Prog: { id class main }
Prog1: { class \epsilon }
Prog2: { id \epsilon }
ClassDecl: { class }
ClassDecl1: { inherits \epsilon }
ClassDecl2: { public private \epsilon }
ClassDecl3: { , \epsilon }
Visibility: { public private }
MemberDecl: { id integer float }
MemberDecl1: { id ( }
FuncHead: { id }
FuncHead1: { :: \epsilon }
FuncDef: { id }
FuncDeclIdless: { ( }
FuncBody: { do local }
FuncBody1: { local \epsilon }
FuncBody2: { id if while read write return \epsilon }
FuncBody3: { id integer float \epsilon }
VarDecl: { id integer float }
VarDecl1: { id }
VarDecl2: { [ \epsilon }
Stat: { id if while read write return }
Stat1: \{ = ; . ( [ \}
Stat2: { ; . }
Stat3: { id = [ }
Stat4: { . ( [ }
Stat5: { ; . }
Stat6: { = ; . }
AssignStat: \{ = \}
StatBlock: { id if while do read write return \epsilon }
StatBlock1: { id if while read write return \epsilon }
Expr: \{ \text{ intnum floatnum id } + - \text{ ( not )} \}
Expr1: \{ <> \Longrightarrow <= >= <> \setminus epsilon \}
RelExpr: \{ \text{ intnum floatnum id } + - \text{ ( not )} \}
ArithExpr: \{ intnum \ floatnum \ id +- \ ( \ not \ \}
Term: \{ \text{ intnum floatnum id } + - \text{ (not } \}
Term1: \{ * / and \backslash epsilon \}
Factor: \{ intnum floatnum id + - (not ) \}
Factor1: { . ( [ \epsilon }
Factor2: { . \epsilon }
Factor3: { id \epsilon }
Factor4: { . ( [ }
Factor5: \{ . \backslash epsilon \}
Var: { id }
Var1: { . ( [ }
Var2: { [ \epsilon }
Var3: { id \epsilon }
Var4: { . }
IdNest: { id }
IdNest1: { ( [ \epsilon }
IdNest2: { [ \epsilon }
IdNest3: \{ ( ) \}
Index: { [
ArraySize: { [ }
Type: { id integer float }
TypeOrVoid: { id integer float void }
FParams: { id integer float \epsilon }
FParams1: { [ \epsilon } FParams2: { , \epsilon }
AParams: { intnum floatnum id +- ( not \epsilon }
AParams1: { , \epsilon }
FParamsTail: { , }
FParamsTail1: { [ \epsilon }
AParamsTail: { , }
```

```
AssignOp: { = }
RelOp: { <> == <=> > }
AddOp: { +- or }
MultOp: { * / and }
IntegerOrFloat: { integer float }
MaybeIntNum: { intnum \epsilon }
```

The follow sets generated are shown in the following code block:

```
Start: { $ }
Prog: { $ }
Prog1: { id main }
Prog2: { main }
ClassDecl: { id class main }
ClassDecl1: { { }
ClassDecl2: { } }
ClassDecl3: { }
Visibility: { id integer float }
MemberDecl: { } public private }
MemberDecl1: { } public private }
FuncHead: { do local }
FuncHead1: { ( }
FuncDef: { id main }
FuncDeclIdless: { } public private }
FuncBody: { ; $ }
FuncBody1: { do }
FuncBody2: { end }
FuncBody3: { do }
VarDecl: { id integer float do }
VarDecl1: { id } integer float do public private }
VarDecl2: { ; } public private }
Stat: { id ; if else while end read write return }
Stat1: { id ; if else while end read write return }
Stat2: { id ; if else while end read write return
Stat3: { id ; if else while end read write return Stat4: { id ; if else while end read write return
Stat5: { id ; if else while end read write return
Stat6: { id ; if else while end read write return }
AssignStat: { ; }
StatBlock: { ; else } StatBlock1: { end }
Expr: { ; , ) }
Expr1: { ; , ) }
RelExpr: { ) }
Sign: \{ \text{ intnum floatnum id } + - \text{ ( not ) } \}
Term: \{ +-<> ; , ) = <=>= <> or \}
Term1: \{+-<>;, \} ] = <=>= < \text{ or } \}

Factor: \{+-*<>/;, \} ] = <=>= < \text{ or and } \}

Factor1: \{+-*<>/;, \} ] = <=>= < \text{ or and } \}
Factor 2: \{ + - * < > / ; , ) = <= >= <> or and <math>\}
Factor3: \{+-*<>/; , \} = <=>= <> or and }
Factor 4: \{ +-*<>/ ; , ) = <=>= <> or and }
Factor5: \{+-*<>/;, \}
Var: { ) }
Var1: { ) [ }
Var2: \{ = \}
Var3: \{ = + - * < > / ; , ) [ ] == <= >= <> or and } Var4: <math>\{ \}
{\tt IdNest: \{ id = +-* <>/\ ; \ , \ ) \ [ \ ] ==<=>= <> \ or \ and \ } \\
IdNest1: \{ = + - * < > / ; , . ) ] == <= >= <> or and \}
{\rm IdNest2}: \ \ {\rm \{} \ = + - \ * \ < > \ / \ \ ; \ \ , \ \ . \ \ ) \ \ ] \ = = < = > = < > \ {\rm or \ and}
IdNest3: \{ = + - * < > / ; , . \} = <= >= <> or and }
Index: \{ = + - * < > / ; , . \} = <= >= <> or and }
Index: { = + - * < > / ; , . ) [ ] == <:
ArraySize: { ; , ) [ } public private }
Type: { id ; do local }
TypeOrVoid: { ; do local }
FParams: { ) }
FParams1: { , ) }
```

```
FParams2: { ) }
AParams: { ) }
AParams1: { ) }
FParamsTail: { , ) }
FParamsTail: { , ) }
AParamsTail: { , ) }
AParamsTail: { , ) }
AssignOp: { intnum floatnum id + - ( not }
RelOp: { intnum floatnum id + - ( not }
AddOp: { intnum floatnum id + - ( not }
MultOp: { intnum floatnum id + - ( not }
IntegerOrFloat: { id ; [ } do public private local }
MaybeIntNum: { ] }
```

Section 3. Design

The overall strategy to designing the parser involved a few steps:

- Implement a first and follow set generator based on a text file containing a DSL based grammar syntax
- Using the generated first and follow sets, implement a paser table generator
- Using the generated table parser, implement a table-driven parser capable of validating a given input file containing source code of the language defined by the grammar and generate a derivation table and list of syntax errors that result
- Add semantic action identifiers to the transformed grammar to direct the parser to perform abstract syntax tree construction steps

The error recovery strategy used was panic mode. Once an empty table cell is found, the implementation attempts to skip the tokens at fault until it finds one that resyncs with the current derivation and continues on parsing the next expressions. This was chosen for its simplicity, even though it does not always identify the true source of the problem in all cases.

Section 4. Use of Tools

The two main tools used for this assignment were Dr. Paquet's grammartool.jar and UCalgary's Grammar Tool. The grammartool.jar was used to quickly remove EBNF repeating/optional notations and left-recursion in the original grammar provided and output the resulting grammar in the notation compatible with the UCalgary Grammar Tool. By having a EBNF notation free grammar without left-recursion, the next step was to use the UCalgary Grammar Tool to iteratively attempt to solve all the ambiguities (as described in Section 1).

Other tools that were used were GraphViz when generating the AST diagram.