

$$\begin{array}{c}
 \textcircled{1} \quad \frac{(\phi), (\psi) \wedge I}{\phi \wedge \psi} \quad \frac{[\sigma]_1}{\sigma} \wedge I \\
 \hline
 \phi \wedge \psi \wedge \sigma \\
 \hline
 \rightarrow I_1 \\
 \sigma \rightarrow (\phi \wedge \psi \wedge \sigma) \\
 \hline
 \rightarrow I_2 \\
 \psi \rightarrow (\sigma \rightarrow \phi \wedge \psi \wedge \sigma) \\
 \hline
 \rightarrow I_3 \\
 \phi \rightarrow (\psi \rightarrow (\sigma \rightarrow (\phi \wedge \psi \wedge \sigma)))
 \end{array}$$

• If $\phi = T$ and $\psi = F$ the formula is false, hence it is not satisfiable.

$$\begin{aligned}
 \textcircled{2} \quad & \neg A \wedge ((\neg B \wedge (C \wedge D) \rightarrow A)) \equiv \\
 & \neg A \wedge \neg B \wedge (\neg C \vee \neg D \vee A) \equiv \\
 & (\neg A \wedge \neg B \wedge \neg C) \vee \\
 & (\neg A \wedge \neg B \wedge \neg D) \vee \\
 & (\neg A \wedge \neg B \wedge A) \equiv (\neg A \wedge \neg B \wedge \neg C) \vee (\neg A \vee \neg B \vee \neg D)
 \end{aligned}$$

$$\textcircled{3} \bullet \psi \equiv (\neg A \vee B \vee C)$$

If $A=T, B=F$ and $C=F$ then

ϕ is T and ψ is F, hence

$$\phi \neq \psi$$

$$\bullet \phi \equiv \neg A \vee C$$

$$\psi \equiv (\neg A \vee B) \rightarrow C \equiv (A \wedge \neg B) \vee C$$

If $A=F, B=T$ and $C=F$ then

ϕ is true and ψ is false, hence

$$\phi \neq \psi$$

$\textcircled{4}$ I_k = idle at time k

C_k = checking " " "

d_k = delivery " " "

$$a) I_k \leftrightarrow (\neg C_k \wedge \neg d_k) \wedge C_k \leftrightarrow (\neg I_k \wedge \neg d_k) \wedge \neg C_k$$

$$\wedge e_k \leftrightarrow (\neg I_k \wedge \neg C_k) \wedge d_k \leftrightarrow (\neg I_k \wedge \neg C_k) \wedge e_k$$

$$b) I_{k-1} \rightarrow (I_k \vee C_k \vee e_k) \wedge C_{k-1} \rightarrow (C_k \vee d_k \vee e_k) \wedge d_{k-1} \rightarrow (d_k \vee I_k \vee e_k)$$

$$c) e_{k-1} \rightarrow e_k$$

⑤ Layer

function: $color(x) = \text{color of node}$
to node x

predicate $edge(x, y)$: there is an edge
between x and y

" node (x) : x is a node

i) Dummy graph:

$$\forall x, y (edge(x, y) \rightarrow color(x) = color(y))$$

ii) Max color

$$\forall x, y (path(x, y) \rightarrow color(x) \neq color(y))$$

when

$$\forall x, y (path(x, y) \leftrightarrow (\exists z (edge(x, z) \wedge path(z, y)) \vee edge(x, y)))$$

⑥ the predicate ~~and~~ count_leaves - prenter
(count for short) is defined as follows.

count (t(-, nil, nil), 0).

count (t(-, t(-, nil, nil), t(-, nil, nil)), 1).

count (t(-, nil, t(-, nil, nil)), 1).

count (t(-, t(-, nil, nil), nil), 1).

~~count (t(-, L, R), N) :-~~

count (t(-, L, R), N) :-

count (L, NL), count (R, NR),

N = NL + NR

⑦ $\mu(x)$

$\mu(x) := -\mu(x)$

$\mu(x) := -\mu(x)$

$\mu(x)$

$\mu(x)$ term hole

$\mu(x)$ slot not
term hole.