

In the following  $A, B, \dots$  are propositional variables,  $a, b, \dots$  constant symbols,  $f, g, \dots$  function symbols,  $X, Y \dots$  variables,  $p, q, \dots$  predicate symbols and  $\phi, \psi, \dots$  formulas (unless differently specified).

1. (4 points) Consider the language of propositional logic. Use natural deduction to prove that the following holds, or find a counter-example to show that it does not hold

- $\vdash \phi \rightarrow (\psi \rightarrow (\sigma \rightarrow (\phi \wedge \psi \wedge \sigma)))$
- $\vdash \phi \rightarrow ((\neg \phi \rightarrow \psi) \rightarrow \psi)$

Hint: Assume that you can use the rule for  $\vee$  introduction

$$\frac{\phi}{\phi \vee \psi}$$

2. (4 points) Transform the following propositional logic formula into an equivalent formula in Disjunctive Normal Form

- $\neg A \wedge (\neg B \wedge ((C \wedge D) \rightarrow A))$

3. (4 points) Prove that that  $\phi \models \psi$  ( $\psi$  is a logical consequence of  $\phi$ ) or that  $\phi \not\models \psi$  for the following formulas:

- $\phi : \neg A \vee \neg B$  and  $\psi : (A \wedge \neg B) \rightarrow (B \vee C)$
- $\phi : A \rightarrow C$  and  $\psi : (A \rightarrow B) \rightarrow C$

4. Define a propositional language which allows to describe the state of an (simplified) ATM at different instants. With the language defined above provide a (set of) formulas which expresses the following facts:

- (a) At each instant the ATM is one (and only one) of these four states: idle, checking the card, delivering the money, error;
- (b) Under normal functioning the ATM switches from idle to checking the card, from checking the card to delivering the money and from delivering the money to idle;
- (c) When the machine enters an error state it remains in this state for all the successive instants.

5. (5 points) Consider the following definitions for a graph  $G$ :

- Adjacent node: a node  $x$  is adjacent to node  $y$  iff there exists an edge between  $x$  and  $y$ .
- Path: a path in a graph is a sequence of edges which joins a sequence of nodes.
- Connected component: A connected component in  $G$  is a subgraph of  $G$  in which any two vertices are connected to each other by paths, and which is connected to no additional vertices in the supergraph  $G$ .
- Dummy graph coloring problem: given a non-oriented graph, associate a color to each of its nodes in such a way that adjacent nodes have the same color.

Provide a FOL language and a set of axioms that formalize the dummy graph coloring problem of a graph with  $k$  connected components, maximizing the number of colors used.

6. (5 points) A binary tree is either empty or it is composed of a root element and two successors, which are binary trees themselves. In Prolog we represent the empty tree by the atom `nil` and the non-empty tree by the term `t(X,L,R)` where  $X$  denotes the root node and  $L$  and  $R$  denote the left and right subtree, respectively. A leaf is a node with no successors. For example, the tree in the Figure is represented by the term `t(a,t(b,nil,nil),t(c,nil,nil))` and  $b$  and  $c$  are leaves.

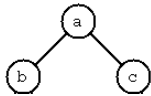


Figure 1: A tree

Write a Prolog program which defines a predicate `count_leaves_parents(T,N)` which counts the number of nodes which have only leafs as children.

```
% count_leaves_parents(T,N) :- the binary tree T has N nodes which  
% have only leafs as children.
```

7. (4 points) Provide a Prolog program  $P$  and a goal  $G$  such that the evaluation of  $G$  in  $P$  successfully terminates, while the evaluation of  $G$  in a program  $P'$ , obtained from  $P$  by changing the order of clauses, does not terminate.
8. (3 points) Describe shortly the difference between symbolic computation and sub-symbolic computation in the context of AI.