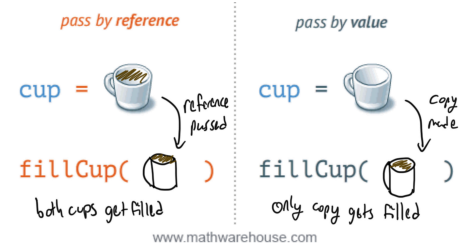


Taylor Uml (tw2)

1. (6 points) Consider the following function written in C:

```
int x=3, y=2, z=1;
void foo(int a, int b) {
    x = x + b;
    a = a - b;
}
```



In each of the cases below, write down the values of x , y and z after the following calls to $\text{foo}()$. If necessary, assume that output arguments are copied back to parameters in the left-to-right order.

- (a)
- $\text{foo}(y,z)$
- where all parameters are passed
- by value

```
int x=3, y=2, z=1;
foo(int y, int z):
    x = x + z => x = 4
    y = y - z => y = 1
```

global variables not updated

Result:
x: 4
y: 2
z: 1

- (b)
- $\text{foo}(y,z)$
- where all parameters are passed
- by reference

```
int x=3, y=2, z=1;
foo(int y, int z):
    x = x + z => x = 4
    y = y - z => y = 1
```

values of x, y updated

Result:
x: 4
y: 1
z: 1

- (c)
- $\text{foo}(y,z)$
- where all parameters are passed
- by value-result

```
int x=3, y=2, z=1;
foo(int y, int z):
    x = x + z => x = 4
    y = y - z => y = 1
```

updates only those values that were parameters, so $y=1$, z same as $z=1$

Result:
x: 4
y: 1
z: 1

- (d)
- $\text{foo}(x,y)$
- where all parameters are passed
- by reference

```
int x=3, y=2, z=1;
foo(int x, int y):
    x = x + y => x = 5
    x = x - y => x = 3
```

value of x updated to 3

Result:
x: 3
y: 2
z: 1

- (e)
- $\text{foo}(x,x)$
- where all parameters are passed
- by reference

```
int x=3, y=2, z=1;
foo(int x, int x):
    x = x + x => x = 6
    x = x - x => x = 0
```

 x is updated to 0

Result:
x: 0
y: 2
z: 1

- (f)
- $\text{foo}(x,x)$
- where all parameters are passed
- by value-result

```
int x=3, y=2, z=1;
foo(int x, int x):
    x = x + x => x = 6
    x = x - x => x = 0
```

only update x

Result:
x: 3
y: 2
z: 1

Stack
 Set x to $x_0 = 6$
 Set x to 0
 Set x to $x_0 = 3$

2. (2 points) Consider the following C++ program, where `&i` means `i` is passed by reference:

```
int bar (int &i) {
    i = i - 2;
    return 2 * i;
}

void foo1 () {
    int x = 3, y = 6, sum;
    sum = x;
    sum = sum + bar(x);
}

void foo2 () {
    int x = 2, y = 7, sum;
    sum = bar(x);
    sum = sum + x;
}
```

`&i` : pass by reference
`i` : pass by value

- (a) What is the value of sum at the end of the function `foo1`? Briefly explain why.

foo1()

1) `int x=3, y=6, sum;` initial declaration Assignment Stack

2) `sum = x;` ~~sum=?~~ ~~x=3~~

3) `sum = sum + bar(x);` ~~sum=3~~ ~~x=1~~

`sum=5`

`bar(&x)` `x=3`, same as aside reference to `x`

`x = x - 2` \Rightarrow `x=1`. update `x` to `x=1`

`return 2 * x`

returns 2

3) `sum = 3 + 2 = 5` `sum=5`

Since `sum` is initially set to 3 on line 2, the next line sets `sum` equal to 3 + the return value from `bar(x=3)`. The value of `x` is passed by reference, so inside `bar`, `x` is updated to `3-2=1`. Then, `bar` returns `2*x` where `x=1`, so returns 2. Therefore, `sum` is equal to `3+2=5`.

- (b) What is the value of `sum` at the end of the function `foo2`? Briefly explain why.

foo2() initial declaration Assignment Stack

1) `int x=2, y=7, sum;` ~~sum=?~~ ~~x=2~~

2) `sum = bar(x);` ~~sum=0~~ ~~x=0~~

`bar(&x)` `x=2`, same as aside reference to `x`

`x = x - 2` \Rightarrow `x=0`. update `x` to `x=0`

`return 2 * x`

returns 0

3) `sum = sum + x = 0` `0 + 0`

`sum=0`

Since `sum` is set to the value of `bar(x)`, and `bar(x)` takes the reference to `x`, after `bar(x=2)` runs, it returns 0, setting `sum=0`. Inside `bar`, `x` is also set to 0, and since `x` is passed by reference, `x=0`. On line 3, `sum` is set to `sum+x`, or `0+0`, making the final value of `sum=0`.

3. (4 points) Consider the Ada program given below.

```

1  procedure Main is
2    A, B, C : Integer;
3    procedure Sub1 (C:Integer) is
4      D, E : Integer;
5      begin — of Sub1
6        ...
7      end; — of Sub1
8    procedure Sub2 (E:Integer) is
9      procedure Sub3 (F:Integer) is
10       C, D: Integer;
11       begin — of Sub3
12         Sub1(100);
13       end; — of Sub3
14     begin — of Sub2
15       Sub3(E);
16     end; — of Sub2
17 begin — of Main
18   Sub2(10);
19 end; — of Main

```

Dynamic Link
points to RA

Static Link points to
static link.

Ada is a statically scoped language. In the above program, the Main function invokes Sub2; Sub2 invokes Sub3; and Sub3 invokes Sub1. Draw the stack of activation records when the program's execution reaches before line 12 and line 6. For each activation record, include local variables, parameters, the dynamic link, the static link, and the return address.

Execution order: Main → Sub2 → Sub3 → Sub1

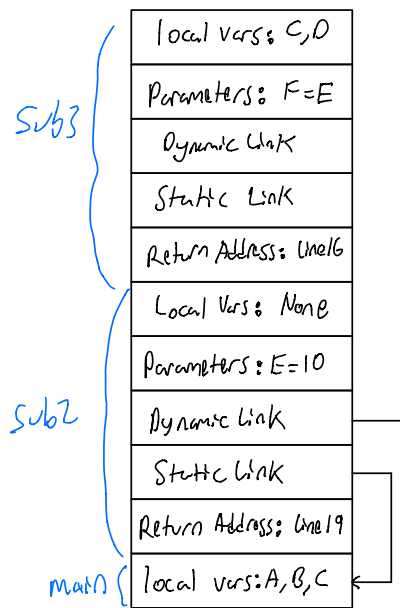
Example Activation Record

Return Value
Local Variables
Parameters
Dynamic Link
Static Link
Source Registers
Return Address

Setup Activation Record as

Sub1
Sub3
Sub2
Main

Before Line 12:



Before Line 6:

