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a) Pseudo Coole for high school optimal

Constant a hashmap to store indices of student prefuences for each school is. What school is preferred at what number for each student. This helps reduce complexity of are some high schools with Available spots, we iterate over a second loop to get each student in the preference list of such school. If the student has not been allotted a school, you assign this one and decrease the spots in the School and increase preference pointer by I (just in case we night have to come back to this school again). Preference pointer helps keep track of how many students night have been covered in the preference list of the school. If the student is already allotted, you check if the preference of student fare that allottment is more than your new school. If the current allotment is preferred, you move to next pointer for this new school, otherwise you switch the two schools of devease available slots for new, increase for current of the tracker of severence. Or well: If the I move the teacher of preference as well. If the current school was filled right now, you also make sure you add it back to the list of schools with available slots.

b) Hashmap to keep teach of preference of school for each student has complexity of O(mn) because first you iterate over each student & them each school. For the main code, we use only doing deletions and insertions inside our loop at indexed locations or using linked lists to add new elements at the end. All of there have a complexity of O(1).

flowerer there are 2 loops seeming, one iterating over schools and other on students, and since they are nested, the complexity becomes O(mn). It is because the operation is upper bound by howing to access each school & then each student which can take m^*n turns.

Since this map owns differently than the main code and they both have complexity of O(mn), hence total complexity is O(mn).

C) Pseudo code foe student optimal.

In this case on well, I made a flashflap to retrieve priorities a student is assigned by different high schools to compare later on. This reduces the complexity of the code a lot. Initial code is sur on students which are left unmatched and for each of these students, making sive that they are unmatched each highschool is considered which is highest in their preference. If that high school has a spot, then the student is assigned to it and spots in the school are decreased. If not, then we check for all the students that school has, if it are decreased. it peepers this new student, it is assigned and the bast preferred one is senered and added to list of students with no netching. If total spots are equal to no. of students, then the ending condition is that all students need to be matched. If they are unequal, then the ending condition is that each student who has not been matched has gone through all the School. School.

d) Harhmap has complexity of (mn), first we iterate on Students and then nested iterations are on schools; just like the case above.

Since, in this case we also have to make suce that we know how each perfect its allocated students, we have to also sum through the spots taken up by each school to get least preferred student for comparison with a new student. Hence that loop also comes in and its iterations will be equal to maximum no- of spots a school has available (ie having to go through all the spots for one school). Hence, the complexity becomes O(mnp) where plenotes maximum no- of Slots in a school.

Special case: When there is only 1 school, and it has n spots for n students, this becomes p = n & $O(1 \times n \times n) = O(n^2)$ just like what we discussed in class.